THE PROGRAMMER'S CP/M® HANDBOOK

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THE PROGRAMMER'S CP/M® HANDBOOK

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Dedication

Several years ago I was told that "Perfection is an English education, an American salary, and a Japanese wife."

Accordingly, I wish to thank the members of Staff at Culford School in England, who gave me the English education, the people who work with me at Johnson-Laird Inc. and Control-C Software and our clients, who give me my American salary, and Mr. and Mrs. Kitagawa, who gave me Kay Kitagawa (who not only married me but took over where my English grammar left off).

A.J-L.

Acknowledgments

Although this book is not authorized or endorsed by Digital Research, I would like to express my thanks to Gary Kildall and Kathy Strutynski of Digital Research, and to Phil Nelson (formerly of Digital Research, now of Victor Technology) for their help in keeping me on the path to truth in this book. I would also like to thank Denise Penrose, Marty McNiff, Mary Borchers, and Ralph Baumgartner at Osborne/McGraw-Hill for their apparently inexhaustible patience.

A.J-L.

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Introduction

This book is a sequel to the *Osborne CP/M® User Guide* by Thom Hogan. It is a technical book written mainly for programmers who require a thorough knowledge of the internal structure of CP/M—how the various pieces of CP/M work, how to use CP/M as an operating system, and finally, how to implement CP/M on different computer systems. This book is written for people who

- Have been working with microcomputers that run Digital Research's CP/M operating system.
- Understand the internals of the microprocessor world—bits, bytes, ports, RAM, ROM, and other jargon of the programmer.
- Know how to write in assembly language for the Intel 8080 or Zilog Z80 Central Processing Unit (CPU) chips.

If you don't have this kind of background, start by getting practical experience on a system running CP/M and by reading the following books from Osborne/ McGraw-Hill:

- *An Introduction to Microcomputers: Volume 1—Basic Concepts*
  This book describes the fundamental concepts and facts that you need to

know about microprocessors in order to program them. If you really need basics, there is a Volume 0 called *The Beginner's Book*.

- *8080A/8085 Assembly Language Programming*
  This book covers all aspects of writing programs in 8080 assembly language, giving many examples.
- *Osborne CP/M® User Guide (2nd Edition)*
  This book introduces the CP/M operating system. It tells you how to use CP/M as a tool to get things done on a computer.

The book you are reading now deals only with CP/M Version 2.2 for the 8080 or Z80 chips. At the time of writing, new versions of CP/M and MP/M (the multi-user, multi-tasking successor to CP/M) were becoming available. CP/M-86 and MP/M-86 for the Intel 8086 CPU chip and MP/M-II for the 8080 or Z80 chips had been released, with CP/M 3.0 (8080 or Z80) in the wings. The 8086, although related architecturally to the 8080, is different enough to make it impossible to cover in detail in this book; and while MP/M-II and MP/M-86 are similar to CP/M, they have many aspects that cannot be adequately discussed within the scope of this book.

Outline of Contents

This book explains topics as if you were starting from the top of a pyramid. Successive "slices" down the pyramid cover the same material but give more detail.

The first chapter includes a brief outline of the notation used in this book for example programs written in Intel 8080 assembly language and in the C programming language.

Chapter 2 deals with the structure of CP/M, describing its major parts, their positions in memory, and their functions.

Chapter 3 discusses CP/M's file system in as much detail as possible, given its proprietary nature. The directory entry, disk parameter block, and file organization are described.

Chapter 4 covers the Console Command Processor (CCP), examining the way in which you enter command lines, the CP/M commands built into the CCP, how the CCP loads programs, and how it transfers control to these programs.

Chapter 5 begins the programming section. It deals with the system calls your programs can make to the high-level part of CP/M, the Basic Disk Operating System (BDOS).

Chapters 6 through 10 deal with the Basic Input/Output System (BIOS). This is the part of CP/M that is unique to each computer system. It is the part that you as a programmer will write and implement for your own computer system.

Chapter 6 describes a standard implementation of the BIOS.

Chapter 7 describes the mechanism for rebuilding CP/M for a different configuration.

Chapter 8 tells you how to write an enhanced BIOS.

Chapter 9 takes a close look at how to handle hardware errors—how to detect and deal with them, and how to make this task easier for the person using the computer.

Chapter 10 discusses the problems you may face when you try to debug your BIOS code. It includes debugging subroutines and describes techniques that will save you time and suffering.

Chapter 11 describes several utility programs, some that work with the features of the enhanced BIOS in Chapter 8 and some that will work with all CP/M 2 implementations.

Chapter 12 concerns error messages and some oddities that you will discover, sometimes painfully, in CP/M. Messages are explained and some probable causes for strange results are documented.

The appendixes contain "ready-reference" information and summaries of information that you need at your side when designing, coding, and testing programs to run under CP/M or your own BIOS routines.

Notation

When you program your computer, you will be sitting in front of your terminal interacting with CP/M and the utility programs that run under it. The sections that follow describe the notation used to represent the dialog that will appear on your terminal and the output that will appear on your printer.

Console Dialog

This book follows the conventions used in the *Osborne CP/M User Guide*, extended slightly to handle more complex dialogs. In this book

- <name> means the ASCII character named between the angle brackets, < and >. For example, <BEL> is the ASCII Bell character, and <HT> is the ASCII Horizontal Tab Character. (Refer to Appendix A for the complete ASCII character set.)
- <cr> means to press the CARRIAGE RETURN key.
- 123 or a number without a suffix means a decimal number.
- 100B or a number followed by B means a binary number.
- 0A5H or a number followed by H means a hexadecimal number. A hexadecimal number starting with a letter is usually shown with a leading 0 to avoid confusion.

- ^x means to hold the CONTROL (CTRL) key down while pressing the x key.
- Underline is keyboard input you type. Output from the computer is shown without underlining.

Assembly Language Program Examples

This book uses Intel 8080 mnemonics throughout as a "lowest common denominator"—the Z80 CPU contains features absent in the 8080, but not vice versa. Output from Digital Research's ASM Assembler is shown so that you can see the generated object code as well as the source.

High-Level Language Examples

The utility programs described in Chapter 11 are written in C, a programming language which lends itself to describing algorithms clearly without becoming entangled in linguistic bureaucracy. Cryptic expressions have been avoided in favor of those that most clearly show how to solve the problem. Ample comments explain the code.

An excellent book for those who do not know how to program in C is *The C Programming Language* by Brian Kernighan and Dennis Ritchie (Prentice-Hall). Appendix A of this book is the C Reference Manual.

Example Programs on Diskette

Example programs in this book have been assembled with ASM and tested with DDT, Digital Research's Dynamic Debugging Tool. C examples were compiled using Leor Zolman's BDS C Compiler (Version 1.50) and tested using the enhanced BIOS described in Chapter 8.

All of the source code shown in this book is available on a single-sided, single-density, 8-inch diskette (IBM 3740 format). Please do *not* contact Osborne/McGraw-Hill to order this diskette. Call or write

Johnson-Laird, Inc.
Attn: The CP/M Programmer's Handbook Diskette
6441 SW Canyon Court
Portland, OR 97221
Tel: (503) 292-6330

The diskette is available for $50 plus shipping costs.

2

The Structure of CP/M

This chapter introduces the pieces that make up CP/M—what they are and what they do. This bird's-eye view of CP/M will establish a framework to which later chapters will add more detailed information.

You may have purchased the standard version of CP/M directly from Digital Research, but it is more likely you received CP/M when you bought your microprocessor system or its disk drive system. Or, you may have purchased CP/M separately from a software distributor. In any case, this distributor or the company that made the system or disk drive will have already modified the standard version of CP/M to work on your specific hardware. Most manufacturers' versions of CP/M have more files on their system diskette than are described here for the standard Digital Research release.

Some manufacturers have rewritten all the documentation so that you may not have received any Digital Research CP/M manuals. If this is the case, you should order the complete set from Digital Research, because as a programmer, you will need to have them for reference.

CP/M from Digital Research

Digital Research provides a standard "vanilla-flavored" version of CP/M that will run only on the Intel Microcomputer Development System (MDS). The CP/M package from Digital Research contains seven manuals and an 8-inch, single-sided, single-density standard IBM 3740 format diskette.

The following manuals come with this CP/M system:

- *An Introduction to CP/M Features and Facilities.* This is a brief description of CP/M and the utility programs you will find on the diskette. It describes only CP/M version 1.4.
- *CP/M 2.0 User's Guide.* Digital Research wrote this manual to describe the new features of CP/M 2.0 and the extensions made to existing CP/M 1.4 features.
- *ED: A Context Editor for the CP/M Disk System.* By today's standards, ED is a primitive line editor, but you can still use it to make changes to files containing ASCII text, such as the BIOS source code.
- *CP/M Assembler (ASM).* ASM is a simple but fast assembler that can be used to translate the BIOS source code on the diskette into machine code. Since ASM is only a bare-bones assembler, many programmers now use its successor, MAC (also from Digital Research).
- *CP/M Dynamic Debugging Tool (DDT).* DDT is an extremely useful program that allows you to load programs in machine code form and then test them, executing the program either one machine instruction at a time or stopping only when the CPU reaches a specific point in the program.
- *CP/M Alteration Guide.* There are two manuals with this title, one for CP/M version 1.4 and the other for 2.0. Both manuals describe, somewhat cryptically, how to modify CP/M.
- *CP/M Interface Guide.* Again, there are two versions, 1.4 and 2.0. These manuals tell you how to write programs that communicate directly with CP/M.

The diskette supplied by Digital Research has the following files:

ASM.COM
The CP/M assembler.

BIOS.ASM
A source code file containing a sample BIOS for the Intel Microcomputer Development System (MDS). Unless you have the MDS, this file is useful only as an example of a BIOS.

CBIOS.ASM
Another source code file for a BIOS. This one is skeletal: There are gaps so that you can insert code for your computer.

DDT.COM
The Dynamic Debugging Tool program.

DEBLOCK.ASM
A source code file that you will need to use in the BIOS if your computer uses sector sizes other than 128 bytes. It is an example of how to block and deblock 128-byte sectors to and from the sector size you need.

DISKDEF.LIB
A library of source text that you will use if you have a copy of Digital Research's advanced assembler, MAC.

DUMP.ASM
The source for an example program. DUMP reads a CP/M disk file and displays it in hexadecimal form on the console.

DUMP.COM
The actual executable program derived from DUMP.ASM.

ED.COM
The source file editor.

LOAD.COM
A program that takes the machine code file output by the assembler, ASM, and creates another file with the data rearranged so that you can execute the program by just typing its name on the keyboard.

MOVCPM.COM
A program that creates versions of CP/M for different memory sizes.

PIP.COM
A program for copying information from one place to another (PIP is short for Peripheral Interchange Program).

STAT.COM
A program that displays statistics about the CP/M and other information that you have stored on disks.

SUBMIT.COM
A program that you use to enter CP/M commands automatically. It helps you avoid repeated typing of long command sequences.

SYSGEN.COM
A program that writes CP/M onto diskettes.

XSUB.COM
An extended version of the SUBMIT program. The files named previously

fall into two groups: One group is used only to rebuild CP/M, while the other set is general-purpose programming tools.

The Pieces of CP/M

CP/M is composed of the Basic Disk Operating System (BDOS), the Console Command Processor (CCP), and the Basic Input/Output System (BIOS).

On occasion you will see references in CP/M manuals to something called the FDOS, which stands for "Floppy Disk Operating System." This name is given to the portion of CP/M consisting of both the BDOS and BIOS and is a relic passed down from the original version. Since it is rarely necessary to refer to the BDOS and the BIOS combined as a single entity, no further references to the FDOS will be made in this book.

The BDOS and the CCP are the proprietary parts of CP/M. Unless you are willing to pay several thousand dollars, you cannot get the source code for them. You do not need to. CP/M is designed so that all of the code that varies from one machine to another is contained in the BIOS, and you do get the BIOS source code from Digital Research. Several companies make specialized BIOSs for different computer systems. In many cases they, as well as some CP/M hardware manufacturers, do not make the source code for their BIOS available; they have put time and effort into building their BIOS, and they wish to preserve the proprietary nature of what they have done.

You may have to build a special configuration of CP/M for a specific computer. This involves no more than the following four steps:

1. Make a version of the BDOS and CCP for the memory size of your computer.
2. Write a modified version of the BIOS that matches the hardware in your computer.
3. Write a small program to load CP/M into memory when you press the RESET button on your computer.
4. Join all of the pieces together and write them out to a diskette.

These steps will be explained in Chapters 7, 8, and 9.

In the third step, you write a small program that loads CP/M into memory when you press the RESET button on your computer. This program is normally called the bootstrap loader. You may also see it called the "boot" or even the "cold start" loader. "Bootstrap" refers to the idea that when the computer is first turned on, there is no program to execute. The task of getting that very first program into the computer is, conceptually, as difficult as attempting to pick yourself up off the ground by pulling on your own bootstraps. In the early days of computing, this operation was performed by entering instructions manually—setting large banks of switches (the computer was built to read the switches as soon as it was turned on). Today, microcomputers contain some small fragment of a program in "non-volatile" read-only memory (ROM)—memory that retains data when the computer is turned off. This stored program, usually a Programmable Read Only Memory (PROM) chip, can load your bootstrap program, which in turn loads CP/M.

CP/M Diskette Format

The standard version of CP/M is formatted on an 8-inch, single-sided diskette. Diskettes other than this type will probably have different layouts; hard disks definitely will be different.

The physical format of the standard 8-inch diskette is shown in Figure 2-1. The

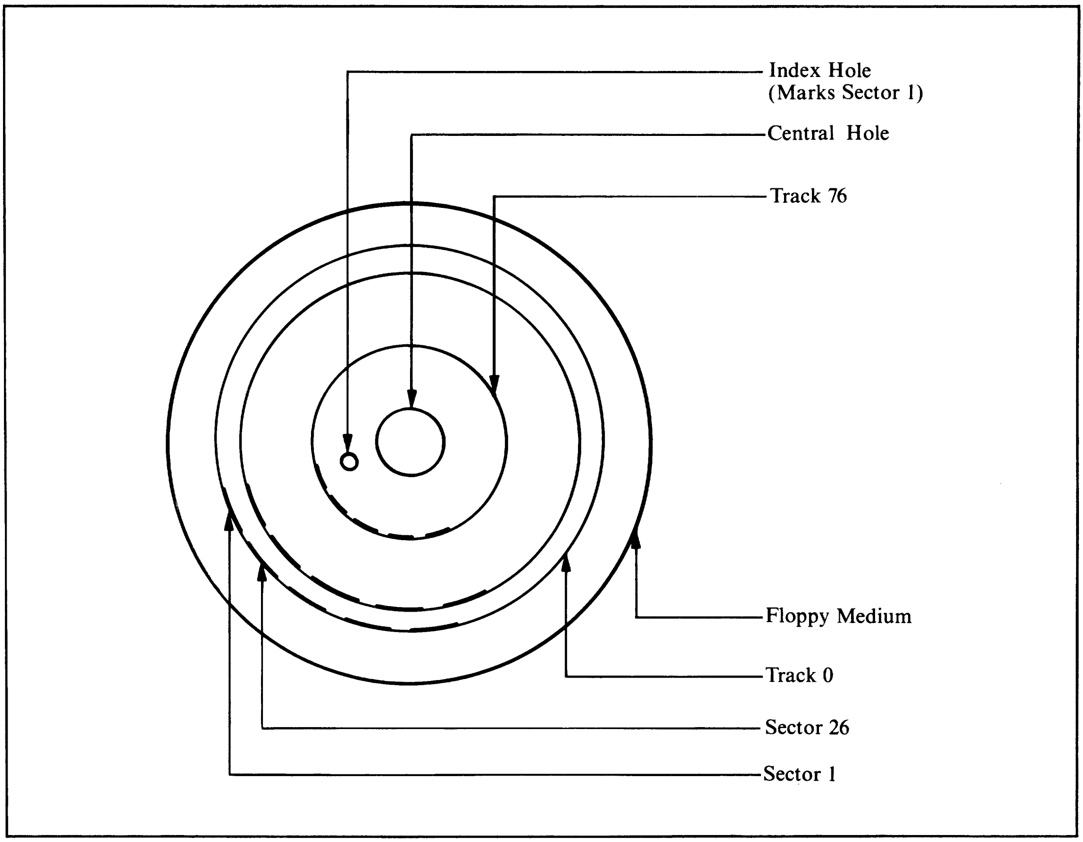


Figure 2-1. Floppy disk layout

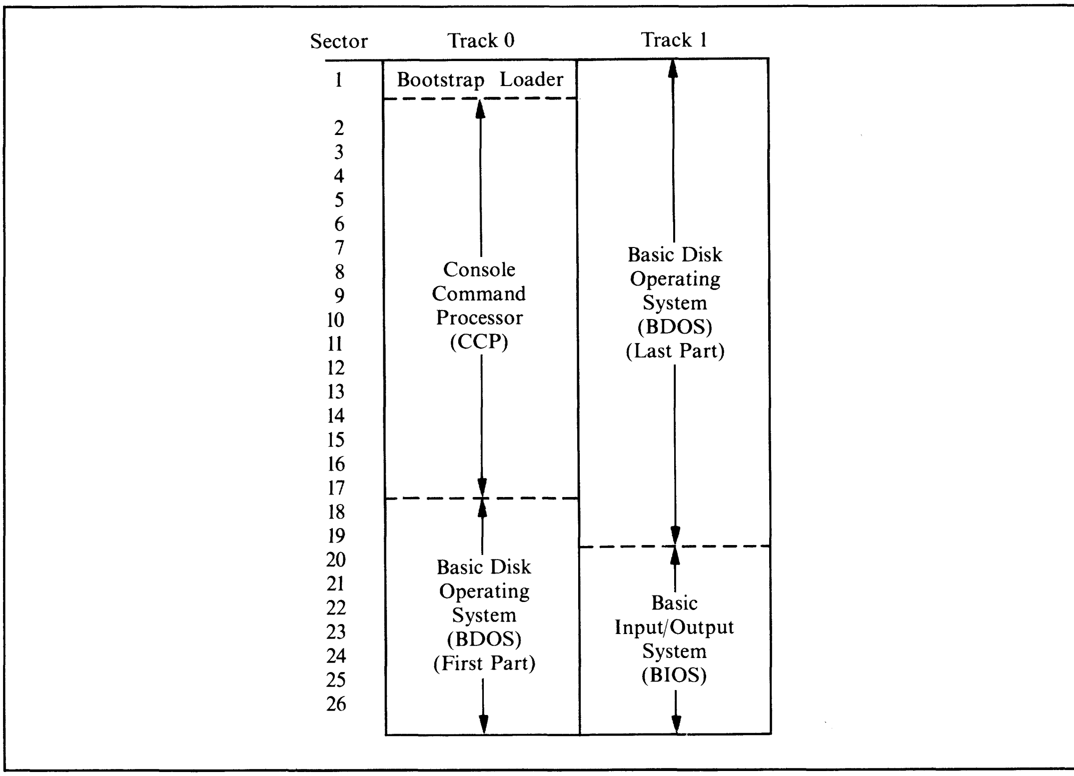


Figure 2-2. Layout of CP/M on tracks 0 and 1 of floppy disk

diskette has a total of 77 concentric tracks numbered from zero (the outermost) to 76 (the innermost). Each of these tracks is divided radially into 26 sectors. These physical sectors are numbered from 1 to 26; physical sector zero does not exist. Each sector has enough space for 128 bytes of data.

Even when CP/M is implemented on a large hard disk with much larger sector sizes, it still works with 128-byte sectors. The BIOS has extra instructions that convert the *real* sectors into CP/M-style 128-byte sectors.

A final note on physical format: The soft-sectored, single-sided, single-density, 8-inch diskette (IBM 3740 format) is the *only* standard format. Any other formats will be unique to the hardware manufacturer that uses them. It is unlikely that you can read a diskette on one manufacturer's computer if it was written on another's, even though the formats appear to be the same. For example, a single-sided, double-density diskette written on an Intel Development System cannot be read on a Digital Microsystems computer even though both use double-density format. If you want to move data from one computer to another, use 8-inch, single-sided, single-density format diskettes, and it *should* work.

In order to see how CP/M is stored on a diskette, consider the first two tracks on the diskette, track 0 and track 1. Figure 2-2 shows how the data is stored on these tracks.

Loading CP/M

The events that occur after you first switch on your computer and put the CP/M diskette into a disk drive are the same as those that occur when you press the RESET button—the computer generates a RESET signal.

The RESET button stops the central processor unit (CPU). All of the internals of the CPU are set to an initial state, and all the registers are cleared to zero. The program counter is also cleared to zero so that when the RESET signal goes away (it only lasts for a few milliseconds), the CPU starts executing instructions at location 0000H in memory.

Memory chips, when they first receive power, cannot be relied upon to contain any particular value. Therefore, hardware designers arrange for some initial instructions to be forced into memory at location 0000H and onward. It is this feat that is like pulling yourself up by your own bootstraps. How can you make the computer obey a particular instruction when there is "nothing" (of any sensible value) inside the machine?

There are two common techniques for placing preliminary instructions into memory:

Force-feeding

With this approach, the hardware engineer assumes that when the RESET signal is applied, some part of the computer system, typically the floppy disk controller, can masquerade as memory. Just before the CPU is unleashed, the floppy disk controller will take control of the computer system and copy a small program into memory at location 0000H and upward. Then the CPU is allowed to start executing instructions at location 0000H. The disk controller preserves the instructions even when power is off because they are stored in nonvolatile PROM-based firmware. These instructions make the disk controller read the first sector of the first track of the system diskette into memory and then transfer control to it.

Shadow ROM

This is a variation of the force-feeding technique. The hardware manufacturer arranges some ROM at location 0000H. There is also some normal read/write memory at location 0000H, but this is electronically disabled when the RESET signal has been activated. The CPU, unleashed at location 0000H, starts to execute the ROM instruction. The first act of the ROM program is to copy itself into read/write memory at some convenient location higher up in memory and transfer control of the machine up to this copy. Then the real memory at location 0000H can be turned on, the ROM turned off, and the first sector on the disk read in.

With either technique, the result is the same. The first sector of the disk is read into memory and control is transferred to the first instruction contained in the sector.

This first sector contains the main CP/M bootstrap program. This program initializes some aspects of the hardware and then reads in the remainder of track 0 and most of the sectors on track 1 (the exact number depends on the overall length of the BIOS itself). The CP/M bootstrap program will contain only the most primitive diskette error handling, trying to read the disk over and over again if the hardware indicates that it is having problems reading a sector.

The bootstrap program loads CP/M to the correct place in memory; the load address is a constant in the bootstrap. If you need to build a version of CP/M that uses more memory, you will need to change this load address inside the bootstrap as well as the address to which the bootstrap will jump when all of CP/M has been read in. This address too is a constant in the bootstrap program.

The bootstrap program transfers control to the first instruction in the BIOS, the cold boot entry point. "Cold" implies that the operation is starting cold from an empty computer.

The cold boot code in the BIOS will set up the hardware in your computer. That is, it programs the various chips that control the speed at which serial ports transmit and receive data. It initializes the serial port chips themselves and generally readies the computer system. Its final act is to transfer control to the first instruction in the BDOS in order to start up CP/M proper.

Once the BDOS receives control, it initializes itself, scans the file directory on the system diskette, and hands over control to the CCP. The CCP then outputs the "A>" prompt to the console and waits for you to enter a command. CP/M is then ready to do your bidding.

At this point, it is worthwhile to review which CP/M parts are in memory, where in memory they are, and what functions they perform.

This overview will look at memory first. Figure 2-3 shows the positions in memory of the Console Command Processor, the Basic Disk Operating System, and the Basic Input/Output System.

By touching upon these major memory components—the CCP, BDOS, and BIOS—this discussion will consider which modules interact with them, how requests for action are passed to them, and what functions they can perform.

Console Command Processor

As you can see in Figure 2-3, the CCP is the first part of CP/M that is encountered going "up" through memory addresses. This is significant when you consider that the CCP is only necessary in between programs. When CP/M is idle, it needs the CCP to interact with you, to accept your next command. Once CP/M has started to execute the command, the CCP is redundant; any console interaction will be handled by the program you are running rather than by the CCP.

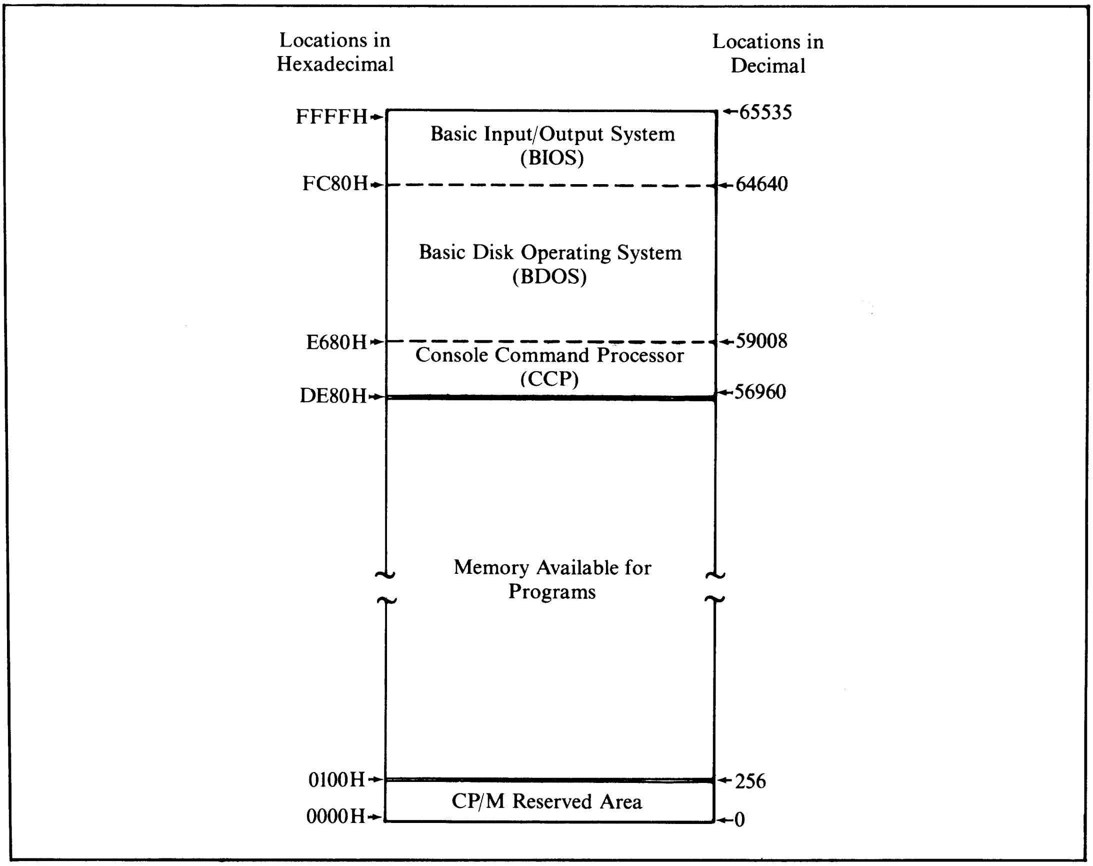


Figure 2-3. Memory layout with CP/M loaded

Therefore, the CCP leads a very jerky existence in memory. It is loaded when you first start CP/M. When you ask CP/M, via the CCP, to execute a program, this program can overwrite the CCP and use the memory occupied by the CCP for its own purposes. When the program you asked for has finished, CP/M needs to reload the CCP, now ready for its interaction with you. This process of reloading the CCP is known as a *warm boot*. In contrast with the cold boot mentioned before, the warm boot is not a complete "start from cold"; it's just a reloading of the CCP. The BDOS and BIOS are not touched.

How does a program tell CP/M that it has finished and that a warm boot must be executed? By jumping to location 0000H. While the BIOS was initializing itself during the cold boot routine, it put an instruction at location 0000H to jump to the warm boot routine, which is also in the BIOS. Once the BIOS warm boot routine has reloaded the CCP from the disk, it will transfer control to the CCP. (The cold and warm boot routines are discussed further in Chapter 6.)

This brief description indicates that every command you enter causes a program to be loaded, the CCP to be overwritten, the program to run, and the CCP to be reloaded when the program jumps to location 0000H on completing its task. This is not completely true. Some frequently needed commands reside in the CCP. Using one of these commands means that CP/M does not have to load anything from a diskette; the programs are already in memory as part of the CCP. These commands, known as "intrinsic" or "resident" commands, are listed here with a brief description of what they do. (All of them are described more thoroughly in Chapter 4.) The "resident" commands are

DIR	Displays which files are on a diskette
ERA	Erases files from a diskette
REN	Changes the names of files on diskette
TYPE	Displays the contents of text files on the console
SAVE	Saves some of memory as a file on diskette
USER	Changes User File Group.

Basic Disk Operating System

The BDOS is the heart of CP/M. The CCP and all of the programs that you run under CP/M talk to the BDOS for all their outside contacts. The BDOS performs such tasks as console input/output, printer output, and file management (creating, deleting, and renaming files and reading and writing sectors).

The BDOS performs all of these things in a rather detached way. It is concerned only with the logical tasks at hand rather than the detailed action of getting a sector from a diskette into memory, for example. These "low-level" operations are done by the BDOS in conjunction with the BIOS.

But how does a program work with the BDOS? By another strategically placed jump instruction in memory. Remember that the cold boot placed the jump to the BIOS warm boot routine in location 0000H. At location 0005H, it puts a jump instruction that transfers control up to the first instruction of the BDOS. Thus, any program that transfers control to location 0005H will find its way into the BDOS. Typically, programs make a CALL instruction to location 0005H so that once the BDOS has performed the task at hand, it can return to the calling program at the correct place. The program enlisting the BDOS's help puts special values into several of the CPU registers before it makes the call to location 0005H. These values tell the BDOS what operation is required and the other values needed for the specific operation.

Basic Input/Output System

As mentioned before, the BDOS deals with the input and output of information in a detached way, unencumbered by the physical details of the computer hardware. It is the BIOS that communicates directly with the hardware, the ports, and the peripheral devices wired to them.

This separation of *logical* input/output in the BDOS from the *physical* input/output in the BIOS is one of the major reasons why CP/M is so popular. It means that the same version of CP/M can be adapted for all types of computers, regardless of the oddities of the hardware design. Digital Research will tell you that there are over 200,000 computers in the world running CP/M. Just about all of them are running *identical* copies of the CCP and BDOS. Only the BIOS is different. If you write a program that plays by the rules and only interacts with the BDOS to get things done, it will run on almost all of those 200,000 computers without your having to change a single line of code.

You probably noticed the word "almost" in the last paragraph. Sometimes programmers make demands of the BIOS directly rather than the BDOS. This leads to trouble. The BIOS should be off limits to your program. You need to know what it is and how it works in order to build a customized version of CP/M, but you must *never* write programs that talk directly to the BIOS if you want them to run on other versions of CP/M.

Now that you understand the perils of talking to the BIOS, it is safe to describe how the BDOS communicates with the BIOS. Unlike the BDOS, which has a single entry point and uses a value in a register to specify the function to be performed, the BIOS has several entry points. The first few instructions in the BIOS are all independent entry points, each taking up three bytes of memory. The BDOS will enter the BIOS at the appropriate instruction, depending on the function to be performed. This group of entry points is similar in function to a railroad marshalling yard. It directs the BDOS to the correct destination in the BIOS for the function it needs to have done. The entry point group consists of a series of JUMP instructions, each one three bytes long. The group as a whole is called the BIOS jump table, or jump vector. Each entry point has a predefined meaning. These points are detailed and will be discussed in Chapter 6.

CCP, BDOS, and BIOS Interactions

Figure 2-4 summarizes the functions that the CCP, BDOS, and BIOS perform, the ways in which these parts of CP/M communicate among themselves, and the way in which one of your programs running under CP/M interacts with the BDOS.

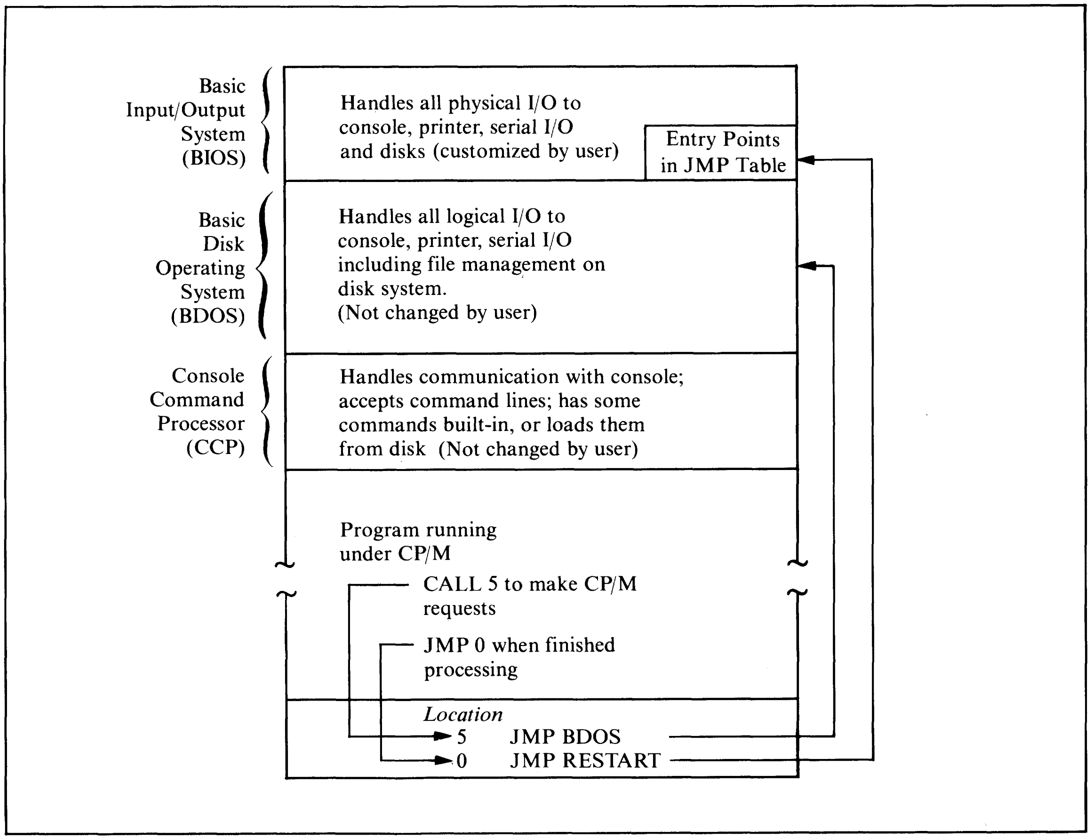


Figure 2-4. CP/M's functional breakdown

How CP/M Views the Disk
The Making of a File
Disk Definition Tables
File Organizations

3

The CP/M File System

This chapter gives you a close look at the CP/M file system. The Basic Disk Operating System (BDOS) is responsible for this file system: It keeps a directory of the files on disk, noting where data are actually stored on the disk. Because the file system automatically keeps track of this information, you can ignore the details of which tracks and sectors on the disk have data for a given file.

How CP/M Views the Disk

To manage files on the disk, CP/M works with the disk in logical terms rather than in physical terms of tracks and sectors. CP/M treats the disk as three major areas.

These are the *reserved area,* which contains the bootstrap program and CP/M itself; the *file directory,* containing one or more entries for each file stored on the disk; and the *data storage area,* which occupies the remainder of the disk. You will be looking at how CP/M allocates the storage to the files as your programs create them.

The Basic Input/Output System (BIOS) has built-in tables that tell CP/M the respective sizes of the three areas. These are the *disk definition tables,* described later in this chapter.

Allocation Blocks

Rather than work with individual 128-byte sectors, CP/M joins several of these sectors logically to form an allocation block. Typically, an allocation block will contain eight 128-byte sectors (which makes it 1024 or 1K bytes long). This makes for easier disk manipulation because the magnitude of the numbers involved is reduced. For example, a standard 8-inch, single-density, single-sided floppy disk has 1950 128-byte sectors; hard disks may have 120,000 or more. By using allocation blocks that view the disk eight sectors at a time, the number of storage units to be managed is substantially reduced. The total number is important because numeric information is handled as 16-bit integers on the 8080 and Z80 microprocessors, and therefore the largest unsigned number possible is 0FFFFH (65,535 or 64K decimal).

Whenever CP/M refers to a specific allocation block, all that is needed is a simple number. The first allocation block is number 0, the next is number 1, and so on, up to the total remaining capacity of the disk.

The typical allocation block contains 1024 (1K) bytes, or eight 128-byte sectors. For the larger hard disks, the allocation block can be 16,384 (16K) bytes, which is 128 128-byte sectors. CP/M is given the allocation via an entry in the disk definition tables in the BIOS.

The size of the allocation block is not arbitrary, but it is a compromise. The originator of the working BIOS for the system—either the manufacturer or the operating system's designer—chooses the size by considering the total storage capacity of the disk. This choice is tempered by the fact that if a file is created with only a single byte of data in it, that file would be given a complete allocation block. Large allocation blocks can waste disk storage if there are many small files, but they can be useful when a few very large files are called for.

This can be seen better by considering the case of a 1K-byte allocation block. If you create a very small file containing just a single byte of data, you will have allocated an entire allocation block. The remaining 1023 bytes will not be used. You can use them by adding to the file, but when you first create this one-byte file, they will be just so much dead space. This is the problem: Each file on the disk will normally have one partly filled allocation block. If these blocks are very large, the amount of wasted (unused) space can be very large. With 16K-byte blocks, a 10-megabyte disk with only 3 megabytes of data on it could become logically full, with all allocation blocks allocated.

On the other hand, when you use large allocation blocks, CP/M's performance is significantly improved because the BDOS refers to the file directory less frequently. For example, it can read a 16K-byte file with only a single directory reference.

Therefore, when considering block allocation, keep the following questions in mind:

How big is the logical disk?

With a larger disk, you can tolerate space wasted by incomplete allocation blocks.

What is the mean file size?

If you anticipate many small files, use small allocation blocks so that you have a larger "supply" of blocks. If you anticipate a smaller number of large files, use larger allocation blocks to get faster file operations.

When a file is first created, it is assigned a single allocation block on the disk. Which block is assigned depends on what other files you already have on the disk and which blocks have already been allocated to them. CP/M maintains a table of which blocks are allocated and which are available. As the file accumulates more data, it will fill up the first allocation block. When this happens, CP/M will extend the file and allocate another block to it. Thus, as the file grows, it occupies more blocks. These blocks need not be adjacent to each other on the disk. The file can exist as a series of allocation blocks scattered all over the disk. However, when you need to see the entire file, CP/M presents the allocation blocks in the correct order. Thus, application programs can ignore allocation blocks. CP/M keeps track of which allocation blocks belong to each file through the file directory.

The File Directory

The *file directory* is sandwiched between the reserved area and the data storage area on the disk. The actual size of the directory is defined in the BIOS's disk definition tables. The directory can have some binary multiple of entries in it, with one or more entries for each file that exists on the disk. For a standard 8-inch floppy diskette, there will be room for 64 directory entries; for a hard disk, 1024 entries would not be unusual. Each directory entry is 32 bytes long.

Simple arithmetic can be used to calculate how much space the directory occupies on a standard floppy diskette. For example, for a floppy disk the formula is $64 \times 32 = 2048$ bytes = 2 allocation blocks of 1024 bytes each.

The directory entry contains the name of the file along with a list of the allocation blocks currently used by the file. Clearly, a single 32-byte directory entry cannot contain all of the allocation blocks necessary for a 5-megabyte file, especially since CP/M uses only 16 bytes of the 32-byte total for storage of allocation block numbers.

Extents

Often CP/M will need to control files that need many allocation blocks. It does this by creating more than one directory entry. Second and subsequent directory entries have the same file name as the first. One of the other bytes of the directory entry is used to indicate the directory entry sequence number. Each new directory entry brings with it a new supply of bytes that can be used to hold more allocation block numbers. In CP/M jargon, each directory entry is called an *extent*. Because the directory entry for each extent has 16 bytes for storing allocation block numbers, it can store either 16 one-byte numbers or 8 two-byte numbers. Therefore, the total number of allocation blocks possible in each extent is either 8 (for disks with more than 255 allocation blocks) or 16 (for smaller disks).

File Control Blocks

Before CP/M can do anything with a file, it has to have some control information in memory. This information is stored in a *file control block,* or FCB. The FCB has been described as a motel for directory entries—a place for them to reside when they are not at home on the disk. When operations on a file are complete, CP/M transforms the FCB back into a directory entry and rewrites it over the original entry. The FCB is discussed in detail at the end of this chapter.

As a summary, Figure 3-1 shows the relationships between disk sectors, allocation blocks, directory entries, and file control blocks.

The Making of a File

To reinforce what you already know about the CP/M file system, this section takes you on a "walk-through" of the events that occur when a program running under CP/M creates a file, writes data to it, and then *closes* the file.

Assume that a program has been loaded in memory and the CPU is about to start executing it. First, the program will declare space in memory for an FCB and will place some preset values there, the most important of which is the file name. The area in the FCB that will hold the allocation block numbers as they are assigned is initially filled with binary 0's. Because the first allocation block that is available for file data is block 1, an allocation block number of 0 will mean that no blocks have been allocated.

The program starts executing. It makes a call to the BDOS (via location 0005H) requesting that CP/M create a file. It transfers to the BDOS the address in memory of the FCB. The BDOS then locates an available entry in the directory, creates a new entry based on the FCB in the program, and returns to the program, ready to write data to the file. Note that CP/M makes no attempt to see if there is already a file of the same name on the disk. Therefore, most real-world programs precede a request to make a file with a request to delete any existing file of the same name.

The program now starts writing data to the file, 128-byte sector by 128-byte sector. CP/ M does not have any provision for writing one byte at a time. It handles data sector-by-sector only, flushing sectors to the disk as they become full.

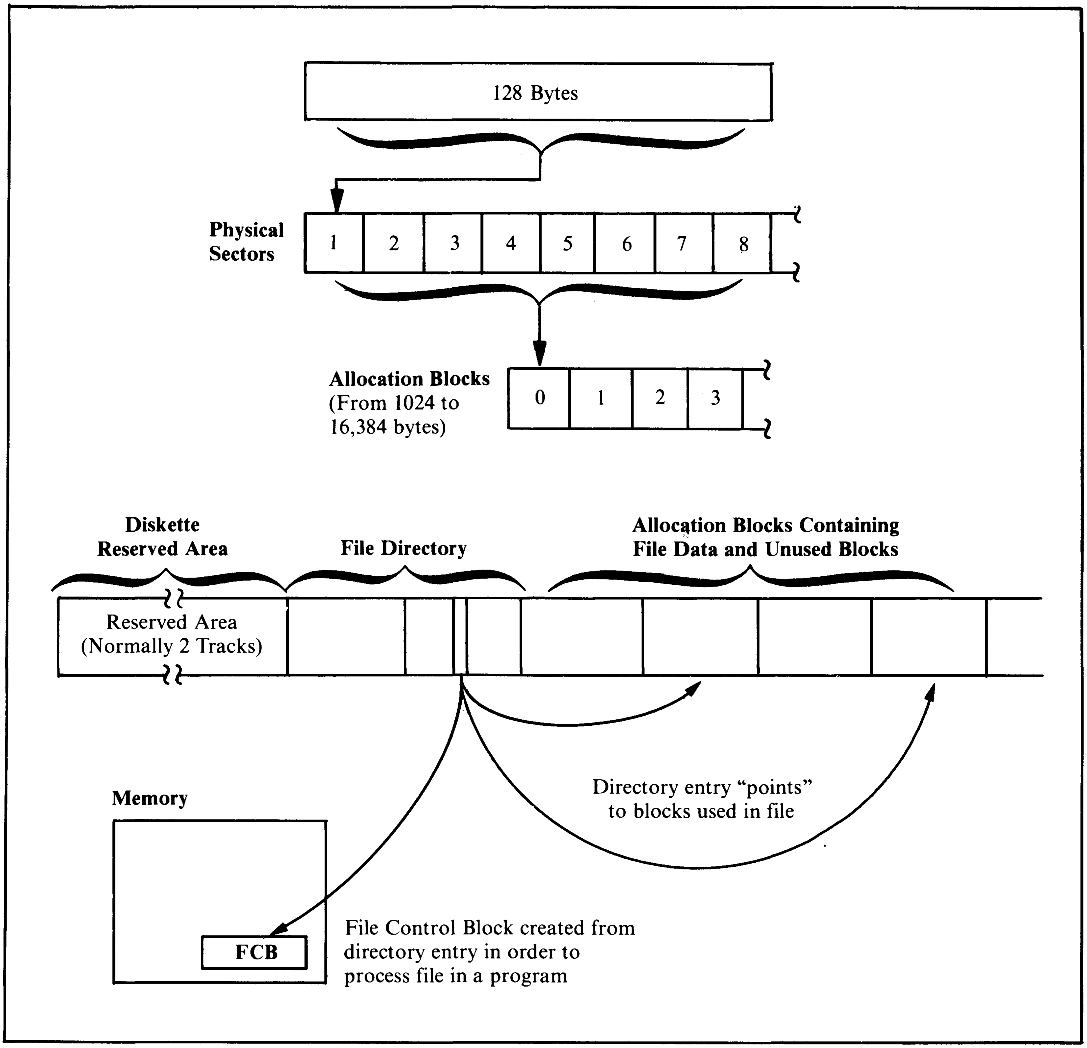


Figure 3-1. The hierarchical relationship between sectors, allocation blocks, directory entires, and FCBs

The first time a program asks CP/M (via a BDOS request) to write a sector onto the file on the disk, the BDOS finds an unused allocation block and assigns it to the file. The number of the allocation block is placed inside the FCB in memory. As each allocation block is filled up, a new allocation block is found and assigned, and its number is added to the list of allocation blocks inside the FCB. Finally, when the FCB has no more room for allocation block numbers, the BDOS

- Writes an updated directory entry out to the disk.

- Seeks out the next spare entry in the directory.
- Resets the FCB in memory to indicate that it is now working on the second extent of the file.
- Clears out the allocation block area in the FCB and waits for the next sector from the program.

Thus the process continues. New extents are automatically opened until the program determines that it is time to finish, writes the last sector out to the disk, and makes a BDOS request to close the file. The BDOS then converts the FCB into a final directory entry and writes to the directory.

Directory Entry

The directory consists of a series of 32-byte entries with one or more entries for each file on the disk. The total number of entries is a binary multiple. The actual number depends on the disk format (it will be 64 for a standard floppy disk and perhaps 2048 for a hard disk).

Figure 3-2 shows the detailed structure of a directory entry. Note that the description is actually Intel 8080 source code for the data definitions you would need in order to manipulate a directory entry. It shows a series of EQU instructions—*equate* instructions, used to assign values or expressions to a label, and in this case used to access an entry. It also shows a series of DS or *define storage* instructions used to declare storage for an entry. The comments on each line describe the function of each of the fields. Where data elements are less than a byte long, the comment identifies which bits are used.

As you study Figure 3-2, you will notice some terminology that as yet has not been discussed. This is described in detail in the sections that follow.

File User Number (Byte 0) The least significant (low order) four bits of byte 0 in the directory entry contain a number in the range 0 to 15. This is the *user number* in which the file belongs. A better name for this field would have been file group number. It works like this: Suppose several users are sharing a computer system with a hard disk that cannot be removed from the system without a lot of trouble. How can each user be sure not to tamper with other users' files? One simple way would be for each to use individual initials as the first characters of any file names. Then each could tell at a glance whether a file was another's and avoid doing anything to anyone else's files. A drawback of this scheme is that valuable character positions would be used in the file name, not to mention the problems resulting if several users had the same initials.

The file user number is prefixed to each file name and can be thought of as part of the name itself. When CP/M is first brought up, User 0 is the default user—the one that will be chosen unless another is designated. Any files created will go into the directory bearing the user number of 0. These files are referred to as being in user area 0. However, with a shared computer system, arrangements must be made for multiple user areas. The USER command makes this possible. User numbers and areas can range from 0 through 15. For example, a user in area 7 would not be able to get a directory of, access, or erase files in user area 5.

This user-number byte serves a second purpose. If this byte is set to a value of 0E5H, CP/M considers that the file directory entry has been deleted and completely ignores the remaining 31 bytes of data. The number 0E5H was not chosen whimsically. When IBM first defined the standard for floppy diskettes, they chose the binary pattern 11100101 (0E5H) as a good test pattern. A new floppy diskette formatted for use has nothing but bytes of 0E5H on it. Thus, the process of erasing a file is a "logical" deletion, where only the first byte of the directory entry is changed to 0E5H. If you accidentally delete a file (and provided that no other directory activity has occurred) it can be resurrected by simply changing this first byte back to a reasonable user number. This process will be explained in Chapter 11.

File Name and Type (Bytes 1 - 8 and 9 - 11) As you can see from Figure 3-2, the file name in a directory entry is eight bytes long; the file type is three. These two fields are used to name a file unambiguously. A file name can be less than eight characters and the file type less than three, but in these cases, the unused character positions are filled with spaces.

Whenever file names and file types are written together, they are separated by a period. You do not need the period if you are not using the file type (which is the same as saying that the file type is all spaces). Some examples of file names are

READ. ME
LONGNAME.TYP
1
1.2

```
0000 =          FDE$USER        EQU     0       ;File user number (LS 4 bits)
0001 =          FDE$NAME        EQU     1       ;File name (8 bytes)
0009 =          FDE$TYP         EQU     9       ;File type
                                                ;Offsets for bits used in type
0009 =          FDE$RO          EQU     9       ;Bit 7 = 1 - Read only
000A =          FDE$SYS         EQU     10      ;Bit 7 = 1 - System status
000B =          FDE$CHANGE      EQU     11      ;Bit 7 = 0 = File Written To
                                                ;
000C =          FDE$EXTENT      EQU     12      ;Extent number
                                                ;13, 14 reserved for CP/M
000F =          FDE$RECUSED     EQU     15      ;Records used in this extent
0010 =          FDE$ABUSED      EQU     16      ;Allocation blocks used
                ;
                ;
                ;
0000            FD$USER:        DS              ;File user number
0001            FD$NAME:        DS      8       ;File name
0009            FD$TYP:         DS      3       ;File type
000C            FD$EXTENT:      DS      1       ;Extent
000D            FD$RESV:        DS      2       ;Reserved for CP/M
000F            FD$RECUSED:     DS      1       ;Records used in this extent
0010            FD$ABUSED:      DS      16      ;Allocation blocks used
```

Figure 3-2. Data declarations for CP/M's file directory entries

A file name and type can contain the characters A through Z, 0 through 9, and some of the so-called "mark" characters such as "/ " and "—". You can also use lowercase letters, but be careful. When you enter commands into the system using the CCP, it converts all lowercases to uppercases, so it will never be able to find files that actually have lowercase letters in their directory entries. Avoid using the "mark" characters excessively. Ones you can use are

! @ # $ % ( ) − + /

Characters that you must not use are

< > . , ; : = ? * [ ]

These characters are used by CP/M in normal command lines, so using them in file names will cause problems.

You can use odd characters in file names to your advantage. For example, if you create files with nongraphic characters in their names or types, the only way you can access these files will be from within programs. You cannot manipulate these files from the keyboard except by using ambiguous file names (described in the next section). This makes it more difficult to erase files accidentally since you cannot specify their names directly from the console.

Ambiguous File Names CP/M has the capability to refer to one or more file names by using special "wild card" characters in the file names. The "?" is the main wildcard character. Whenever you ask CP/M to do something related to files, it will match a "?" with any character it finds in the file name. In the extreme case, a file name and type of "????????.???" will match with any and all file names.

As another example, all the chapters of this book were held in files called "CHAP1.DOC," "CHAP2.DOC," and so on. They were frequently referred to, however, as "CHAP??.DOC." Why two question marks? If only one had been used, for example, "CHAP?.DOC," CP/M would not have been able to match this with "CHAP10.DOC" nor any other chapter with two digits. The matching that CP/M does is strictly character-by-character.

Because typing question marks can be tedious and special attention must be paid to the exact number entered, a convenient shorthand is available. The asterisk character "*" can be used to mean "as many ?'s as you need to fill out the name or the type field." Thus, "????????.???" can be written "*.*" and "CHAP??.DOC" could also be rewritten "CHAP*.DOC."

The use of "*" is allowed only when you are entering file names from the console. The question mark notation, however, can be used for certain BDOS operations, with the file name and type field in the FCB being set to the "?" as needed.

File Type Conventions Although you are at liberty to think up file names without constraint, file types are subject to convention and, in one or two cases, to the mandate of CP/M itself.

The types that will cause problems if you do not use them correctly are

.ASM
: Assembly language source for the ASM program

.MAC
: Macro assembly language

.HEX
: Hexadecimal file output by assemblers

.REL
: Relocatable file output by assemblers

.COM
: Command file executed by entering its name alone

.PRN
: Print file written to disk as a convenience

.LIB
: Library file of programs

.SUB
: Input for CP/M SUBMIT utility program

Examples of conventional file types are

.C
: C source code

.PAS
: Pascal source code

.COB
: COBOL source code

.FTN
: FORTRAN source code

.APL
: APL programs

.TXT
: Text files

.DOC
: Documentation files

.INT
: Intermediate files

.DTA
: Data files

.IDX
Index files

.$$$
Temporary files

The file type is also useful for keeping several copies of the same file, for example, "TEST.001," "TEST.002," and so on.

File Status Each one of the states *Read-Only, System,* and *File Changed* requires only a single bit in the directory entry. To avoid using unnecessary space, they have been slotted into the three bytes used for the file type field. Since these bytes are stored as characters in ASCII (which is a seven-bit code), the most significant bit is not used for the file type and thus is available to show status.

Bit 7 of byte 9 shows Read-Only status. As its name implies, if a file is set to be Read-Only, CP/M will not allow any data to be written to the file or the file to be deleted.

If a file is declared to be System status (bit 7 of byte 10), it will not show up when you display the file directory. Nor can the file be copied from one place to another with standard CP/M utilities such as PIP unless you specifically ask the utility to do so. In normal practice, you should set your standard software tools and application programs to be both Read-Only and System status/Read-Only, so that you cannot accidentally delete them, and System status, so that they do not clutter up the directory display.

The File Changed bit (bit 7 of byte 11) is always set to 0 when you close a file to which you have been writing. This can be useful in conjunction with a file backup utility program that sets this bit to 1 whenever it makes a backup copy. Just by scanning the directory, this utility program can determine which files have changed since it was last run. The utility can be made to back up only those files that have changed. This is much easier than having to remember which files you have changed since you last made backup copies.

With a floppy disk system, there is less need to worry about backing up on a file-by-file basis — it is just as easy to copy the whole diskette. This system is useful, however, with a hard disk system with hundreds of files stored on the disk.

File Extent (Byte 12) Each directory entry represents a file extent. Byte 12 in the directory entry identified the extent number. If you have a file of less than 16,384 bytes, you will need only one extent—number 0. If you write more information to thie file, more extents will be needed. The extent number increases by 1 as each new extent is created.

The extent number is stored in the file directory because the directory entries are in random sequence. The BDOS must do a sequential search from the top of the directory to be sure of finding any given extent of a file. If the directory is large, as it could be on a hard disk system, this search can take several seconds.

Reserved Bytes 13 and 14 These bytes are used by the proprietary parts of CP/M's file system. From your point of view, they will be set to 0.

Record Number (Byte 15) Byte 15 contains a count of the number of records (128-byte sectors) that have been used in the last partially filled allocation block referenced in this directory entry. Since CP/M creates a file sequentially, only the most recently allocated block is not completely full.

Disk Map (Bytes 16 - 31) Bytes 16–31 store the allocation block numbers used by each extent. There are 16 bytes in this area. If the total number of allocation blocks (as defined by you in the BIOS disk tables) is less than 256, this area can hold as many as 16 allocation block numbers. If you have described the disk as having more than 255 allocation blocks, CP/M uses this area to store eight two-byte values. In this case allocation blocks can take on much larger values.

A directory entry can store either 8 or 16 allocation block numbers. If the file has not yet expanded to require this total number of allocation blocks, the unused positions in the entry are filled with zeros. You may think this would create a problem because it appears that several files will have been allocated block 0 over and over. In fact, there is no problem because the file directory itself always occupies block 0 (and depending on its size several of the blocks following). For all practical purposes, block 0 "does not exist," at least for the storage of file data.

Note that if, by accident, the relationship between files and their allocation blocks is scrambled—that is, either the data in a given block is overwritten, or two or more active directory entries contain the same block number—CP/M cannot access information properly and the disk becomes worthless.

Several commercially available utility programs manipulate the directory. You can use them to inspect and change a damaged directory, reviving accidentally erased files if you need to. There are other utilities you can use to logically remove bad sectors on the disk. These utilities find the bad areas, work backward from the track and sector numbers, and compute the allocation block in which the error occurs. Once the block numbers are known, they create a dummy file, either in user area 15 or, in some cases, in an "impossible" user area (one greater than 15), that appears to "own" all the bad allocation blocks.

A good utility program protects the integrity of the directory by verifying that each allocation block is "owned" by only one directory entry.

Disk Definition Tables

As mentioned previously, the BIOS contains tables telling the BDOS how to view the disk storage devices that are part of the computer system. These tables are built *by you*. If you are using standard 8-inch, single-sided, single-density floppy diskettes, you can use the examples in the Digital Research manual *CP/M 2 Alteration Guide*. But if you are using some other, more complex system, you must make some careful judgments. Any mistakes in the *disk definition tables* can create serious problems, especially when you try to correct diskettes created using the erroneous tables. You, as a programmer, must ensure the correctness of the tables by being careful.

One other point before looking at table structures: Because the tables exist and define a particular disk "shape" does not mean that such a disk need necessarily be connected to the system. The tables describe *logical* disks, and there is no way for the physical hardware to check whether your disk tables are correct. You may have a computer system with a single hard disk, yet describe the disk as though it were divided into several *logical* disks. CP/M will view each such "disk" independently, and they should be thought of as separate disks.

Disk Parameter Header Table

This table is the starting point in the disk definition tables. It is the topmost structure and contains nothing but the addresses of other structures. There is one entry in this table for each logical disk that you choose to describe. There is an entry point in the BIOS that returns the address of the parameter header table for a specific logical disk.

An example of the code needed to define a disk parameter header table is shown in Figure 3-3.

Sector Skewing (Skewtable) To define sector *skewing,* also called sector *interlacing,* picture a diskette spinning in a disk drive. The sectors in the track over which the head is positioned are passing by the head one after another—sector 1, sector 2, and so on—until the diskette has turned one complete revolution. Then the sequence repeats. A standard 8-inch diskette has 26 sectors on each track, and the disk spins at 360 rpm. One turn of the diskette takes 60/360 seconds, about 166 milliseconds per track, or 6 milliseconds per sector.

Now imagine CP/M loading a program from such a diskette. The BDOS takes a finite amount of time to read and process each sector since it reads only a single sector at a time. It has to make repeated reads to load a program. By the time the BDOS has read and loaded sector n, it will be too late to read sector n+1. This sector will have already passed by the head and will not come around for another 166 milliseconds. Proceeding in this fashion, almost 4½ seconds are needed to read one complete track.

This problem can be solved by simply numbering the sectors *logically* so that there are several physical sectors between each logical sector. This procedure, called *sector skewing* or *interlace,* is shown in Figure 3-4. Note that unlike physical sectors, logical sectors are numbered from 0 to 25.

Figure 3-4 shows the standard CP/M sector interlace for 8-inch, single-sided, single-density floppy diskettes. You see that logical sector 0 has six sectors between

```
                DPBASE:                         ;Base of the parameter header
                                                ; (used to access the headers)
0000 1000               DW      SKEWTABLE       ;Pointer to logical-to-physical
                                                ; sector conversion table
0002 0000               DW      0               ;Scratch pad areas used by CP/M
0004 0000               DW      0
0006 0000               DW      0
0008 2A00               DW      DIRBUF          ;Pointer to Directory Buffer
                                                ; work area
000A AA00               DW      DPB0            ;Pointer to disk parameter block
000C B900               DW      WACD            ;Pointer to work area (used to
                                                ; check for changed diskettes)
000E C900               DW      ALVEC0          ;Pointer to allocation vector
                ;
                ;
                ;       The following equates would normally be derived from
                ;       values found in the disk parameter Block.
                ;       They are shown here only for the sake of completeness.
                ;
003F =          NODE    EQU     63              ;Number of directory entries 1
00F2 =          NOAB    EQU     242             ;Number of allocation blocks
                ;
                ;       Example data definitions for those objects pointed
                ;       to by the disk parameter header
                ;
                SKEWTABLE:                      ;Sector skew table.
                                                ; Indexed by logical sector
0010 01070D13           DB      01,07,13,19     ;Logical sectors 0,1,2,3
0014 19050B11           DB      25,05,11,17     ;4,5,6,7
0018 1703090F           DB      23,03,09,15     ;8,9,10,11
001C 1502080E           DB      21,02,08,14     ;12,13,14,15
0020 141A060C           DB      20,26,06,12     ;16,17,18,19
0024 1218040A           DB      18,24,04,10     ;20,21,22,23
0028 1016               DB      16,22           ;24,25
                ;
002A            DIRBUF: DS      128             ;Directory buffer
00AA            DPB0:   DS      15              ;Disk parameter block
                                                ;This is normally a table of
                                                ; constants.
                                                ;A dummy definition is shown
                                                ; here
00B9            WACD:   DS      (NODE+1)/4      ;Work area to check directory
                                                ;Only used for removable media
00C9            ALVEC0: DS      (NOAB/8)+1      ;Allocation vector #0
                                                ;Needs 1 bit per allocation
                                                ; block
```

Figure 3-3. Data declarations for a disk parameter header

it and logical sector 1. There is a similar gap between each of the logical sectors, so that there are six "sector times" (about 38 milliseconds) between two adjacent logical sectors. This gives ample time for the software to access each sector. However, several revolutions of the disk are still necessary to read every sector in turn. In Figure 3-4, the vertical columns of logical sectors show which sectors are read on each successive revolution of the diskette.

The wrong interlace can strongly affect performance. It is not a gradual effect, either; if you "miss" the interlace, the perceived performance will be very slow. In the example given here, six turns of the diskette are needed to read the whole track—this lasts one second as opposed to 4½ without any interlacing. But don't imagine that you can change the interlace with impunity; files written with one interlace stay that way. You must be sure to read them back with the same interlace with which they were written.

Some disk controllers can simplify this procedure. When you format the diskette, they can write the sector addresses onto the diskette with the interlace already built in. When CP/M requests sector n, the controller's electronics wait until they see the requested sector's header fly by. They then initiate the read or write operation. In this case you can embed the interlace right into the formatting of the diskette.

Because the wrong interlace gives terrible performance, it is easy to know when you have the right one. Some programmers use the time required to format a diskette as the performance criterion to optimize the interlace. This is not good practice because under normal circumstances you will spend very little time formatting diskettes. The time spent loading a program would be a better arbiter, since far more time is spent doing this. You might argue that doing a file update would be even more representative, but most updates produce slow and sporadic disk activity. This kind of disk usage is not suitable for setting the correct interlace.

Hard disks do not present any problem for sector skewing. They spin at 3600 rpm or faster, and at that speed there simply is no interlace that will help. Some

Physical Sector	**Logical Sector**					
	Pass 1	**Pass 2**	**Pass 3**	**Pass 4**	**Pass 5**	**Pass 6**
1	0					
2				13		
3			9			
4						22
5		5				
6					18	
7	1					
8				14		
9			10			
10						23
11		6				
12					19	
13	2					
14				15		
15			11			
16						24
17		7				
18					20	
19	3					
20				16		
21			12			
22						25
23		8				
24					21	
25	4					
26				17		

NOTE: Additional sector between logical sectors 12 and 13

Figure 3-4. Physical to logical sector skewing tricks can be played to improve the performance of a hard disk—these will be discussed in the section called "Special Considerations for Hard Disks," later in this chapter.

To better understand these theories, study an example of the standard interlace table, or *skewtable*. Bear in mind that the code that will access this table will first be given a *logical* sector. It will then have to return the appropriate *physical* sector.

Figure 3-5 shows the code for the skew table and the code that can be used to access the table. The table is indexed by a logical sector and the corresponding table entry is the physical sector. You can see that the code assumes that the first *logical* sector assigned by CP/M will be sector number 0. Hence there is no need to subtract 1 from the sector number before using it as a table subscript.

Unused Areas in the Disk Parameter Header Table The three words shown as 0's in Figure 3-3 are used by CP/M as temporary variables during disk operations.

Directory Buffer (DIRBUF) The *directory buffer* is a 128-byte area used by CP/M to store a sector from the directory while processing directory entries. You only need one directory buffer; it can be shared by all of the logical disks in the system.

Disk Parameter Block (DPB0) The *disk parameter block* describes the particular characteristics of each logical disk. In general, you will need a separate parameter block for each *type* of logical disk. Logical disks can share a parameter block only if their

```
                        SKEWTABLE:                              ;Logical sector
0000 01070D13                   DB      01,07,13,19             ;0,1,2,3
0004 19050B11                   DB      25,05,11,17             ;4,5,6,7
0008 1703090F                   DB      23,03,09,15             ;8,9,10,11
000C 1502080E                   DB      21,02,08,14             ;12,13,14,15
0010 141A060C                   DB      20,26,06,12             ;16,17,18,19
0014 1218040A                   DB      18,24,04,10             ;20,21,22,23
0018 1016                       DB      16,22                   ;24,25
                        ;
                        ;
                        ;       The code to translate logical sectors to physical
                        ;        sectors is as follows:
                        ;
                        ;       On entry, the logical sector will be transferred from
                        ;       CP/M as a 16-bit value in registers BC.
                        ;       CP/M also transfers the address of the skew table
                        ;       in registers DE (it finds the skew table by looking in
                        ;       the disk parameter header entry).
                        ;
                        ;       On return, the physical sector will be placed
                        ;       in registers HL.
                        ;
                        SECTRAN:
001A EB                         XCHG                    ;HL -> skew table base address
001B 09                         DAD     B               ;HL -> physical sector
                                                        ;       entry in skew table
001C 6E                         MOV     L,M             ;L = physical sector
001D 60                         MOV     H,0             ;HL = Physical Sector
001E C9                         RET                     ;Return to BDOS
```

Figure 3-5. Data declarations for the standard skewtable for standard diskettes characteristics are identical. You can, for example, use a single parameter block to describe all of the single-sided, single-density diskette drives that you have in the system. However, you would need another parameter block to describe double-sided, double-density diskette drives. It is also rare to be able to share parameter blocks when a physical hard disk is split up into several logical disks. You will understand why after looking at the contents of a parameter block, described later in this chapter.

Work Area to Check for Changed Diskettes (WACD) One of the major problems that CP/M faces when working with removable media such as floppy diskettes is that the computer operator, without any warning, can open the diskette drive and substitute a different diskette. On early versions of CP/M, this resulted in the newly inserted diskette being overwritten with data from the original diskette.

With the current version of CP/M, you can request that CP/M check if the diskette has been changed. Given this request, CP/M examines the directory entries whenever it has worked on the directory and, if it detects that the diskette has been changed, declares the whole diskette to be Read-Only status and inhibits any further writing to the diskette. This status will be in effect until the next warm boot operation occurs. A warm boot occurs whenever a program terminates or a CONTROL-C is entered to the CCP, resetting the operating system.

The value of WACD is the address of a buffer, or temporary storage area, that CP/M can use to check the directory. The length of this buffer is defined (somewhat out of place) in the disk parameter block.

Allocation Vector (ALVEC0) CP/M views each disk as a set of allocation blocks, assigning blocks to individual files as those files are created or expanded, and relinquishing blocks as files are deleted.

CP/M needs some mechanism for keeping track of which blocks are used and which are free. It uses the *allocation vector* to form a *bit map,* with each bit in the map corresponding to a specific allocation block. The most significant bit (bit 7) in the first byte corresponds to the first allocation block, number 0. Bit 6 corresponds to block 1, and so on for the entire disk.

Whenever you request CP/M to use a logical disk, CP/M will *log in* the disk. This consists of reading down the file directory and, for each active entry or extent, interacting with the allocation blocks "owned" by that particular file extent. For each block number in the extent, the corresponding bit in the allocation vector is set to 1. At the end of this process, the allocation vector will accurately represent a map of which blocks are in use and which are free.

When CP/M goes looking for an unused allocation block, it tries to find one near the last one used, to keep the file from becoming too fragmented.

In order to reserve enough space for the allocation vector, you need to reserve one bit for each allocation block. Computing the number of allocation blocks is discussed in the section "Maximum Allocation Block Number," later in this chapter.

Disk Parameter Block

The *disk parameter block* in early versions of CP/M was built into the BDOS and was a closely guarded secret of the CP/M file system. To make CP/M adaptable to hard disk systems, Digital Research decided to move the parameter blocks out into the BIOS where everyone could adapt them. Because of the proprietary nature of CP/M's file system, you will still see several odd-looking fields, and you may find the explanation given here somewhat superficial. However, the lack of explanation in no way detracts from your ability to use CP/M as a tool.

Figure 3-6 shows the code necessary to define a parameter block for 8-inch, single-sided diskettes. This table is pointed to by—that is, its address is given in—an entry in the disk parameter header. Each of the entries shown in the disk parameter block is explained in the following sections.

Sectors Per Track This is the number of 128-byte sectors per track. The standard diskette shown in the example has 26 sectors. As you can see, simply telling CP/M that there are 26 sectors per track does not indicate whether the first sector is numbered 0 or 1. CP/M assumes that the first sector is 0; it is left to a sector translate subroutine to decipher which physical sector this corresponds to.

Hard disks normally have sector sizes larger than 128 bytes. This is discussed in the section on considerations for hard disks.

Block Shift, Block Mask, and Extent Mask These mysteriously named fields are used internally by CP/M during disk file operations. The values that you specify for them depend primarily on the size of the allocation block that you want.

Allocation block size can vary from 1024 bytes (1K) to 16,384 bytes (16K). There is a distinct trade-off between these two extremes, as discussed in the section on allocation blocks at the beginning of this chapter.

An allocation block size of 1024 (1K) bytes is suggested for floppy diskettes with capacities up to 1 megabyte, and a block size of 4096 (4K) bytes for larger floppy or hard disks.

```
                      DPB0:
0000 1A00                   DW      26          ;Sectors per track
0002 03                     DB      3           ;Block shift
0003 07                     DB      7           ;Block mask
0004 03                     DB      3           ;Extent mask
0005 F200                   DW      242         ;Max. allocation block number
0007 3F00                   DW      63          ;Number of directory entries 1
0009 C0                     DB      1100$0000B  ;Bit map for allocation blocks
000A 00                     DB      0000$0000B  ; used for directory
000B 1000                   DW      16          ;No. of bytes in dir. check buffer
000D 0200                   DW      2           ;No. of tracks before directory
```

Figure 3-6. Data declarations for the disk parameter block for standard diskettes

If you can define which block size you wish to use, you can now select the values for the block shift and the block mask from Table 3-1.

Table 3-1. Block Shift and Mask Value

Allocation Block Size	Block Shift	Block Mask
1,024	3	7
2,048	4	15
4,096	5	31
8,192	6	63
16,384	7	127

Select your required allocation block size from the left-hand column. This tells you which values of block shift and mask to enter into the disk parameter block.

The last of these three variables, the *extent mask,* depends not only on the block size but also on the total storage capacity of the logical disk. This latter consideration is only important for computing whether or not there will be fewer than 256 allocation blocks on the logical disk. Just divide the chosen allocation block size into the capacity of the logical disk and check whether you will have fewer than 256 blocks.

Keeping this answer and the allocation block size in mind, refer to Table 3-2 for the appropriate value for the extent mask field of the parameter block. Select the appropriate line according to the allocation block size you have chosen. Then, depending on the total number of allocation blocks in the logical disk, select the extent mask from the appropriate column.

Table 3-2. Extent Mask Value

Allocation Block Size	Number of Allocation Blocks	
	1 to 255	256 and Above
1,024	0	(Impossible)
2,048	1	0
4,096	3	1
8,192	7	3
16,384	15	7

Maximum Allocation Block Number This value is the *number* of the last allocation block in the logical disk. As the first block number is 0, this value is *one less* than the total number of allocation blocks on the disk. Where only a partial allocation block exists, the number of blocks is rounded down.

Figure 3-7 has an example for standard 8-inch, single-sided, single-density diskettes. Note that CP/M uses two reserved tracks on this diskette format.

Number of Directory Entries Minus 1 Do not confuse this entry with the number of files that can be stored on the logical disk; it is only the number of *entries* (minus one). Each extent of each file takes one directory entry, so very large files will consume several entries. Also note that the value in the table is *one less* than the number of entries.

On a standard 8-inch diskette, the value is 63 entries. On a hard disk, you may want to use 1023 or even 2047. Remember that CP/M performs a sequential scan down the directory and this takes a noticeable amount of time. Therefore, you should balance the number of logical disks with your estimate of the largest file size that you wish to support.

As a final note, make sure to choose a number of entries that fits evenly into one or more allocation blocks. Each directory entry needs 32 bytes, so you can compute the number of bytes required. Make sure this number can be divided by your chosen allocation block size without a remainder.

Allocation Blocks for the Directory This is a strange value; it is not a number, but a bit map. Looking at Figure 3-6, you see the example value written out in full as a binary value to illustrate how this value is defined. This 16-bit value has a bit set to 1 for each allocation block that is to be used for the file directory.

This value is derived from the number of directory entries you want to have on the disk and the size of the allocation block you want to use. One given, or

Physical characteristics:		**Calculate:**	
77	Tracks/Diskette	77	Tracks/Diskette
26	Sectors/Track	− 2	Tracks Reserved for CP/M
128	Bytes/Sector	75	Tracks for File Storage
2	Tracks Reserved for CP/M	×26	Number of Sectors
1024	Bytes/Allocation Block	1950	Sectors for File Storage
		×128	Bytes per Sector
		249,600	Bytes for File Storage
		÷1024	Bytes/Allocation Block
		243.75	Total Number of Allocation Blocks
		242	Number of the last allocation block (rounded and based on first block being Block 0)

Figure 3-7. Computing the maximum allocation block number for standard diskettes

constant, in this derivation is that the size of each directory entry is 32 bytes.

In the example, 64 entries are required (remember the number shown is one less than the required value). Each entry has 32 bytes. The total number of bytes required for the directory thus is 64 times 32, or 2048 bytes. Dividing this by the allocation block size of 1024 indicates that two allocation blocks must be reserved for the directory. You can see that the example value shows this by setting the two most significant bits of the 16-bit value.

As a word of warning, do not be tempted to declare this value using a DW (define word) pseudo-operation. Doing so will store the value *byte-reversed.*

Size of Buffer for Directory Checking As mentioned before in the discussion of the disk parameter header, CP/M can be requested to check directory entries whenever it is working on the directory. In order to do this, CP/M needs a buffer area, called the *work area to check for changed diskettes,* or WACD, in which it can hold working variables that keep a compressed record of what is on the directory. The length of this buffer area is kept in the disk parameter block; its address is specified in the parameter header. Because CP/M keeps a compressed record of the directory, you need only provide one byte for every four directory entries. You can see in Figure 3-6 that 16 bytes are specified to keep track of the 64 directory entries.

Number of Tracks Before the Directory Figure 3-8 shows the layout of CP/M on a standard floppy diskette. You will see that the first two tracks are reserved, containing the initial bootstrap code and CP/M itself. Hence the example in Figure 3-6, giving the code for a standard floppy disk, shows two reserved tracks (the number of tracks before the directory).

This *track offset value*, as it is sometimes called, provides a convenient method of dividing a physical disk into several logical disks.

Special Considerations for Hard Disks

If you want to run CP/M on a hard disk, you must provide code and build tables that make CP/M work as if it were running on a very large floppy disk. You must even include 128-byte sectors. However, this is not difficult to do.

To adapt hard disks to the 128-byte sector size, you must provide code in the disk driver in your BIOS that will present the illusion of reading and writing 128-byte sectors even though it is really working on sectors of 512 bytes. This code is called the *blocking/deblocking* routine.

If hard disks have sector sizes other than 128 bytes, what of the number of sectors per track, and the number of tracks?

Hard disks come in all sizes. The situation is further confused by the disk controllers, the hardware that controls the disk. In many cases, you can think of the hard disk as just a series of sectors without any tracks at all. The controller, given a *relative* sector number by the BIOS, can translate this sector number into which track, read/write head (if there is more than one platter), and sector are actually being referenced.

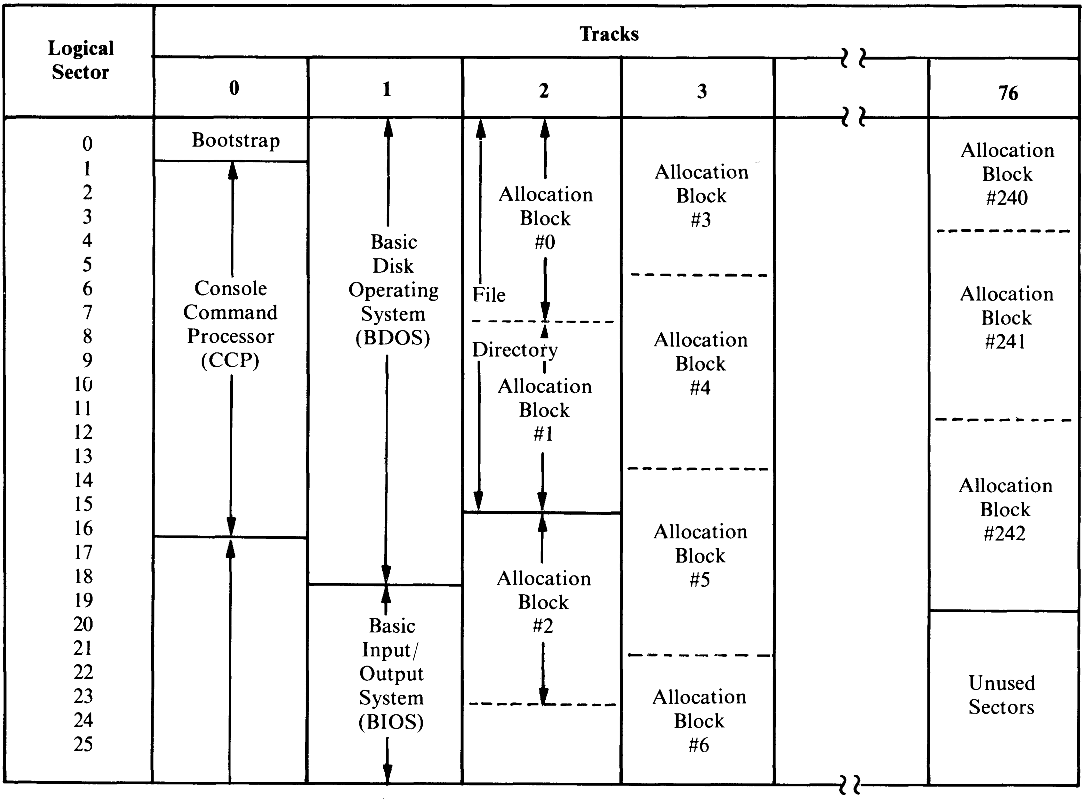


Figure 3-8. Layout of standard diskette

Furthermore, most hard disks rotate so rapidly that there is nothing to be gained by using a sector-skewing algorithm. There is just no way to read more than one physical sector per revolution; there is not enough time.

In many cases it is desirable to divide up a single, physical hard disk into several smaller, logical disks. This is done mainly for performance reasons: Several smaller disks, along with smaller directories, result in faster file operations.

The disk parameter header will have 0's for the skewtable entry and the pointer to the WACD buffer. In general, hard disks *cannot* be changed, at least not without turning off the power and swapping the entire disk drive. If you are using one of the new generation of removable hard disks, you will need to use the directory checking feature of CP/M.

The disk parameter block for a hard disk will be quite different from that used for a floppy diskette. The number of sectors per track needs careful consideration. Remember, this is the number of 128-byte sectors. The conversion from the physical sector size to 128-byte sectors will be done in the disk driver in the BIOS.

If you have a disk controller that works in terms of sectors and tracks, all you need do is compute the number of 128-byte sectors on each track. Multiply the number of physical sectors per track by their size in bytes and then divide the product by 128 to give the result as the number of 128-byte sectors per physical track.

But what of those controllers that view their hard disks as a series of sectors without reference to tracks? They obscure the fact that the sectors are arranged on concentric tracks on the disk's surface. In this case, you can play a trick on CP/M. You can set the "sectors per track" value to the number of 128-byte sectors that will fit into one of the disk's physical sectors. To do this, divide the physical sector size by 128. For example, a 512-byte physical sector size will give an answer of four 128-byte sectors per "track." You can now view the hard disk as having as many "tracks" as there are physical sectors. By using this method, you avoid having to do any kind of arithmetic on CP/M's sector numbers; the "track" number to which CP/M will ask your BIOS to move the disk heads will be the *relative physical sector*. Once the controller has read this physical sector for you, you can look at the 128-byte sector number, which will be 0, 1, 2, or 3 (for a 512-byte physical sector) in order to select which 128 bytes need to be moved in or out of the disk buffer.

The block shift, block mask, and extent mask will be computed as before. Use a 4096-byte allocation block size. This will yield a value of 5 for the block shift, 31 for the block mask, and given that you will have more than 256 allocation blocks for each logical disk, an extent mask value of 1.

The maximum allocation block number will be computed as before. Keep clear in your mind whether you are working with the number of physical sectors (which will be larger than 128 bytes) or with 128-byte sectors when you are computing the storage capacity of each logical disk.

The number of directory entries (less 1) is best set to 511 for logical disks of 1 megabyte and either 1023 or 2047 for larger disks. Remember that under CP/M version 2 you cannot have a logical disk larger than 8 megabytes.

The allocation blocks for the directory are also computed as described for floppy disks.

As a rule, the size of the directory check buffer (WADC) will be set to 0, since there is no need to use this feature on hard disk systems with fixed media.

The number of tracks before the directory (track offset) can be used to divide up the physical disk into smaller logical disks, as shown in Figure 3-9.

There is no rule that says the tracks before a logical disk's directory cannot be used to contain other complete logical disks. You can see this in Figure 3-9. CP/M behaves as if each logical disk starts at track 0 (and indeed they do), but by specifying increasingly larger numbers of tracks before each directory, the logical disks can be staggered across the available space on the physical disk.

Figure 3-10 shows the calculations involved in the first phase of building disk parameter blocks for the hard disk shown in Figure 3-9. The physical characteristics are those imposed by the design of the hard disk. As a programmer, you do not have any control over these; however, you can choose how much of the physical

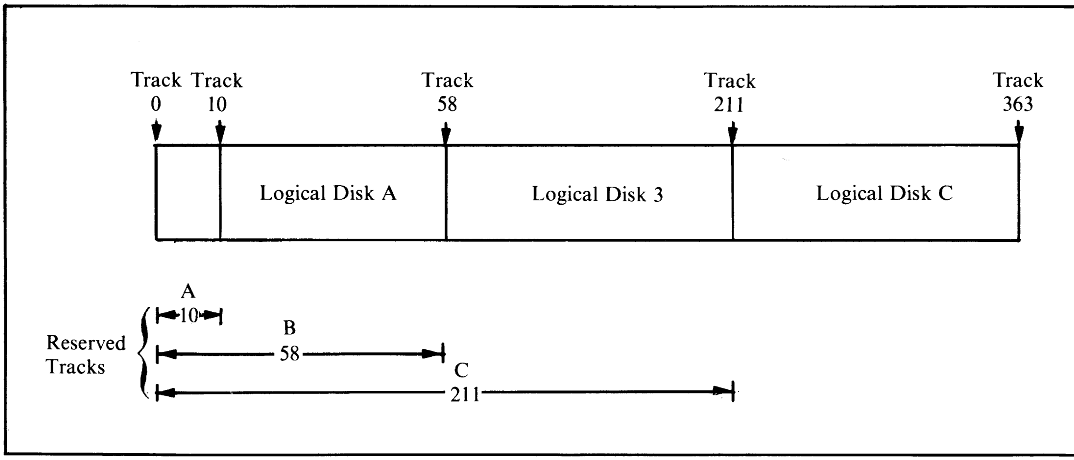


Figure 3-9. Dividing hard disks into logical disks

disk is assigned to each logical disk, the allocation block size, and the number of directory entries. You can see that logical disk A is much smaller than disks B and C, and that B and C are the same size. Disk A will be the systems disk from which most programs will be loaded, so its smaller directory size will make program loading much faster. The allocation block size for disk A is also smaller in order to reduce the amount of space wasted in partially filled allocation blocks.

Figure 3-10 also shows the calculations involved in computing the maximum allocation block number. Again, note that once the total number of allocation blocks has been computed, it is necessary to round it down in the case of any fractional components and then subtract 1 to get the maximum number (the first block being 0).

Figure 3-11 shows the actual values that will be put into the parameter blocks. It is assumed that the disk controller is one of those types that view the physical disk as a series of contiguous sectors and make no reference to tracks; the internal electronics and firmware in the controller take care of these details. For this reason, CP/M is told that each *physical* sector is a "track" in CP/M's terms. Each "track" has 512 bytes and can therefore store four 128-byte sectors. You can see this is the value that is in the sectors/"track" field.

The block shift and mask values are obtained from Table 3-1, using the allocation block size previously chosen. Then, with both the allocation block size and the maximum number of allocation blocks (see Figure 3-10), the extent mask can be obtained from Table 3-2. You can see in Figure 3-11 that extent mask values of 1 were obtained for all three logical disks even though two different allocation block sizes have been chosen, and even though disk A has less than 256 blocks and disks B and C have more.

Physical Characteristics:

364	Tracks/Disk
20	Sectors/Track
512	Bytes/Sector
10,240	Bytes/Track

Chosen Logical Characteristics:

	Tracks	Allocation Block Size
Reserved Area	10	n/a
Disk A:	48	2048
Disk B:	153	4096
Disk C:	153	4096

Calculate:

A:	B: and C:	
48	153	Tracks assigned to Disk
×10,240	×10,240	Bytes/Track
491,520	1,566,720	Bytes/Disk
÷ 2048	÷ 4096	Bytes/Allocation Block
240	382.5	Number of Allocation Blocks
239	381	Maximum Block Number

Figure 3-10. Computing the maximum allocation block number for a hard disk

```
DPBA:       DPBB:       DPBC:
    4           4           4           ;128-byte sectors/"track"
    4           5           5           ;Block shift
    15          31          31          ;Block mask
    1           1           1           ;Extent mask
    239         381         381         ;Max. all. block #
    255         1023        1023        ;No. of directory entries
    11110000B   11111111B   11111111B   ;Bit Map for allocation blocks
    00000000B   00000000B   00000000B   ; used for directory
    0           0           0           ;No. of bytes in dir.check buffer
    (10)        (58)        (211)       ;Actual tracks before directory
    200         1160        4220        ;"Tracks" before directory
```

Figure 3-11. Disk parameter tables for a hard disk

The bit map showing how many allocation blocks are required to hold the file directory is computed by multiplying the number of directory entries by 32 and dividing the product by the allocation block size. This yields results of 4 for disk A and 8 for disks B and C. As you can see, the bit maps have the appropriate number of bits set.

Since most of the hard disks on the market today do not have removable media, the lengths of the directory checking buffer are set to 0.

The number of "tracks" before the directory requires a final touch of skull-duggery. Having already indicated to CP/M that each "track" has four sectors, you need to continue in the same vein and express the number of real tracks before the directories in units of 512-byte physical sectors.

As a final note, if you are specifying these parameter blocks for a disk controller that requires you to communicate with it in terms of physical tracks and 128-byte sectors, then the number of sectors per track must be set to 80 (twenty

512-byte sectors per physical track). You would also have to change the number of tracks before the directory by stating the number of physical tracks (shown in parentheses on Figure 3-11).

Adding Additional Information to the Parameter Block

Normally, some additional information must be associated with each logical disk. For example, in a system that has several physical disks, you need to identify where each *logical* disk resides. You may also want to identify some other *physical* parameters, disk drive types, I/O port numbers, and addresses of driver subroutines.

You may be tempted to extend the disk parameter header entry because there is a separate header entry for each logical disk. But the disk parameter header is exactly 16 bytes long; adding more bytes makes the arithmetic that we need to use in the BIOS awkward. The best place to put these kinds of information is to *prefix* them to the front of each disk parameter block. The label at the front of the block must be left in the same place lest CP/M become confused. Only special additional code that you write will be "smart" enough to look *in front* of the block in order to find the additional parameter information.

File Organizations

CP/M supports two types of files: sequential and random. CP/M views both types as made up of a series of 128-byte *records*. Note that in CP/M's terms, a record is the same as a 128-byte sector. This terminology sometimes gets in the way. It may help to think of 128-byte sectors as *physical* records. Applications programs manipulate *logical* records that bear little or no relation to these physical records. There is code in the applications programs to manipulate logical records.

CP/M does not impose any restrictions on the contents of a file. In many cases, though, certain conventions are used when textual data is stored. Each line of text is terminated by ASCII CARRIAGE RETURN and LINE FEED. The last sector of a text file is filled with ASCII SUB characters; in hexadecimal this is 1AH.

File Control Blocks

In order to get CP/M to work on a file, you need to provide a structure in which both you and the BDOS can keep relevant details about the file, its name and type, and so on. The file control block (FCB) is a derivative of the file directory entry, as you can see in Figure 3-12. This figure shows both a series of equates that can be used to access an entry and a series of DB (define byte) instructions to declare an example.

The first difference you will see between the file directory entry and the FCB is that the very first byte is serving a different purpose. In the FCB, it is used to specify on which disk the file is to be found. You may recall that in the directory, this byte indicates the user number for a given entry. When you are actually processing files, the current user number is set either by the operator in a command from the console or by a BDOS function call; this predefines which subset of files in the directory will be processed. Therefore, the FCB does not need to keep track of the user number.

The disk number in the FCB's first byte is stored in an odd way. A value of 0 indicates to CP/M that it should look for the file on the current default disk. This default disk is selected either by an entry from the console or by making a specific BDOS call from within a program. In general, the default disk should be preset to the disk that contains the set of programs with which you are working. This avoids unnecessary typing on the keyboard when you want to load a program.

A disk number value other than 0 represents a letter of the alphabet based on a simple codification scheme of A = 1, B = 2, and so on.

As you can see from Figure 3-12, the file name and type must be set to the required values, and for sequential file processing, the remainder of the FCB can be set to zeros. Strictly speaking, the last three bytes of the FCB (the random record number and the random record overflow byte) need not even be declared if you are never going to process the file randomly.

This raises a subtle conceptual point. Random files are only random files because *you* process them randomly. Though this sounds like a truism, what it means is that CP/M's files are not intrinsically random or sequential. What they are depends on how you choose to process them at any given point. Therefore,

```
0000 =          FCBE$DISK       EQU     0       ;Disk drive (0 = default, 1=A)
0001 =          FCBE$NAME       EQU     1       ;File name (8 bytes)
0009 =          FCBE$TYP        EQU     9       ;File type
                                                ;Offsets for bits used in type
0009 =          FCBE$RO         EQU     9       ;Bit 7 = 1 - read only
000A =          FCBE$SYS        EQU     10      ;Bit 7 = 1 - system status
000B =          FCBE$CHANGE     EQU     11      ;Bit 7 = 0 - file written to
                                                ;
000C =          FCBE$EXTENT     EQU     12      ;Extent number
                                                ;13, 14 reserved for CP/M
000F =          FCBE$RECUSED    EQU     15      ;Records used in this extent
0010 =          FCBE$ABUSED     EQU     16      ;Allocation blocks used
0020 =          FCBE$SEQREC     EQU     32      ;Sequential rec. to read/write
0021 =          FCBE$RANREC     EQU     33      ;Random rec. to read/write
0023 =          FCBE$RANRECO    EQU     35      ;Random rec. overflow byte (MS)
                ;
                ;
                ;
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;File name
0009 545950     FCB$TYP:        DB      'TYP'   ;File type
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0       ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0       ;Random rec. overflow byte (MS)
```

Figure 3-12. Data declarations for the FCB while the manner in which you process them will be different, there is nothing special built into the file that predicates how it will be used.

Sequential Files

A sequential file begins at the beginning and ends at the end. You can view it as a contiguous series of 128-byte "records."

In order to create a sequential file, you must declare a file control block with the required file name and type and request the BDOS to *create* the file. You can then request the BDOS to write, "record" by "record" (really 128-byte sector by 128-byte sector) into the file. The BDOS will take care of opening up new extents as it needs to. When you have written out all the data, you must make a BDOS request to close the file.

To read an existing file, you also need an FCB with the required file name and type declared. You then make a BDOS request to open the file for processing and a series of Read Sequential requests, each one bringing in the next "record" until either your program detects an end of file condition (by examining the data coming in from the file) or the BDOS discovers that there are no more sectors in the file to read. There is no need to close a file from which you have been reading data—but *do close it*. This is not necessary if you are going to run the program only under CP/M, but it is necessary if you want to run under MP/M (the multiuser version of CP/M).

What if you need to append further information to an existing file? One option is to create a new file, copy the existing file to the new one, and then start adding data to the end of the new file. Fortunately, with CP/M this is not necessary. In the FCB used to read a file, the name and the type were specified, but you can also specify the extent number. If you do, the BDOS will proceed to open (if it can find it) the extent number that you are asking for. If the BDOS opens the extent successfully, all you need do is check if the number of records used in the extent (held in the field FCB$RECUSED) is less than 128 (80H). This indicates the extent is not full. By taking this record number and placing it into the FCB$SEQREC (sequential record number) byte in the FCB, you can make CP/M *jump ahead* and start writing from the effective end of the file.

Random Files

Random files use a simple variation of the technique described above. The main difference is that the random record number must be set in the FCB. The BDOS automatically keeps track of file extents during Read/Write Random requests. (These requests are explained more fully in Chapter 5.)

Conceptually, random files need a small mind-twist. After creating a file as described earlier, you must set the random record number in the FCB before each Write Random request. This is the two-byte value called FCB$RANREC in Figure 3-12. Then, when you give the Write Random request to the BDOS, it will look at the record number; compute in which extent the record must exist; if necessary, create the directory entry for the extent; and finally, write out the data record. Using this scheme, you can dart backward and forward in the file putting records at random throughout the file space, with CP/M creating the necessary directory entries each time you venture into a part of the file that has not yet been written to.

The same technique is used to read a file randomly. You set the random record number in the FCB and then give a system call to the BDOS to open the correct extent and read the data. The BDOS will return an error if it cannot find the required extent or if the particular record is nonexistent.

Problems lie in wait for the unwary. Before starting to do any random reading or writing, you must open up the file at extent 0 even though this extent may not contain any data records. For a new file, this can be done with the Create File request, and for an existing file with the normal Open File request. If you create a *sparse* file, one that has gaps in between the data, you may have some problems manipulating the file. It will appear to have several extents, each one being partially full. This will fool some programs that normally process sequential files; they don't expect to see a partial extent except at the end of a file, and may treat the wrong spot as the end.

4

The Console Command Processor (CCP)

The Console Command Processor processes commands that you enter from the console. As you may recall from the brief overview in Chapter 2, the CCP is loaded into memory immediately below the BDOS. In practice, many programs deliberately overwrite the CCP in order to use the memory it normally occupies. This gives these programs an additional 800H bytes (2K bytes).

When one of these "transient programs" terminates, it relinquishes control to the BIOS, which in turn reloads a fresh copy of the CCP from the system tracks of the disk back into memory and then transfers control to it. Consequently, the CCP leads a sporadic existence—an endless series of being loaded into memory, accepting a command from you at the console, being overwritten by the program you requested to be loaded, and then being brought back into memory when the program terminates.

This chapter discusses what the CCP does for you in those brief periods when it is in memory.

Functions of the CCP

Simply put, once the CCP has control of the machine, so do you. The CCP announces its presence by displaying a prompt of two characters: a letter of the alphabet for the current default disk drive and a "greater than" sign. In the example A>, the A tells you that the default disk drive is currently set to be logical drive A, and the ">," that the message was output by the CCP.

Once you see the prompt, the CCP is ready for you to enter a command line. A command line consists of two major parts: the name of the command and, optionally, some values for the command. This last part is known as the *command tail.*

The command itself can be one of two things: either the name of a file or the name of one of the frequently used commands built into the CCP.

If you enter the name of one of the built-in commands, the CCP does not need to go out to the disk system in order to load the command for execution. The executable code is already inside the CCP.

If the name of the command you entered does not match any of the built-in commands (the CCP has a table of their names), the CCP will search the appropriate logical disk drive for a file with a matching name and a file type of "COM" (which is short for command). You do not enter ".COM" when invoking a command—the CCP assumes a file type of "COM."

If you do not precede the name of the COM file with a logical disk drive specification, the CCP will search the current default drive. If you have prefixed the COM file's name with a specific logical drive, the CCP will look only on that drive for the program. For example, the command MYPROG will cause the CCP to look for a file called "MYPROG.COM" on the current default drive, whereas C:MYPROG would make the CCP search only on drive C.

If you enter a command name that matches neither the CCP's built-in command table nor the name of any COM file on the specified disk, the CCP will output the command name followed by a question mark, indicating it is unable to find the file.

Editing the CCP Command Line

The CCP uses a line buffer to store what you type until you strike either a CARRIAGE RETURN or a LINE FEED. If you make an error or change your mind, you can modify the incomplete command, even to the point of discarding it.

You edit the command line by entering *control characters* from the console. Control characters are designated either by the combination of keys required to generate them from the keyboard or by their official name in the ASCII character set. For example, CONTROL-J is also known as CARRIAGE RETURN or CR.

Whenever CP/M has to represent control characters, the convention is to indicate the "control" aspect of a character with a caret ("^"). For example, CONTROL-A will appear as "^A", CONTROL-Z as "^Z", and so on. But if you press the CONTROL key with the normal shift key and the "6" key, this will produce a CONTROL-^ or "^^". The representation of control keys with the caret is only necessary when outputting to the console or the printer — internally, these characters are held as their appropriate binary values.

CONTROL-C: Warm Boot If you enter a CONTROL-C as the first character of a command line, the CCP will initiate a warm boot operation. This operation resets CP/M completely, including the disk system. A fresh copy of the CCP is loaded into memory and the file directory of the current default disk drive is scanned, rebuilding the allocation bit map held in the BIOS (as discussed in Chapter 3).

The only time you would initiate a warm boot operation is after you have changed a diskette (or a disk, if you have removable media hard disks). Thus, CP/M will reset the disk system.

Note that a CONTROL-C only initiates a warm boot if it is the first character on a command line. If you enter it in any other position, the CCP will just echo it to the screen as "^C". If you have already entered several characters on a command line, use CONTROL-U or CONTROL-X to cancel the line, and then use CONTROL-C to initiate a warm boot. You can tell a warm boot has occurred because there will be a noticeable pause after the CONTROL-C before the next prompt is displayed. The system needs a finite length of time to scan the file directory and rebuild the allocation bit map.

CONTROL-E: Physical End-of-Line The CONTROL-E command is a relic of the days of the teletype and terminals that did not perform an automatic carriage return and line feed when the cursor went off the screen to the right. When you type a CONTROL-E, CP/M sends a CARRIAGE RETURN/LINE FEED command to the console, but does not start to execute the command line you have typed thus far. CONTROL-E is, in effect, a *physical* end-of-line, not a *logical* one.

As you can see, you will need to use this command only if your terminal either overprints (if it is a hard copy device) or does not wrap around when the cursor gets to the right-hand end of the line.

CONTROL-H: Backspace The CONTROL-H command is the ASCII backspace character. When you type it, the CCP will "destructively" backspace the cursor. Use it to correct typing errors you discover before you finish entering the command line. The last character you typed will disappear from the screen. The CCP does this by sending a three-character sequence of backspace, space, backspace to the console.

The CCP ignores attempts to backspace over its own prompt. It also takes care of backspacing over control characters that take two character positions on the line. The CCP sends the character sequence backspace, backspace, space, space, backspace, backspace, erasing both characters.

CONTROL-J: Line Feed/CONTROL-M: Carriage Return The CONTROL-J command is the ASCII LINE FEED character; CONTROL-M is the CARRIAGE RETURN. Both of these characters terminate the command line. The CCP will then execute the command.

CONTROL-P: Printer Echo The CONTROL-P command is used to turn on and off a feature called *printer echo*. When it is turned on, every character sent to the console is also sent to CP/M's list device. You can use this command to get a hard copy of information that normally goes only to the console.

CONTROL-P is a "toggle." The first time you type CONTROL-P it turns on printer echo; the next time you type CONTROL-P it turns off printer echo. Whenever CP/M does a warm boot, printer echo is turned off.

There is no easy way to know whether printer echo is on or off. Try typing a few CARRIAGE RETURNs, and see whether the printer responds; if it does not, type CONTROL-P and try again.

One of the shortcomings in most CP/M implementations is that the printer drivers (the software in the BIOS that controls or "drives" the printer) do not behave very intelligently if the printer is switched off or not ready when you or your program asks it to print. Under these circumstances, the software will wait forever and the system will appear to be dead. So if you "hang" the system in this way when you type a CONTROL-P, check that the printer is turned on and ready. Otherwise, you may have to reset the entire system.

CONTROL-R: Repeat Command Line The CONTROL-R command makes the CCP repeat or retype the current input line. The CCP outputs a "#" character, a CARRIAGE RETURN/LINE FEED, and then the entire contents of the command line buffer. This is a useful feature if you are working on a teletype or other hard copy terminal and have used the RUB or DEL characters. Since these characters do not destructively delete a character, you can get a visually confusing line of text on the terminal. The CONTROL-R character gives you a fresh copy of the line without any of the logically deleted characters cluttering it up. In this way you can see exactly what you have typed into the command line buffer.

See the discussion of the RUB and DEL characters for an example of CONTROL-R in use.

CONTROL-S: Stop Screen Output The CONTROL-S command is the ASCII XOFF (also called DC3) character; XOFF is an abbreviation for "Transmit Off." Typing CONTROL-S will temporarily stop output to the console. In a standard version of

CP/M, the CCP will resume output when *any* character is entered (including another CONTROL-S) from the console. Thus, you can use CONTROL-S as a toggle switch to turn console output on and off.

In some implementations of CP/M, the console driver itself (the low-level code in the BIOS that controls the console) will be maintaining a communication protocol with the console; therefore, a better way of resuming console output after pausing with a CONTROL-S is to use CONTROL-Q, the ASCII XON or "Transmit On" character. Entering a CONTROL-Q instead of relying on the fact that *any* character may be used to continue the output is a fail-safe measure.

The commands CONTROL-S and CONTROL-Q are most useful when you have large amounts of data on the screen. By "riding" the CONTROL-S and CONTROL-Q keys, you can let the data come to the screen in small bursts that you can easily scan.

CONTROL-U or CONTROL-X: Undo Command Line The commands CONTROL-U and CONTROL-X perform the same function: They erase the current partially entered command line so that you can undo any mistakes and start over. The CONTROL-U command was originally intended for hard copy terminals. The CCP outputs a "#" character, then a CARRIAGE RETURN/LINE FEED, and then some blanks to leave the cursor lined up and ready for you to enter the next command line. It leaves what you originally entered in the previous line on the screen. The CONTROL-X command is more suited to screens; the CCP destructively backspaces to the beginning of the command line so that you can reenter it.

RUB or DEL: Delete Last Character The rubout or delete function (keys marked RUB, RUBOUT, DEL, or DELETE) nondestructively deletes the last character that you typed. That is, it deletes the last character from the command line buffer and echoes it back to the console.

Here is an example of a command line with the last few characters deleted using the RUB key:

```
A>RUN PAYROLLLLORYAPSALES
                 ^^^^^^^
                 DELeted
```

You can see that the command line very quickly becomes unreadable. If you lose track of what are data characters and what has been deleted, you can use CONTROL-R to get a fresh copy of what is in the command line buffer.

The example above would then appear as follows:

```
A>RUN PAYROLLLLORYAPSALES#
  RUN SALES_
```

The "#" character is output by the CCP to indicate that the line has been repeated. The "_" represents the position of the cursor, which is now ready to continue with the command line.

Built-In Commands

When you enter a command line and press either CARRIAGE RETURN or LINE FEED, the CCP will check if the command name is one of the set of built-in commands. (It has a small table of command names embedded in it, against which the entered command name is checked.) If the command name matches a built-in one, the CCP executes the command immediately.

The next few sections describe the built-in commands that are available; however, refer to *Osborne CP/M User Guide,* second edition by Thom Hogan (Berkeley: Osborne/McGraw-Hill, 1982) for a more comprehensive discussion with examples of the various forms of each command.

X: — Changing Default Disk Drives The default drive is the currently active drive that CP/M uses for all file access whenever you do not nominate a specific drive. If you wish to change the default drive, simply enter the new default drive's identifying letter followed by a colon. The CCP responds by changing the name of the disk that appears in the prompt line.

On hard disks, this simple operation may take a second or two to complete because the BDOS, requested by the CCP to log in the drive, must read through the disk directory and rebuild the allocation vector for the disk. If you have a diskette or a disk that is removable, changing it and performing a warm boot has the same effect of refreshing CP/M's image of which allocation blocks are used and which are available. It takes longer on a hard disk because, as a rule, the directories are much larger.

DIR — Directory of Files In its simplest form, the DIR command displays a listing of the files set to Directory status in the current user number (or file group) on the current default drive. Therefore, when you do not ask for any files after the DIR command, a file name of "*.*" is assumed. This is a total wildcard, so all files that have not been given System status will be displayed. This is the only built-in command where an omitted file name reference expands to "all file names, all file types."

You can display the directory of a different drive by specifying the drive in the same command line as the DIR command.

You can qualify the files you want displayed by entering a unique or ambiguous file name or extension. Only those files that match the given file name specification will be displayed, and even then, only those files that are not set to System status will appear on the screen. (The standard CP/M utility program STAT can be used to change files from SYS to DIR status.)

Another side effect of the DIR command and files that are SYS status is best illustrated by an example. Imagine that the current logical drive B has two files on it called SYSFILE (which has SYS status) and NONSYS (which does not). Look at the following console dialog, in which user input is underlined:

```
B>DIR<cr>
B: NONSYS                SYSFILE does not show
B>DIR JUNK<cr>
NO FILE                  JUNK does not exist
B>DIR SYSFILE<cr>
B>_
```

Do you see the problem? If a file is not on the disk, the CCP will display NO FILE (or NOT FOUND in earlier versions of CP/M). However, if the file *does* exist but is a SYS file, the CCP does not display it because of its status; nor does the CCP say NO FILE. Instead it quietly returns to the prompt. This can be confusing if you are searching for a file that happens to be set to SYS status. The only safe way to find out if the file does exist is to use the STAT utility.

ERA — Erase a File The ERA command logically removes files from the disk (*logically* because only the file directory is affected; the actual data blocks are not changed).

The logical delete changes the first byte of each directory entry belonging to a file to a value of 0E5H. As you may recall from the discussion on the file directory entry in Chapter 3, this first byte usually contains the file user number. If it is set to 0E5H, it marks the entry as being deleted.

ERA makes a complete pass down the file directory to logically delete all of the extents of the file.

Unlike DIR, the ERA command does not assume "all files, all types" if you omit a file name. If it did, it would be all too easy to erase all of your files by accident. You must enter "*.*" to erase all files, and even then, you must reassure the CCP that you really want to erase all of them from the disk. The actual dialog looks like the following:

```
A>era b:*.*<cr>
ALL (Y/N)?y<cr>
A>_
```

If you change your mind at the last minute, you can press "n" and the CCP will not erase any files.

One flaw in CP/M is that the ERA command only asks for confirmation when you attempt to erase all of your files using a name such as "*.*" or "*.???". Consider the impact of the following command:

```
A>ERA *.C??<cr>
A>_
```

The CCP with no hesitation has wiped out all files that have a file type starting with the letter "C" in the current user number on logical disk A.

If you need to use an ambiguous file name in an ERA command, check which files you will delete by first using a STAT command with exactly the same ambiguous file name. STAT will show you all the files that match the ambiguous name, even those with SYS status that would not be displayed by a DIR command.

There are several utility programs on the market with names like UNERA or WHOOPS, which take an ambiguous file name and reinstate the files that you may have accidentally erased. A design for a version of UNERASE is discussed in Chapter 11.

If you attempt to erase a file that is not on the specified drive, the CCP will respond with a NO FILE message.

REN — Rename a File The REN command renames a file, changing the file name, the file type, or both. In order to rename, you need to enter two file names, the new name and the current file name.

To remember the correct name format, think of the phrase *new = old.* The actual command syntax is

```
A>ren newfile.typ=oldfile.typ<cr>
A>_
```

You can use a logical disk drive letter to specify on which drive the file exists. If you specify the drive, you only need to enter it on one of the file names. If you enter the drive with both file names, it must be the same letter for both.

Unlike the previous built-in command, REN cannot be used with ambiguous file names. If you try, the CCP echoes back the ambiguous names and a question mark, as in the following dialog:

```
A>ren chap*.doc=chapter*.doc<cr>
CHAP*.DOC=CHAPTER*.DOC?
A>_
```

If the REN command cannot find the old file, it will respond NO FILE. If the new file already exists, the message FILE EXISTS will be displayed. If you receive a FILE EXISTS message and want to check that the new file does exist, remember that it is better to use the STAT command than DIR. The extant file may be declared to be SYS status and therefore will not appear if you use the DIR command.

TYPE — Type a Text File The TYPE command copies the specified file to the console. You cannot use ambiguous file names, and you will need to press CONTROL-S if the file has more data than can fill one screen. With the TYPE command, the data in the file will fly past on the screen unless you stop the display by pressing CONTROL-S. Be careful, because if you type any other character, the TYPE command will abort and return control to the CCP.

Once you have had time to see what is displayed on the screen, you can press CONTROL-Q to resume the output of data to the console. With standard CP/M implementations, you will discover that any character can be used to restart the flow of data; however, use CONTROL-Q as a fail-safe measure. CONTROL-S (X-OFF) and CONTROL-Q (X-ON) conform to the standard protocol which should be used.

If you need to get hard copy output of the contents of the file, you should type a CONTROL-P command before you press the CARRIAGE RETURN at the end of the TYPE command line.

As you may have inferred, the TYPE command should only be used to output ASCII text files. If for some reason you use the TYPE command with a file that contains binary information, strange characters will appear on the screen. In fact, you may program your terminal into some state that can only be remedied by turning the power off and then on again. The general rule therefore is *only* use the TYPE command with ASCII text files.

SAVE — Save Memory Image on Disk The SAVE command is the hardest of the CCP's commands to explain. It is more useful to the programmer than to a typical end user. The format of this command is

```
A>SAVE n FILENAME.TYP<cr>
A>_
```

The SAVE command creates a file of the specified name and type (or overwrites an existing file of this name and type), and writes into it the specified number *n* of memory pages. A page in CP/M is 256 (100H) bytes. The SAVE command starts writing out memory from location 100H, the start of the Transient Program Area (TPA). Before you use this command, you will normally have loaded a program into the TPA. The SAVE command does just what its name implies: It saves an image of the program onto a disk file.

More often than not, when you use the SAVE command the file type will be ".COM." With the file saved in this way, the CCP will be able to load and execute the file.

USER — Change User Numbers As mentioned before, the directory of each logical disk consists of several directories that are physically interwoven but logically separated by the user number. When you use a specific user number, those files that were created when you were in another user number are logically not available to you.

The USER command provides a way for you to move from one user number to another. The command format is

```
A>USER n<cr>
A>_
```

where *n* can be any number from 0 to 15. Any other number will provoke the CCP to echoing back your entry, followed by a question mark.

But once you have switched back and forth between user numbers several times, it is easy to become confused about which user number you are in. The STAT command can be used to find the current user number. If you are in a user number that does not make a copy of STAT available to you however, all you can do is use the USER command to set yourself to another user number. You cannot find out which user number you were in; you can only tell the system the user number you want to go to.

In the custom BIOS systems discussed later, there is a way of displaying the current user number each time a warm boot occurs. If you are building a system in which you plan to utilize CP/M's user number features, you should give this display of the current user number serious thought. If you are in the wrong user number and erase files, you can create serious problems.

Some implementations of CP/M have modified the CCP so that the prompt shows the current user number as well as the default drive (similar to the prompt used in MP/M). However, this use of a nonstandard CCP is not a good practice. As a rule, customization should be confined to the BIOS.

Program Loading

The first area to consider when loading a program is the first 100H bytes of memory, called the *base page*. Several fields—units in this area of memory—are set to predetermined values before a program takes control.

To aid in this discussion, imagine a program called COPYFILE that copies one file to another. This program expects you to specify the source and destination file names on the command line. A typical command would read

```
A>copyfile tofile.typ fromfile.typ display
```

Notice the word "display." COPYFILE will, if you specify the "display" option, output the contents of the source file ("fromfile.typ") on the console as the transfer takes place.

When you press the CARRIAGE RETURN key at the end of the command line, the CCP will search the current default drive ("A" in the example) and load a file called COPYFILE.COM into memory starting at location 100H. The CCP then transfers control to location 100H—just past the base page—and COPYFILE starts executing.

Base Page

The base page normally starts from location 0000H in memory, but where there is other material in low memory addresses, it may start at a higher address. Figure 4-1 shows the assembly language code you will need to access the base page. RAM is assumed to start at location 0000H in this example.

```
0000 =      RAM              EQU     0         ;Start of RAM (and the base page)
                                               ;You may need to change this to
                                               ; some other value (e.g. 4300H)

0000        ;                ORG     RAM       ;Set location counter to RAM base
0000        WARMBOOT:        DS      3         ;Contains a JMP to warm boot entry
                                               ; in BIOS Jump vector table
            ;
0002 =      BIOSPAGE         EQU     RAM+2     ;BIOS Jump vector page
            ;
0003        IOBYTE:          DS      1         ;Input/output redirection byte
            ;
0004        CURUSER:         DS      1         ;Current user (bits 7-4)
0004 =      CURDISK          EQU     CURUSER   ;Default logical disk (bits 3-0)
            ;
0005        BDOSE:           DS      3         ;Contains a JMP to BDOS entry
0007 =      TOPRAM           EQU     BDOSE+2   ;Top page of usable RAM
            ;
0005C                        ORG     RAM+5CH   ;Bypass unused locations
            ;
005C        FCB1:            DS      16        ;File control block #1
                                               ;Note: if you use this FCB here
                                               ; you will overwrite FCB2 below.
            ;
006C        FCB2:            DS      16        ;File control block #2
                                               ;You must move this to another
                                               ; place before using it
            ;
0080                         ORG     RAM+80H   ;Bypass unused locations
            ;
            COMTAIL:                           ;Complete command tail
0080        COMTAIL$COUNT:   DS      1         ;Count of the number of chars
                                               ; in command tail (CR not incl.)
0081        COMTAIL$CHARS:   DS      127       ;Characters in command tail
                                               ; converted to uppercase and
                                               ; without trailing carriage ret.
            ;
0080                         ORG     RAM+80H   ;Redefine command tail area
            ;
0080        DMABUFFER:       DS      128       ;Default "DMA" address used
                                               ; as a 128-byte record buffer
            ;
0100                         ORG     RAM+100H  ;Bypass unused locations
            TPA:                               ;Start of transient program area
                                               ; into which programs are loaded.
```

Figure 4-1. Base page data declarations

Some versions of CP/M, such as the early Heathkit/Zenith system, have ROM from location 0000H to 42FFH. Digital Research, responding to market pressure, produced a version of CP/M that assumed RAM starting at 4300H. If you have one of these systems, you must add 4300H to all addresses in the following paragraphs *except* for those that refer to addresses at the top of memory. These will not be affected by the presence of ROM in low memory.

The individual values used in fields in the base page are described in the following sections.

Warmboot The three-byte *warmboot* field contains an instruction to jump up to the high end of RAM. This JMP instruction transfers control into the BIOS and triggers a warm boot operation. As mentioned before, a warm boot causes CP/M to reload the CCP and rebuild the allocation vector for the current default disk. If you need to cause a warm boot from within one of your assembly language programs, code

```
JMP   0             ;Warm Boot
```

BIOSPAGE The BIOS has several different entry points; however, they are all clustered together at the beginning of the BIOS. The first few instructions of the BIOS look like the following:

```
JMP   ENTRY1
JMP   ENTRY2
JMP   ENTRY3        ;and so on
```

Because of the way CP/M is put together, the first jump instruction *always* starts on a page boundary. Remember that a page is 256 (100H) bytes of memory, so a page boundary is an address where the least significant eight bits are zero. For example, the BIOS jump vector (as this set of JMPs is called) may start at an address such as F200H or E600H. The exact address is determined by the size of the BIOS.

By looking at the BIOSPAGE, the most significant byte of the address in the warmboot JMP instruction, the page address of the BIOS jump vector can be determined.

IOBYTE CP/M is based on a philosophy of separating the *physical* world from CP/M's own *logical* view of the world. This philosophy also applies to the character-oriented devices that CP/M supports.

The IOBYTE consists of four two-bit fields that can be used to assign a physical device to each of the logical ones. It is important to understand that the IOBYTE itself is just a passive data structure. Actual assignment occurs only when the physical device drivers examine the IOBYTE, interpreting its contents and selecting the correct physical drive for the cooperation of the BIOS. These device drivers are the low-level (that is, close to machine language) code in the BIOS that actually interfaces and controls the physical device.

The four *logical* devices that CP/M knows about are

1. *The console.* This is the device through which you communicate with CP/M. It is normally a terminal with a screen and a keyboard. The console is a bidirectional device: It can be used as a source for information (input) and a destination to which you can send information (output).

   In CP/M terminology, the console is known by the symbolic name of "CON:". Note the ":"—this differentiates the device name from a disk file that might be called "CON."

2. *The list device.* This is normally a printer of some sort and is used to make hard copy listings. CP/M views the printer as an output device only. This creates problems for printers that need to tell CP/M they are busy, but this

problem can be remedied by adding code to the low-level printer driver.
CP/M's name for this logical device is "LST:".

3. *The paper tape reader.* It is unusual to find a paper tape reader in use today. Originally, CP/M ran on an Intel Microcomputer Development System called the MDS-800, and this system had a paper tape reader. This device can be used only as a source for information.
   CP/M calls this logical device "RDR:".
4. *The paper tape punch.* This, too, is a relic from CP/M's early days and the MDS-800. In this case, the punch can be used only for output.
   The logical device name used by CP/M is "PUN:".

The physical arrangement of the IOBYTE fields is shown in Figure 4-2.

Each two-bit field can take on one of four values: 00, 01, 10, and 11. The particular value can be interpreted by the BIOS to mean a specific physical device, as shown in Table 4-1.

Although the actual interpretation of the IOBYTE is performed by the BIOS, the STAT utility can set the IOBYTE using the logical and physical device names, and PIP (Peripheral Interchange Program) can be used to copy data from one device to another. In addition, you can write a program that simply changes the

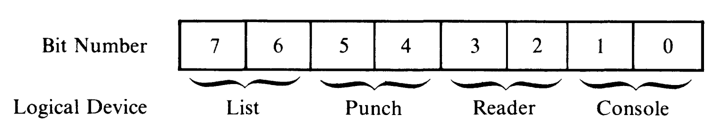


Figure 4-2. Arrangement of the IOBYTE

Table 4-1. IOBYTE Values

Logical Device	Physical Device			
	00	01	10	11
Console (CON:)	TTY:	CRT:	BAT:	UC1:
Reader (RDR:)	TTY:	PTR:	UR1:	UR2:
Punch (PUN:)	TTY:	PTP:	UP1:	UP2:
List (LST:)	TTY:	CRT:	LPT:	UL1:

contents of the IOBYTE. But be careful: Changes in the IOBYTE take effect immediately.

The values in the IOBYTE have the following meanings:

Console (CON:)

00 Teletype driver (TTY:)
This driver is assumed to be connected to a hard copy device being used as the main console.

01 CRT driver (CRT:)
The driver is assumed to be connected to a CRT terminal.

10 Batch mode (BAT:)
This is a rather special case. It is assumed that appropriate drivers will be called so that console input comes from the logical reader (RDR:) and console output is sent to the logical list device (LST:).

11 User defined console (UC1:)
Meaning depends on the individual BIOS implementation. If, for example, you have a high-resolution graphics screen, you could arrange for this setting of the IOBYTE to direct console output to it. You might make console input come in from some graphic tablet, joystick, or other device.

Reader (RDR:)

00 Teletype driver (TTY:)
This refers to the paper tape reader device that was often found on teletype consoles.

01 Paper tape reader (PTR:)
This presumes some kind of high-speed input device connected to the system. Modern systems rarely have such a device, so this setting is often used to connect the logical reader to the input side of a communications line.

10 User defined reader #1 (UR1:)

11 User defined reader #2 (UR2:)
Both of these settings can be used to direct the physical driver to some other specialized devices. These values are included only because they would otherwise have been unassigned. They are rarely used.

Punch (PUN:)

00 Teletype driver (TTY:)
This refers to the paper tape punch that was often found on teletype consoles.

01 Paper tape punch (PTP:)

This presumes that there is some kind of high-speed paper tape punch connected to the system. Again, this is rarely the case, so this setting is often used to connect the logical punch to the output side of a communications line.

10 User defined punch #1 (UP1:)

11 User defined punch #2 (UP2:)
These two settings correspond to the two user defined readers, but they are practically never used.

List (LST:)

00 Teletype driver (TTY:)
Output will be printed on a teletype.

01 CRT driver (CRT:)
Output will be directed to the screen on a CRT terminal.

10 Line printer driver (LPT:)
Output will go to a high-speed printing device. Although the name *line printer* implies a specific type of hardware, it can be any kind of printer.

11 User defined list device (UL1:)
Whoever writes the BIOS can arrange for this setting to cause logical list device output to go to a device other than the main printer.

To repeat: The IOBYTE is not actually used by the main body of CP/M. It is just a passive data structure that can be manipulated by the STAT utility. Whether the IOBYTE has any effect depends entirely on the particular BIOS implementation.

CURUSER The CURUSER field is the most significant four bits (high order nibble) of its byte. It contains the currently selected user number set by the CCP USER command, by a specific call to the BDOS, or by a program setting this nibble to the required value. This last way of changing user numbers may cause compatibility problems with future versions of CP/M, so use it only under controlled conditions.

CURDISK The CURDISK field is the least significant four bits of the byte it shares with CURUSER. It contains a value of 0 if the current disk is A:, 1 if it is B:, and so on.

The CURDISK field can be set from the CCP, by a request to the BDOS, or by a program altering this field. The caveat given for CURUSER regarding compatibility also applies here.

BDOSE This three-byte field contains an instruction to jump to the entry point of the BDOS. Whenever you want the BDOS to do something, you can transfer the request to the BDOS by placing the appropriate values in registers and making a CALL to this JMP instruction. By using a CALL, the return address will be placed on the stack. The subsequent JMP to the BDOS does not put any additional information onto the stack, which operates on a last-in, first-out basis; so when the system returns from the BDOS, it will return directly to your program.

TOPRAM Because the BDOS, like the BIOS, starts on a page boundary, the most significant byte of the address of the BDOS entry tells you in which page the BDOS starts. You must subtract 1 from the value in TOPRAM to get the highest page number that you can use in your program. Note that when you use this technique, you assume that the CCP will be overwritten since it resides in memory just below the BDOS.

FCB1 and FCB2 As a convenience, the CCP takes the first two parameters that appear in the command tail (see next section), attempts to parse them as though they were file names, and places the results in FCB1 and FCB2. The results, in this context, mean that the logical disk letter is converted to its FCB representation, and the file name and type, converted to uppercase, are placed in the FCB in the correct bytes. In addition, any use of "*" in the file name is expanded to one or more question marks. For example, a file name of "abc*.*" will be converted to a name of "ABC?????" and type of "???".

Notice that FCB2 starts only 16 bytes above FCB1, yet a normal FCB is at least 33 bytes long (36 bytes if you want to use random access). In many cases, programs only require a single file name. Therefore, you can proceed to use FCB1 straight away, not caring that FCB2 will be overwritten.

In the case of the COPYFILE program example on previous pages, two file names are required. Before FCB1 can be used, the 16 bytes of FCB2 must be moved into a skeleton FCB that is declared in the body of COPYFILE itself.

COMTAIL The command tail is everything on the command line *other* than the command name itself. For example, the command tail in the COPYFILE command line is shown here:

```
A>copyfile tofile.type fromfile.typ display
```

The CCP takes the command tail (converted to uppercase) and stores it in the COMTAIL area.

COMTAIL$COUNT This is a single-byte binary count of the number of characters in the command tail. The count does *not* include a trailing CARRIAGE RETURN or a blank between the command name and the command tail. For example, if you enter the command line

```
A>PRINT ABC*.*
```

the COMTAIL$COUNT will be six, which is the number of characters in the string "ABC*.*".

COMTAIL$CHARS These are the actual characters in the command tail. This field is not blank-filled, so you must use the COMTAIL$COUNT in order to detect the end of the command tail.

DMA$BUFFER In Figure 4-1, the DMA$BUFFER is actually the same area of memory as the COMTAIL. This is a space-saving trick that works because most programs process the contents of the command tail before they do any disk input or output.

The DMA$BUFFER is a sector buffer (hence it has a length of 128 bytes). The use of the acronym DMA (direct memory access) refers back to the Intel MDS-800. This system had hardware that could move data to and from diskettes by going directly to memory, bypassing the CPU completely. The term is still used even though you may have a computer system that does not use DMA for its disk I/O. You can substitute the idea of "the address to/from which data is read/written" in place of the DMA concept.

You can request CP/M to use a DMA address other than DMA$BUFFER, but whenever the CCP is in control, the DMA address will be set back here.

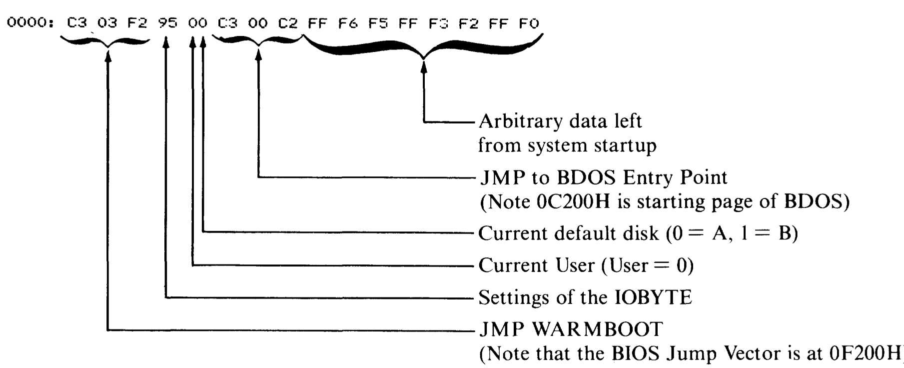
TPA This is the *transient program area* into which the CCP loads programs. The TPA extends up to the base of the BDOS.

The TPA is also the starting address for the memory image that is saved on disk whenever you use the CCP SAVE command.

Memory Dumps of the Base Page

The following are printouts showing the contents of the base page (the first 100H bytes of memory) as the COPYFILE program will see it.

This is an example of the first 16 bytes of memory:



The command line, as you recall, was

```
A>copyfile tofile.typ fromfile.typ display
```

The FCB1 and FCB2 areas will be set by the CCP as follows:

Logical Disk   Logical Disk

```
005C: 00 54 4F 46
      .  T  O  F
0060: 49 4C 45 20 20 54 59 50 00 00 00 00 00 46 52 4F
      I  L  E        T  Y  P  .  .  .  .  .  F  R  O
0070: 4D 46 49 4C 45 54 59 50 00 00 00 00 00 F2 34 F3
      M  F  I  L  E  T  Y  P  .  .  .  .  .  .  4  .
```

Since the logical disks were not specified in the file names in the command line, the CCP has set the disk code in both FCB1 and FCB2 to 00H, meaning "use the default disk." The file name and type have been converted to uppercase, separated, and put into the FCBs in their appointed places.

The complete command tail has been stored in COMTAIL as follows:

31 in decimal

Residue

```
0080: 1F 54 4F 46 49 4C 45 2E 54 59 50 20 46 52 4F 4D
      .  T  O  F  I  L  E  .  T  Y  P     F  R  O  M
0090: 46 49 4C 45 2E 54 59 50 20 44 49 53 50 4C 41 59
      F  I  L  E  .  T  Y  P     D  I  S  P  L  A  Y
00A0: 00 43 52 43 4B 20 20 20 20 43 4F 4D 00 00 00 0A
      .  C  R  C  K              C  O  M  .  .  .  .
00B0: 9B 9C 00 00 00 00 00 00 00 00 00 00 00 00 00 00
      .  .  .  .  .  .  .  .  .  .  .  .  .  .  .  .
00C0: E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5
      .  .  .  .  .  .  .  .  .  .  .  .  .  .  .  .
00D0: E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5
      .  .  .  .  .  .  .  .  .  .  .  .  .  .  .  .
00E0: E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5
      .  .  .  .  .  .  .  .  .  .  .  .  .  .  .  .
00F0: E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5 E5
      .  .  .  .  .  .  .  .  .  .  .  .  .  .  .  .
0100: 01 F9
```

Program Start

You can see that the command tail length is 01FH (31 decimal). This is followed immediately by the command tail characters themselves. Note that the command tail stops at location 9FH. The remainder of the data that you can see is the residue of some previous directory operation by the CCP. You can see the file name CRCK.COM in a directory entry, followed by several 0E5Hs that are unused directory space.

Finally, at location 0100H are the first two bytes of the program.

Processing the Command Tail

One of the first problems facing you if you write a program that can accept parameters from the command tail is to process the command tail itself, isolating each of the parameters. You should use a standard subroutine to do this. This subroutine splits the command line into individual parameters and returns a count of the number of parameters, as well as a pointer to a table of addresses. Each address in this table points in turn to a null-byte-terminated string. Each parameter is placed in a separate string.

Figure 4-3 contains the listing of this subroutine, CTP (Command Tail Processor).

```
0100                             ORG     100H
0100 CD3601      START:          CALL    CTP       ;Test bed for CTP
0103 00                          NOP
                                 ; Remainder of your program

                 ;               This subroutine breaks the command tail apart, placing
                 ;               each value in a separate string area.
                 ;
                 ;               Return parameters:
                 ;                       A = 0 - No error (Z flag set)
                 ;                       B = Count of number of parameters
                 ;                       HL -> Table of addresses
                 ;                               Each address points to a null-byte-
                 ;                               terminated parameter string.
                 ;                       If too many parameters are specified, then A = TMP
                 ;                       If a given parameter is too long, then A = PTL
                 ;                               and D points to the first character of the
                 ;                               offending parameter in the COMTAIL area.
                 ;
0080 =           COMTAIL                 EQU     80H       ;Command tail in base page
0080 =           COMTAIL$COUNT           EQU     COMTAIL   ;Count of chars. in command tail
0001 =           CTP$TMP                 EQU     1         ;Too many parameters error code
0002 =           CTP$PTL                 EQU     2         ;Parameter too long error code
                 ;
                 PTABLE:                                   ;Table of pointers to parameters
0104 0C01                        DW      P1                ; Parameter 1
0106 1A01                        DW      P2                ; Parameter 2
0108 2801                        DW      P3                ; Parameter 3
                                         ; <--- Add more parameter addresses here
010A 0000                        DW      0                 ; Terminator
                 ;
                 ;               Parameter strings.
                 ;               The first byte is 0 so that unused parameters appear
                 ;               to be null strings.
                 ;               The last byte of each is a 0 and is used to detect
                 ;               a parameter that is too long.
010C 0001010101P1:               DB     0,1,1,1,1,1,1,1,1,1,1,1,1,1,0 ;Param. 1 & terminator
011A 0001010101P2:               DB     0,1,1,1,1,1,1,1,1,1,1,1,1,1,0 ;Param. 2 & terminator
0128 0001010101P3:               DB     0,1,1,1,1,1,1,1,1,1,1,1,1,1,0 ;Param. 3 & terminator
                                        ; <--- Add more parameter strings here
                 ;
                 CTP:                                      ;Main entry point <<<<<
0136 210401                      LXI     H,PTABLE          ;HL -> table of addresses
0139 0E00                        MVI     C,0               ;Set parameter count
013B 3A8000                      LDA     COMTAIL$COUNT     ;Character count
013E B7                          ORA     A                 ;Check if any params.
013F C8                          RZ                        ;Exit (return params. already set)
0140 E5                          PUSH    H                 ;Save on top of stack for later
0141 47                          MOV     B,A               ;B = COMTAIL char. count
0142 218100                      LXI     H,COMTAIL+1       ;HL -> Command tail chars.
```

Figure 4-3. Command Tail Processor (CTP)

```
                    CTP$NEXTP:                              ;Next parameter loop
0145 E3                     XTHL                            ;HL -> Table of addresses
                                                            ;Top of stack = COMTAIL ptr.
0146 5E                     MOV     E,M                     ;Get LS byte of param. addr.
0147 23                     INX     H                       ;Update address pointer
0148 56                     MOV     D,M                     ;Get MS byte of param. addr.
                                                            ;DE -> Parameter string (or is 0)
0149 7A                     MOV     A,D                     ;Get copy of MS byte of addr.
014A B3                     ORA     E                       ;Combine MS and LS byte
014B CA8001                 JZ      CTP$TMPX                ;Too many parameters--exit
014E 23                     INX     H                       ;Update pointer to next address
014F E3                     XTHL                            ;HL -> comtail
                                                            ;Top of stack--update addr. ptr.
                                    ;At this point, we have
                                    ; HL -> next byte in command tail
                                    ; DE -> first byte of next parameter string
                    CTP$SKIPB:
0150 7E                     MOV     A,M                     ;Get next parameter byte
0151 23                     INX     H                       ;Update command tail ptr.
0152 05                     DCR     B                       ;Check if characters still remain
0153 FA7301                 MB      CTPX                    ;No, so exit
0156 FE20                   CPI     ' '                     ;Check if blank
0158 CA5001                 JZ      CTP$SKIPB               ;Yes, so skip blanks
015B 0C                     INR     C                       ;Increment parameter counter
                    CTP$NEXTC:
015C 12                     STAX    D                       ;Store in parameter string
015D 13                     INX     D                       ;Update parameter string ptr.
015E 1A                     LDAX    D                       ;Check next byte
015F B7                     ORA     A                       ;Check if terminator
0160 CA7A01                 JZ      CTP$PTLX                ;Parameter too long exit
0163 AF                     XRA     A                       ;Float a 00-byte at end of param.
0164 12                     STAX    D                       ;Store in param. string
0165 7E                     MOV     A,M                     ;Get next character from tail
0166 23                     INX     H                       ;Update command tail pointer
0167 05                     DCR     B                       ;Check if characters still remain
0168 FA7301                 JM      CTPX                    ;No, so exit
016B FE20                   CPI     ' '                     ;Check if parameter terminator
016D CA4501                 JZ      CTP$NEXTP               ;Yes, so move to next parameter
0170 C35C01                 JMP     CTP$NEXTC               ;No, so store it in param. string
                    ;
                    CTPX:                                   ;Normal exit
0173 AF                     XRA     A                       ;A = 0 & Z-flag set
                    ;
                    CTPCX                                   ;Common exit code
0174 E1                     POP     H                       ;Balance stack
0175 210401                 LXI     H,PTABLE                ;Return ptr. to param. addr. table
0178 B7                     ORA     A                       ;Ensure Z-flag set appropriately
0179 C9                     RET
                    ;
                    CTP$PTLX:                               ;Parameter too long exit
017A 3E02                   MVI     A,CTP$PTL               ;Set error code
017C EB                     XCHG                            ;DE -> offending parameter
017D C37401                 JMP     CTPCX                   ;Common exit
                    ;
                    CTP$TMPX:                               ;Too many parameters exit
0180 3E01                   MVI     A,CTP$TMP               ;Set error code
0182 C37401                 JMP     CTPCX                   ;Common exit
                    ;
0185                        END     START
```

Figure 4-3. Command Tail Processor (CTP) (continued)

Available Memory

Many programs need to use all of available memory, and so very early in the program they need to set the stack pointer to the top end of the available RAM. As mentioned before, the CCP can be overwritten as it will be reloaded on the next warm boot.

Figure 4-4 shows the code used to set the stack pointer. This code determines the amount of memory in the TPA and sets the stack pointer to the top of available RAM.

Communicating with the BIOS

If you are writing a utility program to interact with a customized BIOS, there will be occasions where you need to make a *direct* BIOS call. However, if your program ends up on a system running Digital Research's MP/M Operating System, you will have serious problems if you try to call the BIOS directly. Among other things, you will crash the operating system.

If you need to make such a call and you are aware of the dangers of using direct BIOS calls, Figure 4-5 shows you one way to do it.

Remember that the first instructions in the BIOS are the jump vector—a sequence of JMP instructions one after the other. Before you can make a direct call, you need to know the *relative page offset* of the particular JMP instruction you want to go to. The BIOS jump vector always starts on a page boundary, so all you need to know is the least significant byte of its address.

```
0007 =          TOPRAM  EQU     7       ;Most significant byte of
                ;                                   BDOS entry point
0000 3A0700             LDA     TOPRAM  ;Get MS byte of BDOS entry point
0003 3D                 DCR     A       ;Back off one page
0004 2EFF               MVI     L,0FFH  ;Set LS byte of final address
0006 67                 MOV     H,A     ;HL = XXFFH
0007 F9                 SPHL            ;Set stack pointer from HL
```

Figure 4-4. Setting stack pointer to top of available RAM

```
                ;       Use this technique only for CP/M utility programs.
                ;       MP/M programs do not permit this.
                ;
0009 =          CONIN   EQU     09H     ;Get console input character
                                        ; (It's the 4th jump in the vector)
0002 =          BIOSPAGE EQU    2       ;Address of BIOS page
                ;
                ;       At this point you make a direct CONIN
                ;       CALL...
                ;
0000 2E09               MVI     L,CONIN ;Get LS byte of CONIN entry point
0002 CD0500             CALL    BIOS    ;Go to BIOS entry subroutine
                        ;... the rest of your program...
                ;
                ;
                BIOS:
0005 3A0200             LDA     BIOSPAGE;Get BIOS jump vector page
0008 67                 MOV     H,A     ;HL -> entry point
                                        ;(You set LS byte before coming here)
0009 E9                 PCHL            ;"Jump" to BIOS
                                        ;Your return address is already
                                        ; on the stack
```

Figure 4-5. Making a direct BIOS call

```
                        ;       Note: This example assumes you have not
                        ;       overwritten the CCP.
                        ;
0100                            ORG     100H    ;Start at TPA
                        START:
0100 210000                     LXI     H,0     ;Save CCP's stack pointer
0103 39                         DAD     SP      ;By adding it to 0 in HL
0104 220F01                     SHLD    CCP$STACK
0107 314101                     LXI     SP,LOCAL$STACK
                        ;
                        ;       The main body of your program is here
                        ;
                        ;       ... and when you are ready to return
                        ;               to the CCP...
010A 2A0F01                     LHLD    CCP$STACK       ;Get CCP's stack pointer
010D F9                         SPHL                    ;Restore SP
010E C9                         RET                     ;Return to the CCP
                        ;
010F                    CCP$STACK:      DS      2       ;Save area for CCP SP
0111                                    DS      48      ;Local stack
                        LOCAL$STACK:
0141                            END     START
```

Figure 4-6. Returning to CCP at program end

Returning to CP/M

Once your program has run, you will need to return control back to CP/M. If your program has not overwritten the CCP and has left the stack pointer as it was when your program was entered, you can return directly to the CCP using a RET instruction.

Figure 4-6 shows how a normal program would do this if you use a local stack, one within the program. The CCP stack is too small; it has room for only 24 16-bit values.

The advantage of returning directly to the CCP is speed. This is true especially on a hard disk system, where the time needed to perform a warm boot is quite noticeable.

If your program has overwritten the CCP, you have no option but to transfer control to location 0000H and let the warm boot occur. To do this, all you need do is execute

```
EXIT:   JMP   0          ;Warm Boot
```

(As a hint, if you are testing a program and it suddenly exits back to CP/M, the odds are that it has inadvertently blundered to location 0000H and executed a warm boot.)

5

The Basic Disk Operating System

The Basic Disk Operating System is the real heart of CP/M. Unlike the Console Command Processor, it must be in memory all the time. It provides all of the input/output services to CP/M programs, including the CCP.

As a general rule, unless you are writing a system-dependent utility program, you should use the BDOS for *all* of your program's input/output. If you circumvent the BDOS you will probably create problems for yourself later.

What the BDOS Does

The BDOS does all of the system input/output for you. These services can be grouped into two types of functions:

Simple Byte-by-Byte I/O
This is sending and receiving data between the computer system and its logical devices—the console, the "reader" and "punch" (or their substitutes), and the printer.

Disk File I/O
This covers such tasks as creating new files, deleting old files, opening existing files, and reading and writing 128-byte long "records" to and from these files.

The remainder of this chapter explains each of the BDOS functions, shows how to make each operating system request, and gives additional information for each function. You should also refer to Digital Research's manual, *CP/M 2 Interface Guide,* for their standard description of these functions.

BDOS Function Calls

The BDOS function calls are described in the order of their function code numbers. Figure 5-1 summarizes these calls.

Naming Conventions

In practice, whenever you write programs that make BDOS calls, you should include a series of equates for the BDOS function code numbers. We shall be making reference to these values in subsequent examples, so they are shown in Figure 5-2 as they will appear in the programs.

The function names used to define the equates in Figure 5-2 are shorter than those in Figure 5-1 to strike a balance between the abbreviated function names used in Digital Research's documentation and the need for clearer function descriptions.

Making a BDOS Function Request

All BDOS functions are requested by issuing a CALL instruction to location 0005H. You can also request a function by transferring control to location 0005H with the return address on the stack.

In order to tell the BDOS what you need it to do, you must arrange for the internal registers of the CPU to contain the required information before the CALL instruction is executed.

Function Code	Description
	Simple Byte-by-Byte I/O
0	Overall system and BDOS reset
1	Read a byte from the console keyboard
2	Write a byte to the console screen
3	Read a byte from the logical reader device
4	Write a byte to the logical punch device
5	Write a byte to the logical list device
6	Direct console I/O (no CCP-style editing)
7*	Read the current setting of the IOBYTE
8*	Set a new value of the IOBYTE
9	Send a "$"-terminated string to the console
10	Read a string from the console into a buffer
11	Check if a console key is waiting to be read
12	Return the CP/M version number
	Disk File I/O
13	Reset disk system
14	Select specified logical disk drive
15	Open specified file for reading/writing
16	Close specified file after reading/writing
17	Search file directory for first match with filename
18	Search file directory for next match with filename
19	Delete (erase) file
20	Read the next "record" sequentially
21	Write the next "record" sequentially
22	Create a new file with the specified name
23	Rename a file to a new name
24	Indicate which logical disks are active
25	Return the current default disk drive number
26	Set the DMA address (read/write address)
27	Return the address of an allocation vector
28*	Set specified logical disk drive to Read-Only status
29	Indicate which disks are currently Read-Only status
30	Set specified file to System or Read-Only status
31	Return address of disk parameter block (DPB)
32*	Set/Get the current user number
33	Read a "record" randomly
34	Write a "record" randomly
35	Return logical file size (even for random files)
36	Set record number for the next random read/write
37	Reset specified drive
40	Write a "record" randomly with zero fill

*These do not work under MP/M.

Figure 5-1. BDOS function calls

```
0000 =          B$SYSRESET     EQU    0       ;System Reset
0001 =          B$CONIN        EQU    1       ;Read Console Byte
0002 =          B$CONOUT       EQU    2       ;Write Console Byte
0003 =          B$READIN       EQU    3       ;Read "Reader" Byte
0004 =          B$PUNOUT       EQU    4       ;Write "Punch" Byte
0005 =          B$LISTOUT      EQU    5       ;Write Printer Byte
0006 =          B$DIRCONIO     EQU    6       ;Direct Console I/O
0007 =          B$GETIO        EQU    7       ;Get IOBYTE
0008 =          B$SETIO        EQU    8       ;Set IOBYTE
0009 =          B$PRINTS       EQU    9       ;Print Console String
000A =          B$READCONS     EQU    10      ;Read Console String
000B =          B$CONST        EQU    11      ;Read Console Status
000C =          B$GETVER       EQU    12      ;Get CP/M Version Number
000D =          B$DSKRESET     EQU    13      ;Disk System Reset
000E =          B$SELDSK       EQU    14      ;Select Disk
000F =          B$OPEN         EQU    15      ;Open File
0010 =          B$CLOSE        EQU    16      ;Close File
0011 =          B$SEARCHF      EQU    17      ;Search for First Name Match
0012 =          B$SEARCHN      EQU    18      ;Search for Next Name Match
0013 =          B$ERASE        EQU    19      ;Erase (delete) File
0014 =          B$READSEQ      EQU    20      ;Read Sequential
0015 =          B$WRITESEQ     EQU    21      ;Write Sequential
0016 =          B$CREATE       EQU    22      ;Create File
0017 =          B$RENAME       EQU    23      ;Rename File
0018 =          B$GETACTDSK    EQU    24      ;Get Active (Logged-in) Disks
0019 =          B$GETCURDSK    EQU    25      ;Get Current Default Disk
001A =          B$SETDMA       EQU    26      ;Set DMA (Read/Write) Address
001B =          B$GETALVEC     EQU    27      ;Get Allocation Vector Address
001C =          B$SETDSKRO     EQU    28      ;Set Disk to Read Only
001D =          B$GETRODSKS    EQU    29      ;Get Read Only Disks
001E =          B$SETFAT       EQU    30      ;Set File Attributes
001F =          B$GETDPB       EQU    31      ;Get Disk Parameter Block Address
0020 =          B$SETGETUN     EQU    32      ;Set/Get User Number
0021 =          B$READRAN      EQU    33      ;Read Random
0022 =          B$WRITERAN     EQU    34      ;Write Random
0023 =          B$GETFSIZ      EQU    35      ;Get File Size
0024 =          B$SETRANREC    EQU    36      ;Set Random Record Number
0025 =          B$RESETD       EQU    37      ;Reset Drive
0028 =          B$WRITERANZ    EQU    40      ;Write Random with Zero-Fill
```

Figure 5-2. Equates for BDOS function code numbers

The function code number of the specific function call you want performed must be in register C.

If you need to hand a single-byte value to the BDOS, such as a character to be sent to the console, then you must arrange for this value to be in register E. If the value you wish to pass to the BDOS is a 16-bit value, such as the address of a buffer or a file control block (FCB), this value must be in register pair DE.

When the BDOS hands back a single-byte value, such as a keyboard character or a return code indicating the success or failure of the function you requested, it will be returned in register A. When the BDOS returns a 16-bit value, it will be in register pair HL.

On return from the BDOS, registers A and L will contain the same value, as will registers B and H. This odd convention stems from CP/M's origins in PL/M (Programming Language/Microprocessor), a language used by Intel on their MDS system. Thus, PL/M laid the foundations for what are known as "register calling conventions."

The BDOS makes no guarantee about the contents of the other registers. If you need to preserve a value that is in a register, either store the value in memory or push it onto the stack. The BDOS uses its own stack space, so there is no need to worry about it consuming your stack.

To sum up, when you make a function request to the BDOS that requires a byte value, the code and the required entry and exit parameters will be as follows:

```
MVI     C,FUNCTION$CODE         ;C = function code
MVI     E,SINGLE$BYTE           ;E = single byte value
CALL    BDOS                    ;Location 5
                                ;A = return code or value
                                ;or HL = return value
```

For those function requests that need to have an address passed to the BDOS, the calling sequence is

```
MVI     C,FUNCTION$CODE         ;C = function code
LXI     D,ADDRESS               ;DE = address
CALL    BDOS                    ;Location 5
                                ;A = return code or value
                                ;or HL = return value
```

If a function request involves disk files, you will have to tell the BDOS the address of the FCB that you have created for the file. (Refer back to Chapter 3 for descriptions of the FCB.)

Many file processing functions return a value in register A that is either 0FFH, indicating that the file named in the FCB could not be found, or equal to a value of 0, 1, 2, or 3. In the latter case, the BDOS is returning what is called a "directory code." The number is the directory entry number that the BDOS matched to the file name in your FCB. At any given moment, the BDOS has a 128-byte sector from the directory in memory. Each file directory entry is 32 bytes, so four of them (numbered 0, 1, 2, and 3) can be processed at a time. The directory code indicates which one has been matched to your FCB.

References to CP/M "records" in the following descriptions mean 128-byte sectors. Do not confuse them with the logical records used by applications programs. Think of CP/M records as 128-byte sectors throughout.

Function 0: System Reset

Function Code: C = 00H
Entry Parameters: None
Exit Parameters: Does not return

Example

```
0000 =          B$SYSRESET      EQU     0       ;System Reset
0005 =          BDOS            EQU     5       ;BDOS entry point

0000 0E00               MVI     C,B$SYSRESET    ;Set function code
0002 C30500             JMP     BDOS            ;Note: you can use a JMP since
                                                ; you don't get control back
```

Purpose The system reset function makes CP/M do a complete reset, exactly the same as the warm boot function invoked when you transfer control to the WARM-BOOT point (refer to Figure 4-1).

In addition to resetting the BDOS, this function reloads the CCP, rebuilds the allocation vectors for the currently logged disks, sets the DMA address (used by CP/M to address the disk read/write buffer) to 80H, marks all disks as being Read/Write status, and transfers control to the CCP. The CCP then outputs its prompt to the console.

Notes This function is most useful when you are working in a high-level language that does not permit a jump instruction to an absolute address in memory. Use it when your program has finished and you need to return control back to CP/M.

Function 1: Read Console Byte

Function Code: C = 01H
Entry Parameters: None
Exit Parameters: A = Data byte from console

Example

```
0001 =          B$CONIN         EQU     1           ;Console input
0005 =          BDOS            EQU     5           ;BDOS entry

0000 0E01               MVI     C,B$CONIN           ;Get function code
0002 CD0500             CALL    BDOS
```

Purpose This function reads the next byte of data from the console keyboard and puts it into register A. If the character input is a graphic character, it will be echoed back to the console. The only control characters that are echoed are CARRIAGE RETURN, LINE FEED, BACKSPACE, and TAB. In the case of a TAB character, the BDOS outputs as many spaces as are required to move the cursor to the next multiple of eight columns. All of the other control characters, including CONTROL-C, are input but are not echoed.

This function also checks for CONTROL-S (XOFF) to see if console output should be suspended, and for CONTROL-P (printer echo toggle) to see if console output should also be sent to the list device. If CONTROL-S is found, further output will be suspended until you type another character. CONTROL-P will enable the echoing of console output the first time it is pressed and disable it the second time.

If there is no incoming data character, this function will wait until there is one.

Notes This function often hinders rather than helps, because it echoes the input. Whenever you need console input at the byte-by-byte level, you will usually want to suppress this echo back to the console. For instance, you may know that the "console" is actually a communications line such as a modem. You may be trying to accept a password that should not be echoed back. Or you may need to read a cursor control character that would cause an undesirable side effect on the terminal if echoed there.

In addition, if you need more than a single character from the console, your program will be easier to use if the person at the console can take full advantage of the CCP-style line editing. This can best be done by using the Read Console String function (code 10, 0AH).

Read Console String also is more useful for single character input, especially when you are expecting a "Y" or "N" (yes or no) response. If you use the Read Console Byte function, the operator will have only one chance to enter the data. When you use Read Console String, however, users have the chance to type one character, change their minds, backspace, and type another character.

Function 2: Write Console Byte

Function Code: C = 02H
Entry Parameters: E = Data byte to be output
Exit Parameters: None

Example

```
0002 =          B$CONOUT         EQU    2          ;Write Console Byte
0005 =          BDOS             EQU    5          ;BDOS entry

0000 0E02                 MVI    C,B$CONOUT        ;Function code
0002 1E2A                 MVI    E,'*'             ;E = data byte to be output
0004 CD0500               CALL   BDOS
```

Purpose

This function outputs the data byte in register E to the console. As with function 1, if the data byte is a TAB character, it will be expanded by the BDOS to the next column that is a multiple of eight. The BDOS also checks to see if there is an incoming character, and if there is, checks to see if it is a CONTROL-S (in which case console output is suspended) or CONTROL-P (in which case echoing of console output to the printer is toggled on or off).

Notes

You may have problems using this function to output cursor-addressing control sequences to the console. If you try to output a true binary cursor address to position 9, the BDOS will interpret this as a TAB character (ASCII code 9) and dutifully replace it with zero to eight blanks. If you need to output binary values, you must set the most significant bit of the character (use an ORI 80H, for example) so that it will not be taken as the ASCII TAB.

Here are two general-purpose subroutines that you will need for outputting messages. The first one, shown in Figure 5-3, outputs a null-byte-terminated message from a specified address. The second, in Figure 5-4, does essentially the same thing *except* that the message string follows immediately after the call to the subroutine.

```
                        ;MSGOUT (message out)
                        ;Output null-byte-terminated message.

                        ;Calling sequence
                        ;       MESSAGE:        DB      'Message',0
                        ;               :
                        ;       LXI     H,MESSAGE
                        ;       CALL    MSGOUT

                        ;Exit Parameters
                        ;       HL -> Null byte terminator

0002 =                  B$CONOUT        EQU     2       ;Write Console Byte
0005 =                  BDOS            EQU     5       ;BDOS entry point

                        MSGOUT:
0000 7E                         MOV     A,M             ;Get next byte for output
0001 B7                         ORA     A
0002 C8                         RZ                      ;Return when null-byte
0003 23                         INX     H               ;Update message pointer
0004 E5                         PUSH    H               ;Save updated pointer
0005 5F                         MOV     E,A             ;Ready for BDOS
0006 0E02                       MVI     C,B$CONOUT
0008 CD0500                     CALL    BDOS
000B E1                         POP     H               ;Recover message pointer
000C C30000                     JMP     MSGOUT          ;Go back for next character
```

Figure 5-3. Write console byte example, output null-byte terminated message from specified address

```
                        ;MSGOUTI (message out in-line)
                        ;Output null-byte-terminated message that
                        ;follows the CALL to MSGOUTI.

                        ;Calling sequence
                        ;       CALL    MSGOUTI
                        ;       DB      'Message',0
                        ;       ... next instruction

                        ;Exit Parameters
                        ;       HL -> instruction following message

0002 =                  B$CONOUT        EQU     2       ;Write Console Byte
0005 =                  BDOS            EQU     5       ;BDOS entry point

                        MSGOUTI:
0000 E1                         POP     H               ;HL -> message
0001 7E                         MOV     A,M             ;Get next data byte
0002 23                         INX     H               ;Update message pointer
0003 B7                         ORA     A               ;Check if null byte
0004 C20800                     JNZ     MSGOUTIC        ;No, continue
0007 E9                         PCHL                    ;Yes, return to next instruction
                                                        ; after in-line message

                       MSGOUTIC:
0008 E5                         PUSH    H               ;Save message pointer
0009 5F                         MOV     E,A             ;Ready for BDOS
000A 0E02                       MVI     C,B$CONOUT      ;Function code
000C CD0500                     CALL    BDOS
000F C30000                     JMP     MSGOUTI         ;Go back for next char.
```

Figure 5-4. Write console byte example, output null-byte terminated message following call to subroutine

Function 3: Read "Reader" Byte

Function Code: C = 03H
Entry Parameters: None
Exit Parameters: A = Character input

Example

```
0003 =          B$READIN        EQU     3           ;Read "Reader" Byte
0005 =          BDOS            EQU     5           ;BDOS entry

0000 0E03               MVI     C,B$READIN          ;Function code
0002 CD0500             CALL    BDOS                ;A = reader byte
```

Purpose

This function reads the next character from the logical "reader" device into register A. In practice, the physical device that is accessed depends entirely on how your BIOS is configured. In some systems, there is no reader at all; this function will return some arbitrary value such as 1AH (the ASCII CONTROL-Z character, used by CP/M to denote "End of File").

Control is not returned to the calling program until a character has been read.

Notes

Since the physical device (if any) used when you issue this request depends entirely on your particular BIOS, there can be no default standard for all CP/M implementations. This is one of the weaker parts of the BDOS.

You should "connect" the reader device by means of BIOS software to a serial port that can be used for communication with another system. This is only a partial solution to the problem, however, because this function call does not return control to your program until an incoming character has been received. There is no direct way that you can "poll" the reader device to see if an incoming character has been received. Once you make this function call, you lose control until the next character arrives; there is no function corresponding to the Read Console Status (function code 11, 0BH) that will simply read status and return to your program.

One possible solution is to build a timer into the BIOS reader driver that returns control to your program with a dummy value in A if a specified period of time goes by with no incoming character. But this brings up the problem of what dummy value to use. If you ever intend to send and receive files containing pure binary information, there is no character in ASCII that you might not encounter in a legitimate context. Therefore, any dummy character you might choose could also be true data.

The most cunning solution is to arrange for one setting of the IOBYTE (which controls logical-device-to-physical-device mapping) to connect the console to the serial communication line. This done, you can make use of the Read Console Status function, which will return not the physical console status but the serial line status. Your program can then act appropriately if no characters are received within a specified time. Figure 5-11 shows a subroutine that uses this technique in the Set IOBYTE function (code 8, 08H).

Figure 5-5 shows an example subroutine to read lines of data from the reader device. It reads characters from the reader, stacking them in memory until either a LINE FEED or a specified number of characters has been received. Note that CARRIAGE RETURNs are ignored, and the input line is terminated by a byte of 00H. The convention of 00H-byte terminated strings and no CARRIAGE RETURNs is used because it makes for much easier program logic. It also conforms to the conventions of the C language.

```
                ;RL$RDR
                ;Read line from reader device.
                ;Carriage returns are ignored, and input terminates
                ;when specified number of characters have been read
                ;or a line feed is input.

                ;Note: Potential weakness is that there is no
                ;timeout in this subroutine. It will wait forever
                ;if no more characters arrive at the reader device.

                ;Calling sequence
                ;       LXI     H,BUFFER
                ;       LXI     B,MAXCOUNT
                ;       CALL    RL$RDR

                ;Exit Parameters
                ;       HL -> 00H byte terminating string
                ;       BC = residual count (0 if max. chars. read)
                ;       E = last character read

0003 =          B$READIN        EQU     3       ;Reader input
0005 =          BDOS            EQU     5       ;BDOS entry point

000D =          CR              EQU     0DH     ;Carriage return
000A =          LF              EQU     0AH     ;Line feed (terminator)

                RL$RDR:
0000 79                 MOV     A,C             ;Check if count 0
0001 B0                 ORA     B               ;If count 0 on entry, fake
0002 5F                 MOV     E,A             ; last char. read (00H)
0003 CA2000             JZ      RL$RDRX         ;Yes, exit
0006 C5                 PUSH    B               ;Save max. chars. count
0007 E5                 PUSH    H               ;Save buffer pointer

                RL$RDRI:                        ;Loop back here to ignore
0008 0E03               MVI     C,B$READIN
000A CD0500             CALL    BDOS            ;A = character input
000D 5F                 MOV     E,A             ;Preserve copy of chars.
000E FE0D               CPI     CR              ;Check if carriage return
0010 CA0800             JZ      RL$RDRI         ;Yes, ignore it
0013 E1                 POP     H               ;Recover buffer pointer
0014 C1                 POP     B               ;Recover max. Count
0015 FE0A               CPI     LF              ;Check if line feed
0017 CA2000             JZ      RL$RDRX         ;Yes, exit
001A 77                 MOV     M,A             ;No, store char. in buffer
001B 23                 INX     H               ;Update buffer pointer
001C 0B                 DCX     B               ;Downdate count
001D C30000             JMP     RL$RDR          ;Loop back for next char.

                RL$RDRX:
0020 3600               MVI     M,0             ;Null-byte-terminate buffer
0022 C9                 RET
```

Figure 5-5. Read line from reader device

Function 4: Write "Punch" Byte

Function Code: C = 04H
Entry Parameters: E = Byte to be output
Exit Parameters: None

Example

```
0004 =          B$PUNOUT        EQU     4       ;Write "Punch" Byte
0005 =          BDOS            EQU     5

0000 0E04               MVI     C,B$PUNOUT      ;Function code
0002 1E2A               MVI     E,'*'           ;Data byte to output
0004 CD0500             CALL    BDOS
```

Purpose This function is a counterpart to the Read "Reader" Byte described above. It outputs the specified character from register E to the logical punch device. Again, the actual physical device used, if any, is determined by the BIOS. There is no set standard for this device; in some systems the punch device is a "bit bucket," so called because it absorbs all data that you output to it.

Notes The problems and possible solutions discussed under the Read "Reader" Byte function call also apply here. One difference, of course, is that this function outputs data, so the problem of an indefinite loop waiting for the next character is less likely to occur. However, if your punch device is connected to a communications line, and if the output hardware is not ready, the BIOS line driver will wait forever. Unfortunately, there is no legitimate way to deal with this problem since the BDOS does not have a function call that checks whether a logical device is ready for output.

Figure 5-6 shows a useful subroutine that outputs a 00H-byte terminated string to the punch. Wherever it encounters a LINE FEED, it inserts a CARRIAGE RETURN into the output data.

Function 5: Write List Byte

Function Code: C = 05H
Entry Parameters: E = Byte to be output
Exit Parameters: None

Example

```
0005 =          B$LSTOUT        EQU     5       ;Write List Byte
0005 =          BDOS            EQU     5

0000 0E05               MVI     C,B$LSTOUT      ;Function code
0002 1E2A               MVI     E,'*'           ;Data byte to output
0004 CD0500             CALL    BDOS
```

Purpose This function outputs the specified byte in register E to the logical list device. As with the reader and the punch, the physical device used depends entirely on the BIOS.

```
                        ;WL$PUN
                        ;Write line to punch device. Output terminates
                        ;when a 00H byte is encountered.
                        ;A carriage return is output when a line feed is
                        ;encountered.

                        ;Calling sequence
                        ;       LXI     H,BUFFER
                        ;       CALL    WL$PUN

                        ;Exit parameters
                        ;       HL -> 00H byte terminator

0004 =                  B$PUNOUT        EQU     4
0005 =                  BDOS            EQU     5

000D =                  CR              EQU     0DH     ;Carriage return
000A =                  LF              EQU     0AH     ;Line feed

                        WL$PUN:
0000 E5                         PUSH    H               ;Save buffer pointer
0001 7E                         MOV     A,M             ;Get next character
0002 B7                         ORA     A               ;Check if 00H
0003 CA2000                     JZ      WL$PUNX         ;Yes, exit
0006 FE0A                       CPI     LF              ;Check if line feed
0008 CC1600                     CZ      WL$PUNLF        ;Yes, O/P CR
000B 5F                         MOV     E,A             ;Character to be output
000C 0E04                       MVI     C,B$PUNOUT      ;Function code
000E CD0500                     CALL    BDOS            ;Output character
0011 E1                         POP     H               ;Recover buffer pointer
0012 23                         INX     H               ;Increment to next char.
0013 C30000                     JMP     WL$PUN          ;Output next char

                        WL$PUNLF:                       ;Line feed encountered
0016 0E04                       MVI     C,B$PUNOUT      ;Function code
0018 1E0D                       MVI     E,CR            ;Output a CR
001A CD0500                     CALL    BDOS
001D 3E0A                       MVI     A,LF            ;Recreate line feed
001F C9                         RET                     ;Output LF

                        WL$PUNX:                        ;Exit
0020 E1                         POP     H               ;Balance the stack
0021 C9                         RET
```

Figure 5-6. Write line to punch device

Notes One of the major problems associated with this function is that it does not deal with error conditions very intelligently. You cannot be sure which physical device will be used as the logical list device, and most standard BIOS implementations will cause your program to wait forever if the printer is not ready or has run out of paper. The BDOS has no provision to return any kind of error status to indicate that there is a problem with the list device. Therefore, the BIOS will have to be changed in order to handle this situation.

Figure 5-7 is a subroutine which outputs data to the list device. As you can see, this is essentially a repeat of Figure 5-6, which performs the same function for the logical punch device.

```
                        ;WL$LST
                        ;Write line to list device. Output terminates
                        ;when a 00H byte is encountered.
                        ;A carriage return is output when a line feed is
                        ;encountered.

                        ;Calling sequence
                        ;       LXI     H,BUFFER
                        ;       CALL    WL$LST

                        ;Exit parameters
                        ;       HL -> 00H byte terminator

0005 =                  B$LSTOUT        EQU     5
0005 =                  BDOS            EQU     5

000D =                  CR              EQU     0DH     ;Carriage return
000A =                  LF              EQU     0AH     ;Line feed

                        WL$LST:
0000 E5                         PUSH    H               ;Save buffer pointer
0001 7E                         MOV     A,M             ;Get next character
0002 B7                         ORA     A               ;Check if 00H
0003 CA2000                     JZ      WL$LSTX         ;Yes, exit
0006 FE0A                       CPI     LF              ;Check if line feed
0008 CC1600                     CZ      WL$LSTLF        ;Yes, O/P CR
000B 5F                         MOV     E,A             ;Character to be output
000C 0E05                       MVI     C,B$LSTOUT      ;Function code
000E CD0500                     CALL    BDOS            ;Output character
0011 E1                         POP     H               ;Recover buffer pointer
0012 23                         INX     H               ;Update to next char.
0013 C30000                     JMP     WL$LST          ;Output next char.

                        WL$LSTLF:                       ;Line feed encountered
0016 0E05                       MVI     C,B$LSTOUT      ;Function code
0018 1E0D                       MVI     E,CR            ;Output a CR
001A CD0500                     CALL    BDOS
001D 3E0A                       MVI     A,LF            ;Recreate line feed
001F C9                         RET                     ;Output LF

                        WL$LSTX:                        ;Exit
0020 E1                         POP     H               ;Balance the stack
0021 C9                         RET
```

Figure 5-7. Write line to list device

Function 6: Direct Console I/O

Function Code: C = 06H
Entry Parameters: E = 0FFH for Input
E = Other than 0FFH for output
Exit Parameters: A = Input byte or status

Example

```
0006 =          B$DIRCONIO      EQU     6       ;Direct (raw) Console I/O
0005 =          BDOS            EQU     5       ;BDOS entry point

                                                ;Example of console input

0000 0E06               MVI     C,B$DIRCONIO    ;Function code
0002 1EFF               MVI     E,0FFH          ;0FFH means input
0004 CD0500             CALL    BDOS            ;A = 00 if no char. waiting
                                                ;A = NZ if character input
```

```
                                               ;Example of console output

0007 0E06          MVI     C,B$DIRCONIO        ;Function code
0009 1E2A          MVI     E,'*'               ;Not 0FFH means output char.
000B CD0500        CALL    BDOS
```

Purpose This function serves double duty: it both inputs and outputs characters from the console. However, it bypasses the normal control characters and line editing features (such as CONTROL-P and CONTROL-S) normally associated with console I/O. Hence the name "direct" (or "unadorned" as Digital Research describes it). If the value in register E is *not* 0FFH, then E contains a valid ASCII character that is output to the console. The logic used is most easily understood when written in pseudo-code:

```
if this is an input request (E = 0FFH)
     {
     if console status indicates a character is waiting
          {
          read the char from the console and
          return to caller with char in A
          }
     else (no input character waiting) and
          return to caller with A = 00
     }
else (output request)
     {
     output the char in E to the console and
     return to caller
     }
```

Notes This function works well provided you never have to send a value of 0FFH or expect to receive a value of 00H. If you do need to send or receive pure binary data, you cannot use this function, since these values are likely to be part of the data stream.

To understand why you might want to send and receive binary data, remember that the logical "reader" does not have any method for you to check its status to see if an incoming character has arrived. All you can do is attempt to read a character (Read Reader Byte, function code 3). However, the BDOS will not give control back to you until a character arrives (which could be a very long time). One possibility is to logically assign the console to a communications line by the use of the IOBYTE (or some similar means) and then use this Direct I/O call to send and receive data to and from the line. Then you could indeed "poll" the communications line and avoid having your program go into an indefinite wait for an incoming character. An example subroutine using this technique is shown in Figure 5-11 under Set IOBYTE (function code 8).

Figure 5-8 shows a subroutine that uses the Direct Console Input and Output. Because this example is more complex than any shown so far, the code used to check the subroutine has also been included.

Function 7: Get IOBYTE Setting

Function Code: C = 07H
Entry Parameters: None
Exit Parameters: A = IOBYTE current value

```
                ;--------------------------------------------------
                ;TESTBED CODE
                ;Because of the complexity of this subroutine, the
                ; actual testbed code has been left in this example.
                ; It assumes that DDT or ZSID
                ; will be used for checkout.
                ;--------------------------------------------------

                        IF      1                       ;Change to IF 0 to disable testbed
0100                    ORG     100H
0100 C31101             JMP     START                   ;Bypass "variables" setup by DDT

0103 00         OPTIONS:        DB      0               ;Option flags
0104 41454900   TERMS:          DB      'A','E','I',0   ;Terminators
0108 05         BUFFER          DB      5               ;Max. characters in buffer
0109 00                         DB      0               ;Actual count
010A 6363636363                 DB      99,99,99,99,99  ;Data bytes
010F 6363                       DB      99,99

                START:
0111 210801             LXI     H,BUFFER        ;Get address of buffer
0114 110401             LXI     D,TERMS         ;Address of terminator table
0117 3A0301             LDA     OPTIONS         ;Get options set by DDT
011A 47                 MOV     B,A             ;Put in correct register
011B CD2B01             CALL    RCS             ;Enter subroutine
011E CD3800             CALL    38H             ;Force DDT breakpoint
0121 C31101             JMP     START           ;Test again
                        ENDIF                   ;End of testbed

                ;RCS: Read console string (using raw input)
                ;Reads a string of characters into a memory
                ; buffer using raw input.

                ;Supports options:
                ;       o to echo characters or not (when echoing,
                ;         a carriage return will be echoed followed
                ;         by line feed)
                ;       o warm boot on input of control-C or not
                ;       o terminating input either on:
                ;               o max. no of chars input
                ;               o matching terminator character

                ; Calling Sequence
                ;       LXI     H,BUFFER
                ;               Buffer has structure:
                ;                       BUFFER: DB      10      Max. size
                ;                               DB      0       Actual Read
                ;                               DS      10+1    Buffer area
                ;       MVI     B,OPTIONS       Options required
                ;                               (see equates)
                ;       LXI     D,TERMS         Pointer to 00H-byte
                ;                               terminated Chars,
                ;                               any one of which is a
                ;                               terminator.
                ;       CALL    RCS

                ; Exit Parameters
                ;       BUFFER: Updated with data bytes and actual
                ;               character count input.
                ;               (Does not include the terminator).
                ;       A = Terminating Code
                ;               0 =     Maximum number of characters input.
                ;               NZ =    Terminator character found.

0001 =          RCS$ECHO        EQU     0000$0001B      ;Input characters to be echoed
0002 =          RCS$ABORT       EQU     0000$0010B      ;Abort on Control-C
0004 =          RCS$FOLD        EQU     0000$0100B      ;Fold lowercase to uppercase
0008 =          RCS$TERM        EQU     0000$1000B      ;DE -> term. char. set

0006 =          B$DIRCONIO      EQU     6       ;Direct console I/O
0005 =          BDOS            EQU     5       ;BDOS entry point

0003 =          CTL$C           EQU     03H     ;Control-C
000D =          CR              EQU     0DH     ;Carriage return
```

Figure 5-8. Read/write string from/to console using raw I/O

```
000A =          LF              EQU     0AH     ;Line feed
0008 =          BS              EQU     08H     ;Backspace

                RCS$ST:                         ;Internal standard terminator table
0124 0D                 DB      0DH             ;Carriage return
0125 0A                 DB      0AH             ;Line feed
0126 00                 DB      0               ;End of table

                RCS$BSS:                        ;Destructive backspace sequence
0127 08200800           DB      BS,' ',BS,0

                RCS:                            ;<<<<< Main entry
012B 23                 INX     H               ;HL -> actual count
012C 3600               MVI     M,0             ;Reset to initial state
012E 2B                 DCX     H               ;HL -> max. count

                RCS$L:
012F E5                 PUSH    H               ;Save buffer pointer
0130 CD9201             CALL    RCS$GC          ;Get character and execute:
                                                ; ECHO, ABORT, and FOLD options
                                                ;C = character input
0133 E1                 POP     H               ;Recover buffer pointer
0134 3E08               MVI     A,RCS$TERM      ;Check if user-specified terminator
0136 A0                 ANA     B               ;B = options
0137 C23D01             JNZ     RCS$UST         ;User specified terminators
013A 112401             LXI     D,RCS$ST        ;Standard terminators

                RCS$UST:
013D CDD401             CALL    RCS$CT          ;Check for terminator
0140 CA4C01             JZ      RCS$NOTT        ;Not terminator
0143 47                 MOV     B,A             ;Preserve terminating char.

                RCS$MCI:                        ;(Max. char. input shares this code)
0144 0E00               MVI     C,0             ;Terminate buffer
0146 CD7F01             CALL    RCS$SC          ;Save character
0149 78                 MOV     A,B             ;Recover terminating char.
014A B7                 ORA     A               ;Set flags
014B C9                 RET

                RCS$NOTT:                       ;Not a terminator
014C 3E08               MVI     A,BS            ;Check for backspace
014E B9                 CMP     C
014F CA6001             JZ      RCS$BS          ;Backspace entered
0152 CD7F01             CALL    RCS$SC          ;Save character in buffer
0155 CD8B01             CALL    RCS$UC          ;Update count
0158 C22F01             JNZ     RCS$L           ;Not max. so get another char.
015B 0600               MVI     B,0             ;Fake terminating char.
015D C34401             JMP     RCS$MCI         ;A = 0 for max. chars. input

                RCS$BS:                         ;Backspace entered
0160 E5                 PUSH    H               ;Save buffer pointer
0161 23                 INX     H               ;HL -> actual count
0162 35                 DCR     M               ;Back up one
0163 FA7A01             JM      RCS$NBS         ;Check if count negative
0166 212701             LXI     H,RCS$BSS       ;HL -> backspacing sequence
0169 3E01               MVI     A,RCS$ECHO      ;No, check if echoing
016B A0                 ANA     B               ;BS will have been echoed if so
016C CA7001             JZ      RCS$BSNE        ;No, input BS not echoed
016F 23                 INX     H               ;Bypass initial backspace

                RCS$BSNE:
0170 C5                 PUSH    B               ;Save options and character
0171 D5                 PUSH    D               ;Save terminator table pointer
0172 CDF601             CALL    WCS             ;Write console string
0175 D1                 POP     D               ;Recover terminator table pointer
0176 C1                 POP     B               ;Recover options and character
0177 C37B01             JMP     RCS$BSX         ;Exit from backspace logic

                RCS$NBS:
017A 34                 INR     M               ;Reset count to 0

                RCS$BSX:
017B E1                 POP     H               ;Recover buffer pointer
017C C32F01             JMP     RCS$L           ;Get next character
```

Figure 5-8. (Continued)

```
                    RCS$SC:                               ;Save character in C in buffer
                                                          ;HL -> buffer pointer
0017F D5                    PUSH     D                    ;Save terminator table pointer
0180 E5                     PUSH     H                    ;Save buffer pointer
0181 23                     INX      H                    ;HL -> actual count in buffer
0182 5E                     MOV      E,M                  ;Get actual count
0183 1C                     INR      E                    ;Count of 0 points to first data byte
0184 1600                   MVI      D,0                  ;Make word value of actual count
0186 19                     DAD      D                    ;HL -> next free data byte
0187 71                     MOV      M,C                  ;Save data byte away
0188 E1                     POP      H                    ;Recover buffer pointer
0189 D1                     POP      D                    ;Recover terminator table
                                                          ; pointer
018A C9                     RET

                    RCS$UC:                               ;Update buffer count and check for max.
                                                          ;Return Z set if = to max., NZ
                                                          ; if not HL -> buffer on entry
018B E5                     PUSH     H                    ;Save buffer pointer
018C 7E                     MOV      A,M                  ;Get max. count
018D 23                     INX      H                    ;HL -> actual count
018E 34                     INR      M                    ;Increase actual count
018F BE                     CMP      M                    ;Compare max. to actual
0190 E1                     POP      H                    ;Recover buffer pointer
0191 C9                     RET                           ;Z-flag set


                    RCS$GC:                               ;Get character and execute
                                                          ; ECHO, ABORT and FOLD options
0192 D5                     PUSH     D                    ;Save terminator table pointer
0193 E5                     PUSH     H                    ;Save buffer pointer
0194 C5                     PUSH     B                    ;Save option flags

                    RCS$WT:
0195 0E06                   MVI      C,B$DIRCONIO         ;Function code
0197 1EFF                   MVI      E,0FFH               ;Specify input
0199 CD0500                 CALL     BDOS
019C B7                     ORA      A                    ;Check if data waiting
019D CA9501                 JZ       RCS$WT               ;Go back and wait
01A0 C1                     POP      B                    ;Recover option flags
01A1 4F                     MOV      C,A                  ;Save data byte
01A2 3E02                   MVI      A,RCS$ABORT          ;Check if abort option enabled
01A4 A0                     ANA      B
01A5 CAAE01                 JZ       RCS$NA               ;No abort
01A8 3E03                   MVI      A,CTL$C              ;Check for control-C
01AA B9                     CMP      C
01AB CA0000                 JZ       0                    ;Warm boot

                    RCS$NA:
01AE 3E04                   MVI      A,RCS$FOLD           ;Check if folding enabled
01B0 A0                     ANA      B
01B1 C4E501                 CNZ      TOUPPER              ;Convert to uppercase
01B4 3E01                   MVI      A,RCS$ECHO           ;Check if echo required
01B6 A0                     ANA      B
01B7 CAD101                 JZ       RCS$NE               ;No echo required
01BA C5                     PUSH     B                    ;Save options and character
01BB 59                     MOV      E,C                  ;Move character for output
01BC 0E06                   MVI      C,B$DIRCONIO         ;Function code
01BE CD0500                 CALL     BDOS                 ;Echo character
01C1 C1                     POP      B                    ;Recover options and character
01C2 3E0D                   MVI      A,CR                 ;Check if carriage return
01C4 B9                     CMP      C
01C5 C2D101                 JNZ      RCS$NE               ;No
01C8 C5                     PUSH     B                    ;Save options and character
01C9 0E06                   MVI      C,B$DIRCONIO         ;Function code
01CB 1E0A                   MVI      E,LF                 ;Output line feed
01CD CD0500                 CALL     BDOS
01D0 C1                     POP      B                    ;Recover options and character

                    RCS$NE:
01D1 E1                     POP      H                    ;Recover buffer pointer
01D2 D1                     POP      D                    ;Recover terminator table
01D3 C9                     RET                           ;Character in C
```

Figure 5-8. (Continued)

```
                RCS$CT:                         ;Check for terminator
                                                ;C = character just input
                                                ;DE -> 00-byte character
                                                ; string of term. chars.
                                                ;Returns Z status if no
                                                ; match found, NZ if found
                                                ; (with A = C = terminating
                                                ; character)
01D4 D5                 PUSH    D               ;Save table pointer

                RCS$CTL:
01D5 1A                 LDAX    D               ;Get next terminator character
01D6 B7                 ORA     A               ;Check for end of table
01D7 CAE201             JZ      RCS$CTX         ;No terminator matched
01DA B9                 CMP     C               ;Compare to input character
01DB CAE201             JZ      RCS$CTX         ;Terminator matched
01DE 13                 INX     D               ;Move to next terminator
01DF C3D501             JMP     RCS$CTL         ; loop to try next character in table

                RCS$CTX:                        ;Check terminator exit
01E2 B7                 ORA     A               ;At this point, A will either
                                                ; be 0 if the end of the
                                                ; table has been reached, or
                                                ; NZ if a match has been
                                                ; found. The Z-flag will be
                                                ; set.
01E3 D1                 POP     D               ;Recover table pointer
01E4 C9                 RET

                ;TOUPPER - Fold lowercase letters to upper
                ;       C = Character on entry and exit

                TOUPPER:
01E5 3E60               MVI     A,'a'-1         ;Check if folding needed
01E7 B9                 CMP     C               ;Compare to input char.
01E8 D2F501             JNC     TOUPX           ;No, char. is < or =  "a"-1
01EB 3E7A               MVI     A,'z'           ;Maybe, char. is = or > "a"
01ED B9                 CMP     C
01EE DAF501             JC      TOUPX           ;No, char. is > "z"
01F1 3EDF               MVI     A,0DFH          ;Fold character
01F3 A1                 ANA     C
01F4 4F                 MOV     C,A             ;Return folded character

                TOUPX:
01F5 C9                 RET

                ;WCS - Write console string (using raw I/O)
                ;Output terminates when a 00H byte is encountered.
                ;A carriage return is output when a line feed is
                ;encountered.

                ;Calling sequence
                ;       LXI     H,BUFFER
                ;       CALL    WCS

                ;Exit parameters
                ;       HL -> 00H byte terminator

                WCS:
01F6 E5                 PUSH    H               ;Save buffer pointer
01F7 7E                 MOV     A,M             ;Get next character
01F8 B7                 ORA     A               ;Check if 00H
01F9 CA1602             JZ      WCSX            ;Yes, exit
01FC FE0A               CPI     LF              ;Check if line feed
01FE CC0C02             CZ      WCSLF           ;Yes, output a carriage return
0201 5F                 MOV     E,A             ;Character to be output
0202 0E06               MVI     C,B$DIRCONIO    ;Function code
0204 CD0500             CALL    BDOS            ;Output character
0207 E1                 POP     H               ;Recover buffer pointer
0208 23                 INX     H               ;Update to next char.
0209 C3F601             JMP     WCS             ;Output next char.

                WCSLF:                          ;Line feed encountered
020C 0E06               MVI     C,B$DIRCONIO    ;Function code
```

Figure 5-8. (Continued)

```
020E 1E0D               MVI   E,CR          ;Output a CR
0210 CD0500             CALL  BDOS
0213 3E0A               MVI   A,LF          ;Recreate line feed
0215 C9                 RET                 ;Output LF

              WCSX:                         ;Exit
0216 E1                 POP   H             ;Balance the stack
0217 C9                 RET
```

Figure 5-8. (Continued)

Example

```
0007 =         B$GETIO        EQU    7       ;Get IOBYTE
0005 =         BDOS           EQU    5       ;BDOS entry point

0000 0E07             MVI    C,B$GETIO      ;Function code
0002 CD0500           CALL   BDOS           ;A = IOBYTE
```

Purpose This function places the current value of the IOBYTE in register A.

Notes As we saw in Chapter 4, the IOBYTE is a means of associating CP/M's logical devices (console, reader, punch, and list) with the physical devices supported by a particular BIOS. Use of the IOBYTE is completely optional. CP/M, to quote from the Digital Research *CP/M 2.0 Alteration Guide,* "...tolerate[s] the existence of the IOBYTE at location 0003H."

In practice, the STAT utility provided by Digital Research does have some features that set the IOBYTE to different values from the system console.

Figure 5-9 summarizes the IOBYTE structure. A more detailed description was given in Chapter 4.

Each two-bit field can take on one of four values: 00, 01, 10, and 11. The value can be interpreted by the BIOS to mean a specific physical device, as shown in Table 4-1.

Figure 5-10 has equates that are used to refer to the IOBYTE. You can see that the values shown are declared using the SHL (shift left) operator in the Digital Research Assembler. This is just a reminder that the values are structured this way in the IOBYTE itself.

```
                    +-------+-------+-------+-------+
        Bit No.     | 7 : 6 | 5 : 4 | 3 : 2 | 1 : 0 |
                    +-------+-------+-------+-------+
 Logical Device        List   Punch  Reader Console
```

Figure 5-9. The IOBYTE structure

```
                        ;IOBYTE equates
                        ;These are for accessing the IOBYTE.

                        ;Mask values to isolate specific devices.
                        ;(These can also be inverted to preserve all BUT the
                        ; specific device)
0003 =                  IO$CONM EQU     0000$0011B      ;Console mask
000C =                  IO$RDRM EQU     0000$1100B      ;Reader mask
0030 =                  IO$PUNM EQU     0011$0000B      ;Punch mask
00C0 =                  IO$LSTM EQU     1100$0000B      ;List mask

                                                        ;Console values
0000 =                  IO$CTTY EQU     0               ;Console -> TTY:
0001 =                  IO$CCRT EQU     1               ;Console -> CRT:
0002 =                  IO$CBAT EQU     2               ;Console input <- RDR:
                                                        ;Console output -> LST:
0003 =                  IO$CUC1 EQU     3               ;Console -> UC1: (user console 1)

                                                        ;Reader values
0000 =                  IO$RTTY EQU     0 SHL 2         ;Reader <- TTY:
0004 =                  IO$RRDR EQU     1 SHL 2         ;Reader <- RDR:
0008 =                  IO$RUR1 EQU     2 SHL 2         ;Reader <- UR1: (user reader 1)
000C =                  IO$RUR2 EQU     3 SHL 2         ;Reader <- UR2: (user reader 2)

                                                        ;Punch values
0000 =                  IO$PTTY EQU     0 SHL 4         ;Punch -> TTY:
0010 =                  IO$PPUN EQU     1 SHL 4         ;Punch -> PUN:
0020 =                  IO$PUP1 EQU     2 SHL 4         ;Punch -> UP1: (user punch 1)
0030 =                  IO$PUP2 EQU     3 SHL 4         ;Punch -> UP2: (user punch 2)

                                                        ;List values
0000 =                  IO$LTTY EQU     0 SHL 6         ;List -> TTY:
0040 =                  IO$LCRT EQU     1 SHL 6         ;List -> CRT:
0080 =                  IO$LLPT EQU     2 SHL 6         ;List -> LPT: (physical line printer)
00C0 =                  IO$LUL1 EQU     3 SHL 6         ;List -> UL1: (user list 1)
```

Figure 5-10. IOBYTE equates

Function 8: Set IOBYTE

Function Code: C = 08H
Entry Parameters: E = New IOBYTE value
Exit Parameters: None

Example This listing shows you how to assign the logical reader device to the BIOS's console driver. It makes use of some equates from Figure 5-10.

```
0007 =          B$GETIO         EQU     7       ;Get IOBYTE
0008 =          B$SETIO         EQU     8       ;Set IOBYTE
0005 =          BDOS            EQU     5       ;BDOS entry point

000C =          IO$RDRM         EQU     0000$1100B      ;Reader bit mask
0008 =          IO$RUR1         EQU     2 SHL 2         ;User reader select

                ;This example shows how to assign the logical
                ;reader to the user-defined reader #1 (UR1:)

0100                    ORG     100H
0100 0E07               MVI     C,B$GETIO       ;First, get current IOBYTE
```

```
0102 CD0500         CALL    BDOS
0105 E6F3           ANI     (NOT IO$RDRM) AND 0FFH ;Preserve all but
                                            ; reader bits
0107 F608           ORI     IO$RUR1         ;OR in new setting
0109 5F             MOV     E,A             ;Ready for set IOBYTE
010A 0E08           MVI     C,B$SETIO       ;Set new value
010C CD0500         CALL    BDOS
```

Purpose This function sets the IOBYTE to a new value which is given in register E. Because of the individual bit fields in the IOBYTE, you will normally use the Get IOBYTE function, change some bits in the current value, and then call the Set IOBYTE function.

Notes You can use the Set IOBYTE, Get IOBYTE, and Direct Console I/O functions together to create a small program that transforms your computer system into a "smart" terminal. Any data that you type on your keyboard can be sent out of a serial communications line to another computer, and any data received on the line can be sent to the screen.

Figure 5-11 shows this program and illustrates the use of all of these functions.

For this program to function correctly, your BIOS must check the IOBYTE and detect whether the logical console is connected to the physical console (with the IOBYTE set to TTY:) or to the input side of the serial communications line (with the IOBYTE set to RDR:).

Figure 5-11 shows how to use the Get and Set IOBYTE functions to make a simple terminal emulator. For this example to work, the BIOS must detect the Console Value as 3 (IO$CUC1) and connect Console Status, Input, and Output functions to the communications line.

```
0006 =          B$DIRCONIO      EQU     6           ;Direct console input/output
0007 =          B$GETIO         EQU     7           ;Get IOBYTE
0008 =          B$SETIO         EQU     8           ;Set IOBYTE
000B =          B$CONST         EQU     11          ;Get console status (sneak preview)
0005 =          BDOS            EQU     5           ;BDOS entry point

0003 =          IO$CONM EQU     0000$0011B          ;Console mask for IOBYTE
0001 =          IO$CCRT EQU     1                   ;Console -> CRT:
0003 =          IO$CUC1 EQU     3                   ;Console -> user console #1

                TERM:
0000 CD2A00             CALL    SETCRT              ;Connect console -> CRT:

                TERM$CKS:
0003 CD5200             CALL    CONST               ;Get CRT status
0006 CA2400             JZ      TERM$NOKI           ;No console input
0009 CD4B00             CALL    CONIN               ;Get keyboard character
000C CD3000             CALL    SETCOMM             ;Connect console -> comm. line
000F CD4500             CALL    CONOUT              ;Output to comm. line

                TERM$CCS:                           ;Check comm. status
0012 CD5200             CALL    CONST               ;Get "console" status
0015 CA0000             JZ      TERM                ;No incoming comm. character
0018 CD4B00             CALL    CONIN               ;Get incoming comm. character
```

Figure 5-11. Simple terminal emulator

```
001B CD2A00         CALL    SETCRT          ;Connect console -> CRT:
001E CD4500         CALL    CONOUT          ;Output to CRT
0021 C30300         JMP     TERM$CKS        ;Loop back to check keyboard status

            TERM$NOKI:
0024 CD3000         CALL    SETCOMM         ;Connect console -> comm. line
0027 C31200         JMP     TERM$CCS        ;Loop back to check comm. status

            SETCRT:                         ;Connect console -> CRT:
002A F5             PUSH    PSW             ;Save possible data character
002B 0601           MVI     B,IO$CCRT       ;Connect console -> CRT:
002D C33300         JMP     SETCON          ;Common code

            SETCOMM:                        ;Connect console -> comm. line
0030 F5             PUSH    PSW             ;Save possible data character
0031 0603           MVI     B,IO$CUC1       ;Connect console -> comm. line
                                            ;Drop into SETCON

            SETCON:                         ;Set console device
                                            ;New code in B (in bits 1,0)
0033 C5             PUSH    B               ;Save code
0034 0E07           MVI     C,B$GETIO       ;Get current IOBYTE
0036 CD0500         CALL    BDOS
0039 E6FC           ANI     (NOT IO$CONM) AND 0FFH  ;Preserve all but console
003B C1             POP     B               ;Recover required code
003C B0             ORA     B               ;OR in new bits
003D 5F             MOV     E,A             ;Ready for setting
003E 0E08           MVI     C,B$SETIO       ;Function code
0040 CD0500         CALL    BDOS
0043 F1             POP     PSW             ;Recover possible data character
0044 C9             RET

            CONOUT:
0045 5F             MOV     E,A             ;Get data byte for output
0046 0E06           MVI     C,B$DIRCONIO    ;Function code
0048 C30500         JMP     BDOS            ;BDOS returns to CONOUT's caller

            CONIN:
004B 0E06           MVI     C,B$DIRCONIO    ;Function code
004D 1EFF           MVI     E,0FFH          ;Indicate console input
004F C30500         JMP     BDOS            ;BDOS returns to CONIN's caller

            CONST:
0052 0E0B           MVI     C,B$CONST       ;Function code
0054 CD0500         CALL    BDOS
0057 B7             ORA     A               ;Set Z-flag to result
0058 C9             RET
```

Figure 5-11. (Continued)

Function 9: Display "$"-Terminated String

Function Code: C = 09H
Entry Parameters: DE = Address of first byte of string
Exit Parameters: None

Example

```
0009 =        B$PRINTS    EQU    9       ;Print $-Terminated String
0005 =        BDOS        EQU    5       ;BDOS entry point

000D =        CR          EQU    0DH     ;Carriage return
000A =        LF          EQU    0AH     ;Line feed
0009 =        TAB         EQU    09H     ;Horizontal tab
```

```
0000 0D0A09546８MESSAGE:        DB      CR,LF,TAB,'This is a message',CR,LF,'$'

0017 0E09              MVI     C,B$PRINTS         ;Function code
0019 110000            LXI     D,MESSAGE          ;Pointer to message
001C CD0500            CALL    BDOS
```

Purpose This function outputs a string of characters to the console device. The address of this string is in registers DE. You must make sure that the last character of the string is "$"; the BDOS uses this character as a marker for the end of the string. The "$" itself does not get output to the console.

While the BDOS is outputting the string, it expands tabs as previously described, checks to see if there is an incoming character, and checks for CONTROL-S (XOFF, which stops the output until another character is entered) or CONTROL-P (which turns on or off echoing of console characters to the printer).

Notes One of the biggest drawbacks of this function is its use of "$" as a terminating character. As a result, you cannot output a string with a "$" in it. To be truly general-purpose, it would be better to use a subroutine that used an ASCII NUL (00H) character as a terminator, and simply make repetitive calls to the BDOS CONOUT function (code 2). Figure 5-3 is an example of such a subroutine.

Figure 5-12 shows an example of a subroutine that outputs one of several messages. It selects the message based on a message code that you give it as a parameter. Therefore, it is useful for handling error messages; the calling code can pass it an 8-bit error code. You may find it more flexible to convert this subroutine to using 00H-byte-terminated messages using the techniques shown in Figure 5-3.

```
;OM (Output message)
;This subroutine selects one of several messages based on
; the contents of the A register on entry. It then displays
; this message on the console.

;Each message is declared with a "$" as its last character.
; If the A register contains a value larger than the number
; of messages declared, OM will output "Unknown Message".

;As an option, OM can output carriage return / line feed
; prior to outputting the message text.

;Entry parameters
;       HL -> message table
;               This has the form :
;                 DB      3       ;Number of messages in table
;                 DW      MSG0    ;Address of text (A = 0)
;                 DW      MSG1    ;(A = 1)
;                 DW      MSG2    ;(A = 2)
;
;         MSG0:   DB      'Message text$'
;                         ...etc.
;                 A = Message code (from 0 on up)
;                 B = Output CR/LF if non-zero
```

Figure 5-12. Display $-terminated message on console

```
                ;       Calling sequence
                ;               LXI     H,MSG$TABLE
                ;               LDA     MSGCODE
                ;               MVI     B,0     ;Suppress CR/LF
                ;               CALL    OM

0009 =          B$PRINTS        EQU     9       ;Print $-terminated string
0005 =          BDOS            EQU     5       ;BDOS entry point

000D =          CR              EQU     0DH     ;Carriage return
000A =          LF              EQU     0AH     ;Line feed

0000 0D0A24     OM$CRLF:        DB      CR,LF,'$'
0003 556E6B6E6FOM$UM:           DB      'Unknown Message$'

                OM:
0013 F5                 PUSH    PSW             ;Save message code
0014 E5                 PUSH    H               ;Save message table pointer
0015 78                 MOV     A,B             ;Check if CR/LF required
0016 B7                 ORA     A
0017 CA2200             JZ      OM$NOCR         ;No
001A 110000             LXI     D,OM$CRLF       ;Output CR/LF
001D 0E09               MVI     C,B$PRINTS
001F CD0500             CALL    BDOS

                OM$NOCR:
0022 E1                 POP     H               ;Recover message table pointer
0023 F1                 POP     PSW             ;Recover message code
0024 BE                 CMP     M               ;Compare message to max. value
0025 D23700             JNC     OM$ERR          ;Error-code not <= max.
0028 23                 INX     H               ;Bypass max. value in table
0029 87                 ADD     A               ;Message code * 2
002A 5F                 MOV     E,A             ;Make (code * 2) a word value
002B 1600               MVI     D,0
002D 19                 DAD     D               ;HL -> address of message text
002E 5E                 MOV     E,M             ;Get LS byte
002F 23                 INX     H               ;HL -> MS byte
0030 56                 MOV     D,M             ;Get MS byte
                                                ;DE -> message text itself

                OM$PS:                          ;Print string entry point
0031 0E09               MVI     C,B$PRINTS      ;Function code
0033 CD0500             CALL    BDOS
0036 C9                 RET                     ;Return to caller

                OM$ERR:                         ;Error
0037 110300             LXI     D,OM$UM         ;Point to "Unknown Message"
003A C33100             JMP     OM$PS           ;Print string
```

Figure 5-12. (Continued)

Function 10: Read Console String

Function Code: C = 0AH
Entry Parameters: DE = Address of string buffer
Exit Parameters: String buffer with console bytes in it

Example

```
000A =          B$READCONS      EQU     10      ;Read Console String
0005 =          BDOS            EQU     5       ;BDOS entry point
```

```
0050 =          BUFLEN          EQU     80      ;Buffer length

                BUFFER:                         ;Console input buffer
0000 50         BUFMAXCH:       DB      BUFLEN  ;Max. no. of characters in
                                                ; buffer
0001 00         BUFACTCH:       DB      0       ;Actual no. of characters input
0002            BUFCH:          DS      BUFLEN  ;Buffer characters

0052 0E0A               MVI     C,B$READCONS    ;Function code
0054 110000             LXI     D,BUFFER        ;Pointer to buffer
0057 CD0500             CALL    BDOS
```

Purpose This function reads a string of characters from the console device and stores them in a buffer (address in DE) that you define. Full line editing is possible: the operator can backspace, cancel the line and start over, and use all the normal control functions. What you will ultimately see in the buffer is the final version of the character string entered, without any of the errors or control characters used to do the line editing.

The buffer that you define has a special format. The first byte in the buffer tells the BDOS the maximum number of characters to be accepted. The second byte is reserved for the BDOS to tell you how many characters were actually placed in the buffer. The following bytes contain the characters of the string.

Character input will cease either when a CARRIAGE RETURN is entered or when the maximum number of characters, as specified in the buffer, has been received. The CARRIAGE RETURN is not stored in the buffer as a character—it just serves as a terminator.

If the first character entered is a CARRIAGE RETURN, then the BDOS sets the "characters input" byte to 0. If you attempt to input more than the maximum number of characters, the "characters input" count will be the same as the maximum value allowed.

Notes This function is useful for accepting console input, especially because of the line editing that it allows. It should be used even for single-character responses, such as "Y/N" (yes or no), because the operator can type "Y", backspace, and overtype with "N". This makes for more "forgiving" programs, tolerant of humans who change their minds.

Figure 5-13 shows an example subroutine that uses this function. It accepts console input, matches the input against a table, and transfers control to the appropriate subroutine. Many interactive programs need to do this; they accept an operator command and then transfer control to the appropriate command processor to deal with that command.

This example also includes two other subroutines that are useful in their own right. One compares null-byte-terminated strings (FSCMP), and the other converts, or "folds," lowercase letters to uppercase (FOLD).

```
                        ;RSA
                        ;Return subprocessor address
                        ;This subroutine returns one of several addresses selected
                        ; from a table by matching keyboard input against specified
                        ; strings. It is normally used to switch control to a
                        ; particular subprocessor according to an option entered
                        ; by the operator from the keyboard.
                        ;
                        ;Character string comparisons are performed with case-folding;
                        ; that is, lowercase letters are converted to uppercase.
                        ;
                        ;If the operator input fails to match any of the specified
                        ; strings, then the carry flag is set. Otherwise, it is
                        ; cleared.

                        ;Entry parameters
                        ;       HL -> Subprocessor select table
                        ;               This has the form :
                        ;               DW      TEXT0,SUBPROC0
                        ;               DW      TEXT1,SUBPROC1
                        ;               DW      0       ;Terminator
                        ;       TEXT0:  DB      'add',0 ;00H-byte terminated
                        ;       TEXT1:  DB      'subtract',0
                        ;       SUBPROC0:
                        ;               Code for processing ADD function.
                        ;       SUBPROC1:
                        ;               Code for processing SUBTRACT function.

                        ;Exit parameters
                        ;       DE -> operator input string (00H-terminated
                        ;             input string).
                        ;       Carry Clear, HL -> subprocessor.
                        ;       Carry Set, HL = 0000H.

                        ;Calling sequence
                        ;       LXI     H,SUBPROCTAB    ;Subprocessor table
                        ;       CALL    RSA
                        ;       JC      ERROR           ;Carry set only on error
                        ;       LXI     D,RETURN        ;Fake CALL instruction
                        ;       PUSH    D               ;Push return address on stack
                        ;       PCHL                    ;"CALL" to subprocessor
                        ;       RETURN:

000A =                  B$READCONS      EQU     10      ;Read console string into buffer
0005 =                  BDOS            EQU     5       ;BDOS entry point

0050 =                  RSA$BL          EQU     80      ;Buffer length
0000 50                 RSA$BUF:        DB      RSA$BL  ;Max. no. of characters
0001 00                 RSA$ACTC:       DB      0       ;Actual no. of characters
0002                    RSA$BUFC:       DS      RSA$BL  ;Buffer characters
0052 00                                 DB      0       ;Safety terminator

                        RSA:
0053 2B                         DCX     H               ;Adjust Subprocessor pointer
0054 2B                         DCX     H               ; for code below
0055 E5                         PUSH    H               ;Top of stack (TOS) -> subproc. table - 2
0056 0E0A                       MVI     C,B$READCONS    ;Function code
0058 110000                     LXI     D,RSA$BUF       ;DE -> buffer
005B CD0500                     CALL    BDOS            ;Read operator input and
                                                        ; Convert to 00H-terminated
005E 210100                     LXI     H,RSA$ACTC      ;HL -> actual no. of chars. input
0061 5E                         MOV     E,M             ;Get actual no. of chars. input
0062 1600                       MVI     D,0             ;Make into word value
0064 23                         INX     H               ;HL -> first data character
0065 19                         DAD     D               ;HL -> first UNUSED character in buffer
0066 3600                       MVI     M,0             ;Make input buffer 00H terminated

                        RSA$ML:                         ;Compare input to specified values
                                                        ; Main loop
0068 E1                         POP     H               ;Recover subprocessor table pointer
0069 23                         INX     H               ;Move to top of next entry
006A 23                         INX     H               ;HL -> text address
006B 5E                         MOV     E,M             ;Get text address
```

Figure 5-13. Read console string for keyboard options

```
006C 23                 INX     H
006D 56                 MOV     D,M             ;DE -> text

006E 7A                 MOV     A,D             ;Check if at end of subprocessor table
006F B3                 ORA     E
0070 CA8500             JZ      RSA$NFND        ;Match not found

0073 23                 INX     H               ;HL -> subprocessor address
0074 E5                 PUSH    H               ;Save ptr. to subprocessor table
0075 210200             LXI     H,RSA$BUFC      ;HL -> input characters
0078 CD8A00             CALL    FSCMP           ;Folded string compare
007B C26800             JNZ     RSA$ML          ;No match, move to next entry
007E E1                 POP     H               ;Match found, recover subprocessor ptr.
007F 5E                 MOV     E,M             ;Get actual subprocessor address
0080 23                 INX     H
0081 56                 MOV     D,M             ;DE -> Subprocessor code
0082 EB                 XCHG                    ;HL -> Subprocessor code
0083 B7                 ORA     A               ;Clear carry (match found)
0084 C9                 RET

                RSA$NFND:
0085 210000             LXI     H,0             ;Indicate no match found
0088 37                 STC                     ;Set carry
0089 C9                 RET

                ;FSCMP
                ;Compare folded (lowercase to upper) string.
                ;This subroutine compares two 00H-byte terminated
                ;strings and returns with the condition flags set
                ;to indicate their relationship.

                ;Entry parameters
                ;       DE -> string 1
                ;       HL -> string 2

                ;Exit parameters
                ;       Flags set (based on string 1 - string 2, on a
                ;       character-by-character basis)

                FSCMP:
008A 1A                 LDAX    D               ;Get string 1 character
008B CD9E00             CALL    FOLD            ;Fold to uppercase
008E F5                 PUSH    PSW             ;Save string 1 character
008F 7E                 MOV     A,M             ;Get string 2 character
0090 CD9E00             CALL    FOLD            ;Fold to uppercase
0093 47                 MOV     B,A             ;Save string 2 character
0094 F1                 POP     PSW             ;Recover string 1 character
0095 B8                 CMP     B               ;String 1 - string 2
0096 C0                 RNZ                     ;Return if not equal
0097 B7                 ORA     A               ;Equal, so check if end of strings
0098 C8                 RZ                      ;Yes
0099 13                 INX     D               ;No, update string 1 pointer
009A 23                 INX     H               ; and string 2 pointer
009B C38A00             JMP     FSCMP           ;Check next character

                ;FOLD
                ;Folds a lowercase letter (a-z) to uppercase (A-Z)
                ;The character to be folded is in A on entry and on exit.

                FOLD:
009E 4F                 MOV     C,A             ;Preserve input character
009F 3E60               MVI     A,'a'-1         ;Check if folding needed
00A1 B9                 CMP     C               ;Compare to input character
00A2 D2AF00             JNC     FOLDX           ;No, char. is <= "a"
00A5 3E7A               MVI     A,'z'           ;Check if < "z"
00A7 B9                 CMP     C
00A8 DAAF00             JC      FOLDX           ;No, char. is > "z"
00AB 3EDF               MVI     A,0DFH          ;Fold character
00AD A1                 ANA     C
00AE C9                 RET

                FOLDX:
00AF 79                 MOV     A,C             ;Recover original input char.
00B0 C9                 RET
```

Figure 5-13. (Continued)

Function 11: Read Console Status

Function Code: C = 0BH
Entry Parameters: None
Exit Parameters: A = 00H if no incoming data byte
A = 0FFH if incoming data byte

Example

```
000B =          B$CONST         EQU     11          ;Get Console Status
0005 =          BDOS            EQU     5           ;BDOS entry point

0000 0E0B               MVI     C,B$CONST           ;Function code
0002 CD0500             CALL    BDOS                ;A = 00 if no character waiting
                                                    ;A = 0FFH if character waiting
```

Purpose This function tells you whether a console input character is waiting to be processed. Unlike the Console Input functions, which will wait until there is input, this function simply checks and returns immediately.

Notes Use this function wherever you want to interrupt an executing program if a console keyboard character is entered. Just put a Console Status call in the main loop of the program. Then, if the program detects that keyboard data is waiting, it can take the appropriate action. Normally this would be to jump to location 0000H, thereby aborting the current program and initiating a warm boot.

Figure 5-11 is an example subroutine that shows how to use this function.

Function 12: Get CP/M Number

Function Code: C = 0CH
Entry Parameters: None
Exit Parameters: HL = Version number code

Example

```
000C =          B$GETVER        EQU     12          ;Get CP/M Version Number
0005 =          BDOS            EQU     5           ;BDOS entry point

0000 0E0C               MVI     C,B$GETVER          ;Function code
0002 CD0500             CALL    BDOS                ;H = 00 for CP/M
                                                    ;L = version (e.g. 22H for 2.2)
```

Purpose This function tells you which version of CP/M you are currently running. A two-byte value is returned:

H = 00H for CP/M, H = 01H for MP/M

L = 00H for all releases before CP/M 2.0

L = 20H for CP/M 2.0, 21H for 2.1, 22H for 2.2, and so on for any subsequent releases.

This information is of interest only if your program has some version-specific logic built into it. For example, CP/M version 1.4 does not support the same Random File Input/Output operations that CP/M 2.2 does. Therefore, if your program uses Random I/O, put this check at the beginning to ensure that it is indeed running under the appropriate version of CP/M.

Notes Figure 5-14 is a subroutine that checks the current CP/M version number, and, if it is not CP/M 2.2, displays an explanatory message on the console and does a warm boot by jumping to location 0000H.

Function 13: Reset Disk System

Function Code: C = 0DH
Entry Parameters: None
Exit Parameters: None

```
                    ;CCPM
                    ;Check if CP/M
                    ;This subroutine determines the version number of the
                    ;operating system and, if not CP/M version 2, displays
                    ;an error message and executes a warm boot.

                    ;Entry and exit parameters
                    ;       None

                    ;Calling sequence
                    ;       CALL    CCPM    ;Warm boots if not CP/M 2

0009 =              B$PRINTS        EQU     9       ;Display $-terminated string
000C =              B$GETVER        EQU     12      ;Get version number
0005 =              BDOS            EQU     5       ;BDOS entry point

000D =              CR              EQU     0DH     ;Carriage return
000A =              LF              EQU     0AH     ;Line feed

0000 0D0A           CCPMM:  DB      CR,LF
0002 5468697320             DB      'This program can only run under CP/M version 2.'
0031 0D0A24                 DB      CR,LF,'$'

                    CCPM:
0034 0E0C                   MVI     C,B$GETVER      ;Get version number
0036 CD0500                 CALL    BDOS
0039 7C                     MOV     A,H             ;H must be 0 for CP/M
003A B7                     ORA     A
003B C24700                 JNZ     CCPME           ;Must be MP/M
003E 7D                     MOV     A,L             ;L = version number of CP/M
003F E6F0                   ANI     0F0H            ;Version number in MS nibble
0041 FE20                   CPI     20H             ;Check if version 2
0043 C24700                 JNZ     CCPME           ;Must be an earlier version
0046 C9                     RET                     ;Yes, CP/M version 2

                    CCPME:                          ;Error
0047 0E09                   MVI     C,B$PRINTS      ;Display error message
0049 110000                 LXI     D,CCPMM
004C CD0500                 CALL    BDOS
004F C30000                 JMP     0               ;Warm boot
```

Figure 5-14. Determine the CP/M version number

Example

```
000D =          B$DSKRESET      EQU     13              ;Reset Disk System
0005 =          BDOS            EQU     5               ;BDOS entry point

0000 0E0D               MVI     C,B$DSKRESET            ;Function code
0002 CD0500             CALL    BDOS
```

Purpose This function requests CP/M to completely reset the disk file system. CP/M then resets its internal tables, selects logical disk A as the default disk, resets the DMA address back to 0080H (the address of the buffer used by the BDOS to read and write to the disk), and marks all logical disks as having Read/Write status.

The BDOS will then have to log in each logical disk as each disk is accessed. This involves reading the entire file directory for the disk and rebuilding the allocation vectors (which keep track of which allocation blocks are free and which are used for file storage).

Notes This function lets you change the diskettes under program control. If the operator were to simply change diskettes, without CP/M knowing about it, the next access to the (now different) diskette would force CP/M to declare the disk Read-Only, thwarting any further attempts to write on the diskette. If you need to reset one or two disks, rather than the entire disk system, look ahead to the Reset Disk function (code 37) described at the end of this chapter.

Figure 5-15 shows a simple subroutine that outputs a message on the console, requesting that the diskette in a specified drive be changed. It then issues a Reset Disk function call to make sure that CP/M will log in the diskette on the next access to the drive.

```
                ;CDISK
                ;Change disk
                ;This subroutine displays a message requesting the
                ;user to change the specified logical disk, then waits
                ;for a carriage return to be pressed. It then issues
                ;a Disk Reset and returns to the caller.

                ;Entry parameters
                ;       A = Logical disk to be changed (A = 0, B = 1)

                ;Exit parameters
                ;       None

                ;Calling sequence
                ;       MVI     A,0             ;Change drive A:
                ;       CALL    CDISK

000D =          B$DSKRESET      EQU     13      ;Disk Reset function code
0009 =          B$PRINTS        EQU     9       ;Print $-terminated string
0001 =          B$CONIN         EQU     1       ;Get console input
0005 =          BDOS            EQU     5       ;BDOS entry point
```

Figure 5-15. Reset requested disk drive

```
000D =          CR              EQU     0DH
000A =          LF              EQU     0AH

0000 0D0A436861CDISKM:          DB      CR,LF,'Change logical disk '
0016 00         CDISKD:         DB      0
0017 3A20616E64                 DB      ': and press Carriage Return to continue$'

                CDISK:
003F C640               ADI     'A'-1           ;Convert to letter
0041 321600             STA     CDISKD          ;Store in message
0044 0E09               MVI     C,B$PRINTS      ;Display message
0046 110000             LXI     D,CDISKM
0049 CD0500             CALL    BDOS

                CDISKW:
004C 0E01               MVI     C,B$CONIN       ;Get keyboard character
004E CD0500             CALL    BDOS
0051 FE0D               CPI     CR
0053 C24C00             JNZ     CDISKW
0056 0E0D               MVI     C,B$DSKRESET    ;Now reset disk system
0058 CD0500             CALL    BDOS
005B C9                 RET
```

Figure 5-15. Reset requested disk drive (continued)

Function 14: Select Logical Disk

Function Code: C = 0EH
Entry Parameters: E = Logical Disk Code
00H = Drive A
01H = Drive B and so on
Exit Parameters: None

Example

```
000E =          B$SELDSK        EQU     14      ;Select Logical Disk
0005 =          BDOS            EQU     5       ;BDOS entry point

0000 0E0E               MVI     C,B$SELDSK      ;Function code
0002 1E00               MVI     E,0             ;E = 0 for A:, 1 for B: etc.
0004 CD0500             CALL    BDOS
```

Purpose This function makes the logical disk named in register E the default disk. All subsequent references to disk files that do not specify the disk will use this default.

When you reference a disk file that *does* have an explicit logical disk in its name you do not have to issue another Select Disk function; the BDOS will take care of that for you.

Notes Notice the way in which the logical disk is specified in register E. It is not the same as the disk drive specification in the first byte of the file control block. In the FCB, a value of 00H is used to mean "use the current default disk" (as specified in the last Select Disk call or by the operator on the console). With this function, a value of 00H in register A means that A is the selected drive, a value of 01H means drive B, and so on to 0FH for drive P, allowing 16 drives in the system.

If you select a logical disk that does not exist in your computer system, the BDOS will display the following message:

```
BDOS Err on J: Select
```

If you type a CARRIAGE RETURN in order to proceed, the BDOS will do a warm boot and transfer control back to the CCP. To avoid this, you must rely on the computer operator not to specify nonexistent disks or build into your program the knowledge of how many logical disk drives are on the system.

Another problem with this function is that you cannot distinguish a logical disk for which the appropriate tables have been built into the BIOS, but for which there is no physical disk drive. The BDOS does not check to see if the drive is physically present when you make the Select Disk call. It merely sets up some internal values ready to access the logical disk. If you then attempt to access this nonexistent drive, the BIOS will detect the error. What happens next is completely up to the BIOS. The standard BIOS will return control to the BDOS, indicating an error condition. The BDOS will output the message

```
BDOS Err on C: Bad Sector
```

You then have a choice. You can press CARRIAGE RETURN, in which case the BDOS will ignore the error and attempt to continue with whatever appears to have been read in. Or you can enter a CONTROL-C, causing the program to abort and CP/M to perform a warm boot.

Note that the Select Disk function does not return any values. If your program gets control back, you can assume that the logical disk you asked for at least has tables declared for it.

Function 15: Open File

Function Code: C = 0FH
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
000F =          B$OPEN          EQU     15      ;Open File
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;File name
0009 545950     FCB$TYP:        DB      'TYP'   ;File type
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
```

```
0021 0000    FCB$RANRECz:     DW      0          ;Random rec. to read/write
0023 00      FCB$RANRECO:     DB      0          ;Random rec. overflow byte (MS)

0024 0E0F             MVI     C,B$OPEN          ;Function code
0026 110000           LXI     D,FCB             ;DE -> File control block
0029 CD0500           CALL    BDOS              ;A = 0FFH if file not found
```

Purpose This function opens a specified file for reading or writing. The FCB, whose address must be in register DE, tells CP/M the user number, the logical disk, the file name, and the file type. All other bytes of the FCB will normally be set to 0.

The code returned by the BDOS in register A indicates whether the file has been opened successfully. If A contains 0FFH, then the BDOS was unable to find the correct entry in the directory. If A = 0, 1, 2, or 3, then the file has been opened.

Notes The Open File function searches the entire file directory on the specified logical disk looking for the file name, type, and extent specified in the FCB; that is, it is looking for an exact match for bytes 1 through 14 of the FCB. The file name and type may be ambiguous; that is, they may contain "?" characters. In this case, the BDOS will open the first file in the directory that matches the ambiguous name in the FCB. If the file name or type is shorter than eight or three characters respectively, then the remaining characters must be filled with blanks.

When the BDOS searches the file directory, it expects to find an *exact* match with each character of the file name and type, including lowercase letters or nongraphic characters. However, the BDOS uses only the least significant seven bits of each character—the most significant bit is used to indicate special file status characteristics, or *attributes.*

By matching the file extent as well as the name and type, you can, if you wish, open the file at some point other than its beginning. For normal sequential access, you would not usually want to do this, but if your program can predict which file extent is required, this is a method of moving directly to it.

It is also possible to open the same file more than once. Each instance requires a separate FCB. The BDOS is not aware that this is happening. It is really only safe to do this when you are reading the file. Each FCB can be used to read the file independently.

Once the file has been found in the directory, the number of records and the allocation blocks used are copied from the directory entry into the FCB (bytes 16 through 31). If the file is to be accessed sequentially from the beginning of the file, the current record (byte 32) must be set to zero by your program.

The value returned in register A is the relative directory entry number of the entry that matched the FCB. As previously explained, the buffer that CP/M uses holds a 128-byte record from the directory with four directory entries numbered 0, 1, 2, and 3. This *directory code* is returned by almost all of the file-related BDOS functions, but under normal circumstances you will be concerned only with whether the value returned in A is 0FFH or not.

Figure 5-16 shows a subroutine that takes a 00H-byte terminated character string, creates a valid FCB, and then opens the specified file. Shown as part of this example is the subroutine BF (Build FCB). It performs the brunt of the work of converting a string of ASCII characters into an FCB-style disk, file name, and type.

```
                        ;OPENF
                        ;Open File

                        ;Given a pointer to a 00H-byte-terminated file name,
                        ;and an area that can be used for a file control
                        ;block, this subroutine builds a valid file control
                        ;block and attempts to open the file.

                        ;If the file is opened, it returns with the carry flag clear.
                        ;If the file cannot be opened, this subroutine returns
                        ;with the carry flag set.

                        ;Entry parameters
                        ;       DE -> 36-byte area for file control block
                        ;       HL -> 00H-byte terminated file name of the
                        ;               form {disk:} Name {.typ}
                        ;               (disk and typ are optional)

                        ;Exit parameters
                        ;       Carry clear : File opened correctly.
                        ;       Carry set   : File not opened.

                        ;Calling Sequence
                        ;       LXI     D,FCB
                        ;       LXI     H,FNAME
                        ;       CALL    OPENF
                        ;       JC      ERROR
                        ;where
                        ;FCB:   DS      36              ;Space for file control block
                        ;FNAME: DB      'A:TESTFILE.DAT',0

000F =                  B$OPEN          EQU     15      ;File Open function code
0005 =                  BDOS            EQU     5       ;BDOS entry point

                        OPENF:
0000 D5                         PUSH    D               ;Preserve pointer to FCB
0001 CD0C00                     CALL    BF              ;Build file control block
0004 0E0F                       MVI     C,B$OPEN
0006 D1                         POP     D               ;Recover pointer to FCB
0007 CD0500                     CALL    BDOS
000A 17                         RAL                     ;If A=0FFH, carry set
                                                        ;otherwise carry clear
000B C9                         RET

                        ;BF
                        ;Build file control block
                        ;This subroutine formats a 00H-byte-terminated string
                        ;(presumed to be a file name) into an FCB, setting
                        ;the disk and file name and type and clearing the
                        ;remainder of the FCB to 0's.

                        ;Entry parameters
                        ;       DE -> file control block (36 Bytes)
                        ;       HL -> file name string (00H-byte-terminated)

                        ;Exit parameters
                        ;       The built file control block
                        ;Calling sequence
                        ;       LXI     D,FCB
                        ;       LXI     H,FILENAME
                        ;       CALL    BF

                        BF:
```

Figure 5-16. Open file request

```
000C 23              INX     H              ;Check if 2nd char. is ":"
000D 7E              MOV     A,M            ;Get character from file name
000E 2B              DCX     H              ;HL -> now back at 1st char.
000F FE3A            CPI     ':'            ;If ":", then disk specified
0011 C21C00          JNZ     BF$ND          ;No disk
0014 7E              MOV     A,M            ;Get disk letter
0015 E61F            ANI     0001$1111B     ;A (41H) -> 1, B (42H) -> 2 ...
0017 23              INX     H              ;Bypass disk letter
0018 23              INX     H              ;Bypass ":"
0019 C31D00          JMP     BF$SD          ;Store disk in FCB

                BF$ND:                      ;No disk present
001C AF              XRA     A              ;Indicate default disk

                BF$SD:
001D 12              STAX    D              ;Store disk in FCB
001E 13              INX     D              ;DE -> 1st char. of name in FCB
001F 0E08            MVI     C,8            ;File name length
0021 CD3700          CALL    BF$GT          ;Get token
                                            ;Note -- at this point, BF$GT
                                            ;will have advanced the string
                                            ;pointer to either a "." or
                                            ;00H byte
0024 FE2E            CPI     '.'            ;Check terminating character
0026 C22A00          JNZ     BF$NT          ;No file type specified
0029 23              INX     H              ;Bypass "." in file name

                BF$NT:
002A 0E03            MVI     C,3            ;File type length
002C CD3700          CALL    BF$GT          ;Get token
                                            ;Note -- if no file type is
                                            ;present BF$GT will merely
                                            ;spacefill the FCB
002F 0600            MVI     B,0            ;0-fill the remainder of the FCB
0031 0E18            MVI     C,24           ;36 - 12 (disk, name, type = 12 chars.)
0033 CD6400          CALL    BF$FT          ;Re-use fill token S/R
0036 C9              RET

                ;BF$GT
                ;Build FCB -- get token

                ;This subroutine scans a file name string,
                ;placing characters into a file control block.
                ;On encountering a terminator character ("." or 00H),
                ;the remainder of the token is space filled.
                ;If an "*" is encountered, the remainder of the token
                ;is filled with "?".

                ;Entry parameters
                ;       DE -> Into file control block
                ;       HL -> Into file name string
                ;       C = Maximum no. of characters in token


                ;Exit parameters
                ;       File control block contains next token
                ;       A = Terminating character

                BF$GT:
0037 7E              MOV     A,M            ;Get next string character
0038 B7              ORA     A              ;Check if end of string
0039 CA5700          JZ      BF$SFT         ;Yes, space fill token
003C FE2A            CPI     '*'            ;Check if ?-fill required
003E CA5C00          JZ      BF$QFT         ;Yes, fill with ?
0041 FE2E            CPI     '.'            ;Assume current token is file
                                            ;name
                                            ;Check if file type coming up
                                            ;(If current token is file
                                            ;type this check is
                                            ;benignly redundant)
0043 CA5700          JZ      BF$SFT         ;Yes, space fill token
0046 12              STAX    D              ;None of the above, so store
                                            ;in FCB
0047 13              INX     D              ;Update FCB pointer
0048 23              INX     H              ;Update string pointer
```

Figure 5-16. (Continued)

```
0049 0D                 DCR     C               ;Countdown on token length
004A C23700             JNZ     BF$GT           ;Still more characters to go

               BF$SKIP:                         ;Skip chars. until "." or 00H
004D 7E                 MOV     A,M             ;Get next string character
004E B7                 ORA     A               ;Check if 00H
004F C8                 RZ                      ;Yes
0050 FE2E               CPI     '.'             ;Check if "."
0052 C8                 RZ                      ;Yes
0053 23                 INX     H               ;Update string pointer (only)
0054 C34D00             JMP     BF$SKIP         ;Try next character

               BF$SFT:                          ;Space fill token
0057 0620               MVI     B,' '
0059 C36400             JMP     BF$FT           ;Common fill token code
                                                ;BF$FT returns to caller

               BF$QFT:                          ;Question mark fill token
005C 063F               MVI     B,'?'
005E CD6400             CALL    BF$FT           ;Common fill token code
0061 C34D00             JMP     BF$SKIP         ;Bypass multiple "*" etc.

               BF$FT:                           ;Fill token
0064 F5                 PUSH    PSW             ;Save terminating character
0065 78                 MOV     A,B             ;Get fill characer

               BF$FTL:                          ;Inner loop
0066 12                 STAX    D               ;Store in FCB
0067 13                 INX     D               ;Update FCB Pointer
0068 0D                 DCR     C               ;Downdate residual count
0069 C26600             JNZ     BF$FTL          ;Keep going
006C F1                 POP     PSW             ;Recover terminating character
006D C9                 RET
```

Figure 5-16. (Continued)

Function 16: Close File

Function Code: C = 10H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
0010 =          B$CLOSE         EQU     16      ;Close File
0005 =          BDOS            EQU     5       ;BDOS entry point

0000            FCB:            DS      36      ;File control block

0024 0E10               MVI     C,B$CLOSE       ;Function code
0026 110000             LXI     D,FCB           ;DE -> File control block
0029 CD0500             CALL    BDOS            ;A = 0,1,2,3 if successful
                                                ;A = 0FFH if file name not
                                                ; in directory
```

Purpose This function terminates the processing of a file to which you have written information. Under CP/M you do not need to close a file that you have been reading. However, if you ever intend for your program to function correctly under MP/M (the multi-user version of CP/M) you should close all files regardless of their use.

The Close File function, like Open File, returns a directory code in the A register. Register A will contain 0FFH if the BDOS could not close the file successfully. If A is 0, 1, 2, or 3, then the file has been closed.

Notes When the BDOS closes a file to which data has been written, it writes the current contents of the FCB out to the disk directory, updating an existing directory entry by matching the disk, name, type, and extent number in the same manner that the Open File function does.

Note that the BDOS does not transfer the last record of the file to the disk during the close operation. It merely updates the file directory. You must arrange to flush any partly filled record to the disk. If the file that you have created is a standard CP/M ASCII text file, you must arrange to fill the unused portion of the record with the standard 1AH end-of-file characters as CP/M expects, as explained in the section on the Write Sequential function (code 21).

Function 17: Search for First Name Match

Function Code: C = 11H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
0011 =          B$SEARCHF       EQU     17      ;Search First
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C453FFCB$NAME:        DB      'FILE????'      ;Ambiguous file name
0009 543F50     FCB$TYP:        DB      'T?P'   ;Ambiguous file type
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0       ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0       ;Random rec. overflow byte (MS)

0024 0E11               MVI     C,B$SEARCHF     ;Function code
0026 110000             LXI     D,FCB           ;DE -> File control block
0029 CD0500             CALL    BDOS            ;A = 0,1,2,3.
                                                ;(A * 32) + DMA -> directory
                                                ; entry
                                                ;A = 0FFH if file name not
                                                ; found
```

Purpose This function scans down the file directory for the first entry that matches the file name, type, and extent in the FCB addressed by DE. The file name, type, and extent may contain a "?" (ASCII 3FH) in one or more character positions. Where a "?" occurs, the BDOS will match *any* character in the corresponding position in the file directory. This is known as ambiguous file name matching.

The first byte of an FCB normally contains the logical disk number code. A value of 0 indicates the default disk, while 1 means disk A, 2 is B, and so on up to a possible maximum of 16 for disk P. However, if this byte contains a "?", the BDOS will search the default logical disk and will match the file name and type regardless of the user number. This function is normally used in conjunction with the Search Next function (which is described immediately after this function). Search First, in the process of matching a file, leaves certain variables in the BDOS set, ready for a subsequent Search Next.

Both Search First and Search Next return a directory code in the A register. With Search First, A = 0FFH when no files match the FCB; if a file match is found, A will have a value of 0, 1, 2, or 3.

Notes To locate the particular directory entry that either the Search First or Search Next function matched, multiply the directory code returned in A by the length of a directory entry (32 bytes). This is easily done by adding the A register to itself five times (see the code in Figure 5-17 near the label GNFC). Then add the DMA address to get the actual address where the matched directory entry is stored.

There are many occasions when you may need to write a program that will accept an ambiguous file name and operate on all of the file names that match it. (The DIR and ERA commands built into the CCP are examples that use ambiguous file names.) To do this, you must use several BDOS functions: the Set DMA Address function (code 26, described later in this chapter), this function (Search First), and Search Next (code 18). All of this is shown in the subroutine given in Figure 5-17.

```
;GNF
;This subroutine returns an FCB setup with either the
;first file matched by an ambiguous file name, or (if
;specified by entry parameter) the next file name.

;Note : this subroutine is context sensitive. You must
;       not have more than one ambiguous file name
;       sequence in process at any given time.

;>>>    Warning : This subroutine changes the DMA address
;>>>              inside the BDOS.

;Entry parameters
;       DE -> Possibly ambiguous file name
;               (00-byte terminated)
;               (Only needed for FIRST request)
;       HL -> File control block
;       A = 0  : Return FIRST file name that matches
;         = NZ : Return NEXT file name that matches

;Exit parameters
;Carry set : A = FF, no file name matches
;           A not = OFFH, error in input file name
;Carry clear : FCB setup with next name
;           HL -> Directory entry returned
;               by Search First/Next

;Calling sequence
;       LXI     D,FILENAME
;       LXI     H,FCB
```

Figure 5-17. Search first/next calls for ambiguous file name

```
                        ;       MVI     A,0     ;or MVI A,1 for NEXT
                        ;       CALL    GNF

0011 =                  B$SEARCHF       EQU     17      ;Search for first file name
0012 =                  B$SEARCHN       EQU     18      ;Search for next file name
001A =                  B$SETDMA        EQU     26      ;Set up DMA address
0005 =                  BDOS            EQU     5       ;BDOS entry point

0080 =                  GNFDMA  EQU     80H             ;Default DMA address
000D =                  GNFSVL  EQU     13              ;Save length (no. of chars to move)
0024 =                  GNFFCL  EQU     36              ;File control block length
0000                    GNFSV:  DS      GNFSVL          ;Save area for file name/type

                        GNF:
000D E5                         PUSH    H               ;Save FCB pointer
000E D5                         PUSH    D               ;Save file name pointer
000F F5                         PUSH    PSW             ;Save first/next flag

0010 118000                     LXI     D,GNFDMA        ;Set DMA to known address
0013 0E1A                       MVI     C,B$SETDMA      ;Function code
0015 CD0500                     CALL    BDOS
0018 F1                         POP     PSW             ;Recover first/next flag
0019 E1                         POP     H               ;Recover file name pointer
001A D1                         POP     D               ;Recover FCB pointer
001B D5                         PUSH    D               ;Resave FCB pointer

001C B7                         ORA     A               ;Check if FIRST or NEXT
001D C23E00                     JNZ     GNFN            ;NEXT
0020 CD9300                     CALL    BF              ;Build file control block
0023 E1                         POP     H               ;Recover FCB pointer (to balance stack)
0024 D8                         RC                      ;Return if error in file name
0025 E5                         PUSH    H               ;Resave FCB pointer

                                                        ;Move ambiguous file name to
                                                        ;save area
                                                        ;HL -> FCB
0026 110000                     LXI     D,GNFSV         ;DE -> save area
0029 0E0D                       MVI     C,GNFSVL        ;Get save length
002B CD8A00                     CALL    MOVE
002E D1                         POP     D               ;Recover FCB pointer
002F D5                         PUSH    D               ;and resave

0030 0E11                       MVI     C,B$SEARCHF     ;Search FIRST
0032 CD0500                     CALL    BDOS
0035 E1                         POP     H               ;Recover FCB pointer
0036 FEFF                       CPI     0FFH            ;Check for error
0038 CA7D00                     JZ      GNFEX           ;Error exit
003B C35D00                     JMP     GNFC            ;Common code

                        GNFN:                           ;Execute search FIRST to re-
                                                        ;establish contact with
                                                        ;previous file
                                                        ;User's FCB still has
                                                        ;name/type in it
003E CD7F00                     CALL    GNFZF           ;Zero-fill all but file name/type
0041 D1                         POP     D               ;Recover FCB address
0042 D5                         PUSH    D               ;and resave
0043 0E11                       MVI     C,B$SEARCHF     ;Re-find the file
0045 CD0500                     CALL    BDOS
0048 D1                         POP     D               ;Recover FCB pointer
0049 D5                         PUSH    D               ;and resave
004A 210000                     LXI     H,GNFSV         ;Move file name from save area
                                                        ;into FCB
004D 0E0D                       MVI     C,GNFSVL        ;Save area length
004F CD8A00                     CALL    MOVE

0052 0E12                       MVI     C,B$SEARCHN     ;Search NEXT
0054 CD0500                     CALL    BDOS
0057 E1                         POP     H               ;Recover FCB address
0058 FEFF                       CPI     0FFH            ;Check for error
005A CA7D00                     JZ      GNFEX           ;Error exit

                        GNFC:
005D E5                         PUSH    H               ;Save FCB address
005E 87                         ADD     A               ;Multiply BDOS return code * 32
```

Figure 5-17. (Continued)

```
005F 87             ADD     A               ;* 4
0060 87             ADD     A               ;* 8
0061 87             ADD     A               ;* 16
0062 87             ADD     A               ;* 32
0063 218000         LXI     H,GNFDMA        ;HL -> DMA address
0066 5F             MOV     E,A             ;Make (code * 32) a word value
                                            ;in DE
0067 1600           MVI     D,0
0069 19             DAD     D               ;HL -> file's directory entry

                                            ;Move file name into FCB
006A D1             POP     D               ;Recover FCB address
006B E5             PUSH    H               ;Save directory entry pointer
006C D5             PUSH    D               ;and resave
006D 0E0D           MVI     C,GNFSVL        ;Length of save area
006F CD8A00         CALL    MOVE
0072 3A0000         LDA     GNFSV           ;Get disk from save area
0075 D1             POP     D               ;Recover FCB address
0076 12             STAX    D               ;Overwrite user number in FCB

                                            ;Set up to zero-fill tail end
                                            ;of FCB
0077 CD7F00         CALL    GNFZF           ;Zero-fill
007A E1             POP     H               ;Recover directory entry
                                            ;pointer
007B AF             XRA     A               ;Clear carry
007C C9             RET

                GNFEX:
007D 37             STC                     ;Set carry to indicate error
007E C9             RET

                ;GNFZF
                ;Get next file -- zero fill
                ;This subroutine zero-fills the bytes that follow the
                ;file name and type in an FCB.

                ;Entry parameters
                ;       DE -> file control block

                GNFZF:
007F 210D00         LXI     H,GNFSVL        ;Bypass area that holds file name
0082 19             DAD     D               ;HL -> FCB + GNFSVL
0083 54             MOV     D,H             ;DE -> FCB + GNFSVL
0084 5D             MOV     E,L
0085 13             INX     D               ;DE -> FCB + GNFSVL + 1
0086 3600           MVI     M,0             ;FCB + GNFSVL = 0
0088 0E17           MVI     C,GNFFCL-GNFSVL ;Remainder of file control block

                ;Drop into MOVE
                ;Spread 0's through remainder
                ;of FCB

                ;MOVE
                ;This subroutine moves C bytes from HL to DE.

                MOVE:
008A 7E             MOV     A,M             ;Get source byte
008B 12             STAX    D               ;Save destination byte
008C 13             INX     D               ;Increment destination pointer
008D 23             INX     H               ;Increment source pointer
008E 0D             DCR     C               ;Decrement count
008F C28A00         JNZ     MOVE            ;Go back for more
0092 C9             RET

                ;BF
                ;Build file control block

                ;This subroutine formats a 00H-byte terminated string

                ;(presumed to be a file name) into an FCB, setting the
                ;disk and file name and type, and clearing the
                ;remainder of the FCB to 0's.
```

Figure 5-17. (Continued)

```
                    ;Entry parameters
                    ;       DE -> File control block (36 bytes)
                    ;       HL -> File name string (00H-byte-terminated)

                    ;Exit parameters
                    ;       The built file control block

                    ;This subroutine is shown in full in Figure 5-16
0093 C9             BF:     RET                     ;Dummy subroutine for this example
```

Figure 5-17. (Continued)

Function 18: Search for Next Name Match

Function Code: C = 12H
Entry Parameters: None (assumes previous Search First call)
Exit Parameters: A = Directory code

Example

```
0012 =          B$SEARCHN       EQU     18              ;Search Next
0005 =          BDOS            EQU     5               ;BDOS entry point

0000 0E12               MVI     C,B$SEARCHN     ;Function code
                                                ;Note: No FCB pointer
                                                ;You must precede this call
                                                ; with a call to Search First
0002 CD0500             CALL    BDOS            ;A = 0,1,2,3
                                                ;(A * 32) + DMA -> directory
                                                ; entry
                                                ;A = 0FFH if file name not
                                                ; found
```

Purpose This function searches down the file directory for the *next* file name, type, and extent that match the FCB specified in a previous Search First function call.

Search First and Search Next are the only BDOS functions that must be used together. As you can see, the Search Next function does not require an FCB address as an input parameter—all the necessary information will have been left in the BDOS on the Search First call.

Like Search First, Search Next returns a directory code in the A register; in this case, if A = 0FFH, it means that there are no *more* files that match the file control block. If A is not 0FFH, it will be a value of 0, 1, 2, or 3, indicating the relative directory entry number.

Notes There are two ways of using the Search First/Next calls. Consider a simple file copying program that takes as input an ambiguous file name. You could scan the file directory, matching all of the possible file names, possibly displaying them on the console, and storing the names of the files to be copied in a table inside your program. This would have the advantage of enabling you to present the file names to the operator before any copying occurred. You could even arrange for the operator to select which files to copy on a file-by-file basis. One disadvantage would be that you could not accurately predict how many files might be selected. On some hard disk systems you might have to accommodate several thousand file names.

The alternative way of handling the problem would be to match one file name, copy it, then match the next file name, copy it, and so on. If you gave the operator the choice of selecting which files to copy, this person would have to wait at the terminal as each file was being copied, but the program would not need to have large table areas set aside to hold file names. This solution to the problem is slightly more complicated, as you can see from the logic in Figure 5-17.

The subroutine in Figure 5-17, Get Next File (GNF), contains all of the necessary logic to search down a directory for both alternatives described. It does require that you indicate *on entry* whether it should search for the first or next file match, by setting A to zero or some nonzero value respectively.

You can see from Figure 5-17 that whenever the subroutine is called to get the *next* file, you must execute a Search First function to re-find the previous file. Only then can a Search Next be issued.

As with all functions that return a directory code in A, if this value is not 0FFH, it will be the relative directory entry number in the directory record currently in memory. This directory record will have been read into memory at whatever address was specified at the last Set DMA Address function call (code 26, 1AH). Notwithstanding its odd name, the DMA Address is simply the address into which any record input from disk will be placed. If the Set DMA Address function has not been used to change the value, then the CP/M default DMA address, location 0080H, will be used to hold the directory record.

The actual code for locating the address of the particular directory entry matched by the Search First/Next functions is shown in Figure 5-17 near the label GNFC. The method involves multiplying the directory code by 32 and then adding this product to the current DMA address.

Function 19: Erase (Delete) File

Function Code: C = 13H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
0013 =          B$ERASE         EQU     19      ;Erase File
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 3F3F4C454EFCB$NAME:        DB      '??LENAME'      ;Ambiguous file name
0009 3F5950     FCB$TYP:        DB      '?YP'   ;Ambiguous file type
000C 00         FCB$EXTENT:     DB      0       ;Extent
```

```
000D 0000         FCB$RESV:       DB      0,0          ;Reserved for CP/M
000F 00           FCB$RECUSED:    DB      0            ;Records used in this extent
0010 0000000000FCB$ABUSED:        DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                   DB      0,0,0,0,0,0,0,0
0020 00           FCB$SEQREC:     DB      0            ;Sequential rec. to read/write
0021 0000         FCB$RANREC:     DW      0            ;Random rec. to read/write
0023 00           FCB$RANRECO:    DB      0            ;Random rec. overflow byte (MS)

0024 0E13                 MVI     C,B$ERASE            ;Function code
0026 110000               LXI     D,FCB                ;DE -> file control block
0029 CD0500               CALL    BDOS                 ;A = 0FFH if file not found
```

Purpose This function logically deletes from the file directory files that match the FCB addressed by DE. It does so by replacing the first byte of each relevant directory entry (remember, a single file can have several entries, one for each extent) by the value 0E5H. This flags the directory entry as being available for use.

Notes Like the previous two functions, Search First and Search Next, this function can take an ambiguous file name and type as part of the file control block, but unlike those functions, the logical disk select code cannot be a "?".

This function returns a directory code in A in the same way as the previous file operations.

Function 20: Read Sequential

Function Code: C = 14H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
0014 =            B$READSEQ       EQU     20           ;Read Sequential
0005 =            BDOS            EQU     5            ;BDOS entry point

                  FCB:                                 ;File control block
0000 00           FCB$DISK:       DB      0            ;Search on default disk drive.
0001 46494C454EFCB$NAME:          DB      'FILENAME'       ;file name
0009 545950       FCB$TYP:        DB      'TYP'        ;File type
000C                              DS      24           ;Set by file open


                                                       ;Record will be read into
                                                       ; address set by prior SETDMA
                                                       ; call
0024 0E14                 MVI     C,B$READSEQ          ;Function code
0026 110000               LXI     D,FCB                ;DE -> File control block
0029 CD0500               CALL    BDOS                 ;A = 00 if operation successful
                                                       ;A  = nonzero if no data in
                                                       ; file
```

Purpose This function reads the next record (128-byte sector) from the designated file into memory at the address set by the last Set DMA function call (code 26, 1AH). The record read is specified by the FCB's sequential record field (FCB$SEQREC in the example listing for the Open File function, code 15). This field is incremented by 1 so that a subsequent call to Read Sequential will get the next record from the file. If the end of the current extent is reached, then the BDOS will

```
                        ;GETC
                        ;This subroutine gets the next character from a
                        ;sequential disk file. It assumes that the file has
                        ;already been opened.

                        ;>>>     Note : this subroutine changes CP/M's DMA address.

                        ;Entry parameters
                        ;        DE -> file control block

                        ;Exit parameters
                        ;        A = next character from file
                        ;            (= OFFH on physical end of file)
                        ;            Note : 1AH is normal EOF character for
                        ;                   ASCII Files.

                        ;Calling sequence
                        ;        LXI     DE,FCB
                        ;        CALL    GETC
                        ;        CPI     1AH
                        ;        JZ      EOFCHAR
                        ;        CPI     OFFH
                        ;        JZ      ACTUALEOF

0014 =                  B$READSEQ       EQU     20      ;Read sequential
001A =                  B$SETDMA        EQU     26      ;Set DMA address
0005 =                  BDOS            EQU     5       ;BDOS entry point

0080 =                  GETCBS  EQU     128             ;Buffer size
0000                    GETCBF: DS      GETCBS          ;Declare buffer
0080 00                 GETCCC: DB      0               ;Char. count (initially
                                                        ;"empty")

                        GETC:
0081 3A8000                     LDA     GETCCC          ;Check if buffer is empty
0084 B7                         ORA     A
0085 CA9900                     JZ      GETCFB          ;Yes, fill buffer

                        GETCRE:                         ;Re-entry point after buffer filled
0088 3D                         DCR     A               ;No, downdate count
0089 328000                     STA     GETCCC          ;Save downdated count

008C 47                         MOV     B,A             ;Compute offset of next
                                                        ;character
008D 3E7F                       MVI     A,GETCBS-1      ;By subtracting
008F 90                         SUB     B               ;(buffer size -- downdated count)
0090 5F                         MOV     E,A             ;Make result into word value
0091 1600                       MVI     D,0
0093 210000                     LXI     H,GETCBF        ;HL -> base of buffer
0096 19                         DAD     D               ;HL -> next character in buffer
0097 7E                         MOV     A,M             ;Get next character
0098 C9                         RET

                        GETCFB:                         ;Fill buffer
0099 D5                         PUSH    D               ;Save FCB pointer
009A 110000                     LXI     D,GETCBF        ;Set DMA address to buffer
009D 0E1A                       MVI     C,B$SETDMA      ;function code
009F CD0500                     CALL    BDOS
00A2 D1                         POP     D               ;Recover FCB pointer
00A3 0E14                       MVI     C,B$READSEQ     ;Read sequential "record" (sector)
00A5 CD0500                     CALL    BDOS
00A8 B7                         ORA     A               ;Check if read unsuccessful (A = NZ)
00A9 C2B400                     JNZ     GETCX           ;Yes
00AC 3E80                       MVI     A,GETCBS        ;Reset count
00AE 328000                     STA     GETCCC
00B1 C38800                     JMP     GETCRE          ;Re-enter subroutine

                        GETCX:                          ;Physical end of file
00B4 3EFF                       MVI     A,OFFH          ;Indicate such
00B6 C9                         RET
```

Figure 5-18. Read next character from sequential disk file automatically open the next extent and reset the sequential record field to 0, ready for the next Read function call.

The file specified in the FCB must have been readied for input by issuing an Open File (code 15, 0FH) or a Create File (code 22, 16H) BDOS call.

The value 00H is returned in A to indicate a successful Read Sequential operation, while a nonzero value shows that the Read could not be completed because there was no data in the next record, as at the end of file.

Notes

Although it is not immediately obvious, you can change the sequential record number, FCB$SEQREC, and within a given extent, read a record at random. If you want to access any given record within a file, you must compute which extent that record would be in and set the extent field in the file control block (FCB$EXTENT) before you open the file. Thus, although the function name implies sequential access, in practice you can use it to perform a simple type of random access. If you need to do true random access, look ahead to the Random Read function (code 33), which takes care of opening the correct extent automatically.

Figure 5-18 shows an example of a subroutine that returns the data from a sequential file byte-by-byte, reading in records from the file as necessary. This subroutine, GETC, is useful as a low-level "primitive" on which you can build more sophisticated functions, such as those that read a fixed number of characters or read characters up to a CARRIAGE RETURN/LINE FEED combination.

When you read data from a CP/M text file, the normal convention is to fill the last record of the file with 1AH characters (CONTROL-Z). Therefore, two possible conditions can indicate end-of-file: either encountering a 1AH, or receiving a return code from the BDOS function (in the A register) of 0FFH. However, if the file that you are reading is not an ASCII text file, then a 1AH character has no special meaning—it is just a normal data byte in the body of the file.

Function 21: Write Sequential

Function Code: C = 15H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
0015 =          B$WRITESEQ      EQU     21          ;Write Sequential
0005 =          BDOS            EQU     5           ;BDOS entry point

                FCB:                                ;File control block
0000 00         FCB$DISK:       DB      0           ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'        ;file name
0009 545950     FCB$TYP:        DB      'TYP'       ;File type
000C                            DS      24          ;Set by Open or Create File

                                                    ;Record must be in address
                                                    ; set by prior SETDMA call
0024 0E15               MVI     C,B$WRITESEQ        ;Function code
0026 110000             LXI     D,FCB               ;DE -> File control block
0029 CD0500             CALL    BDOS                ;A = 00H if operation
                                                    ; successful
                                                    ;A = nonzero if disk full
```

Purpose This function writes a record from the address specified in the last Set DMA (code 26, 1AH) function call to the file defined in the FCB. The sequential record number in the FCB (FCB$SEQREC) is updated by 1 so that the next call to Write Sequential will write to the next record position in the file. If necessary, a new extent will be opened to receive the new record.

This function is directly analogous to the Read Sequential function, writing instead of reading. The file specified in the FCB must first be activated by an Open File (code 15, 0FH) or create File call (code 22, 16H).

A directory code of 00H is returned in A to indicate that the Write was successful; a nonzero value is returned if the Write could not be completed because the disk was full.

Notes As with the Read Sequential function (code 20, 14H), you can achieve a simple form of random writing to the file by manipulating the sequential record number (FCB$SEQREC). However, you can only overwrite *existing* records in the file, and if you want to move to another extent, you must close the file and reopen it with the FCB$EXTENT field set to the correct value. For true random writing to the file, look ahead to the Write Random function (code 34, 22H). This takes care of opening or creating the correct extent of the file automatically.

The only logical error condition that can occur when writing to a file is insufficient room on the disk to accommodate the next extent of the file. Any hardware errors detected will be handled by the disk driver built into the BIOS or BDOS.

Figure 5-19 shows a subroutine, PUTC, to which you can pass data a byte at a time. It assembles this data into a buffer, making a call to Write Sequential whenever the buffer becomes full. You can see that provision is made in the entry parameters (by setting register B to a nonzero value) for the subroutine to fill the remaining unused characters of the buffer with 1AH characters. You must do this to denote the end of an ASCII text file.

Function 22: Create (Make) File

Function Code: C = 16H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
0016 =          B$CREATE        EQU     22      ;File Create
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;file name
0009 545950     FCB$TYP:        DB      'TYP'   ;File type
000C 00         FCB$EXTENT:     DB      0       ;Extent
```

```
000D 0000       FCB$RESV:       DB      0,0       ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0         ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0         ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0         ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0         ;Random rec. overflow byte (MS)

                                                  ;Note : file to be created
                                                  ;must not already exist....
0024 0E16               MVI     C,B$CREATE        ;Function code
0026 110000             LXI     D,FCB             ;DE -> file control block
0029 CD0500             CALL    BDOS              ;A = 0,1,2,3 if operation
                                                  ; successful
                                                  ;A = 0FFH if directory full
```

```
                ;PUTC
                ;This subroutine either puts the next chararacter out
                ;to a sequential file, writing out completed "records"
                ;(128-byte sectors) or, if requested to, will fill the
                ;remainder of the current "record" with 1AH's to
                ;indicate end of file to CP/M.

                ;Entry parameters
                ;       DE -> File control block
                ;       B = 0, A = next data character to be output
                ;       B /= 0, fill the current "record" with 1AH's

                ;Exit parameters
                ;       none.

                ;Calling sequence
                ;       LXI     D,FCB
                ;       MVI     B,0     ;Not end of file
                ;       LDA     CHAR
                ;       CALL    PUTC
                ;     or
                ;       LXI     D,FCB
                ;       MVI     B,1     ;Indicate end of file
                ;       CALL    PUTC

0015 =          B$WRITESEQ      EQU     21      ;Write sequential
001A =          B$SETDMA        EQU     26      ;Set DMA address
0005 =          BDOS            EQU     5       ;BDOS entry point

0080 =          PUTCBS  EQU     128             ;Buffer size
0000            PUTCBF: DS      PUTCBS          ;Declare buffer
0080 00         PUTCCC: DB      0               ;Char. count (initially "empty")

                PUTC:
0081 D5                 PUSH    D               ;Save FCB address
0082 F5                 PUSH    PSW             ;Save data character
0083 78                 MOV     A,B             ;Check if end of file requested
0084 B7                 ORA     A
0085 C29900             JNZ     PUTCEF          ;Yes
0088 CDC300             CALL    PUTCGA          ;No, get address of next free byte
                                                ;HL -> next free byte
                                                ;E = Current char. count (as
                                                ;well as A)
008B F1                 POP     PSW             ;Recover data character
008C 77                 MOV     M,A             ;Save in buffer
008D 7B                 MOV     A,E             ;Get current character count
008E 3C                 INR     A               ;Update character count
008F FE80               CPI     PUTCBS          ;Check if buffer full
0091 CAA900             JZ      PUTCWB          ;Yes, write buffer
0094 328000             STA     PUTCCC          ;No, save updated count
0097 D1                 POP     D               ;Dump FCB address for return
0098 C9                 RET
```

Figure 5-19. Write next character to sequential disk file

```
                        PUTCEF:                                     ;End of file
0099 F1                         POP     PSW                         ;Dump data character
009A CDC300                     CALL    PUTCGA                      ;HL -> next free byte
                                                                    ;A = current character count

                        PUTCCE:                                     ;Copy EOF character
009D FE80                       CPI     PUTCBS                      ;Check for end of buffer
009F CAA900                     JZ      PUTCWB                      ;Yes, write out the buffer
00A2 361A                       MVI     M,1AH                       ;No, store EOF in buffer
00A4 3C                         INR     A                           ;Update count
00A5 23                         INX     H                           ;Update buffer pointer
00A6 C39D00                     JMP     PUTCCE                      ;Continue until end of buffer

                        PUTCWB:                                     ;Write buffer
00A9 AF                         XRA     A                           ;Reset character count to 0
00AA 328000                     STA     PUTCCC
00AD 110000                     LXI     D,PUTCBF                    ;DE -> buffer
00B0 0E1A                       MVI     C,B$SETDMA                  ;Set DMA address -> buffer
00B2 CD0500                     CALL    BDOS
00B5 D1                         POP     D                           ;Recover FCB address
00B6 0E15                       MVI     C,B$WRITESEQ                ;Write sequential record
00B8 CD0500                     CALL    BDOS
00BB B7                         ORA     A                           ;Check if error
00BC C2C000                     JNZ     PUTCX                       ;Yes if A = NZ
00BF C9                         RET                                 ;No, return to caller

                        PUTCX:                                      ;Error exit
00C0 3EFF                       MVI     A,0FFH                      ;Indicate such
00C2 C9                         RET

                        PUTCGA:                                     ;Return with HL -> next free char.
                                                                    ;and A = current char. count
00C3 3A8000                     LDA     PUTCCC                      ;Get current character count
00C6 5F                         MOV     E,A                         ;Make word value in DE
00C7 1600                       MVI     D,0
00C9 210000                     LXI     H,PUTCBF                    ;HL -> Base of buffer
00CC 19                         DAD     D                           ;HL -> next free character
00CD C9                         RET
```

Figure 5-19. Write next character to sequential disk file (continued)

Purpose

This function creates a new file of the specified name and type. You must first ensure that no file of the same name and type already exists on the same logical disk, either by trying to open the file (if this succeeds, the file already exists) or by unconditionally erasing the file.

In addition to creating the file and its associated file directory entry, this function also effectively opens the file so that it is ready for records to be written to it.

This function returns a normal directory code if the file creation has completed successfully or a value of 0FFH if there is insufficient disk or directory space.

Notes

Under some circumstances, you may want to create a file that is slightly more "secure" than normal CP/M files. You can do this by using either lowercase letters or nongraphic ASCII characters such as ASCII NUL (00H) in the file name or type. Neither of these classes of characters can be generated from the keyboard; in the first case, the CCP changes all lowercase characters to uppercase, and in the second, it rejects names with odd characters in them. Thus, computer operators cannot erase such a file because there is no way that they can create the same file name from the CCP.

The converse is also true; the only way that you can erase these files is by using a program that *can* set the exact file name into an FCB and then issue an Erase File function call.

Note that this function cannot accept an ambiguous file name in the FCB.

Figure 5-20 shows a subroutine that creates a file only after it has erased any existing files of the same name.

Function 23: Rename File

Function Code: C = 17H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Directory code

Example

```
0017 =          B$RENAME        EQU     23      ;Rename file
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00                         DB      0       ;Search on default disk drive
0001 4F4C444E41                 DB      'OLDNAME '      ;File name
0009 545950                     DB      'TYP'   ;File type
000C 00000000                   DB      0,0,0,0
```

```
                ;CF
                ;Create file
                ;This subroutine creates a file. It erases any
                ;previous file before creating the new one.

                ;Entry parameters
                ;       DE -> File control block for new file

                ;Exit parameters
                ;       Carry clear if operation successful
                ;               (A = 0,1,2,3)
                ;       Carry set if error (A = 0FFH)

                ;Calling sequence
                ;       LXI     D,FCB
                ;       CALL    CF
                ;       JC      ERROR

0013 =          B$ERASE         EQU     19      ;Erase file
0016 =          B$CREATE        EQU     22      ;Create file
0005 =          BDOS            EQU     5       ;BDOS entry point

                CF:
0000 D5                 PUSH    D               ;Preserve FCB pointer
0001 0E13               MVI     C,B$ERASE       ;Erase any existing file
0003 CD0500             CALL    BDOS
0006 D1                 POP     D               ;Recover FCB pointer
0007 0E16               MVI     C,B$CREATE      ;Create (and open new file)
0009 CD0500             CALL    BDOS
000C FEFF               CPI     0FFH            ;Carry set if OK, clear if error
000E 3F                 CMC                     ;Complete to use Carry set if Error
000F C9                 RET
```

Figure 5-20. Create file request

```
0010 00                          DB      0            ;FCB + 16
0011 4E45574E41                  DB      'NEWNAME '          ;File name
0019 545950                      DB      'TYP'       ;File type
001C 00000000                    DB      0,0,0,0

0020 0E17                MVI     C,B$RENAME          ;Function code
0022 110000              LXI     D,FCB               ;DE -> file control block
0025 CD0500              CALL    BDOS                ;A = 00H if operation succesful
                                                     ;A = 0FFH if file not found
```

Purpose This function renames an existing file name and type to a new name and type. It is unusual in that it uses a single FCB to store both the old file name and type (in the first 16 bytes) and the new file name and type (in the second 16 bytes).

This function returns a normal directory code if the file rename was completed successfully or a value of 0FFH if the old file name could not be found.

Notes The Rename File function only checks that the old file name and type exist; it makes no check to ensure that the new name and type combination does not already exist. Therefore, you should try to open the new file name and type. If you succeed, do not attempt the rename operation. CP/M will create more than one file of the same name and type, and you stand to lose the information in both files as you attempt to sort out the problem.

For security, you can also use lowercase letters and nongraphic characters in the file name and type, as described under the File Create function (code 22, 16H) above.

Never use ambiguous file names in a rename operation; it produces strange effects and may result in files being irreparably damaged. This function will change *all* occurrences of the old file name to the new name.

Figure 5-21 shows a subroutine that will accept an existing file name and type and a new name and type and rename the old to the new. It checks to make sure that the new file name does not already exist, returning an error code if it does.

Function 24: Get Active Disks (Login Vector)

Function Code: C = 18H
Entry Parameters: None
Exit Parameters: HL = Active disk map (login vector)

Example

```
0018 =          B$GETACTDSK     EQU     24      ;Get Active Disks
0005 =          BDOS            EQU     5       ;BDOS entry point

                                                ;Example of getting active
0000 0E18             MVI     C,B$GETACTDSK     ; disk function code
0002 CD0500           CALL    BDOS              ;HL = active disk bit map
                                                ;Bits are = 1 if disk active
                                                ;Bits 15 14 13 ... 2 1 0
                                                ;Disk  P  O  N ... C B A
```

Purpose This function returns a bit map, called the *login vector,* in register pair HL, indicating which logical disk drives have been selected since the last warm boot or

```
                    ;RF
                    ;Rename file
                    ;This subroutine renames a file.
                    ;It uses the BF (build FCB) subroutine shown in Figure 5.16

                    ;Entry parameters
                    ;       *** No case-folding of file names occurs ***
                    ;       HL -> old file name (00-byte terminated)
                    ;       DE -> new file name (00-byte terminated)

                    ;Exit parameters
                    ;       Carry clear if operation successful
                    ;               (A = 0,1,2,3)
                    ;       Carry set if error
                    ;               A = OFEH if new file name already exists
                    ;               A = OFFH if old file name does not exist

                    ;Calling sequence
                    ;       LXI     H,OLDNAME       ;HL -> old name
                    ;       LXI     D,NEWNAME       ;DE -> new name
                    ;       CALL    RF
                    ;       JC      ERROR

000F =              B$OPEN          EQU     15      ;Open file
0017 =              B$RENAME        EQU     23      ;Rename file
0005 =              BDOS            EQU     5       ;BDOS entry point

0000 0000000000RFFCB:   DW      0,0,0,0,0,0,0,0 ;1 1/2 FCB's long
0010 0000000000         DW      0,0,0,0,0,0,0,0
0020 0000000000         DW      0,0,0,0,0,0,0,0
0030 000000             DW      0,0,0


                    RF:
0036 D5                 PUSH    D               ;Save new name pointer
0037 110000             LXI     D,RFFCB         ;Build old name FCB
                                                ;HL already -> old name
003A CD5D00             CALL    BF

003D E1                 POP     H               ;Recover new name pointer
003E 111000             LXI     D,RFFCB+16      ;Build new name in second part of file
0041 CD5D00             CALL    BF              ;control block

0044 111000             LXI     D,RFFCB+16      ;Experimentally try
0047 0E0F               MVI     C,B$OPEN        ;to open the new file
0049 CD0500             CALL    BDOS            ;to ensure it does
004C FEFF               CPI     OFFH            ;not already exist
004E 3EFE               MVI     A,OFEH          ;Assume error (flags unchanged)
0050 D8                 RC                      ;Carry set if A was 0,1,2,3

0051 110000             LXI     D,RFFCB         ;Rename the file
0054 0E17               MVI     C,B$RENAME
0056 CD0500             CALL    BDOS
0059 FEFF               CPI     OFFH            ;Carry set if OK, clear if error
005B 3F                 CMC                     ;Invert to use carry, set if error
005C C9                 RET

                    ;BF
                    ;Build file control block
                    ;This subroutine formats a 00H-byte terminated string
                    ;(presumed to be a file name) into an FCB, setting the
                    ;disk and the file name and type, and clearing the
                    ;remainder of the FCB to 0's.

                    ;Entry parameters
                    ;       DE -> file control block (36 bytes)
                    ;       HL -> file name string (00H-byte terminated)

                    ;Exit parameters
                    ;       The built file control block.

                    ;Calling sequence
                    ;       LXI     D,FCB
                    ;       LXI     H,FILENAME
                    ;       CALL    BF

                    BF:
005D C9                 RET                     ;Dummy subroutine : see Figure 5.16.
```

Figure 5-21. Rename file request

Reset Disk function (code 13, 0DH). The least significant bit of L corresponds to disk A, while the highest order bit in H maps disk P. The bit corresponding to the specific logical disk is set to 1 if the disk has been selected or to 0 if the disk is not currently on-line.

Logical disks can be selected programmatically through any file operation that sets the drive field to a nonzero value, through the Select Disk function (code 14, 0EH), or by the operator entering an "X:" command where "X" is equal to A, B, ..., P.

Notes This function is intended for programs that need to know which logical disks are currently active in the system—that is, those logical disks which have been selected.

Function 25: Get Current Default Disk

Function Code: C = 19H
Entry Parameters: None
Exit Parameters: A = Current disk
(0 = A, 1 = B, ..., F=P)

Example

```
0019 =          B$GETCURDSK     EQU     25          ;Get Current Disk
0005 =          BDOS            EQU     5           ;BDOS entry point

0000 0E19               MVI     C,B$GETCURDSK       ;Function code
0002 CD0500             CALL    BDOS                ;A = 0 if A:, 1 if B: ...
```

Purpose This function returns the current default disk set by the last Select Disk function call (code 14, 0EH) or by the operator entering the "X:" command (where "X" is A, B, ..., P) to the CCP.

Notes This function returns the current default disk in coded form. Register A = 0 if drive A is the current drive, 1 if drive B, and so on. If you need to convert this to the corresponding ASCII character, simply add 41H to register A.

Use this function when you convert a file name and type in an FCB to an ASCII string in order to display it. If the first byte of the FCB is 00H, the current default drive is to be used. You must therefore use this function to determine the logical disk letter for the default drive.

Function 26: Set DMA (Read/Write) Address

Function Code: C = 1AH
Entry Parameters: DE = DMA (read/write) address
Exit Parameters: None

Example

```
001A =          B$SETDMA        EQU     26          ;Set DMA Address
0005 =          BDOS            EQU     5           ;BDOS entry point
```

```
0000            SECBUFF:        DS      128     ;Sector buffer

0080 0E1A               MVI     C,B$SETDMA      ;Function code
0082 110000             LXI     D,SECBUFF       ;Pointer to buffer
0085 CD0500             CALL    BDOS
```

Purpose This function sets the BDOS's direct memory access (DMA) address to a new value. The name is an historic relic dating back to the Intel Development System on which CP/M was originally developed. This machine, by virtue of its hardware, could read data from a diskette directly into memory or write data to a diskette directly from memory. The name *DMA address* now applies to the address of the buffer to and from which data is transferred whenever a diskette Read, Write, or directory operation is performed.

Whenever CP/M first starts up (cold boot) or a warm boot or Reset Disk operation occurs, the DMA address is reset to its default value of 0080H.

Notes No function call can tell you the current value of the DMA address. All you can do is make a Set DMA function call to ensure that it is where you want it.

Once you have set the DMA address to the correct place for your program, it will remain set there until another Set DMA call, Reset Disk, or warm boot occurs.

The Read and Write Sequential and Random operations use the current setting of the DMA address, as do the directory operations Search First and Search Next.

Function 27: Get Allocation Vector

Function Code: C = 1BH
Entry Parameters: None
Exit Parameters: HL = Address of allocation vector

Example

```
001B =          B$GETALVEC      EQU     27      ;Get Allocation Vector Address
0005 =          BDOS            EQU     5       ;BDOS entry point

0000 0E1B               MVI     C,B$GETALVEC    ;Function code
0002 CD0500             CALL    BDOS            ;HL -> Base address of
                                                ;       allocation vector
```

Purpose This function returns the base, or starting, address of the allocation vector for the currently selected logical disk. This information, indicating which parts of the disk are assigned, is used by utility programs and the BDOS itself to determine how much unused space is on the logical disk, to locate an unused allocation block in order to extend a file, or to relinquish an allocation block when a file is deleted.

Notes Digital Research considers the actual layout of the allocation vector to be proprietary information.

Function 28: Set Logical Disk to Read-Only Status

Function Code: C = 1CH
Entry Parameters: None
Exit Parameters: None

Example

```
001C =          B$SETDSKRO      EQU     28              ;Set disk to Read Only
                                                        ; function code
0005 =          BDOS            EQU     5               ;BDOS entry point

                                                        ;Sets disk selected by prior
                                                        ;Select disk function call
0000 0E1C               MVI     C,B$SETDSKRO            ;Function code
0002 CD0500             CALL    BDOS
```

Purpose This function logically sets the currently selected disk to a Read-Only state. Any attempts to execute a Write Sequential or Write Random function to the selected disk will be intercepted by the BDOS, and the following message will appear on the console:

```
BDOS Err on X: R/O
```

where X: is the selected disk.

Notes Once you have requested Read-Only status for the currently selected logical disk, this status will persist even if you proceed to select other logical disks. In fact, it will remain in force until the next warm boot or Reset Disk System function call.

Digital Research documentation refers to this function code as Disk Write Protect. The Read-Only description is used here because it corresponds to the error message produced if your program attempts to write on the disk.

Function 29: Get Read-Only Disks

Function Code: C = 1DH
Entry Parameters: None
Exit Parameters: HL = Read-Only disk map

Example

```
001D =          B$GETRODSKS     EQU     29              ;Get Read Only disks
0005 =          BDOS            EQU     5               ;BDOS entry point

0000 0E19               MVI     C,B$GETRODSKS           ;Function code
0002 CD0500             CALL    BDOS                    ;HL = Read Only disk bit map
                                                        ;Bits are = 1 if disk Read Only
                                                        ;Bits 15 14 13 ... 2 1 0
                                                        ;Disk  P  O  N ... C B A
```

Purpose This function returns a bit map in registers H and L showing which logical disks in the system have been set to Read-Only status, either by the Set Logical

Disk to Read-Only function call (code 28, 1CH), or by the BDOS itself, because it detected that a diskette had been changed.

The least significant bit of L corresponds to logical disk A, while the most significant bit of H corresponds to disk P. The bit corresponding to the specific logical disk is set to 1 if the disk has been set to Read-Only status.

Function 30: Set File Attributes

Function Code: C = 1EH
Entry Parameters: DE = Address of FCB
Exit Parameters: A = Directory code

Example

```
001E =          B$SETFAT        EQU     30      ;Set File Attribute
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;File name
0009 D4         FCB$TYP:        DB      'T'+80H         ;Type with R/O
                                                        ; attribute
000A 5950                       DB      'YP'
000C 0000000000                 DW      0,0,0,0,0,0,0,0,0,0,0

0022 0E1E               MVI     C,B$SETFAT      ;Function code
0024 110000             LXI     D,FCB           ;DE -> file control block
                                                ;MS bits set in file name/type
0027 CD0500             CALL    BDOS            ;A = 0FFH if file not found
```

Purpose

This function sets the bits that describe attributes of a file in the relevant directory entries for the specified file. Each file can be assigned up to 11 file attributes. Of these 11, two have predefined meanings, four others are available for you to use, and the remaining five are reserved for future use by CP/M.

Each attribute consists of a single bit. The most significant bit of each byte of the file name and type is used to store the attributes. The file attributes are known by a code consisting of the letter "f" (for file name) or "t" (for file type), followed by the number of the character position and a single quotation mark. For example, the Read-Only attribute is t1′.

The significance of the attributes is as follows:

- f1′ to f4′ Available for you to use
- f5′ to f8′ Reserved for future CP/M use
- t1′ Read-Only File attribute
- t2′ System File attribute
- t3′ Reserved for future CP/M use

Attributes are set by presenting this function with an FCB in which the unambiguous file name has been preset with the most significant bits set appropriately. This function then searches the directory for a match and changes the matched entries to contain the attributes which have been set in the FCB.

The BDOS will intercept any attempt to write on a file that has the Read-Only attribute set. The DIR command in the CCP does not display any file with System status.

Notes You can use the four attributes available to you to set up a file security system, or perhaps to flag certain files that must be backed up to other disks. The Search First and Search Next functions allow you to view the complete file directory entry, so your programs can test the attributes easily.

The example subroutines in Figures 5-22 and 5-23 show how to set file attributes (SFA) and get file attributes (GFA), respectively. They both use a bit map in which the most significant 11 bits of the HL register pair are used to indicate the corresponding high bits of the 11 characters of the file name/type combination. You will also see some equates that have been declared to make it easier to manipulate the attributes in this bit map.

```
                        ;SFA
                        ;Set file attributes
                        ;This subroutine takes a compressed bit map of all the
                        ;file attribute bits, expands them into an existing
                        ;file control block and then requests CP/M to set
                        ;the attributes in the file directory.

                        ;Entry parameters
                        ;       DE -> file control block
                        ;       HL = bit map. Only the most significant 11
                        ;            bits are used. These correspond directly
                        ;            with the possible attribute bytes.

                        ;Exit parameters
                        ;       Carry clear if operation successful (A = 0,1,2,3)
                        ;       Carry set if error (A = 0FFH)

                        ;Calling sequence
                        ;       LXI     D,FCB
                        ;       LXI     H,0000$0000$1100$0000B  ;Bit Map
                        ;       CALL    SFA
                        ;       JC      ERROR

                                                        ;File Attribute Equates

8000 =                  FA$F1   EQU     1000$0000$0000$0000B    ;F1' - F4'
4000 =                  FA$F2   EQU     0100$0000$0000$0000B    ;Available for use by
2000 =                  FA$F3   EQU     0010$0000$0000$0000B    ; application programs
1000 =                  FA$F4   EQU     0001$0000$0000$0000B

0800 =                  FA$F5   EQU     0000$1000$0000$0000B    ;F5' - F8'
0400 =                  FA$F6   EQU     0000$0100$0000$0000B    ;Reserved for CP/M
0200 =                  FA$F7   EQU     0000$0010$0000$0000B
0100 =                  FA$F8   EQU     0000$0001$0000$0000B

0080 =                  FA$T1   EQU     0000$0000$1000$0000B    ;T1' -- read/only file
0080 =                  FA$RO   EQU     FA$T1
0040 =                  FA$T2   EQU     0000$0000$0100$0000B    ;T2' -- system files
0040 =                  FA$SYS  EQU     FA$T2
0020 =                  FA$T3   EQU     0000$0000$0010$0000B    ;T3' -- reserved for CP/M

001E =                  B$SETFAT        EQU     30      ;Set file attributes
0005 =                  BDOS            EQU     5       ;BDOS entry point
```

Figure 5-22. Set file attributes

```
                         SFA:
0000 D5                         PUSH    D               ;Save FCB pointer
0001 13                         INX     D               ;HL -> 1st character of file name
0002 0E0B                       MVI     C,8+3           ;Loop count for file name and type

                         SFAL:                          ;Main processing loop
0004 AF                         XRA     A               ;Clear carry and A
0005 29                         DAD     H               ;Shift next MS bit into carry
0006 CE00                       ACI     0               ;A = 0 or 1 depending on carry
0008 0F                         RRC                     ;Rotate LS bit of A into MS bit
0009 47                         MOV     B,A             ;Save result (00H or 80H)
000A EB                         XCHG                    ;HL -> FCB character
000B 7E                         MOV     A,M             ;Get FCB character
000C E67F                       ANI     7FH             ;Isolate all but attribute bit
000E B0                         ORA     B               ;Set attribute with result
000F 77                         MOV     M,A             ;and store back into FCB
0010 EB                         XCHG                    ;DE -> FCB, HL = remaining bit map
0011 13                         INX     D               ;DE -> next character in FCB
0012 0D                         DCR     C               ;Downdate character count
0013 C20400                     JNZ     SFAL            ;Loop back for next character
0016 0E1E                       MVI     C,B$SETFAT      ;Set file attribute function code
0018 D1                         POP     D               ;Recover FCB pointer
0019 CD0500                     CALL    BDOS
001C FEFF                       CPI     0FFH            ;Carry set if OK, clear if error
001E 3F                         CMC                     ;Invert to use carry set if error
001F C9                         RET
```

Figure 5-22. Set file attributes (continued)

```
                        ;GFA
                        ;Get file attributes
                        ;This subroutine finds the appropriate file using a
                        ;search for First Name Match function rather than opening
                        ;the file. It then builds a bit map of the file attribute
                        ;bits in the file name and type. This bit map is then ANDed
                        ;with the input bit map, and the result is returned in the
                        ;zero flag. The actual bit map built is also returned in case
                        ;more complex check is required.

                        ;>>>    Note: This subroutine changes the CP/M DMA address.

                        ;Entry parameters
                        ;       DE -> File control block
                        ;       HL = Bit map mask to be ANDed with attribute
                        ;            results

                        ;Exit parameters
                        ;       Carry clear, operation successful
                        ;               Nonzero status set to result of AND between
                        ;               input mask and attribute bits set.
                        ;               HL = Unmasked attribute bytes set.
                        ;       Carry set, file could not be found

001A =                  B$SETDMA        EQU     26      ;Set DMA address
0011 =                  B$SEARCHF       EQU     17      ;Search for first entry to match
0005 =                  BDOS            EQU     5       ;BDOS entry point
0080 =                  GFADMA          EQU     80H     ;Default DMA address

                        ;Calling sequence
                        ;       LXI     D,FCB
                        ;       LXI     H,0000$0000$1100$0000B  ;Bit map
                        ;       CALL    GFA
                        ;       JC      ERROR

                                                        ;File attribute equates

8000 =                  FA$F1   EQU     1000$0000$0000$0000B    ;F1' - F5'
4000 =                  FA$F2   EQU     0100$0000$0000$0000B    ;Available for use by
```

Figure 5-23. Get file attributes

```
2000 =          FA$F3   EQU     0010$0000$0000$0000B     ;Application programs
1000 =          FA$F4   EQU     0001$0000$0000$0000B

0800 =          FA$F5   EQU     0000$1000$0000$0000B     ;F6' - F8'
0400 =          FA$F6   EQU     0000$0100$0000$0000B     ;Reserved for CP/M
0200 =          FA$F7   EQU     0000$0010$0000$0000B
0100 =          FA$F8   EQU     0000$0001$0000$0000B

0080 =          FA$T1   EQU     0000$0000$1000$0000B     ;T1' -- read/only file
0080 =          FA$RO   EQU     FA$T1
0040 =          FA$T2   EQU     0000$0000$0100$0000B     ;T2' -- system files
0040 =          FA$SYS  EQU     FA$T2
0020 =          FA$T3   EQU     0000$0000$0010$0000B     ;T3' -- reserved for CP/M

                GFA:
0000 E5                 PUSH    H               ;Save AND-mask
0001 D5                 PUSH    D               ;Save FCB pointer
0002 0E1A               MVI     C,B$SETDMA      ;Set DMA to default address
0004 118000             LXI     D,GFADMA        ;DE -> DMA address
0007 CD0500             CALL    BDOS

000A D1                 POP     D               ;Recover FCB pointer
000B 0E11               MVI     C,B$SEARCHF     ;Search for match with name
000D CD0500             CALL    BDOS
0010 FEFF               CPI     0FFH            ;Carry set if OK, clear if error
0012 3F                 CMC                     ;Invert to use set carry if error
0013 DA4100             JC      GFAX            ;Return if error
                                                ;Multiply by 32 to get offset into DMA buffer
0016 87                 ADD     A               ;* 2
0017 87                 ADD     A               ;* 4
0018 87                 ADD     A               ;* 8
0019 87                 ADD     A               ;* 16
001A 87                 ADD     A               ;* 32
001B 5F                 MOV     E,A             ;Make into a word value
001C 1600               MVI     D,0
001E 218000             LXI     H,GFADMA        ;HL -> DMA address
0021 19                 DAD     D               ;HL -> Directory entry in DMA buffer
0022 23                 INX     H               ;HL -> 1st character of file name
0023 EB                 XCHG                    ;DE -> 1st character of file name

0024 0E0B               MVI     C,8+3           ;Count of characters in file name and type
0026 210000             LXI     H,0             ;Clear bit map

                GFAL:                           ;Main loop
0029 1A                 LDAX    D               ;Get next character of file name
002A E680               ANI     80H             ;Isolate attribute bit
002C 07                 RLC                     ;Move MS bit into LS bit
002D B5                 ORA     L               ;OR in any previously set bits
002E 6F                 MOV     L,A             ;Save result
002F 29                 DAD     H               ;Shift HL left one bit for next time
0030 13                 INX     D               ;DE -> next character in file name, type
0031 0D                 DCR     C               ;Downdate count
0032 C22900             JNZ     GFAL            ;Go back for next character

0035 29                 DAD     H               ;Left justify attribute bits in HL
0036 29                 DAD     H               ;MS attribute bit will already be in
0037 29                 DAD     H               ;bit 11 of HL, so only 4 shifts are
0038 29                 DAD     H               ;necessary

0039 D1                 POP     D               ;Recover AND-mask
003A 7A                 MOV     A,D             ;Get MS byte of mask
003B A4                 ANA     H               ;AND with MS byte of result
003C 47                 MOV     B,A             ;Save interim result
003D 7B                 MOV     A,E             ;Get LS byte of mask
003E A5                 ANA     L               ;AND with LS byte of result
003F B0                 ORA     B               ;Combine two results to set Z flag

0040 C9                 RET

                GFAX:                           ;Error exit
0041 E1                 POP     H               ;Balance stack
0042 C9                 RET
```

Figure 5-23. Get file attributes (continued)

Function 31: Get Disk Parameter Block Address

Function Code: C = 1FH
Entry Parameters: None
Exit Parameters: HL = Address of DPB

Example

```
001F =          B$GETDPB        EQU     31      ;Get Disk Parameter Block
                                                ; Address
0005 =          BDOS            EQU     5       ;BDOS entry point

                                                ;Returns DPB address of
                                                ; logical disk previously
                                                ; selected with a Select
                                                ; Disk function.
0000 0E1F               MVI     C,B$GETDPB      ;Function code
0002 CD0500             CALL    BDOS            ;HL -> Base address of current
                                                ; disk's parameter block
```

Purpose This function returns the address of the disk parameter block (DPB) for the last selected logical disk. The DPB, explained in Chapter 3, describes the physical characteristics of a specific logical disk—information mainly of interest for system utility programs.

Notes The subroutines shown in Figure 5-24 deal with two major problems. First, given a track and sector number, what allocation block will they fall into? Converseley, given an allocation block, what is its starting track and sector?

These subroutines are normally used by system utilities. They first get the DPB address using this BDOS function. Then they switch to using direct BIOS calls to perform their other functions, such as selecting disks, tracks, and sectors and reading and writing the disk.

The first subroutine, GTAS (Get Track and Sector), in Figure 5-24, takes an allocation block number and converts it to give you the starting track and sector number. GMTAS (Get Maximum Track and Sector) returns the maximum track and sector number for the specified disk. GDTAS (Get Directory Track and Sector) tells you not only the starting track and sector for the file directory, but also the number of 128-byte sectors in the directory.

Note that whenever a track number is used as an entry or an exit parameter, it is an absolute track number. That is, the number of reserved tracks on the disk before the directory has already been added to it.

GNTAS (Get Next Track and Sector) helps you read sectors sequentially. It adds 1 to the sector number, and when you reach the end of a track, updates the track number by 1 and resets the sector number to 1.

GAB (Get Allocation Block) is the converse of GTAS (Get Track and Sector). It returns the allocation block number, given a track and sector.

Finally, Figure 5-24 includes several useful 16-bit subroutines to divide the HL register pair by DE (DIVHL), to multiply HL by DE (MULHL), to subtract DE from HL (SUBHL —this can also be used as a 16-bit compare), and to shift HL right one bit (SHLR). The divide and multiply subroutines are somewhat primitive, using iterative subtraction and addition, respectively. Nevertheless, they do perform their role as supporting subroutines.

```
                ;Useful subroutines for accessing the data in the
                ;disk parameter block

000E =          B$SELDSK        EQU     14              ;Select Disk function code
001F =          B$GETDPB        EQU     31              ;Get DPB address
0005 =          BDOS            EQU     5               ;BDOS entry point

                ;It makes for easier, more compact code to copy the
                ;specific disk parameter block into local variables
                ;while manipulating the information.
                ;Here are those variables --

                DPB:                                    ;Disk parameter block
0000 0000       DPBSPT: DW      0                       ;128-byte sectors per track
0002 00         DPBBS:  DB      0                       ;Block shift
0003 00         DPBBM:  DB      0                       ;Block mask
0004 00         DPBEM:  DB      0                       ;Extent mask
0005 0000       DPBMAB: DW      0                       ;Maximum allocation block number
0007 0000       DPBNOD: DW      0                       ;Number of directory entries - 1
0009 0000       DPBDAB: DW      0                       ;Directory allocation blocks
000B 0000       DPBCBS: DW      0                       ;Check buffer size
000D 0000       DPBTBD: DW      0                       ;Tracks before directory (reserved tracks)

000F =          DPBSZ   EQU     $-DPB                   ;Disk parameter block size

                ;GETDPB
                ;Gets disk parameter block
                ;This subroutine copies the DPB for the specified
                ;logical disk into the local DPB variables above.

                ;Entry parameters
                ;       A = Logical disk number (A: = 0, B: = 1...)

                ;Exit parameters
                ;       Local variables contain DPB

                GETDPB:
000F 5F                 MOV     E,A             ;Get disk code for select disk
0010 0E0E               MVI     C,B$SELDSK      ;Select the disk
0012 CD0500             CALL    BDOS
0015 0E1F               MVI     C,B$GETDPB      ;Get the disk parameter base address
0017 CD0500             CALL    BDOS            ;HL -> DPB
001A 0E0F               MVI     C,DPBSZ         ;Set count
001C 110000             LXI     D,DPB           ;Get base address of local variables

                GDPBL:                          ;Copy DPB into local variables
001F 7E                 MOV     A,M             ;Get byte from DPB
0020 12                 STAX    D               ;Store into local variable
0021 13                 INX     D               ;Update local variable pointer
0022 23                 INX     H               ;Update DPB pointer
0023 0D                 DCR     C               ;Downdate count
0024 C21F00             JNZ     GDPBL           ;Loop back for next byte
0027 C9                 RET

                ;GTAS
                ;Get track and sector (given allocation block number)

                ;This subroutine converts an allocation block into a
                ;track and sector number -- note that this is based on
                ;128-byte sectors.

                ;>>>>>  Note: You must call GETDPB before
                ;>>>>>        you call this subroutine

                ;Entry parameters
                ;       HL = allocation block number

                ;Exit parameters
                ;       HL = track number
                ;       DE = sector number

                ;Method :
                ;In mathematical terms, the track can be derived from:
                ;Trk = ((allocation block * sec. per all. block) / sec. per trk)
                ;       + tracks before directory
```

Figure 5-24. Accessing disk parameter block data

```
                    ;The sector is derived from:
                    ;Sec = ((allocation block * sec. per all. block) modulo/
                    ;       sec. per trk) + 1

                    GTAS:
0028 3A0200                 LDA     DPBBS           ;Get block shift -- this will be 3 to
                                                    ;7 depending on allocation block size
                                                    ;It will be used as a count for shifting

                    GTASS:
002B 29                     DAD     H               ;Shift allocation block left one place
002C 3D                     DCR     A               ;Decrement block shift count
002D C22B00                 JNZ     GTASS           ;More shifts required
0030 EB                     XCHG                    ;DE = all. block * sec. per block
                                                    ;i.e. DE = total number of sectors
0031 2A0000                 LHLD    DPBSPT          ;Get sectors per track
0034 EB                     XCHG                    ;HL = sec. per trk, DE = tot. no. of sec.
0035 CD8F00                 CALL    DIVHL           ;BC = HL/DE, HL = remainder
                                                    ;BC = track, HL = sector
0038 23                     INX     H               ;Sector numbering starts from 1
0039 EB                     XCHG                    ;DE = sector, HL = track
003A 2A0D00                 LHLD    DPBTBD          ;Tracks before directory
003D 09                     DAD     B               ;DE = sector, HL = absolute track
003E C9                     RET

                    ;GMTAS
                    ;Get maximum track and sector

                    ;This is just a call to GTAS with the maximum
                    ;allocation block as the input parameter

                    ;>>>>>  Note: You must call GETDPB before
                    ;>>>>>        you call this subroutine

                    ;Entry parameters: none

                    ;Exit parameters:
                    ;       HL = maximum track number
                    ;       DE = maximum sector

                    GMTAS:
003F 2A0500                 LHLD    DPBMAB          ;Get maximum allocation block
0042 C32800                 JMP     GTAS            ;Return from GTAS with parameters in HL and DE

                    ;GDTAS
                    ;Get directory track and sector

                    ;This returns the START track and sector for the
                    ;file directory, along with the number of sectors
                    ;in the directory.

                    ;>>>>>  Note: You must call GETDPB before
                    ;>>>>>        you call this subroutine

                    ;Entry parameters: none

                    ;Exit parameters:
                    ;       BC = number of sectors in directory
                    ;       DE = directory start sector
                    ;       HL = directory start track

                    GDTAS:
0045 2A0700                 LHLD    DPBNOD          ;Get number of directory entries - 1
0048 23                     INX     H               ;Make true number of entries
                                                    ;Each entry is 32 bytes long, so to
                                                    ;convert to 128 byte sectors, divide by 4
0049 CDD000                 CALL    SHLR            ;/ 2 (by shifting HL right one bit)
004C CDD000                 CALL    SHLR            ;/ 4
004F E5                     PUSH    H               ;Save number of sectors
0050 210000                 LXI     H,0             ;Directory starts in allocation block 0
0053 CD2800                 CALL    GTAS            ;HL = track, DE = sector
0056 C1                     POP     B               ;Recover number of sectors
0057 C9                     RET
```

Figure 5-24. (Continued)

```
                         ;GNTAS
                         ;Get NEXT track and sector

                         ;This subroutine updates the input track and sector
                         ;by one, incrementing the track and resetting the
                         ;sector number as required.

                         ;>>>>>  Note: You must call GETDPB before
                         ;>>>>>        you call this subroutine

                         ; Note: you must check for end of disk by comparing
                         ;       the track number returned by this subroutine
                         ;       to that returned by by GMTAS + 1. When
                         ;       equality occurs, the end of disk has been reached.

                         ;Entry parameters
                         ;       HL = current track number
                         ;       DE = current sector number

                         ;Exit parameters
                         ;       HL = updated track number
                         ;       DE = updated sector number

                         GNTAS:
0058 E5                          PUSH    H               ;Save track
0059 13                          INX     D               ;Update sector
005A 2A0000                      LHLD    DPBSPT          ;Get sectors per track
005D CDC900                      CALL    SUBHL           ;HL = HL - DE
0060 E1                          POP     H               ;Recover current track
0061 D0                          RNC                     ;Return if updated sector <= sec. per trk.
0062 23                          INX     H               ;Update track if upd. sec > sec. per trk.
0063 110100                      LXI     D,1             ;Reset sector to 1
0066 C9                          RET

                         ;GAB
                         ;Get allocation block

                         ;This subroutine returns an allocation block number
                         ;given a specific track and sector. It also returns
                         ;the offset down the allocation block at which the
                         ;sector will be found. This offset is in units of
                         ;128-byte sectors.

                         ;>>>>>   Note: You must call GETDPB before
                         ;>>>>>         you call this subroutine

                         ;Entry parameters
                         ;       HL = track number
                         ;       DE = sector number

                         ;Exit parameters
                         ;       HL = allocation block number

                         ;Method
                         ;The allocation block is formed from:
                         ;AB = (sector + ((track - tracks before directory)
                         ;       * sectors per track)) / log2 (sectors per all. block)

                         ;The sector offset within allocation block is formed from:
                         ;Offset = (sector + ((track - tracks before directory)
                         ;       * sectors per track)) / AND (sectors per all. block - 1)

                         GAB:
0067 D5                          PUSH    D               ;Save sector
0068 EB                          XCHG                    ;DE = track
0069 2A0D00                      LHLD    DPBTBD          ;Get no. of tracks before directory
006C EB                          XCHG                    ;DE = no. of tracks before dir. HL = track
006D CDC900                      CALL    SUBHL           ;HL = HL - DE
                                                         ;HL = relative track within logical disk
0070 EB                          XCHG                    ;DE = relative track
0071 2A0000                      LHLD    DPBSPT          ;Get sectors per track
0074 CDA400                      CALL    MULHL           ;HL = HL * DE
                                                         ;HL = number of sectors
0077 EB                          XCHG                    ;DE = number of sectors
```

Figure 5-24. (Continued)

```
0078 E1                 POP     H               ;Recover sector
0079 2B                 DCX     H               ;Make relative to 0
007A 19                 DAD     D               ;HL = relative sector
007B 3A0300             LDA     DPBBM           ;Get block mask
007E 47                 MOV     B,A             ;Ready for AND operation
007F 7D                 MOV     A,L             ;Get LS byte of relative sector
0080 A0                 ANA     B               ;AND with block mask
0081 F5                 PUSH    PSW             ;A = sector displacement
0082 3A0200             LDA     DPBBS           ;Get block shift
0085 4F                 MOV     C,A             ;Make into counter

                GABS:                           ;Shift loop
0086 CDD000             CALL    SHLR            ;HL shifted right (divided by 2)
0089 0D                 DCR     C               ;Count down
008A C28600             JNZ     GABS            ;Shift again if necessary
008D F1                 POP     PSW             ;Recover offset
008E C9                 RET

                ;Utility subroutines
                ;These perform 16-bit arithmetic on the HL register pair.

                ;DIVHL
                ;Divides HL by DE using an iterative subtract.
                ;In practice, it uses an iterative ADD of the complemented divisor.

                ;Entry parameters
                ;       HL = dividend
                ;       DE = divisor

                ;Exit parameters
                ;       BC = quotient
                ;       HL = remainder

                DIVHL:
008F D5                 PUSH    D               ;Save divisor
                                                ;Note : 2's complement is formed by
                                                ;inverting all bits and adding 1.
0090 7B                 MOV     A,E             ;Complement divisor (for iterative
0091 2F                 CMA                     ;ADD later on)
0092 5F                 MOV     E,A
0093 7A                 MOV     A,D             ;Get MS byte
0094 2F                 CMA                     ;Complement it
0095 57                 MOV     D,A
0096 13                 INX     D               ;Make 2's complement
                                                ;Now, subtract negative divisor until
                                                ;dividend goes negative, counting the number
                                                ;of times the subtract occurs
0097 010000             LXI     B,0             ;Initialize quotient
                DIVHLS:                         ;Subtract loop
009A 03                 INX     B               ;Add 1 to quotient
009B 19                 DAD     D               ;"Subtract" divisor
009C DA9A00             JC      DIVHLS          ;Dividend not yet negative
                                                ;Dividend now negative, quotient 1 too large
009F 0B                 DCX     B               ;Correct quotient
                                                ;Compute correct remainder
00A0 EB                 XCHG                    ;DE = remainder - divisor
00A1 E1                 POP     H               ;Recover positive divisor
00A2 19                 DAD     D               ;HL = remainder
00A3 C9                 RET                     ;BC = quotient, HL = remainder

                ;MULHL
                ;Multiply HL * DE using iterative ADD.

                ;Entry parameters
                ;       HL = multiplicand
                ;       DE = multiplier

                ;Exit parameters
                ;       HL = product
                ;       DE = multiplier

                MULHL:
00A4 C5                 PUSH    B               ;Save user register
                                                ;Check if either multiplicand
                                                ; or multiplier is 0
```

Figure 5-24. (Continued)

```
00A5 7C               MOV     A,H
00A6 B5               ORA     L
00A7 CAC400           JZ      MULHLZ          ;Yes, fake product
00AA 7A               MOV     A,D
00AB B3               ORA     E
00AC CAC400           JZ      MULHLZ          ;Yes, fake product

                                              ;This routine will be faster if
                                              ; the smaller value is in DE
00AF 7A               MOV     A,D             ;Get MS byte of current DE value
00B0 BC               CMP     H               ;Check which is smaller
00B1 DAB500           JC      MULHLN          ;C set if D < H, so no exchange
00B4 EB               XCHG

              MULHLN:
00B5 42               MOV     B,D             ;BC = multiplier
00B6 4B               MOV     C,E
00B7 54               MOV     D,H             ;DE = HL = multiplicand
00B8 5D               MOV     E,L
00B9 0B               DCX     B               ;Adjust count as
                                              ;1 * multiplicand = multiplicand

              MULHLA:                         ;ADD loop
00BA 78               MOV     A,B             ;Check if all iterations completed
00BB B1               ORA     C
00BC CAC700           JZ      MULHLX          ;Yes, exit
00BF 19               DAD     D               ;HL = multiplicand + multiplicand
00C0 0B               DCX     B               ;Countdown on multiplier - 1
00C1 C3BA00           JMP     MULHLA          ;Loop back until all ADDs done

              MULHLZ:
00C4 210000           LXI     H,0             ;Fake product as either multiplicand
                                              ; or multiplier is 0

              MULHLX:
00C7 C1               POP     B               ;Recover user register
00C8 C9               RET

              ;SUBHL
              ;Subtract HL - DE

              ;Entry parameters
              ;       HL = subtrahend
              ;       DE = subtractor

              ;Exit parameters
              ;       HL = difference

              SUBHL:
00C9 7D               MOV     A,L             ;Get LS byte
00CA 93               SUB     E               ;Subtract without regard to carry
00CB 6F               MOV     L,A             ;Put back into difference
00CC 7C               MOV     A,H             ;Get MS byte
00CD 9A               SBB     D               ;Subtract including carry
00CE 67               MOV     H,A             ;Move back into difference
00CF C9               RET

              ;SHLR
              ;Shift HL right one place (dividing HL by 2)

              ;Entry parameters
              ;       HL = value to be shifted

              ;Exit parameters
              ;       HL = value/2

              SHLR:
00D0 B7               ORA     A               ;Clear carry
00D1 7C               MOV     A,H             ;Get MS byte
00D2 1F               RAR                     ;Bit 7 set from previous carry,
                                              ; bit 0 goes into carry
00D3 67               MOV     H,A             ;Put shift MS byte back
00D4 7D               MOV     A,L             ;Get LS byte
00D5 1F               RAR                     ;Bit 7 = bit 0 of MS byte
00D6 6F               MOV     L,A             ;Put back into result
00D7 C9               RET
```

Figure 5-24. (Continued)

Function 32: Set/Get User Number

Function Code: C = 20H
Entry Parameters: E = 0FFH to get user number, or
E = 0 to 15 to set user number
Exit Parameters: A = Current user number if E was 0FFH

Example

```
0020 =           B$SETGETUN        EQU     32       ;Set/Get User Number
0005 =           BDOS              EQU     5        ;BDOS entry point

                                                    ;To set user number
0000 0E20                  MVI     C,B$SETGETUN     ;Function code
0002 1E0F                  MVI     E,15             ;Required user number
0004 CD0500                CALL    BDOS             ;To get user number
0007 0E20                  MVI     C,B$SETGETUN     ;Function code
0009 1EFF                  MVI     E,0FFH           ;Indicate request to GET
000B CD0500                CALL    BDOS             ;A = Current user no. (0 -- 15)
```

Purpose This subroutine either sets or gets the current user number. The current user number determines which file directory entries are matched during all disk file operations.

When you call this function, the contents of the E register specify what action is to be taken. If E = 0FFH, then the function will return the current user number in the A register. If you set E to a number in the range 0 to 15 (that is, a valid user number), the function will set the current user number to this value.

Notes You can use this function to share files with other users. You can locate a file by attempting to open a file and switching through all of the user numbers. Or you can share a file in another user number by setting to that number, operating on the file, and then reverting back to the original user number.

If you do change the current user number, make provisions in your program to return to the original number before your program terminates. It is disconcerting for computer operators to find that they are in a different user number after a program. Files can easily be damaged or accidentally erased this way.

Function 33: Read Random

Function Code: C = 21H
Entry Parameters: DE = Address of FCB
Exit Parameters: A = Return code

Example

```
0021 =           B$READRAN         EQU     33       ;Read Random
0005 =           BDOS              EQU     5        ;BDOS entry point

                 FCB:                               ;File control block
0000 00          FCB$DISK:         DB      0        ;Search on default disk drive
0001 46494C454EFCB$NAME:           DB      'FILENAME'       ;File name
0009 545950      FCB$TYP:          DB      'TYP'    ;File type
```

```
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0       ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0       ;Random rec. overflow byte (MS)

0024 D204       RANRECNO:       DW      1234    ;Example random record number

                                                ;Record will be read into
                                                ; address set by prior
                                                ; SETDMA call
0026 2A2400             LHLD    RANRECNO        ;Get random record number
0029 222100             SHLD    FCB$RANREC      ;Set up file control block
002C 0E21               MVI     C,B$READRAN     ;Function code
002E 110000             LXI     D,FCB           ;DE -> file control block
0031 CD0500             CALL    BDOS            ;A = 00 if operation successful
                                                ;A  = nonzero if no data in
                                                ; file specifically:
                                                ;A = 01 -- attempt to read
                                                ;         unwritten record
                                                ;    03 -- CP/M could not
                                                ;         close current extent
                                                ;    04 -- attempt to read
                                                ;         unwritten extent
                                                ;    06 -- attempt to read
                                                ;         beyond end of disk
```

Purpose This function reads a specific CP/M record (128 bytes) from a random file—that is, a file in which records can be accessed directly. It assumes that you have already opened the file, set the DMA address using the BDOS Set DMA function, and set the specific record to be read into the random record number in the FCB. This function computes the extent of the specified record number and attempts to open it and read the correct CP/M record into the DMA address.

The random record number in the FCB is three bytes long (at relative bytes 33, 34, and 35). Byte 33 is the least significant byte, 34 is the middle byte, and 35 the most significant. CP/M uses only the most significant byte (35) for computing the overall file size (function 35). You must set this byte to 0 when setting up the FCB. Bytes 33 and 34 are used together for the Read Random, so you can access from record 0 to 65535 (a maximum file size of 8,388,480 bytes).

This function returns with A set to 0 to indicate that the operation has been completed successfully, or A set to a nonzero value if an error has occurred. The error codes are as follows:

A = 01 (attempt to read unwritten record)

A = 03 (CP/M could not close current extent)

A = 04 (attempt to read unwritten extent)

A = 06 (attempt to read beyond end of disk)

Unlike the Read Sequential BDOS function (code 20, 14H), which updates the current (sequential) record number in the FCB, the Read Random function leaves the record number unchanged, so that a subsequent Write Random will replace the record just read.

You can follow a Read Random with a Write Sequential (code 21, 15H). This will rewrite the record just read, but will then update the sequential record number. Or you may choose to use a Read Sequential after the Read Random. In this case, the same record will be reread and the sequential record number will be incremented. In short, the file can be sequentially read or written once the Read Random has been used to position to the required place in the file.

Notes

To use the Read Random function, you must first open the *base extent* of the file, that is, extent 0. Even though there may be no actual data records in this extent, opening permits the file to be processed correctly.

One problem that is not immediately obvious with random files is that they can easily be created with gaps in the file. If you were to create the file with record number 0 and record number 5000, there would be no intervening file extents. Should you attempt to read or copy the file sequentially, even using CP/M's file copy utility, only the first extent (and in this case, record 0) would get copied. A Read Sequential function would return an "end of file" error after reading record 0. You must therefore be conscious of the type of the file that you try and read.

See Figure 5-26 for an example subroutine that performs Random File Reads and Writes. It reads or writes records of sizes other than 128 bytes, where necessary reading or writing several CP/M records, prereading them into its own buffer when the record being written occupies only part of a CP/M record. It also contains subroutines to produce a 32-bit product from multiplying HL by DE (MLDL—Multiply double length) and a right bit shift for DE, HL (SDLR—Shift double length right).

Function 34: Write Random

Function Code: C = 22H
Entry Parameters: DE = Address of file control block
Exit Parameters: A = Return code

Example

```
0022 =          B$WRITERAN      EQU     34      ;Write Random
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;File name
0009 545950     FCB$TYP:        DB      'TYP'   ;File type
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0       ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0       ;Random rec. overflow byte (MS)

0024 D204       RANRECNO:       DW      1234    ;Example random record number

                                                ;Record will be written from
                                                ; address set by prior
                                                ; SETDMA call
```

```
0026 2A2400        LHLD    RANRECNO         ;Get random record number
0029 222100        SHLD    FCB$RANREC       ;Set up file control block
002C 0E22          MVI     C,B$WRITERAN     ;Function code
002E 110000        LXI     D,FCB            ;DE -> file control block
0031 CD0500        CALL    BDOS             ;A = 00 if operation successful
                                            ;A = nonzero if no data in file
                                            ; specifically:
                                            ;A = 03 -- CP/M could not
                                            ;        close current extent
                                            ;     05 -- directory full
                                            ;     06 -- attempt to write
                                            ;        beyond end of disk
```

Purpose

This function writes a specific CP/M record (128 bytes) into a random file. It is initiated in much the same way as the companion function, Read Random (code 33, 21H). It assumes that you have already opened the file, set the DMA address to the address in memory containing the record to be written to disk, and set the random record number in the FCB to the specified record being written. This function also computes the extent in which the specified record number lies and opens the extent (creating it if it does not already exist). The error codes returned in A by this call are the same as those for Read Random, with the addition of error code 05, which indicates a full directory.

Like the Read Random (but unlike the Write Sequential), this function does not update the logical extent and sequential (current) record number in the FCB. Therefore, any subsequent sequential operation will access the record just written by the Read Random call, but these functions will update the sequential record number. The Write Random can therefore be used to position to the required place in the file, which can then be accessed sequentially.

Notes

In order to use the Write Random, you must first open the base extent (extent 0) of the file. Even though there may be no data records in this extent, opening permits the file to be processed correctly.

As explained in the notes for the Read Random function, you can easily create a random file with gaps in it. If you were to create a file with record number 0 and record number 5000, there would be no intervening file extents.

Figure 5-25 shows an example subroutine that creates a random file (CRF) but avoids this problem. You specify the number of 128-byte CP/M records in the file. The subroutine creates the file and then writes zero-filled records throughout. This makes it easier to process the file and permits standard CP/M utility programs to copy the file because there is a data record in every logical record position in the file. It is no longer a "sparse" file.

Figure 5-26 shows a subroutine that ties the Read and Write Random functions together. It performs Random Operations (RO). Unlike the standard BDOS functions that operate on 128-byte CP/M records, RO can handle arbitrary record size from one to several thousand bytes. You specify the relative record number of your record, not the CP/M record number (RO computes this). RO also prereads a CP/M record when your logical record occupies part of a 128-byte record, either because your record is less than 128 bytes or because it spans more than one

```
                        ;CRF
                        ;Create random file
                        ;This subroutine creates a random file. It erases any previous
                        ;file before creating the new one, and then writes 0-filled
                        ;records throughout the entire file.

                        ;Entry parameters
                        ;       DE -> file control block for new file
                        ;       HL = Number of 128-byte CP/M records to be
                        ;               zero-filled.

                        ;Exit parameters
                        ;       Carry clear if operation successful (A = 0,1,2,3)
                        ;       Carry set if error (A = 0FFH)

                        ;Calling sequence
                        ;       LXI     D,FCB
                        ;       CALL    CRF
                        ;       JC      ERROR

0013 =                  B$ERASE         EQU     19      ;Erase file
0016 =                  B$CREATE        EQU     22      ;Create file
001A =                  B$SETDMA        EQU     26      ;Set DMA address
0015 =                  B$WRITESEQ      EQU     21      ;Write sequential record
0005 =                  BDOS            EQU     5       ;BDOS entry point

                        CRFBUF:                         ;Zero-filled buffer
0000 0000000000                 DW      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,
                                          0,0,0
0032 0000000000                 DW      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,
                                          0,0,0
0064 0000000000                 DW      0,0,0,0,0,0,0,0,0,0,0,0,0,0

0080 0000               CRFRC:  DW      0               ;Record count

                        CRF:
0082 228000                     SHLD    CRFRC           ;Save record count
0085 D5                         PUSH    D               ;Preserve FCB pointer
0086 0E13                       MVI     C,B$ERASE       ;Erase any existing file
0088 CD0500                     CALL    BDOS
008B D1                         POP     D               ;Recover FCB pointer
008C D5                         PUSH    D               ; and resave
008D 0E16                       MVI     C,B$CREATE      ;Create (and open new file)
008F CD0500                     CALL    BDOS
0092 FEFF                       CPI     0FFH            ;Carry set if OK, clear if error
0094 3F                         CMC                     ;Complete to use carry set if error
0095 D1                         POP     D               ;Recover FCB address
0096 D8                         RC                      ;Return if error
0097 D5                         PUSH    D               ;Resave FCB pointer

0098 0E1A                       MVI     C,B$SETDMA      ;Set DMA address to 0-buffer
009A 110000                     LXI     D,CRFBUF
009D CD0500                     CALL    BDOS
00A0 D1                         POP     D               ;Recover FCB pointer

                        CRFL:
00A1 2A8000                     LHLD    CRFRC           ;Get record count
00A4 7D                         MOV     A,L
00A5 B4                         ORA     H               ;Check if count now zero
00A6 C8                         RZ                      ;Yes, exit
00A7 2B                         DCX     H               ;Downdate count
00A8 228000                     SHLD    CRFRC           ;Save count
00AB D5                         PUSH    D               ;Resave FCB address
00AC 0E15                       MVI     C,B$WRITESEQ    ;Write sequentially
00AE CD0500                     CALL    BDOS

00B1 D1                         POP     D               ;Recover FCB
00B2 C3A100                     JMP     CRFL            ;Write next record
```

Figure 5-25. Create random file

128-byte sector. The subroutine suppresses this preread if you happen to use a record size that is some multiple of 128 bytes. In this case, your records will fit exactly onto a 128-byte record, so there will never be some partially occupied 128-byte sector.

This example also contains subroutines to produce a 32-bit product from multiplying HL by DE (MLDL—Multiply double length) and a right bit shift for DE, HL (SDLR—Shift double length right).

```
                        ;RO
                        ;Random operation (read or write)

                        ;This subroutine reads or writes a random record from a file.
                        ;The record length can be other than 128-bytes. This
                        ;subroutine computes the start CP/M record (which
                        ;is 128 bytes), and, if reading, performs a random read
                        ;and moves the user-specified record into a user buffer.
                        ;If necessary, more CP/M records will be read until the complete
                        ;user-specified record has been input.
                        ;For writing, if the size of the user-specified record is not an exact
                        ;multiple of CP/M records, the appropriate sectors will be preread.
                        ;It is not necessary to preread when the user-specified record
                        ;is an exact CP/M record, nor when subroutine is processing
                        ;CP/M records entirely spanned by a user-specified record.

                        ;Entry parameters
                        ;       HL -> parameter block of the form:
                        ;               DB      0       ;OFFH when reading, 00H for write
                        ;               DW      FCB     ;Pointer to FCB
                        ;               DW      RECNO   ;User record number
                        ;               DW      RECSZ   ;User record size
                        ;               DW      BUFFER  ;Pointer to buffer of
                        ;                               ; RECSZ bytes in length

                        ;Exit parameters
                        ;       A = 0 if operation completed (and user record
                        ;               copied into user buffer)
                        ;           1 if attempt to read unwritten CP/M record
                        ;           3 if CP/M could not close an extent
                        ;           4 if attempt to read unwritten extent
                        ;           5 if CP/M could not create a new extent
                        ;           6 if attempt to read beyond end of disk

                        ;Calling sequence
                        ;       LXI     H,PARAMS        ;HL -> parameter block
                        ;       CALL    RO
                        ;       ORA     A               ;Check if error
                        ;       JNZ     ERROR
0021 =                  FCBE$RANREC     EQU     33      ;Offset of random record no. in FCB
001A =                  B$SETDMA        EQU     26      ;Set the DMA address
0021 =                  B$READRAN       EQU     33      ;Read random record
0028 =                  B$WRITERANZ     EQU     40      ;Write random record with zero-fill
                                                        ; previously unallocated allocation
                                                        ; blocks
0005 =                  BDOS            EQU     5       ;BDOS entry point

                        ROPB:                           ;Parameter block image
0000 00                 ROREAD: DB      0               ;NZ when reading, Z when writing
0001 0000               ROFCB:  DW      0               ;Pointer to FCB
0003 0000               ROURN:  DW      0               ;User record number
0005 0000               ROURL:  DW      0               ;User record length
0007 0000               ROUB:   DW      0               ;Pointer to user buffer
0009 =                  ROPBL   EQU     $-ROPB          ;Parameter block length

0009 0000               ROFRP:  DW      0               ;Pointer to start of user record fragment
                                                        ; in first CP/M-record read in
```

Figure 5-26. Read/Write variable length records randomly

```
000B 00         ROFRL:  DB      0               ;Fragment length
000C 0000       RORNP:  DW      0               ;Record number pointer (in user FCB)
000E 00         ROWECR: DB      0               ;NZ when writing user records that are an
                                                ; exact super-multiple of CP/M-record (and
                                                ; therefore no preread is required)

000F            ROBUF:  DS      128             ;Buffer for CP/M record

                RO:
008F 110000             LXI     D,ROPB          ;DE -> local parameter block
0092 0E09               MVI     C,ROPBL         ;Parameter block length
0094 CDFE01             CALL    MOVE            ;Move C bytes from HL to DE

                        ;To compute offset of user record in CP/M record,
                        ; compute the relative BYTE offset of the start
                        ; of the user record within the file (i.e.
                        ; user record number * record size). The least
                        ; significant 7 bits of this product give the
                        ; byte offset of the start of the user record.
                        ;The product / 128 (shifted left 7 bits) gives the
                        ;CP/M record number of the start of the user record.

0097 2A0500             LHLD    ROURL           ;Get user record length
009A 7D                 MOV     A,L             ;Get LS bytes of user rec. length
009B E67F               ANI     7FH             ;Check if exact multiple of 128
009D B7                 ORA     A               ;(i.e. exact CP/M records)
009E 3E00               MVI     A,0             ;A = 0, flags unchanged
00A0 C2A400             JNZ     RONE            ;Not exact CP/M records
00A3 3D                 DCR     A               ;A =FF

                RONE:
00A4 320E00             STA     ROWECR          ;Set write-exact-CP/M-records flag
00A7 EB                 XCHG                    ;DE = user record length
00A8 2A0300             LHLD    ROURN           ;Get user record number
00AB CDB801             CALL    MLDL            ;DE,HL = HL * DE
                                                ;DE,HL = user-record byte offset in file
00AE D5                 PUSH    D               ;Save user-record byte offset
00AF E5                 PUSH    H
00B0 7D                 MOV     A,L             ;Get LS byte of product
00B1 E67F               ANI     7FH             ;Isolate byte offset within

00B3 4F                 MOV     C,A             ;CP/M record
00B4 0600               MVI     B,0             ;Make into word value
00B6 210F00             LXI     H,ROBUF         ;Get base address of local buffer
00B9 09                 DAD     B               ;HL -> Start of fragment in buffer
00BA 220900             SHLD    ROFRP           ;Save fragment pointer

                        ;Compute maximum fragment length that could reside in
                        ;remainder of CP/M record, based on the offset in the
                        ;CP/M record where the fragment starts.

00BD 47                 MOV     B,A             ;Take copy of offset in CP/M record
00BE 3E80               MVI     A,128           ;CP/M record size
00C0 90                 SUB     B               ;Compute 128 - offset
00C1 320B00             STA     ROFRL           ;Assume this is the fragment length

                        ;If the user record length is less than the assumed
                        ; fragment length, use it in place of the result above

00C4 47                 MOV     B,A             ;Get copy of assume frag. length
00C5 3A0600             LDA     ROURL+1         ;Get MS byte of user record length
00C8 B7                 ORA     A               ;If NZ, rec. len. must be > 128
00C9 C2D600             JNZ     ROFLOK          ;So fragment length is OK
00CC 3A0500             LDA     ROURL           ;Still a chance that rec. len.
00CF B8                 CMP     B               ; less than fragment len.
00D0 D2D600             JNC     ROFLOK          ;NC if user rec. len. => frag. len.
00D3 320B00             STA     ROFRL           ;User rec. len. < frag. len. so
                                                ; reset fragment length to smaller

                ROFLOK:
00D6 3A0E00             LDA     ROWECR          ;Get exact CP/M record flag
00D9 47                 MOV     B,A             ;for ANDing with READ flag
00DA 3A0000             LDA     ROREAD          ;Get read operation flag
00DD 2F                 CMA                     ;Invert so NZ when writing
```

Figure 5-26. (Continued)

```
00DE A0                 ANA     B               ;Form logical AND
00DF 320E00             STA     ROWECR          ;Save back in flag

                        ;Recover the double length byte offset within the file
                        ;of the start of the user record. Shift 7 places right
                        ;to divide by 128 and get the CP/M record number for
                        ;the start of the user record.

00E2 E1                 POP     H               ;Recover user rec. byte offset
00E3 D1                 POP     D
00E4 0E07               MVI     C,7             ;Count for shift right

                ROS:
00E6 CDF101             CALL    SDLR            ;DE,HL = DE,HL / 2
00E9 0D                 DCR     C
00EA C2E600             JNZ     ROS

00ED 7A                 MOV     A,D             ;Error if DE still NZ after
00EE B3                 ORA     E               ; division by 128.
00EF C2AC01             JNZ     ROERO

                                                ;Set CP/M record number in FCB
00F2 EB                 XCHG                    ;DE = CP/M record number
00F3 2A0100             LHLD    ROFCB           ;Get pointer to FCB
00F6 012100             LXI     B,FCBE$RANREC   ;Offset of random record no. in FCB
00F9 09                 DAD     B               ;HL -> ran. rec. no. in FCB
00FA 220C00             SHLD    RORNP           ;Save record number pointer
00FD 73                 MOV     M,E             ;Store LS byte
00FE 23                 INX     H
00FF 72                 MOV     M,D             ;Store MS byte

0100 0E1A               MVI     C,B$SETDMA      ;Set DMA address to local buffer
0102 110F00             LXI     D,ROBUF
0105 CD0500             CALL    BDOS

0108 3A0E00             LDA     'ROWECR         ;Bypass preread if exact sector write
010B B7                 ORA     A
010C C21F01             JNZ     ROMNF

010F 2A0100             LHLD    ROFCB           ;Get pointer to FCB
0112 EB                 XCHG                    ;DE -> FCB
0113 0E21               MVI     C,B$READRAN     ;Read random function
0115 CD0500             CALL    BDOS

0118 FE05               CPI     5               ;Check if error code < 5
011A DCAF01             CC      ROCIE           ;Yes, check if ignorable error
                                                ; (i.e. error reading unwritten part
                                                ; of file for write operation preread)
011D B7                 ORA     A               ;Check if error
011E C0                 RNZ                     ;Yes

                ROMNF:                          ;Move next fragment
011F 2A0700             LHLD    ROUB            ;Get pointer to user buffer
0122 EB                 XCHG                    ;DE -> user buffer
0123 2A0900             LHLD    ROFRP           ;HL -> start of user rec. in local buffer
0126 3A0B00             LDA     ROFRL           ;Get fragment length
0129 4F                 MOV     C,A             ;Ready for MOVE

012A 3A0000             LDA     ROREAD          ;Check if reading
012D B7                 ORA     A
012E C23201             JNZ     RORD1           ;Yes, so leave DE, HL unchanged
0131 EB                 XCHG                    ;Writing, so swap source and destination
                                                ;DE -> start of user rec. in local buffer
                                                ;HL -> user buffer

                RORD1:
0132 CDFE01             CALL    MOVE            ;Reading - fragment local -> user buffer
                                                ;Writing - fragment user -> local buffer
0135 3A0000             LDA     ROREAD          ;Check if writing
0138 B7                 ORA     A
0139 CA3D01             JZ      ROWR1           ;Writing, so leave HL -> user buffer
013C EB                 XCHG                    ;HL -> next byte in user buffer

                ROWR1:
013D 220700             SHLD    ROUB            ;Save updated user buffer pointer
0140 3A0000             LDA     ROREAD          ;Check if reading
```

Figure 5-26. (Continued)

```
0143 B7                 ORA     A
0144 C25001             JNZ     RORD3           ;Yes, bypass write code

0147 0E28               MVI     C,B$WRITERANZ   ;Write random
0149 2A0100             LHLD    ROFCB           ;Get address of FCB
014C EB                 XCHG                    ;DE -> FCB
014D CD0500             CALL    BDOS

                RORD3:  ;Compute residual length of user record as yet unmoved.
                        ;If necessary (because more data needs to be transferred)
                        ;more CP/M records will be read. In this case
                        ;the start of the fragment will be offset 0. The fragment
                        ;length depends on whether the user record finishes within
                        ;the next sector or spans it. If the residual length  of the
                        ;user record is > 128, the fragment length will be set to
                        ;128.

0150 2A0500             LHLD    ROURL           ;Get residual user rec. length
0153 3A0B00             LDA     ROFRL           ;Get fragment length just moved
0156 5F                 MOV     E,A             ;Make into a word value
0157 1600               MVI     D,0
0159 CDEA01             CALL    SUBHL           ;Compute ROURL - ROFRL
015C 7C                 MOV     A,H             ;Check if result 0
015D B5                 ORA     L
015E C8                 RZ                      ;Return when complete USER
                                                ; record has been transferred
015F 220500             SHLD    ROURL           ;Save downdated residual rec. length
0162 4D                 MOV     C,L             ;Assume residual length < 128
0163 118000             LXI     D,128           ;Check if residual length is < 128
0166 CDEA01             CALL    SUBHL           ;HL = HL - DE
0169 FA6E01             JM      ROLT128         ;negative if < 128
016C 0E80               MVI     C,128           ;=> 128, so set frag.length to 128

                ROLT128:
016E 79                 MOV     A,C
016F 320B00             STA     ROFRL           ;Fragment length now is either 128
                                                ; if more than 128 bytes left to input
                                                ; in user record, or just the right
                                                ; number of bytes (< 128) to complete
                                                ; the user record.
0172 210F00             LXI     H,ROBUF         ;All subsequent CP/M records will start
0175 220900             SHLD    ROFRP           ; at beginning of buffer

                                                ;Update random record number in FCB
0178 2A0C00             LHLD    RORNP           ;HL -> random record number in  user FCB
017B 5E                 MOV     E,M             ;Increment the random record number
017C 23                 INX     H               ;HL -> MS byte of record number
017D 56                 MOV     D,M             ;Get MS byte
017E 13                 INX     D               ;Update record number itself
017F 7A                 MOV     A,D             ;Check if record now 0
0180 B3                 ORA     E
0181 C28701             JNZ     ROSRN           ;No, so save record number
0184 3E06               MVI     A,6             ;Indicate "seek past end of disk"
0186 C9                 RET                     ;Return to user
                ROSRN:
0187 72                 MOV     M,D             ;Save record number
0188 2B                 DCX     H               ;HL -> LS byte
0189 73                 MOV     M,E

                                                ;If writing, check if preread required
018A 3A0E00             LDA     ROWECR          ;Check if exact CP/M record write
018D B7                 ORA     A
018E C21F01             JNZ     ROMNF           ;Yes, go move next fragment

0191 3A0000             LDA     ROREAD          ;If reading, perform read unconditionally
0194 B7                 ORA     A
0195 C2A001             JNZ     RORD2

0198 3A0B00             LDA     ROFRL           ;For writes, bypass preread if
019B FE80               CPI     128             ; whole CP/M-record is to be overwritten
019D CA1F01             JZ      ROMNF           ; (fragment length = 128)

                RORD2:
01A0 0E21               MVI     C,B$READRAN     ;Read the next CP/M record
01A2 2A0100             LHLD    ROFCB           ; in sequence
```

Figure 5-26. (Continued)

```
01A5 EB                 XCHG                    ;DE -> FCB
01A6 CD0500             CALL    BDOS
01A9 C31F01             JMP     ROMNF           ;Go back to move next fragment

                ROERO:                          ;Error because user record number
                                                ; * User record length / 128 gives
                                                ; a CP/M record number > 65535.
01AC 3E04               MVI     A,4             ;Indicate "attempt to read unwritten
01AE C9                 RET                     ; extent"

                ROCIE:                          ;Check ignorable error (preread
                                                ; for write operation)
01AF 47                 MOV     B,A             ;Save original error code
01B0 3A0000             LDA     ROREAD          ;Check if read operation
01B3 B7                 ORA     A
01B4 78                 MOV     A,B             ;Restore original error code but
                                                ; leave flags unchanged
01B5 C0                 RNZ                     ;Return if reading
01B6 AF                 XRA     A               ;Fake "no error" indicator
01B7 C9                 RET

                ;MLDL
                ;Multiply HL * DE using iterative ADD with product
                ;returned in DE,HL.

                ;Entry parameters
                ;       HL = multiplicand
                ;       DE = multiplier

                ;Exit parameters
                ;       DE,HL = product
                ;       DE = multiplier

                MLDL:
01B8 010000             LXI     B,0             ;Put 0 on top of stack
01BB C5                 PUSH    B               ; to act as MS byte of product
                                                ;Check if either multiplicand
                                                ; or multiplier is 0
01BC 7C                 MOV     A,H
01BD B5                 ORA     L
01BE CAE501             JZ      MLDLZ           ;Yes, fake product
01C1 7A                 MOV     A,D
01C2 B3                 ORA     E
01C3 CAE501             JZ      MLDLZ           ;Yes, fake product

                                                ;This routine will be faster if
                                                ; the smaller value is in DE
01C6 7A                 MOV     A,D             ;Get MS byte of current DE value
01C7 BC                 CMP     H               ;Check which is smaller
01C8 DACC01             JC      MLDLNX          ;C set if D < H, so no exchange
01CB EB                 XCHG
                MLDLNX:
01CC 42                 MOV     B,D             ;BC = multiplier
01CD 4B                 MOV     C,E

01CE 54                 MOV     D,H             ;DE = HL = multiplicand
01CF 5D                 MOV     E,L

01D0 0B                 DCX     B               ;Adjust count as
                                                ; 1 * multiplicand = multiplicand
                MLDLA:                          ;ADD loop
01D1 78                 MOV     A,B             ;Check if all iterations completed
01D2 B1                 ORA     C
01D3 CAE801             JZ      MLDLX           ;Yes, exit
01D6 19                 DAD     D               ;HL = multiplicand + multiplicand
01D7 E3                 XTHL                    ;HL = MS bytes of result, TOS = part prod.
01D8 7D                 MOV     A,L             ;Get LS byte of top half of product
01D9 CE00               ACI     0               ;Add one if carry set
01DB 6F                 MOV     L,A             ;Replace
01DC 7C                 MOV     A,H             ;Repeat for MS byte
01DD CE00               ACI     0
01DF 67                 MOV     H,A
01E0 E3                 XTHL
01E1 0B                 DCX     B               ;Countdown on multiplier - 1
01E2 C3D101             JMP     MLDLA           ;Loop back until all ADDs done
```

Figure 5-26. (Continued)

```
                MLDLZ:
01E5 210000             LXI     H,0             ;Fake product as either multiplicand
                                                ; or multiplier is 0

                MLDLX:
01E8 D1                 POP     D               ;Recover MS part of product
01E9 C9                 RET

                ;SUBHL
                ;Subtract HL - DE.

                ;Entry parameters
                ;       HL = subtrahend
                ;       DE = subtractor

                ;Exit parameters
                ;       HL = difference

                SUBHL:
01EA 7D                 MOV     A,L             ;Get LS byte
01EB 93                 SUB     E               ;Subtract without regard to carry
01EC 6F                 MOV     L,A             ;Put back into difference
01ED 7C                 MOV     A,H             ;Get MS byte
01EE 9A                 SBB     D               ;Subtract including carry
01EF 67                 MOV     H,A             ;Move back into difference
01F0 C9                 RET

                ;SDLR
                ;Shift DE,HL right one place (dividing DE,HL by 2)

                ;Entry parameters
                ;       DE,HL = value to be shifted

                ;Exit parameters
                ;       DE,HL = value / 2

                SDLR:
01F1 B7                 ORA     A               ;Clear carry
01F2 EB                 XCHG                    ;Shift DE first
01F3 CDF701             CALL    SDLR2
01F6 EB                 XCHG                    ;Now shift HL

                                                ;Drop into SDLR2 with carry
                                                ; set correctly from LS bit
                                                ; of DE
                SDLR2:                          ;Shift HL right one place
01F7 7C                 MOV     A,H             ;Get MS byte
01F8 1F                 RAR                     ;Bit 7 set from previous carry,
                                                ;Bit 0 goes into carry
01F9 67                 MOV     H,A             ;Put shift MS byte back
01FA 7D                 MOV     A,L             ;Get LS byte
01FB 1F                 RAR                     ;Bit 7 = bit 0 of MS byte
01FC 6F                 MOV     L,A             ;Put back into result
01FD C9                 RET

                ;MOVE
                ;Moves C bytes from HL to DE

                MOVE:
01FE 7E                 MOV     A,M             ;Get source byte
01FF 12                 STAX    D               ;Store in destination
0200 13                 INX     D               ;Update destination pointer
0201 23                 INX     H               ;Update source pointer
0202 0D                 DCR     C               ;Downdate count
0203 C2FE01             JNZ     MOVE            ;Get next byte
0206 C9                 RET
```

Figure 5-26. (Continued)

Function 35: Get File Size

Function Code: C = 23H
Entry Parameters: DE = Address of FCB
Exit Parameters: Random record field set in FCB

Example

```
0023 =          B$GETFSIZ       EQU     35      ;Get Random File LOGICAL size
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;File name
0009 545950     FCB$TYP:        DB      'TYP'   ;File type
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0       ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0       ;Random rec. overflow byte (MS)

0024 0E23               MVI     C,B$GETFSIZ     ;Function code
0026 110000             LXI     D,FCB           ;DE -> file control block
0029 CD0500             CALL    BDOS
002C 2A2100             LHLD    FCB$RANREC      ;Get random record number
                                                ;HL = LOGICAL file size
                                                ; i.e. the record number of the
                                                ; last record
```

Purpose This function returns the virtual size of the specified file. It does so by setting the random record number (bytes 33-35) in the specified FCB to the maximum 128-byte record number in the file. The virtual file size is calculated from the record address of the record following the end of the file. Bytes 33 and 34 form a 16-bit value that contains the record number, with overflow indicated in byte 35. If byte 35 is 01, this means that the file has the maximum record count of 65,536.

If the function cannot find the file specified by the FCB, it returns with the random record field set to 0.

You can use this function when you want to add data to the end of an existing file. By calling this function first, the random record bytes will be set to the end of file. Subsequent Write Random calls will write out records to this preset address.

Notes Do not confuse the virtual file size with the actual file size. In a random file, if you write just a single CP/M record to record number 1000 and then call this function, it will return with the random record number field set in the FCB to 1000—even though only a single record exists in the file.

For sequential files, this function returns the number of records in the file. In this case, the virtual and actual file sizes coincide.

Function 36: Set Random Record Number

Function Code: C = 24H
Entry Parameters: DE = Address of FCB
Exit Parameters: Random record field set in FCB

Example

```
0024 =          B$SETRANREC     EQU     36      ;Set Random Record Number
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;File name
0009 545950     FCB$TYP:        DB      'TYP'   ;File type
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0       ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0       ;Random rec. overflow byte (MS)

                                                ;... file opened and read
                                                ; or written sequentially...

0024 0E24               MVI     C,B$SETRANREC   ;Function code
0026 110000             LXI     D,FCB           ;DE -> file control block
0029 CD0500             CALL    BDOS
002C 2A2100             LHLD    FCB$RANREC      ;Get random record number
                                                ;HL = random record number
                                                ; that corresponds to the
                                                ; sequential progress down
                                                ; the file.
```

Purpose This function sets the random record number in the FCB to the correct value for the last record read or written sequentially to the file.

Notes This function provides you with a convenient way to build an index file so that you can randomly access a sequential file. Open the sequential file, and as you read each record, extract the appropriate key field from the data record. Make the BDOS Set Random Record request and create a new data record with just the key field and the random record number. Write the new data record out to the index file.

Once you have done this for each record in the file, your index file provides a convenient method, given a search key value, of finding the appropriate CP/M record in which the data lies.

You can also use this function as a means of finding out where you are currently positioned in a sequential file—either to relate a CP/M record number to the position, or simply as a place-marker to allow a repositioning to the same place later.

Function 37: Reset Logical Disk Drive

Function Code: C = 25H
Entry Parameters: DE = Logical drive bit map
Exit Parameters: A = 00H

Example

```
0025 =          B$RESETD        EQU     37      ;Reset Logical Disks
0005 =          BDOS            EQU     5       ;BDOS entry point
```

```
                                                ;DE = Bit map of disks to be
                                                ; reset
                                                ;Bits are = 1 if disk to be
                                                ; reset
                                                ;Bits 15 14 13 ... 2 1 0
                                                ;Disk  P  O  N ... C B A
0000 110200         LXI     D,0000$0000$0000$0010B  ;Reset drive B:
0003 0E25           MVI     C,B$RESETD          ;Function code
0005 CD0500         CALL    BDOS
```

Purpose This function resets individual disk drives. It is a more precise version of the Reset Disk System function (code 13,0DH), in that you can set specific logical disks rather than all of them.

The bit map in DE shows which disks are to be reset. The least significant bit of E represents disk A, and the most significant bit of D, disk P. The bits set to 1 indicate the disks to be reset.

Note that this function returns a zero value in A in order to maintain compatibility with MP/M.

Notes Use this function when only specific diskettes need to be changed. Changing a diskette without requesting CP/M to log it in will cause the BDOS to assume that an error has occurred and to set the new diskette to Read-Only status as a protective measure.

Function 40: Write Random with Zero-fill

Function Code: C = 28H
Entry Parameters: DE = Address of FCB
Exit Parameters: A = Return Code

Example

```
0028 =          B$WRITERANZ     EQU     40      ;Write Random with Zero-Fill
0005 =          BDOS            EQU     5       ;BDOS entry point

                FCB:                            ;File control block
0000 00         FCB$DISK:       DB      0       ;Search on default disk drive
0001 46494C454EFCB$NAME:        DB      'FILENAME'      ;File name
0009 545950     FCB$TYP:        DB      'TYP'   ;File type
000C 00         FCB$EXTENT:     DB      0       ;Extent
000D 0000       FCB$RESV:       DB      0,0     ;Reserved for CP/M
000F 00         FCB$RECUSED:    DB      0       ;Records used in this extent
0010 0000000000FCB$ABUSED:      DB      0,0,0,0,0,0,0,0 ;Allocation blocks used
0018 0000000000                 DB      0,0,0,0,0,0,0,0
0020 00         FCB$SEQREC:     DB      0       ;Sequential rec. to read/write
0021 0000       FCB$RANREC:     DW      0       ;Random rec. to read/write
0023 00         FCB$RANRECO:    DB      0       ;Random rec. overflow byte (MS)

0024 D204       RANRECNO:       DW      1234    ;Example random record number

                                                ;Record will be written from
                                                ; address set by prior
                                                ; SETDMA call
0026 2A2400         LHLD    RANRECNO            ;Get random record number
0029 222100         SHLD    FCB$RANREC          ;Set up file control block
002C 0E28           MVI     C,B$WRITERANZ       ;Function code
002E 110000         LXI     D,FCB               ;DE -> file control block
0031 CD0500         CALL    BDOS                ;A = 00 if operation successful
```

```
;A = nonzero if no data in file
; specifically :
;A = 03 -- CP/M could not
;          close current extent
;     05 -- directory full
;     06 -- attempt to write
;          beyond end of disk
```

Purpose This function is an extension to the Write Random function described previously. In addition to performing the Write Random, it will also fill each new allocation block with 00H's. Digital Research added this function to assist Microsoft with the production of its COBOL compiler—it makes the logic of the file handling code easier. It also is an economical way to completely fill a random file with 00H's. You need only write one record per allocation block; the BDOS will clear the rest of the block for you.

Notes Refer to the description of the Write Random function (code 34).

The BIOS Components
The BIOS Entry Points
Bootstrap Functions
Character Input/Output Functions
Disk Functions
Calling the BIOS Functions Directly
Example BIOS

6

The Basic Input/Output System

This chapter takes a closer look at the Basic Input/Output System (BIOS). The BIOS provides the software link between the Console Command Processor (CCP), the Basic Disk Operating System (BDOS), and the physical hardware of your computer system. The CCP and BDOS interact with the parts of your computer system only as logical devices. They can therefore remain unchanged from one computer system to the next. The BIOS, however, is customized for your particular type of computer and disk drives. The only predictable part of the BIOS is the way in which it interfaces to the CCP and BDOS. This must remain the same no matter what special features are built into the BIOS.

The BIOS Components

A standard BIOS consists of low-level subroutines that drive four types of physical devices:

- Console: CP/M communicates with the outside world via the console. Normally this will be a video terminal or a hard-copy terminal.
- "Reader" and "punch": These devices are normally used to communicate between computer systems — the names "reader" and "punch" are just historical relics from the early days of CP/M.
- List: This is a hard-copy printer, either letter-quality or dot-matrix.
- Disk drives: These can be anything from the industry standard single-sided, single-density, 8-inch floppy diskette drives to hard disk drives with capacities of several hundred megabytes.

The BIOS Entry Points

The first few instructions of the BIOS are all jump (JMP) instructions. They transfer control to the 17 different subroutines in the BIOS. The CCP and the BDOS, when making a specific request of the BIOS, do so by transferring control to the appropriate JMP instruction in this BIOS *jump table* or *jump vector.* The BIOS jump vector always starts at the beginning of a 256-byte page, so the address of the first jump instruction is always of the form xx00H, where "xx" is the page address. Location 0000H to 0002H has a jump instruction to the second entry of the BIOS jump vector — so you can always find the page address of the jump vector by looking in location 0002H.

Figure 6-1 shows the contents of the BIOS jump vector along with the page-relative address of each jump. The labels used in the jump instructions have been adopted by convention.

The following sections describe the functions of each of the BIOS's main subroutines. You should also refer to Digital Research's manual *CP/M 2.0 Alteration Guide* for their description of the BIOS routines.

Bootstrap Functions

There are two bootstrap functions. The cold bootstrap loads the entire CP/M operating system when the system is either first turned on or reset. The warm bootstrap reloads the CCP whenever a program branches to location 0000H.

```
xx00H   JMP  BOOT     ;"Cold" (first time) bootstrap
xx03H   JMP  WBOOT    ;"Warm" bootstrap
xx06H   JMP  CONST    ;Console input status
xx09H   JMP  CONIN    ;Console input
xx0CH   JMP  CONOUT   ;Console output
xx0FH   JMP  LIST     ;List output
xx12H   JMP  PUNCH    ;"Punch" output
xx15H   JMP  READER   ;"Reader" input
xx18H   JMP  HOME     ;Home disk heads (to track 0)
xx1BH   JMP  SELDSK   ;Select logical disk
xx1EH   JMP  SETTRK   ;Set track number
xx21H   JMP  SETSEC   ;Set sector number
xx24H   JMP  SETDMA   ;Set DMA address
xx27H   JMP  READ     ;Read (128-byte) sector
xx2AH   JMP  WRITE    ;Write (128-byte) sector
xx2DH   JMP  LISTST   ;List device output status
xx30H   JMP  SECTRAN  ;Sector translate
```

Figure 6-1. Layout of the standard BIOS jump vector

BOOT: "Cold" Bootstrap

The BOOT jump instruction is the first instruction executed in CP/M. The bootstrap sequence must transfer control to the BOOT entry point in order to bring up CP/M. In general, a PROM receives control either when power is first applied or after you press the RESET button on the computer. This reads in the CP/M loader on the first sector of the physical disk drive chosen to be logical disk A. This CP/M loader program reads the binary image of the CCP, BDOS, and BIOS into memory at some predetermined address. Then it transfers control to the BOOT entry point in the BIOS jump vector.

This BOOT routine must initialize all of the required computer hardware. It sets up the baud rates for the physical console (if this has not already been done during the bootstrap sequence), the "reader," "punch," and list devices, and the disk controller. It must also set up the base page of memory so that there is a jump at location 0000H to the warm boot entry point in the BIOS jump vector (at xx03H) and a jump at location 0005H to the BDOS entry point.

Most BOOT routines sign on by displaying a short message on the console, indicating the current version of CP/M and the computer hardware that this BIOS can support.

The BOOT routine terminates by transferring control to the start of the CCP + 6 bytes (the CCP has its own small jump vector at the beginning). Just before the BOOT routine jumps into the CCP, it sets the C register to 0 to indicate that logical disk A is to be the default disk drive. This is what causes "A>" to be the CCP's initial prompt.

The actual CCP entry point is derived from the base address of the BIOS. The CCP and BDOS together require 1E00H bytes of code, so the first instruction of the CCP starts at BIOS − 1E00H.

WBOOT: "Warm" Bootstrap

Unlike the "cold" bootstrap entry point, which executes only once, the WBOOT or warm boot routine will be executed every time a program terminates by jumping to location 0000H, or whenever you type a CONTROL-C on the console as the first character of an input line.

The WBOOT routine is responsible for reloading the CCP into memory. Programs often use all of memory up to the starting point of the BDOS, overwriting the CCP in the process. The underlying philosophy is that while a program is executing, the CCP is not needed, so the program can use the memory previously occupied by the CCP. The CCP occupies 800H (2048) bytes of memory—and this is frequently just enough to make the difference between a program that cannot run and one that can.

A few programs that are self-contained and do not require the BDOS's facilities will also overwrite the BDOS to get another 1600H (5632) bytes of memory. Therefore, to be really safe, the WBOOT routine should read in both the CCP and the BDOS. It also needs to set up the two JMPs at location 0000H (to WBOOT itself) and at location 0005H (to the BDOS). Location 0003H should be set to the initial value of the IOBYTE if this is implemented in the BIOS.

As its last act, the WBOOT routine sets register C to indicate which logical disk is to be selected (C = 0 for A, 1 for B, and so on). It then transfers control into the CCP at the first instruction in order to restart the CCP. Again, the actual address is computed based on the knowledge that the CCP starts 1E00H bytes lower in memory than the base address of the BIOS.

Character Input/Output Functions

Character input/output functions deal with logical devices: the console, "reader," "punch," and list devices. Because these logical devices can in practice be connected by software to one of several physical character I/O devices, many BIOS's use CP/M's IOBYTE features to assign logical devices to physical ones.

In this case, each of the BIOS functions must check the appropriate bit fields of the IOBYTE (see Figure 4-2 and Table 4-1) to transfer control to the correct physical device *driver* (program that controls a physical device).

CONST: Console Input Status

CONST simply returns an indicator showing whether there is an incoming character from the console device. The convention is that A = 0FFH if a character is waiting to be processed, A = 0 if one is not. Note that the zero flag need not be set to reflect the contents of the A register—it is the contents that are important.

CONST is called by the CCP whenever the CCP is in the middle of an operation that can be interrupted by pressing a keyboard character.

The BDOS will call CONST if a program makes a Read Console Status function call (B$CONST, code 11, 0BH). It is also called by the console input BIOS routine, CONIN (described next).

CONIN: Console Input

CONIN reads the next character from the console to the A register and sets the most significant (parity) bit to 0.

Normally, CONIN will call the CONST routine until it detects A = 0FFH. Only then will it input the data character and mask off the parity bit.

CONIN is called by the CCP and by the BDOS when a program executes a Read Console Byte function (B$CONIN, code 1).

CONOUT: Console Output

CONOUT outputs the character (in ASCII) in register C to the console. The most significant (parity) bit of the character will always be 0.

CONOUT must first check that the console device is ready to receive more data, delaying if necessary until it is, and only then sending the character to the device.

CONOUT is called by the CCP and by the BDOS when a program executes a Write Console Byte function (B$CONOUT, code 2).

LIST: List Output

LIST is similar to CONOUT except that it sends the character in register C to the list device. It too checks first that the list device is ready to receive the character.

LIST is called by the CCP in response to the CONTROL-P toggle for printer echo of console output, and by the BDOS when a program makes a Write Printer Byte or Display String call (B$LISTOUT and B$PRINTS, codes 5 and 9).

PUNCH: "Punch" Output

PUNCH sends the character in register C to the "punch" device. As mentioned earlier, the "punch" is rarely a real paper tape punch. In most BIOS's, the PUNCH entry point either returns immediately and is effectively a null routine, or it outputs the character to a communications device, such as a modem, on your computer.

PUNCH must check that the "punch" device is indeed ready to accept another character for output, and must wait if it is not.

Digital Research's documentation states that the character to be output will always have its most significant bit set to 0. This is not true. The BDOS simply transfers control over to the PUNCH entry point in the BIOS; the setting of the most significant bit will be determined by the program making the BDOS function request (B$PUNOUT, code 4). This is important because the requirement of a zero would preclude being able to send pure binary data via the BIOS PUNCH function.

READER: "Reader" Input

As with the PUNCH entry point, the READER entry point rarely connects to a real paper tape reader.

The READER function must return the next character from the reader device in the A register, waiting, if need be, until there is a character.

Digital Research's documentation again says that the most significant bit of the A register must be 0, but this is not the case if you wish to receive pure binary information via this function.

READER is called whenever a program makes a Read "Reader" Byte function request (B$READIN, code 3).

Disk Functions

All of the disk functions that follow were originally designed to operate on the 128-byte sectors used on single-sided, single-density, 8-inch floppy diskettes that were standard in the industry at the time. Now that CP/M runs on many different types of disks, some of the BIOS disk functions seem strange because most of the new disk drives use sector sizes other than 128 bytes.

To handle larger sector sizes, the BIOS has some additional code that makes the BDOS respond as if it were still handling 128-byte sectors. This code is referred to as the *blocking/deblocking* code. As its name implies, it blocks together several 128-byte "sectors" and only writes to the disk when a complete *physical* sector has been assembled. When reading, it reads in a physical sector and then deblocks it, handing back several 128-byte "sectors" to the BDOS.

To do all of this, the blocking/deblocking code uses a special buffer area of the same size as the physical sectors on the disk. This is known as the host disk buffer or HSTBUF. Physical sectors are read into this buffer and written to the disk from it.

In order to optimize this blocking/deblocking routine, the BIOS has code in it to reduce the number of times that an actual disk read or write occurs. A side effect is that at any given moment, several 128-byte "sectors" may be stored in the HSTBUF, waiting to be written out to the disk when HSTBUF becomes full. This sometimes complicates the logic of the BIOS disk functions. You cannot simply select a new disk drive, for example, when the HSTBUF contains data destined for another disk drive. You will see this complication in the BIOS only in the form of added logical operations; the BIOS disk functions rarely trigger immediate physical operations. It is easier to understand these BIOS functions if you consider that they make *requests*—and that these requests are satisfied only when it makes sense to do so, taking into account the blocking/deblocking logic.

HOME: Home Disk

HOME sets the requested track and sector to 0.

SELDSK: Select Disk

SELDSK does not do what its name implies. It does not (and must not) physically select a logical disk. Instead, it returns a pointer in the HL register pair to the disk parameter header for the logical disk specified in register C on entry. C = 0 for drive A, 1 for drive B, and so on. SELDSK also stores this code for the requested disk to be used later in the READ and WRITE functions.

If the logical disk code in register C refers to a nonexistent disk or to one for which no disk parameter header exists, then SELDSK must return with HL set to 0000H. Then the BDOS will output a message of the form

```
"BDOS Err on X: Select"
```

Note that SELDSK not only does not select the disk, but also does not indicate whether or not the requested disk is physically present—merely whether or not there are disk tables present for the disk.

SELDSK is called by the BDOS either during disk file operations or by a program issuing a Select Disk request (B$SELDSK, code 14).

SETTRK: Set Track

SETTRK saves the requested disk track that is in the BC register pair when SETTRK gets control. Note that this is an absolute track number; that is, the number of reserved tracks before the file directory will have been added to the track number relative to the start of the logical disk.

The number of the requested track will be used in the next BIOS READ or WRITE function (described later in this chapter).

SETTRK is called by the BDOS when it needs to read or write a 128-byte sector. Legitimate track numbers are from 0 to 0FFFFH (65,535).

SETSEC: Set Sector

SETSEC is similar to SETTRK in that it stores the requested sector number for later use in BIOS READ or WRITE functions. The requested sector number is handed to SETSEC in the A register; legitimate values are from 0 to 0FFH (255).

The sector number is a logical sector number. It does not take into account any sector skewing that might be used to improve disk performance.

SETSEC is called by the BDOS when it needs to read or write a 128-byte sector.

SETDMA: Set DMA Address

SETDMA saves the address in the BC register pair in the requested DMA address. The next BIOS READ or WRITE function will use the DMA address as a pointer to the 128-byte sector buffer into which data will be read or from which data will be written.

The default DMA address is 0080H. SETDMA is called by the BDOS when it needs to READ or WRITE a 128-byte sector.

READ: Read Sector

READ reads in a 128-byte sector provided that there have been previous BIOS function calls to

SELDSK—"select" the disk

SETDMA—set the DMA address

SETTRK—set the track number

SETSEC—set the sector number.

Because of the blocking/deblocking code in the BIOS, there are frequent occasions when the requested sector will already be in the host buffer (HSTBUF), so that a physical disk read is not required. All that is then required is for the BIOS to move the appropriate 128 bytes from the HSTBUF into the buffer pointed at by the DMA address.

Only during the READ function will the BIOS normally communicate with the physical disk drive, selecting it and seeking to read the requested track and sector. During this process, the READ function must also handle any hardware errors that occur, trying an operation again if a "soft," or recoverable, error occurs.

The READ function must return with the A register set to 00H if the read operation is completed successfully. If the READ function returns with the A register set to 01H, the BDOS will display an error message of the form

```
BDOS Err on X: Bad Sector
```

Under these circumstances, you have only two choices. You can enter a CARRIAGE RETURN, ignore the fact that there was an error, and attempt to make sense of the data in the DMA buffer. Or you can type a CONTROL-C to abort the operation, perform a warm boot, and return control to the CCP.

As you can see, CP/M's error handling is not particularly helpful, so most BIOS writers add more sophisticated error recovery right in the disk driver. This can include some interaction with the console so that a more determined effort can be made to correct errors or, if nothing else, give you more information as to what has gone wrong. Such error handling is discussed in Chapter 9.

If you are working with a hard disk system, the BIOS driver must also handle the management of bad sectors. You cannot simply replace a hard disk drive if one or two sectors become unreadable. This bad sector management normally requires that a directory of "spare" sectors be put on the hard disk before it is used to store data. Then, when a sector is found to be bad, one of the spare sectors is substituted in its place. This is also discussed in Chapter 9.

WRITE: Write Sector

WRITE is similar to READ but with the obvious difference that data is transferred from the DMA buffer to the specified 128-byte sector. Like READ, this function requires that the following function calls have already been made:

SELDSK—"select" the disk

SETDMA—set the DMA address

SETTRK—set the track number

SETSEC—set the sector number.

Again, it is only in the WRITE routine that the driver will start to talk directly to the physical hardware, selecting the disk unit, track, and sector, and transferring the data to the disk.

With the blocking/deblocking code, the BDOS optimizes the number of disk writes that are needed by indicating in register C the type of disk write that is to be performed:

0 = normal sector write

1 = write to file directory sector

2 = write to sector of previously unused allocation block.

Type 0 occurs whenever the BDOS is writing to a data sector in an already used allocation block. Under these circumstances, the disk driver must preread the appropriate host sector because there may be previously stored information on it.

Type 1 occurs whenever the BDOS is writing to a file directory sector—in this case, the BIOS must not defer writing the sector to the disk, as the information is too valuable to hold in memory until the HSTBUF is full. The longer the information resides in the HSTBUF, the greater the chance of a power failure or glitch, making file data already physically written to the disk inaccessible because the file directory is out of date.

Type 2 occurs whenever the BDOS needs to write to the first sector of a previously unused allocation block. Unused, in this context, includes an allocation block that has become available as a result of a file being erased. In this case, there is no need for the disk driver to preread an entire host-sized sector into the HSTBUF, as there is no data of value in the physical sector.

As with the READ routine, the WRITE function returns with A set to 00H if the operation has been completed successfully. If the WRITE function returns with A set to 01H, then the BDOS will display the *same* message as for READ:

```
BDOS Err on X: Bad Sector
```

You can see now why most BIOS writers add extensive error-recovery and user-interaction routines to their disk drivers.

For hard disk systems, some disk drivers are written so that they automatically "spare out" a failing sector, writing the data to one of the spare sectors on the disk.

LISTST: List Status

As you can tell from its position in the list of BIOS functions, the LISTST function was a latecomer. It was added when CP/M was upgraded from version 1.4 to version 2.0.

This function returns the current status of the list device, using the IOBYTE if necessary to select the correct physical device. It sets the A register to 0FFH if the list device can accept another character for output or to 00H if it is not ready.

Digital Research's documentation states that this function is used by the DESPOOL utility program (which allows you to print a file "simultaneously" with other operations) to improve console response during its operation, and that it is acceptable for the routine always to return 00H if you choose not to implement it fully.

Unfortunately, this statement is wrong. Many other programs use the LISTST function to "poll" the list device to make sure it is ready, and if it fails to come ready after a predetermined time, to output a message to the console indicating that the printer is not ready. If you ever make a call to the BDOS list output functions, Write Printer Byte and Print String (codes 5 and 9), and the printer is not ready, then CP/M will wait forever—and your program will have lost control so it cannot even detect that the problem has occurred. If LISTST always returns a 00H, then the printer will always appear not to be ready. Not only does this make nonsense out of the LISTST function, but it also causes a stream of false "Printer not Ready" error messages to appear on the console.

SECTRAN: Sector Translate

SECTRAN, given a logical sector number, locates the correct physical sector number in the sector translate table for the previously selected (via SELDSK) logical disk drive.

Note that both logical and physical sector numbers are 128-byte sectors, so if you are working with a hard disk system, it is not too efficient to impose a sector interlace at the 128-byte sector level. It is better to impose the sector interlace right inside the hard disk driver, if at all; in general, hard disks spin so rapidly that CP/M simply cannot take advantage of sector interlace.

The BDOS hands over the logical sector number in the BC register pair, with the address of the sector translate table in the DE register pair. SECTRAN must return the physical sector number in HL.

If SECTRAN is to be a null routine, it must move the contents of BC to HL and return.

Calling the BIOS Functions Directly

As a general rule, you should not make direct calls to the BIOS. To do so makes your programs less transportable from one CP/M system to the next. It precludes being able to run these programs under MP/M, which has a different form of BIOS called an extended I/O system, or XIOS.

There are one or two problems, however, that can only be solved by making direct BIOS calls. These occur in utility programs that, for example, need to make direct access to the CP/M file directory, or need to access some "private" jump instructions which have been added to the standard BIOS jump vector.

If you really do need direct access to the BIOS, Figure 6-2 shows an example subroutine that does this. It requires that the A register contain a BIOS function code indicating the offset in the jump vector of the jump instruction to which control is to be passed.

```
                ;       Equates for use with BIOS subroutine
                ;
0003 =          WBOOT   EQU     03H     ;Warm boot
0006 =          CONST   EQU     06H     ;Console status
0009 =          CONIN   EQU     09H     ;Console input
000C =          CONOUT  EQU     0CH     ;Console output
000F =          LIST    EQU     0FH     ;Output to list device
0012 =          PUNCH   EQU     12H     ;Output to punch device
0015 =          READER  EQU     15H     ;Input from reader
0018 =          HOME    EQU     18H     ;Home selected disk to track 0
001B =          SELDSK  EQU     1BH     ;Select disk
001E =          SETTRK  EQU     1EH     ;Set track
0021 =          SETSEC  EQU     21H     ;Set sector
0024 =          SETDMA  EQU     24H     ;Set DMA address
0027 =          READ    EQU     27H     ;Read 128-byte sector
002A =          WRITE   EQU     2AH     ;Write 128-byte sector
002D =          LISTST  EQU     2DH     ;Return list status
0030 =          SECTRAN EQU     30H     ;Sector translate
                ;
                                        ;Add further "private" BIOS codes here
                ;
                ;       BIOS
                ;       This subroutine transfers control to the appropriate
                ;       entry in the BIOS Jump Vector, based on a code number
                ;       handed to it in the L register.
                ;
                ;       Entry parameters
                ;
                ;       L = Code number (which is in fact the page-relative
                ;               address of the correct JMP instruction within
                ;               the jump vector)
                ;       All other registers are preserved and handed over to
                ;               the BIOS routine intact.
                ;
                ;       Exit parameters
                ;
```

Figure 6-2. BIOS equates

```
                        ;       This routine does not CALL the BIOS routine, therefore
                        ;       when the BIOS routine RETurns, it will do so directly
                        ;       to this routine's caller.
                        ;
                        ;       Calling sequence
                        ;
                        ;               MVI     L,Code$Number
                        ;               CALL    BIOS
                        ;
                        BIOS:
0000 F5                         PUSH    PSW     ;Save user's A register
0001 3A0200                     LDA     0002H   ;Get BIOS JMP vector page from
                                                ;  warm boot JMP
0004 67                         MOV     H,A     ;HL -> BIOS JMP vector entry
0005 F1                         POP     PSW     ;Recover user's A register
0006 E9                         PCHL            ;Transfer control into the BIOS routine
```

Figure 6-2. BIOS equates (continued)

Line Numbers	Functional Component or Routine
0072-0116	BIOS Jump Vector
0120-0270	Initialization Code
0275-0286	Display Message
0289-0310	Enter CP/M
0333-0364	CONST - Console Status
0369-0393	CONIN - Console Input
0397-0410	CONOUT - Console Output
0414-0451	LISTST - List Status
0456-0471	LIST - List Output
0476-0492	PUNCH - Punch Output
0496-0511	READER - Reader Input
0516-0536	IOBYTE Driver Select
0540-0584	Device Control Tables
0589-0744	Low-level Drivers for Console, List,etc.
0769-0824	Disk Parameter Header Tables
0831-0878	Disk Parameter Blocks
0881-0907	Other Disk data areas
0910-0955	SELDSK - Select Disk
0958-0964	SETTRK - Set Track
0967-0973	SETSEC - Set Sector
0978-0984	SETDMA - Set DMA Address
0987-1025	Sector Skew Tables
1028-1037	SECTRAN - Logical to Physical Sector translation
1041-1056	HOME - Home to Track 0
1059-1154	Deblocking Algorithm data areas
1157-1183	READ - Read 128-byte sector
1185-1204	WRITE - Write 128-byte sector
1206-1378	Deblocking Algorithm
1381-1432	Buffer Move
1435-1478	Deblocking subroutines
1481-1590	8" Floppy Physical Read/Write
1595-1681	5 1/4" Floppy Physical Read/Write
1685-1764	WBOOT - Warm Boot

Figure 6-3. Functional Index to Figure 6-4

Example BIOS

The remainder of this chapter is devoted to an example BIOS listing. This actual working BIOS shows the overall structure and interface to the individual BIOS subroutines.

Unlike most BIOS's, this one has been written specifically to be understood easily. The variable names are uncharacteristically long and descriptive, and each block of code has commentary to put it into context.

Each source line has been sequentially numbered (an infrequently used option that Digital Research's Assembler, ASM, permits). Figure 6-3 contains a functional index to the BIOS as a whole so that you can find particular functions in the listing in Figure 6-4 by line number.

```
0001 <-- Line Number ;  Figure 6-4.
0002                 ;
0003                 ;*****************************************************
0004                 ;*                                                   *
0005                 ;*          Simple BIOS Listing                      *
0006                 ;*                                                   *
0007                 ;*****************************************************
0008                 ;
0009                 ;
0010  3030 =         VERSION         EQU     '00'    ;Equates used in the sign on message
0011  3730 =         MONTH           EQU     '07'
0012  3531 =         DAY             EQU     '15'
0013  3238 =         YEAR            EQU     '82'
0014                 ;
0015                 ;*******************************************************************
0016                 ;*                                                                 *
0017                 ;* This BIOS is for a computer system with the following           *
0018                 ;* hardware configuration :                                        *
0019                 ;*                                                                 *
0020                 ;*         - 8080 CPU                                              *
0021                 ;*         - 64KBytes of RAM                                       *
0022                 ;*         - CRT/keyboard controller that transfers data           *
0023                 ;*           as though it were a serial port (but requires         *
0024                 ;*           no baud rate generator or USART programming)          *
0025                 ;*         - A serial port, used for both list and "reader"/       *
0026                 ;*           "punch" devices. The serial port chip is an           *
0027                 ;*           Intel 8251A with an 8253 baud rate generator.         *
0028                 ;*         - Two 5 1/4" mini-floppy, double-sided, double-         *
0029                 ;*           density drives. These drives use 512-byte sectors.    *
0030                 ;*           These are used as logical disks A: and B:.            *
0031                 ;*         - Two 8" standard diskette drives (128-byte sectors).   *
0032                 ;*           These are used as logical disks C: and D:.            *
0033                 ;*                                                                 *
0034                 ;*           Two intelligent disk controllers are used, one for    *
0035                 ;*           each diskette type. These controllers access memory   *
0036                 ;*           directly, both to read the details of the             *
0037                 ;*           operations they are to perform and also to read       *
0038                 ;*           and write data from and to the diskettes.             *
0039                 ;*                                                                 *
0040                 ;*                                                                 *
0041                 ;*******************************************************************
0042
0043                 ;
0044                 ;  Equates for defining memory size and the base address and
0045                 ;  length of the system components.
```

Figure 6-4. Simple BIOS listing

```
                  ;
0040 =            Memory$Size          EQU    64      ;Number of Kbytes of RAM
                  ;
                  ;  The BIOS Length must be determined by inspection.
                  ;  Comment out the ORG BIOS$Entry line below by changing the first
                  ;  character to a semicolon. (This will make the Assembler start
                  ;  the BIOS at location 0.) Then assemble the BIOS and round up to
                  ;  the nearest 100H the address displayed on the console at the end
                  ;  of the assembly.
                  ;
0900 =            BIOS$Length          EQU    0900H
                  ;
0800 =            CCP$Length EQU       0800H   ;Constant
0E00 =            BDOS$Length          EQU    0E00H   ;Constant
                  ;
0008 =            Overall$Length       EQU    ((CCP$Length + BDOS$Length + BIOS$Length) / 1024) + 1
                  ;
E000 =            CCP$Entry  EQU       (Memory$Size - Overall$Length) * 1024
E806 =            BDOS$Entry EQU       CCP$Entry + CCP$Length + 6
F600 =            BIOS$Entry EQU       CCP$Entry + CCP$Length + BDOS$Length
                  ;
                  ;
                  ;

F600                 ORG     BIOS$Entry      ;Assemble code at BIOS address
                  ;
                  ;  BIOS jump vector
                  ;  Control will be transferred to the appropriate entry point
                  ;  from the CCP or the BDOS, both of which compute the relative
                  ;  address of the BIOS jump vector in order to locate it.
                  ;  Transient programs can also make direct BIOS calls transferring
                  ;  control to location xx00H, where xx is the value in location
                  ;  0002H.
                  ;
F600 C3F9F6          JMP     BOOT    ;Cold boot -- entered from CP/M bootstrap loader
                  Warm$Boot$Entry:   ;  Labelled so that the initialization code can
                                     ;  put the warm boot entry address down in location
                                     ;  0001H and 0002H of the base page
F603 C329FE          JMP     WBOOT   ;Warm boot -- entered by jumping to location 0000H.
                                     ;  Reloads the CCP which could have been
                                     ;  overwritten by previous program in transient
                                     ;  program area
F606 C362F8          JMP     CONST   ;Console status -- returns A = 0FFH if there is a
                                     ;  console keyboard character waiting
F609 C378F8          JMP     CONIN   ;Console input -- returns the next console keyboard
                                     ;  character in A
F60C C386F8          JMP     CONOUT  ;Console output -- outputs the character in C to
                                     ;  the console device
F60F C3ACF8          JMP     LIST    ;List output -- outputs the character in C to the
                                     ;  list device
F612 C3BCF8          JMP     PUNCH   ;Punch output -- outputs the character in C to the
                                     ;  logical punch device
F615 C3CDF8          JMP     READER  ;Reader input -- returns the next input character from
                                     ;  the logical reader device in A
F618 C3D3FB          JMP     HOME    ;Homes the currently selected disk to track 0
F61B C32BFB          JMP     SELDSK  ;Selects the disk drive specified in register C and
                                     ;  returns the address of the disk parameter header
F61E C358FB          JMP     SETTRK  ;Sets the track for the next read or write operation
                                     ;  from the BC register pair
F621 C35EFB          JMP     SETSEC  ;Sets the sector for the next read or write operation
                                     ;  from the A register
F624 C365FB          JMP     SETDMA  ;Sets the direct memory address (disk read/write)
                                     ;  address for the next read or write operation
                                     ;  from the DE register pair
F627 C3FBFB          JMP     READ    ;Reads the previously specified track and sector from
                                     ;  the selected disk into the DMA address
F62A C315FC          JMP     WRITE   ;Writes the previously specified track and sector onto
                                     ;  the selected disk from the DMA address
F62D C394F8          JMP     LISTST  ;Returns A = 0FFH if the list device can accept
                                     ;  another output character
F630 C3CDFB          JMP     SECTRAN ;Translates a logical sector into a physical one
                  ;
                  ;
                  ;
                  ;  The cold boot initialization code is only needed once.
```

Figure 6-4. (Continued)

```
                     ;  It can be overwritten once it has been executed.
                     ;  Therefore, it is "hidden" inside the main disk buffer.
                     ;  When control is transferred to the BOOT entry point, this
                     ;  code will be executed, only being overwritten by data from
                     ;  the disk once the initialization procedure is complete.
                     ;
                     ;  To hide code in the buffer, the buffer is first declared
                     ;  normally. Then the value of the location counter following
                     ;  the buffer is noted. Then, using an ORG (ORiGin) statement, the
                     ;  location counter is "wound back" to the start of the buffer
                     ;  again and the initialization code written normally.
                     ;  At the end of this code, another ORG statement is used to
                     ;  set the location counter back as it was after the buffer had
                     ;  been declared.
                     ;
                     ;
0200 =               Physical$Sector$Size        EQU     512     ;This is the actual sector size
                                                         ;for the 5 1/4" mini-floppy diskettes.
                                                         ;The 8" diskettes use 128-byte sectors.
                                                         ;Declare the physical disk buffer for the
                                                         ;5 1/4" diskettes
F633                 Disk$buffer:        DS      Physical$Sector$Size
                     ;
                                                         ;Save the location counter
F833 =               After$Disk$Buffer   EQU     $       ;$ = Current value of location counter
                     ;
F633                                     ORG     Disk$Buffer     ;Wind the location counter back
                     ;
                     Initialize$Stream: ;This stream of data is used by the
                                        ;initialize subroutine. It has the following
                                        ;format:
                                        ;
                                        ;       DB      Port number to be initialized
                                        ;       DB      Number of bytes to be output
                                        ;       DB      xx,xx,xx,xx data to be output
                                        ;       :
                                        ;       :
                                        ;       DB      Port number of 00H terminator
                                        ;
                                        ;Note : On this machine, the console port does
                                        ;       not need to be initialized. This has
                                        ;       already been done by the PROM bootstrap code.
                                        ;
                                                ;Initialize the 8251A USART used for
                                                ;  the list and communications devices.
F633 ED                 DB      Communication$Status$Port       ;Port number
F634 06                 DB      6                               ;Number of bytes
F635 00                 DB      0               ;Get chip ready to be programmed by
F636 00                 DB      0               ;  sending dummy data out to it
F637 00                 DB      0
F638 42                 DB      0100$0010B      ;Reset and raise data terminal ready
F639 6E                 DB      01$10$11$10B    ;1 stop bit, no parity, 8 bits per character
                                                ;  baud rate divide factor of 16.
F63A 25                 DB      0010$0101B      ;Raise request to send, and enable
                                                ;  transmit and receive.
                     ;
                                                ;Initialize the 8253 programmable interval
                                                ;  timer used to generate the baud rate for
                                                ;  the 8251A USART
F63B DF                 DB      Communication$Baud$Mode         ;Port number
F63C 01                 DB      1                               ;Number of bytes
F63D B6                 DB      10$11$011$0B    ;Select counter 2, load LS byte first,
                                                ;  Mode 3 (for baud rates), binary count.
                     ;
F63E DE                 DB      Communication$Baud$Rate         ;Port number
F63F 02                 DB      2                               ;Number of bytes
F640 3800               DW      0038H           ;1200 baud (based on 16X divide-down selected
                                                ;  in the 8251A USART)
                     ;
F642 00                 DB      0               ;Port number of 0 terminates
                     ;
                     ;
                     ;  Equates for the sign-on message
                     ;
000D =               CR EQU     0DH             ;Carriage return
```

Figure 6-4. (Continued)

```
000A =          LF EQU     0AH              ;Line feed
                ;
                Signon$Message:             ;Main sign-on message
F643 43502F4D20    DB      'CP/M 2.2.'
F64C 3030          DW      VERSION          ;Current version number
F64E 20            DB      ' '
F64F 3037          DW      MONTH            ;Current date
F651 2F            DB      '/'
F652 3135          DW      DAY
F654 2F            DB      '/'
F655 3832          DW      YEAR
F657 0D0A0A        DB      CR,LF,LF
F65A 53696D706C    DB      'Simple BIOS',CR,LF,LF
F668 4469736B20    DB      'Disk configuration :',CR,LF,LF
F67F 2020202020    DB      '     A: 0.35 Mbyte 5" Floppy',CR,LF
F69D 2020202020    DB      '     B: 0.35 Mbyte 5" Floppy',CR,LF,LF
F6BC 2020202020    DB      '     C: 0.24 Mbyte 8" Floppy',CR,LF
F6DA 2020202020    DB      '     D: 0.24 Mbyte 8" Floppy',CR,LF
                ;
F6F8 00            DB      0
                ;
0004 =          Default$Disk        EQU     0004H   ;Default disk in base page
                ;
                BOOT:      ;Entered directly from the BIOS JMP vector.
                           ;Control will be transferred here by the CP/M
                           ;  bootstrap loader.
                           ;The initialization state of the computer system
                           ;  will be determined by the
                           ;  PROM bootstrap and the CP/M loader setup.
                           ;
                                            ;Initialize system.
                                            ;This routine uses the Initialize$Stream
                                            ;  declared above.
F6F9 F3            DI                       ;Disable interrupts to prevent any
                                            ;  side effects during initialization.
F6FA 2133F6        LXI     H,Initialize$Stream     ;HL -> Data stream
                ;
                Initialize$Loop:
F6FD 7E            MOV     A,M              ;Get port number
F6FE B7            ORA     A                ;If 00H, then initialization complete
F6FF CA13F7        JZ      Initialize$Complete
F702 320AF7        STA     Initialize$Port ;Set up OUT instruction
F705 23            INX     H                ;HL -> Count of number of bytes to output
F706 4E            MOV     C,M              ;Get byte count
                ;
                Initialize$Next$Byte:
F707 23            INX     H                ;HL -> Next data byte
F708 7E            MOV     A,M              ;Get next data byte
F709 D3            DB      OUT              ;Output to correct port
                Initialize$Port:
F70A 00            DB      0                ;<- Set above
F70B 0D            DCR     C                ;Count down
F70C C207F7        JNZ     Initialize$Next$Byte    ;Go back if more bytes
F70F 23            INX     H                ;HL -> Next port number
F710 C3FDF6        JMP     Initialize$Loop ;Go back for next port initialization
                ;
                Initialize$Complete:
                ;

F713 3E01          MVI     A,00$00$00$01B          ;Set IOBYTE to indicate terminal
F715 320300        STA     IOBYTE                  ;  is to act as console

F718 2143F6        LXI     H,Signon$Message        ;Display sign-on message on console
F71B CD33F8        CALL    Display$Message
                ;

F71E AF            XRA     A                ;Set default disk drive to A:
F71F 320400        STA     Default$Disk
F722 FB            EI                       ;Interrupts can now be enabled
                ;
F723 C340F8        JMP     Enter$CPM        ;Complete initialization and enter
                                            ;  CP/M by going to the Console Command
                                            ;  Processor.
                ;
                ;  End of cold boot initialization code
                ;
```

Figure 6-4. (Continued)

```
F833                ORG     After$Disk$Buffer       ;Reset location counter
                 ;
                 ;
                 Display$Message:   ;Displays the specified message on the console.
                                    ;On entry, HL points to a stream of bytes to be
                                    ;  output. A 00H-byte terminates the message.
F833 7E             MOV     A,M                ;Get next message byte
F834 B7             ORA     A                  ;Check if terminator
F835 C8             RZ                         ;Yes, return to caller
F836 4F             MOV     C,A                ;Prepare for output
F837 E5             PUSH    H                  ;Save message pointer
F838 CD86F8         CALL    CONOUT             ;Go to main console output routine
F83B E1             POP     H                  ;Recover message pointer
F83C 23             INX     H                  ;Move to next byte of message
F83D C333F8         JMP     Display$Message ;Loop until complete message output
                 ;
                 ;
                 Enter$CPM: ;This routine is entered either from the cold or warm
                            ;  boot code. It sets up the JMP instructions in the
                            ;  base page, and also sets the high-level disk driver's
                            ;  input/output address (also known as the DMA address).
                 ;
F840 3EC3           MVI     A,JMP              ;Get machine code for JMP
F842 320000         STA     0000H              ;Set up JMP at location 0000H
F845 320500         STA     0005H              ;  and at location 0005H
                 ;
F848 2103F6         LXI     H,Warm$Boot$Entry       ;Get BIOS vector address
F84B 220100         SHLD    0001H              ;Put address at location 0001H

F84E 2106E8         LXI     H,BDOS$Entry       ;Get BDOS entry point address
F851 220600         SHLD    6                  ;Put address at location 0005H
                 ;
F854 018000         LXI     B,80H              ;Set disk I/O address to default
F857 CD65FB         CALL    SETDMA             ;Use normal BIOS routine
                 ;
F85A FB             EI                         ;Ensure interrupts are enabled
F85B 3A0400         LDA     Default$Disk       ;Transfer current default disk to
F85E 4F             MOV     C,A                ;  Console Command Processor
F85F C300E0         JMP     CCP$Entry          ;Transfer to CCP
                 ;
                 ;
                 ;  Serial input/output drivers
                 ;
                 ;  These drivers all look at the IOBYTE at location
                 ;  0003H, which will have been set by the cold boot routine.
                 ;  The IOBYTE can be modified by the STAT utility, by
                 ;  BDOS calls, or by a program that puts a value directly
                 ;  into location 0003H.
                 ;
                 ;  All of the routines make use of a subroutine, Select$Routine,
                 ;  that takes the least significant two bits of the A register
                 ;  and uses them to transfer control to one of the routines whose
                 ;  address immediately follows the call to Select$Routine.
                 ;  A second entry point, Select$Routine$21, uses bits
                 ;  2 and 1 to do the same job -- this saves some space
                 ;  by avoiding an unnecessary instruction.
                 ;
0003 =           IOBYTE      EQU     0003H   ;I/O redirection byte
                 ;
                 ;
                 ;
                 CONST:             ;Get console status
                                    ;Entered directly from the BIOS JMP vector
                                    ;  and returns a parameter that reflects whether
                                    ;  there is incoming data from the console.
                                    ;
                                    ;A = 00H (zero flag set) if no data
                                    ;A = 0FFH (zero flag clear) if data
                                    ;
                                    ;CONST will be called by programs that
                                    ;  make periodic checks to see if the computer
                                    ;  operator has pressed any keys -- for example,
                                    ;  to interrupt an executing program.
                                    ;
F862 CD6AF8         CALL    Get$Console$Status      ;Return A = zero or nonzero
                                               ;According to status, then convert
```

Figure 6-4. (Continued)

```
                                                 ;  to return parameter convention.
F865 B7            ORA     A                     ;Set flags to reflect status
F866 C8            RZ                            ;If 0, no incoming data
F867 3EFF          MVI     A,0FFH                ;Otherwise return A = 0FFH to
F869 C9            RET                           ;  indicate incoming data
                ;
                Get$Console$Status:
F86A 3A0300        LDA     IOBYTE                ;Get I/O redirection byte
                                                 ;Console is selected according to
                                                 ;  bits 1,0 of IOBYTE
F86D CDDCF8        CALL    Select$Routine        ;Select appropriate routine
                                                 ;These routines return to the caller
                                                 ;  of Get$Console$Status.
F870 F6F8          DW      Teletype$In$Status        ;00 <- IOBYTE bits 1,0
F872 FCF8          DW      Terminal$In$Status        ;01
F874 02F9          DW      Communication$In$Status   ;10
F876 08F9          DW      Dummy$In$Status           ;11
                ;
                ;
                ;
                ;
                CONIN:                  ;Get console input character
                                        ;Entered directly from the BIOS JMP vector;
                                        ;  returns the next data character from the
                                        ;  Console in the A register. The most significant
                                        ;  bit of the data character will be 0, except
                                        ;  when "reader" (communication port) input has
                                        ;  been selected. In this case, the full eight bits
                                        ;  of data are returned to permit binary data to be
                                        ;  received.
                                        ;
                                        ;Normally, this routine will be called after
                                        ;  a call to CONST has indicated that a data character
                                        ;  is ready, but whenever the CCP or the BDOS can
                                        ;  proceed no further until console input occurs,
                                        ;  then CONIN will be called without a preceding
                                        ;  CONST call.
                                        ;
F878 3A0300        LDA     IOBYTE                ;Get I/O redirection byte
F87B CDDCF8        CALL    Select$Routine        ;Select correct CONIN routine
                                                 ;These routines return directly
                                                 ; to CONIN's caller.
F87E 20F9          DW      Teletype$Input            ;00 <- IOBYTE bits 1,0
F880 26F9          DW      Terminal$Input            ;01
F882 2FF9          DW      Communication$Input       ;10
F884 35F9          DW      Dummy$Input               ;11
                ;
                ;
                ;
                CONOUT:                 ;Console output
                                        ;Entered directly from BIOS JMP vector;
                                        ; outputs the data character in the C register
                                        ; to the appropriate device according to bits
                                        ; 1,0 of IOBYTE
                                        ;
F886 3A0300        LDA     IOBYTE                ;Get I/O redirection byte
F889 CDDCF8        CALL    Select$Routine        ;Select correct CONOUT routine
                                                 ;These routines return directly
                                                 ; to CONOUT's caller.
F88C 38F9          DW      Teletype$Output           ;00 <- IOBYTE bits 1,0
F88E 3EF9          DW      Terminal$Output           ;01
F890 44F9          DW      Communication$Output      ;10
F892 4AF9          DW      Dummy$Output              ;11
                ;
                ;
                ;
                LISTST:                 ;List device (output) status
                                        ;Entered directly from the BIOS JMP vector;
                                        ; returns in A list device status that
                                        ; indicates whether the list device can accept
                                        ; another output character. The IOBYTE's bits
                                        ; 7,6 determine the physical device used.
                                        ;
                                        ;A = 00H (zero flag set): cannot accept data
                                        ;A = 0FFH (zero flag clear): can accept data
                                        ;
```

Figure 6-4. (Continued)

```
                                        ;Digital Research's documentation indicates
                                        ; that you can always return with A = 00H
                                        ; ("Cannot accept data") if you do not wish to
                                        ; implement the LISTST routine. This is NOT TRUE.
                                        ;If you do not wish to implement the LISTST routine
                                        ; always return with A = 0FFH ("Can accept data").
                                        ;The LIST driver will then take care of things rather
                                        ; than potentially hanging the system.
                                        ;
F894 CD9CF8         CALL    Get$List$Status ;Return A = zero or nonzero
                                        ; according to status, then convert
                                        ; to return parameter convention
F897 B7             ORA     A           ;Set flags to reflect status
F898 C8             RZ                  ;If 0, cannot accept data for output
F899 3EFF           MVI     A,0FFH      ;Otherwise return A = 0FFH to
F89B C9             RET                 ; indicate can accept data for output
                ;
                Get$List$Status:
F89C 3A0300         LDA     IOBYTE      ;Get I/O redirection byte
F89F 07             RLC                 ;Move bits 7,6 to 1,0
F8A0 07             RLC
F8A1 CDDCF8         CALL    Select$Routine  ;Select appropriate routine
                                                ;These routines return directly
                                                ; to Get$List$Status's caller.
F8A4 0BF9           DW      Teletype$Out$Status         ;00 <- IOBYTE bits 1,0
F8A6 11F9           DW      Terminal$Out$Status         ;01
F8A8 17F9           DW      Communication$Out$Status    ;10
F8AA 1DF9           DW      Dummy$Out$Status            ;11

                ;
                ;
                ;
                LIST:               ;List output
                                    ;Entered directly from BIOS JMP vector;
                                    ; outputs the data character in the C register
                                    ; to the appropriate device according to bits
                                    ; 7,6 of IOBYTE
                                    ;
F8AC 3A0300         LDA     IOBYTE          ;Get I/O redirection byte
F8AF 07             RLC                     ;Move bits 7,6 to 1,0
F8B0 07             RLC
F8B1 CDDCF8         CALL    Select$Routine          ;Select correct LIST routine
                                                    ;These routines return directly
                                                    ; to LIST's caller.
F8B4 38F9           DW      Teletype$Output         ;00 <- IOBYTE bits 1,0
F8B6 3EF9           DW      Terminal$Output         ;01
F8B8 44F9           DW      Communication$Output    ;10
F8BA 4AF9           DW      Dummy$Output            ;11

                ;
                ;
                ;
                PUNCH:              ;Punch output
                                    ;Entered directly from BIOS JMP vector;
                                    ; outputs the data character in the C register
                                    ; to the appropriate device according to bits
                                    ; 5,4 of IOBYTE
                                    ;
F8BC 3A0300         LDA     IOBYTE          ;Get I/O redirection byte
F8BF 0F             RRC                     ;Move bits 5,4 to 2,1
F8C0 0F             RRC
F8C1 0F             RRC
F8C2 CDDDF8         CALL    Select$Routine$21       ;Select correct PUNCH routine
                                                    ;These routines return directly
                                                    ; to PUNCH's caller.
F8C5 38F9           DW      Teletype$Output         ;00 <- IOBYTE bits 1,0
F8C7 4AF9           DW      Dummy$Output            ;01
F8C9 44F9           DW      Communication$Output    ;10
F8CB 3EF9           DW      Terminal$Output         ;11
                ;
                ;
                ;
                READER:             ;Reader input
                                    ;Entered directly from BIOS JMP vector;
                                    ; inputs the next data character from the
                                    ; reader device into the A register
```

Figure 6-4. (Continued)

```
                                        ;The appropriate device is selected according
                                        ; to bits 3,2 of IOBYTE.
                                        ;
F8CD 3A0300       LDA     IOBYTE                  ;Get I/O redirection byte
F8D0 0F           RRC                             ;Move bits 3,2 to 2,1
F8D1 CDDDF8       CALL    Select$Routine$21       ;Select correct READER routine
                                                  ;These routines return directly
                                                  ; to READER's caller.
F8D4 38F9         DW      Teletype$Output         ;00 <- IOBYTE bits 1,0
F8D6 4AF9         DW      Dummy$Output            ;01
F8D8 44F9         DW      Communication$Output    ;10
F8DA 3EF9         DW      Terminal$Output         ;11
              ;
              ;
              ;
              Select$Routine:             ;Transfers control to a specified address
                                          ; following its calling address according to
                                          ; the value of bits 1,0 in A.
F8DC 07           RLC                     ;Shift select values into bits 2,1
                                          ; in order to do word arithmetic
              ;
              Select$Routine$21:          ;Entry point to select routine selection bits
                                          ; are already in bits 2,1
F8DD E606         ANI     0000$0110B      ;Isolate just bits 2,1
F8DF E3           XTHL                    ;HL -> first word of addresses after
                                          ; CALL instruction
F8E0 5F           MOV     E,A             ;Add on selection value to address table
F8E1 1600         MVI     D,0             ; base
F8E3 19           DAD     D               ;HL -> selected routine address
                                          ;Get routine address into HL
F8E4 7E           MOV     A,M             ;LS byte
F8E5 23           INX     H               ;HL -> MS byte
F8E6 66           MOV     H,M             ;MS byte
F8E7 6F           MOV     L,A             ;HL -> routine
F8E8 E3           XTHL                    ;Top of stack -> routine
F8E9 C9           RET                     ;Transfer to selected routine
              ;
              ;
              ;
              ;   Input/Output Equates
              ;
00ED =        Teletype$Status$Port                EQU     0EDH
00EC =        Teletype$Data$Port          EQU     0ECH
0001 =        Teletype$Output$Ready               EQU     0000$0001B      ;Status mask
0002 =        Teletype$Input$Ready                EQU     0000$0010B      ;Status mask
              ;
0001 =        Terminal$Status$Port                EQU     01H
0002 =        Terminal$Data$Port          EQU     02H
0001 =        Terminal$Output$Ready               EQU     0000$0001B      ;Status mask
0002 =        Terminal$Input$Ready                EQU     0000$0010B      ;Status mask
              ;
00ED =        Communication$Status$Port   EQU     0EDH
00EC =        Communication$Data$Port             EQU     0ECH
0001 =        Communication$Output$Ready  EQU     0000$0001B      ;Status mask
0002 =        Communication$Input$Ready   EQU     0000$0010B      ;Status mask
              ;
00DF =        Communication$Baud$Mode             EQU     0DFH            ;Mode Select
00DE =        Communication$Baud$Rate             EQU     0DEH            ;Rate Select
              ;
              ;
              ;   Serial device control tables
              ;
              ;   In order to reduce the amount of executable code,
              ;   the same low-level driver code is used for all serial ports.
              ;   On entry to the low-level driver, HL points to the
              ;   appropriate control table.
              ;
              Teletype$Table:
F8EA ED           DB      Teletype$Status$Port
F8EB EC           DB      Teletype$Data$Port
F8EC 01           DB      Teletype$Output$Ready
F8ED 02           DB      Teletype$Input$Ready
              ;
              Terminal$Table:
F8EE 01           DB      Terminal$Status$Port
```

Figure 6-4. (Continued)

```
F8EF 02            DB      Terminal$Data$Port
F8F0 01            DB      Terminal$Output$Ready
F8F1 02            DB      Terminal$Input$Ready
               ;
               Communication$Table:
F8F2 ED            DB      Communication$Status$Port
F8F3 EC            DB      Communication$Data$Port
F8F4 01            DB      Communication$Output$Ready
F8F5 02            DB      Communication$Input$Ready
               ;
               ;
               ;
               ;
               ;   The following routines are "called" by Select$Routine
               ;   to perform the low-level input/output
               ;
               Teletype$In$Status:
F8F6 21EAF8        LXI     H,Teletype$Table         ;HL -> control table
F8F9 C34BF9        JMP     Input$Status             ;Note use of JMP. Input$Status
                                                    ; will execute the RETurn.
               ;
               Terminal$In$Status:
F8FC 21EEF8        LXI     H,Terminal$Table         ;HL -> control table
F8FF C34BF9        JMP     Input$Status             ;Note use of JMP. Input$Status
                                                    ; will execute the RETurn.
               ;
               Communication$In$Status:
F902 21F2F8        LXI     H,Communication$Table    ;HL -> control table
F905 C34BF9        JMP     Input$Status             ;Note use of JMP. Input$Status
                                                    ; will execute the RETurn.
               ;
               Dummy$In$Status:                     ;Dummy status, always returns
F908 3EFF          MVI     A,0FFH                   ; indicating incoming data is ready
F90A C9            RET
               ;
               ;
               Teletype$Out$Status:
F90B 21EAF8        LXI     H,Teletype$Table         ;HL -> control table
F90E C356F9        JMP     Output$Status            ;Note use of JMP. Output$Status
                                                    ; will execute the RETurn.
               ;
               Terminal$Out$Status:
F911 21EEF8        LXI     H,Terminal$Table         ;HL -> control table
F914 C356F9        JMP     Output$Status            ;Note use of JMP. Output$Status
                                                    ; will execute the RETurn.
               ;
               Communication$Out$Status:
F917 21F2F8        LXI     H,Communication$Table    ;HL -> control table
F91A C356F9        JMP     Output$Status            ;Note use of JMP. Output$Status
                                                    ; will execute the RETurn.
               ;
               Dummy$Out$Status:                    ;Dummy status, always returns
F91D 3EFF          MVI     A,0FFH                   ; indicating ready for output
F91F C9            RET
               ;
               ;
               Teletype$Input:
F920 21EAF8        LXI     H,Teletype$Table         ;HL -> control table
F923 C360F9        JMP     Input$Data               ;Note use of JMP. Input$Data
                                                    ; will execute the RETurn.
               ;
               Terminal$Input:
F926 21EEF8        LXI     H,Terminal$Table         ;HL -> control table
                                                    ; will execute the RETurn.
F929 CD60F9        CALL    Input$Data               ;** Special case **
                                                    ;Input$Data will return here
F92C E67F          ANI     7FH                      ; so that parity bit can be set 0
F92E C9            RET
               ;
               Communication$Input:
F92F 21F2F8        LXI     H,Communication$Table    ;HL -> control table
F932 C360F9        JMP     Input$Data               ;Note use of JMP. Input$Data
                                                    ; will execute the RETurn.
               ;
               Dummy$Input:                         ;Dummy input, always returns
F935 3E1A          MVI     A,1AH                    ; indicating CP/M end of file
```

Figure 6-4. (Continued)

```
F937 C9            RET
                ;
                ;
                ;
                ;
                Teletype$Output:
F938 21EAF8        LXI      H,Teletype$Table         ;HL -> control table
F93B C370F9        JMP      Output$Data              ;Note use of JMP. Output$Data
                                                     ; will execute the RETurn.
                ;
                Terminal$Output:
F93E 21EEF8        LXI      H,Terminal$Table         ;HL -> control table
                                                     ; will execute the RETurn.
F941 C370F9        JMP      Output$Data              ;Note use of JMP. Output$Data
                                                     ; will execute the RETurn.
                ;
                Communication$Output:
F944 21F2F8        LXI      H,Communication$Table    ;HL -> control table
F947 C370F9        JMP      Output$Data              ;Note use of JMP. Output$Data
                                                     ; will execute the RETurn.
                ;
                Dummy$Output:                        ;Dummy output, always discards
F94A C9            RET                               ; the output character
                ;
                ;
                ;
                ;
                ;  These are the general purpose low-level drivers.
                ;  On entry, HL points to the appropriate control table.
                ;  For output, the C register contains the data to be output.
                ;
                Input$Status:                ;Return with A = 00H if no incoming data,
                                             ; otherwise A = nonzero.
F94B 7E            MOV      A,M              ;Get status port
F94C 3250F9        STA      Input$Status$Port        ;*** Self-modifying code ***
F94F DB            DB       IN               ;Input to A from correct status port
                ;
                Input$Status$Port:
F950 00            DB       00               ;<- Set above
F951 23            INX      H                ;Move HL to point to input data mask
F952 23            INX      H
F953 23            INX      H
F954 A6            ANA      M                ;Mask with input status
F955 C9            RET
                ;
                ;
                Output$Status:               ;Return with A = 00H if not ready for output
                                             ; otherwise A = nonzero.
F956 7E            MOV      A,M              ;Get status port
F957 325BF9        STA      Output$Status$Port       ;*** Self-modifying code ***
F95A DB            DB       IN               ;Input to A from correct status port
                ;
                Output$Status$Port:
F95B 00            DB       00               ;<- Set above
F95C 23            INX      H                ;Move HL to point to output data mask
F95D 23            INX      H
F95E A6            ANA      M                ;Mask with output status
F95F C9            RET
                ;
                ;
                Input$Data:                  ;Return with next data character in A.
                                             ;Wait for status routine to indicate
                                             ; incoming data.
F960 E5            PUSH     H                ;Save control table pointer
F961 CD4BF9        CALL     Input$Status     ;Get input status in zero flag
F964 E1            POP      H                ;Recover control table pointer
F965 CA60F9        JZ       Input$Data       ;Wait until incoming data
F968 23            INX      H                ;HL -> data port
F969 7E            MOV      A,M              ;Get data port
F96A 326EF9        STA      Input$Data$Port  ;*** Self-modifying code ***
F96D DB            DB       IN               ;Input to A from correct data port
                ;
                Input$Data$Port:
F96E 00            DB       0                ;<- Set above
F96F C9            RET
                ;
```

Figure 6-4. (Continued)

```
                      ;
                      Output$Data:                    ;Output the data character in the C register.
                                                      ;Wait for status routine to indicate device
                                                      ; ready to accept another character
F970 E5                  PUSH   H                     ;Save control table pointer
F971 CD56F9              CALL   Output$Status         ;Get output status in zero flag
F974 E1                  POP    H                     ;Recover control table pointer
F975 CA70F9              JZ     Output$Data           ;Wait until ready for output
F978 23                  INX    H                     ;HL -> output port
F979 7E                  MOV    A,M                   ;Get output port
F97A 327FF9              STA    Output$Data$Port              ;*** Self-modifying code ***
F97D 79                  MOV    A,C                   ;Get data character to be output
F97E D3                  DB     OUT                   ;Output data to correct port
                      ;
                      Output$Data$Port:
F97F 00                  DB     0                     ;<- Set above
F980 C9                  RET
                      ;
                      ;
                      ;  High level diskette drivers
                      ;
                      ;  These drivers perform the following functions:
                      ;
                      ;  SELDSK  Select a specified disk and return the address of
                      ;          the appropriate disk parameter header
                      ;  SETTRK  Set the track number for the next read or write
                      ;  SETSEC  Set the sector number for the next read or write
                      ;  SETDMA  Set the DMA (read/write) address for the next read or write.
                      ;  SECTRAN Translate a logical sector number into a physical
                      ;  HOME    Set the track to 0 so that the next read or write will
                      ;          be on Track 0
                      ;
                      ;  In addition, the high-level drivers are responsible for making
                      ;  the 5 1/4" floppy diskettes that use a 512-byte sector appear
                      ;  to CP/M as though they used a 128-byte sector. They do this
                      ;  by using what is called blocking/deblocking code,
                      ;  described in more detail later in this listing,
                      ;  just prior to the code itself.
                      ;
                      ;
                      ;
                      ;  Disk parameter tables
                      ;
                      ;  As discussed in Chapter 3, these describe the physical
                      ;  characteristics of the disk drives. In this example BIOS,
                      ;  there are two types of disk drives; standard single-sided,
                      ;  single-density 8", and double-sided, double-density 5 1/4"
                      ;  diskettes.
                      ;
                      ;  The standard 8" diskettes do not need to use the blocking/
                      ;  deblocking code, but the 5 1/4" drives do. Therefore an additional
                      ;  byte has been prefixed to the disk parameter block to
                      ;  tell the disk drivers each logical disk's physical
                      ;  diskette type, and whether or not it needs deblocking.
                      ;
                      ;
                      ;  Disk definition tables
                      ;
                      ;  These consist of disk parameter headers, with one entry
                      ;  per logical disk driver, and disk parameter blocks, with
                      ;  either one parameter block per logical disk or the same
                      ;  parameter block for several logical disks.
                      ;
                      ;
                      Disk$Parameter$Headers:                       ;Described in Chapter 3
                      ;
                                         ;Logical Disk A: (5 1/4" Diskette)
F981 6BFB                DW     Floppy$5$Skewtable                  ;5 1/4" skew table
F983 0000000000          DW     0,0,0                               ;Reserved for CP/M
F989 C1F9                DW     Directory$Buffer
F98B 42FA                DW     Floppy$5$Parameter$Block
F98D 61FA                DW     Disk$A$Workarea
F98F C1FA                DW     Disk$A$Allocation$Vector
                      ;
                                         ;Logical Disk B: (5 1/4" Diskette)
F991 6BFB                DW     Floppy$5$Skewtable                  ;Shares same skew table as A:
```

Figure 6-4. (Continued)

```
F993 0000000000   DW      0,0,0                           ;Reserved for CP/M
F999 C1F9         DW      Directory$Buffer                ;Share same buffer as A:
F99B 42FA         DW      Floppy$5$Parameter$Block        ;Same DPB as A:
F99D 81FA         DW      Disk$B$Workarea                 ;Private work area
F99F D7FA         DW      Disk$B$Allocation$Vector        ;Private allocation vector
                ;
                              ;Logical Disk C: (8" Floppy)
F9A1 B3FB         DW      Floppy$8$Skewtable              ;8" skew table
F9A3 0000000000   DW      0,0,0                           ;Reserved for CP/M
F9A9 C1F9         DW      Directory$Buffer                ;Share same buffer as A:
F9AB 52FA         DW      Floppy$8$Parameter$Block
F9AD A1FA         DW      Disk$C$Workarea                 ;Private work area
F9AF EDFA         DW      Disk$C$Allocation$Vector        ;Private allocation vector
                ;
                              ;Logical Disk D: (8" Floppy)
F9B1 6BFB         DW      Floppy$5$Skewtable              ;Shares same skew table as A:
F9B3 0000000000   DW      0,0,0                           ;Reserved for CP/M
F9B9 C1F9         DW      Directory$Buffer                ;Share same buffer as A:
F9BB 52FA         DW      Floppy$8$Parameter$Block        ;Same DPB as C:
F9BD B1FA         DW      Disk$D$Workarea                 ;Private work area
F9BF 0CFB         DW      Disk$D$Allocation$Vector        ;Private allocation vector

                ;
                ;
F9C1            Directory$Buffer:  DS      128
                ;
                ;
                ;
                ;  Disk Types
                ;
0001 =          Floppy$5    EQU     1          ;5 1/4" mini floppy
0002 =          Floppy$8    EQU     2          ;8" floppy (SS SD)
                ;
                ;  Blocking/deblocking indicator
                ;
0080 =          Need$Deblocking    EQU     1000$0000B      ;Sector size > 128 bytes
                ;
                ;
                ;  Disk parameter blocks
                ;
                ;  5 1/4" mini floppy
                ;
                                           ;Extra byte prefixed to indicate
                                           ; disk type and blocking required
FA41 81           DB      Floppy$5 + Need$Deblocking
                Floppy$5$Parameter$Block:
FA42 4800         DW      72               ;128-byte sectors per track
FA44 04           DB      4                ;Block shift
FA45 0F           DB      15               ;Block mask
FA46 01           DB      1                ;Extent mask
FA47 AE00         DW      174              ;Maximum allocation block number
FA49 7F00         DW      127              ;Number of directory entries - 1
FA4B C0           DB      1100$0000B       ;Bit map for reserving 1 alloc. block
FA4C 00           DB      0000$0000B       ;  for file directory
FA4D 2000         DW      32               ;Disk changed work area size
FA4F 0100         DW      1                ;Number of tracks before directory
                ;
                ;
                ;  Standard 8" Floppy
                                           ;Extra byte prefixed to DPB for
                                           ; this version of the BIOS
FA51 02           DB      Floppy$8         ;Indicates disk type and the fact
                                           ; that no deblocking is required
                Floppy$8$Parameter$Block:
FA52 1A00         DW      26               ;Sectors per track
FA54 03           DB      3                ;Block shift
FA55 07           DB      7                ;Block mask
FA56 00           DB      0                ;Extent mask
FA57 F200         DW      242              ;Maximum allocation block number
FA59 3F00         DW      63               ;Number of directory entries - 1
FA5B C0           DB      1100$0000B       ;Bit map for reserving 2 alloc. blocks
FA5C 00           DB      0000$0000B       ;  for file directory
FA5D 1000         DW      16               ;Disk changed work area size
FA5F 0200         DW      2                ;Number of tracks before directory
                ;
                ;
```

Figure 6-4. (Continued)

```
                        ;  Disk work areas
                        ;
                        ;  These are used by the BDOS to detect any unexpected
                        ;  change of diskettes. The BDOS will automatically set
                        ;  such a changed diskette to read-only status.
                        ;
FA61                    Disk$A$Workarea:   DS      32      ; A:
FA81                    Disk$B$Workarea:   DS      32      ; B:
FAA1                    Disk$C$Workarea:   DS      16      ; C:
FAB1                    Disk$D$Workarea:   DS      16      ; D:
                        ;
                        ;
                        ;  Disk allocation vectors
                        ;
                        ;  These are used by the BDOS to maintain a bit map of
                        ;  which allocation blocks are used and which are free.
                        ;  One byte is used for eight allocation blocks, hence the
                        ;  expression of the form (allocation blocks/8)+1.
                        ;
FAC1                    Disk$A$Allocation$Vector   DS      (174/8)+1       ; A:
FAD7                    Disk$B$Allocation$Vector   DS      (174/8)+1       ; B:
                        ;
FAED                    Disk$C$Allocation$Vector   DS      (242/8)+1       ; C:
FB0C                    Disk$D$Allocation$Vector   DS      (242/8)+1       ; D:
                        ;
                        ;
0004 =                  Number$of$Logical$Disks            EQU     4
                        ;
                        ;
                        SELDSK:                            ;Select disk in C
                                                           ;C = 0 for drive A, 1 for B, etc.
                                                           ;Return the address of the appropriate
                                                           ; disk parameter header in HL, or 0000H
                                                           ; if the selected disk does not exist.
                                                           ;
FB2B 210000                LXI     H,0                     ;Assume an error
FB2E 79                    MOV     A,C                     ;Check if requested disk valid
FB2F FE04                  CPI     Number$of$Logical$Disks
FB31 D0                    RNC                             ;Return if > maximum number of disks
                        ;
FB32 32EAFB                STA     Selected$Disk           ;Save selected disk number
                                                           ;Set up to return DPH address
FB35 6F                    MOV     L,A                     ;Make disk into word value
FB36 2600                  MVI     H,0
                                                           ;Compute offset down disk parameter
                                                           ; header table by multiplying by
                                                           ; parameter header length (16 bytes)
FB38 29                    DAD     H                       ; *2
FB39 29                    DAD     H                       ; *4
FB3A 29                    DAD     H                       ; *8
FB3B 29                    DAD     H                       ; *16
FB3C 1181F9                LXI     D,Disk$Parameter$Headers        ;Get base address
FB3F 19                    DAD     D                       ;DE -> Appropriate DPH
FB40 E5                    PUSH    H                       ;Save DPH address
                        ;
                                                           ;Access disk parameter block
                                                           ; to extract special prefix byte that
                                                           ; identifies disk type and whether
                                                           ; deblocking is required
                                                           ;
FB41 110A00                LXI     D,10                    ;Get DPB pointer offset in DPH
FB44 19                    DAD     D                       ;DE -> DPB address in DPH
FB45 5E                    MOV     E,M                     ;Get DPB address in DE
FB46 23                    INX     H
FB47 56                    MOV     D,M
FB48 EB                    XCHG                            ;DE -> DPB
FB49 2B                    DCX     H                       ;DE -> prefix byte
FB4A 7E                    MOV     A,M                     ;Get prefix byte
FB4B E60F                  ANI     0FH                     ;Isolate disk type
FB4D 32FAFB                STA     Disk$Type               ;Save for use in low-level driver
FB50 7E                    MOV     A,M                     ;Get another copy of prefix byte
FB51 E680                  ANI     Need$Deblocking                 ;Isolate deblocking flag
FB53 32F9FB                STA     Deblocking$Required     ;Save for use in low-level driver
FB56 E1                    POP     H                       ;Recover DPH pointer
FB57 C9                    RET
                        ;
```

Figure 6-4. (Continued)

```
                         ;
                         ;  Set logical track for next read or write
                         ;
                         SETTRK:
FB58 60                     MOV     H,B              ;Selected track in BC on entry
FB59 69                     MOV     L,C
FB5A 22EBFB                 SHLD    Selected$Track   ;Save for low-level driver
FB5D C9                     RET
                         ;
                         ;
                         ;  Set logical sector for next read or write
                         ;
                         ;
                         SETSEC:                      ;Logical sector in C on entry
FB5E 79                     MOV     A,C
FB5F 32EDFB                 STA     Selected$Sector  ;Save for low-level driver
FB62 C9                     RET
                         ;
                         ;
                         ;  Set disk DMA (input/output) address for next read or write
                         ;
FB63 0000                DMA$Address:       DW       0        ;DMA address
                         ;
                         SETDMA:                      ;Address in BC on entry
FB65 69                     MOV     L,C              ;Move to HL to save
FB66 60                     MOV     H,B
FB67 2263FB                 SHLD    DMA$Address      ;Save for low-level driver
FB6A C9                     RET
                         ;
                         ;
                         ;  Translate logical sector number to physical
                         ;
                         ;  Sector translation tables
                         ;  These tables are indexed using the logical sector number,
                         ;  and contain the corresponding physical sector number.
                         ;
                         Floppy$5$Skewtable:          ;Each physical sector contains four
                                                      ; 128-byte sectors.
                            ;      Physical 128b   Logical 128b       Physical 512-byte
FB6B 00010203               DB     00,01,02,03     ;00,01,02,03          0  )
FB6F 10111213               DB     16,17,18,19     ;04,05,06,07          4  )
FB73 20212223               DB     32,33,34,35     ;08,09,10,11          8  )
FB77 0C0D0E0F               DB     12,13,14,15     ;12,13,14,15          3  ) Head
FB7B 1C1D1E1F               DB     28,29,30,31     ;16,17,18,19          7  )  0
FB7F 08090A0B               DB     08,09,10,11     ;20,21,22,23          2  )
FB83 18191A1B               DB     24,25,26,27     ;24,25,26,27          6  )
FB87 04050607               DB     04,05,06,07     ;28,29,30,31          1  )
FB8B 14151617               DB     20,21,22,23     ;32,33,34,35          5  )
                         ;
FB8F 24252627               DB     36,37,38,39     ;36,37,38,39          0  ]
FB93 34353637               DB     52,53,54,55     ;40,41,42,43          4  ]
FB97 44454647               DB     68,69,70,71     ;44,45,46,47          8  ]
FB9B 30313233               DB     48,49,50,51     ;48,49,50,51          3  ] Head
FB9F 40414243               DB     64,65,66,67     ;52,53,54,55          7  ]  1
FBA3 2C2D2E2F               DB     44,45,46,47     ;56,57,58,59          2  ]
FBA7 3C3D3E3F               DB     60,61,62,63     ;60,61,62,63          6  ]
FBAB 28292A2B               DB     40,41,42,43     ;64,65,66,67          1  ]
FBAF 38393A3B               DB     56,57,58,59     ;68,69,70,71          5  ]
                         ;
                         ;
                         Floppy$8$Skewtable:          ;Standard 8" Driver
                            ;       01,02,03,04,05,06,07,08,09,10    Logical sectors
FBB3 01070D1319             DB      01,07,13,19,25,05,11,17,23,03   ;Physical sectors
                            ;
                            ;       11,12,13,14,15,16,17,18,19,20    Logical sectors
FBBD 090F150208             DB      09,15,21,02,08,14,20,26,06,12   ;Physical sectors
                            ;
                            ;       21,22,23,24,25,26        Logical sectors
FBC7 1218040A10             DB      18,24,04,10,16,22       ;Physical sectors
                         ;
                         ;
                         SECTRAN:                     ;Translate logical sector into physical
                                                      ;On entry, BC = logical sector number
                                                      ;         DE -> appropriate skew table
                                                      ;
                                                      ;on exit, HL = physical sector number
```

Figure 6-4. (Continued)

```
FBCD EB              XCHG                         ;HL -> skew table base
FBCE 09              DAD     B                    ;Add on logical sector number
FBCF 6E              MOV     L,M                  ;Get physical sector number
FBD0 2600            MVI     H,0                  ;Make into a 16-bit value
FBD2 C9              RET
                 ;
                 ;
                 ;
                 HOME:                            ;Home the selected logical disk to track 0.
                                                  ;Before doing this, a check must be made to see
                                                  ; if the physical disk buffer has information
                                                  ; that must be written out. This is indicated by
                                                  ; a flag, Must$Write$Buffer, set in the
                                                  ; deblocking code.
                                                  ;
FBD3 3AE9FB          LDA     Must$Write$Buffer           ;Check if physical buffer must
FBD6 B7              ORA     A                           ;  be written out to disk
FBD7 C2DDFB          JNZ     HOME$No$Write
FBDA 32E8FB          STA     Data$In$Disk$Buffer         ;No, so indicate that buffer
                                                         ;  is now unoccupied.
                 HOME$No$Write:
FBDD 0E00            MVI     C,0                         ;Set to track 0 (logically --
FBDF CD58FB          CALL    SETTRK                      ;  no actual disk operation occurs)
FBE2 C9              RET

                 ;
                 ;  Data written to or read from the mini-floppy drive is transferred
                 ;  via a physical buffer that is actually 512 bytes long (it was
                 ;  declared at the front of the BIOS and holds the "one-time"
                 ;  initialization code used for the cold boot procedure).
                 ;
                 ;  The blocking/deblocking code attempts to minimize the amount
                 ;  of actual disk I/O by storing the disk, track, and physical sector
                 ;  currently residing in the Physical Buffer. If a read request is for
                 ;  a 128-byte CP/M "sector" that already is in the physical buffer,
                 ;  then no disk access occurs.
                 ;
                 ;
0800 =           Allocation$Block$Size      EQU     2048
0012 =           Physical$Sec$Per$Track     EQU     18
0004 =           CPM$Sec$Per$Physical       EQU     Physical$Sector$Size/128
0048 =           CPM$Sec$Per$Track          EQU     CPM$Sec$Per$Physical*Physical$Sec$Per$Track
0003 =           Sector$Mask                EQU     CPM$Sec$Per$Physical-1
0002 =           Sector$Bit$Shift           EQU     2         ;LOG2(CPM$Sec$Per$Physical)
                 ;
                                   ;These are the values handed over by the BDOS
                                   ;  when it calls the WRITE operation.
                                   ;The allocated/unallocated indicates whether the
                                   ;  BDOS is set to write to an unallocated allocation
                                   ;  block (it only indicates this for the first
                                   ;  128-byte sector write)  or to an allocation block
                                   ;  that has already been allocated to a file.
                                   ;The BDOS also indicates if it is set to write to
                                   ;  the file directory.
                                   ;
0000 =           Write$Allocated            EQU     0
0001 =           Write$Directory            EQU     1
0002 =           Write$Unallocated          EQU     2
                 ;
FBE3 00          Write$Type:                DB      0         ;Contains the type of write
                                                              ;  indicated by the BDOS.
                 ;
                 ;
                 In$Buffer$Dk$Trk$Sec:                        ;Variables for physical sector
                                                              ; currently in Disk$Buffer in memory
FBE4 00          In$Buffer$Disk:            DB      0         ; These are moved and compared
FBE5 0000        In$Buffer$Track:           DW      0         ; as a group, so do not alter
FBE7 00          In$Buffer$Sector:          DB      0         ; these lines.
                 ;
FBE8 00          Data$In$Disk$Buffer:       DB      0         ;When nonzero, the disk buffer has
                                                              ;  data from the disk in it.
FBE9 00          Must$Write$Buffer:         DB      0         ;Nonzero when data has been
                                                              ;  written into Disk$Buffer but
                                                              ;  not yet written out to disk
                 ;
                 Selected$Dk$Trk$Sec:       ;Variables for selected disk, track, and sector
```

Figure 6-4. (Continued)

```
                                                      ; (Selected by SELDSK, SETTRK,and SETSEC)
FBEA 00         Selected$Disk:                  DB      0       ; These are moved and
FBEB 0000       Selected$Track:                 DW      0       ; compared as a group so
FBED 00         Selected$Sector:                DB      0       ; do not alter order.

FBEE 00         Selected$Physical$Sector:       DB      0       ;Selected physical sector derived
                                                                ;  from selected (CP/M) sector by
                                                                ;  shifting it right the number of
                                                                ;  of bits specified by
                                                                ;  Sector$Bit$Shift
                ;
FBEF 00         Selected$Disk$Type:             DB      0       ;Set by SELDSK to indicate either
                                                                ;  8" or 5 1/4" floppy
FBF0 00         Selected$Disk$Deblock:          DB      0       ;Set by SELDSK to indicate whether
                                                                ;  deblocking is required.


                Unallocated$Dk$Trk$Sec:                 ;Parameters for writing to a previously
                                                                ;  unallocated allocation block.
FBF1 00         Unallocated$Disk:               DB      0       ; These are moved and compared
FBF2 0000       Unallocated$Track:              DW      0       ; as a group so do not alter
FBF4 00         Unallocated$Sector:             DB      0       ; these lines.

FBF5 00         Unallocated$Record$Count:       DB      0       ;Number of unallocated "records"
                                                                ; in current previously unallocated
                                                                ; allocation block.

FBF6 00         Disk$Error$Flag:                DB      0       ;Nonzero to indicate an error
                                                                ;  that could not be recovered
                                                                ;  by the disk drivers. BDOS will
                                                                ;  output a "bad sector" message.
                ;
                ;Flags used inside the deblocking code

FBF7 00         Must$Preread$Sector:            DB      0       ;Nonzero if a physical sector must
                                                                ;  be read into the disk buffer
                                                                ;  either before a write to an
                                                                ;  allocated block can occur, or
                                                                ;  for a normal CP/M 128-byte
                                                                ;  sector read
FBF8 00         Read$Operation:                 DB      0       ;Nonzero when a CP/M 128-byte
                                                                ;  sector is to be read
FBF9 00         Deblocking$Required:            DB      0       ;Nonzero when the selected disk
                                                                ;  needs deblocking (set in SELDSK)
FBFA 00         Disk$Type:                      DB      0       ;Indicates 8" or 5 1/4" floppy
                                                                ;  selected (set in SELDSK).
                ;
                ;
                ;   Read in the 128-byte CP/M sector specified by previous calls
                ;   to select disk and to set track and sector. The sector will be read
                ;   into the address specified in the previous call to set DMA address.
                ;
                ;   If reading from a disk drive using sectors larger than 128 bytes,
                ;   deblocking code will be used to "unpack" a 128-byte sector from
                ;   the physical sector.
                READ:
FBFB 3AF9FB         LDA     Deblocking$Required     ;Check if deblocking needed
FBFE B7             ORA     A                       ;(flag was set in SELDSK call)
FBFF CA52FD         JZ      Read$No$Deblock         ;No, use normal nondeblocked

                                ;The deblocking algorithm used is such
                                ;  that a read operation can be viewed
                                ;  up until the actual data transfer as
                                ;  though it was the first write to an
                                ;  unallocated allocation block.
FC02 AF             XRA     A                       ;Set the record count to 0
FC03 32F5FB         STA     Unallocated$Record$Count ;  for first "write"
FC06 3C             INR     A                       ;Indicate that it is really a read
FC07 32F8FB         STA     Read$Operation          ;  that is to be performed
FC0A 32F7FB         STA     Must$Preread$Sector     ;  and force a preread of the sector
                                                    ;  to get it into the disk buffer
FC0D 3E02           MVI     A,Write$Unallocated     ;Fake deblocking code into responding
FC0F 32E3FB         STA     Write$Type              ;  as if this is the first write to an
                                                    ;  unallocated allocation block.
FC12 C36EFC         JMP     Perform$Read$Write      ;Use common code to execute read
```

Figure 6-4. (Continued)

```
                        ;
                        ;  Write a 128-byte sector from the current DMA address to
                        ;  the previously selected disk, track, and sector.
                        ;
                        ;  On arrival here, the BDOS will have set register C to indicate
                        ;  whether this write operation is to an already allocated allocation
                        ;  block (which means a preread of the sector may be needed),
                        ;  to the directory (in which case the data will be written to the
                        ;  disk immediately), or to the first 128-byte sector of a previously
                        ;  unallocated allocation block (in which case no preread is required).
                        ;
                        ;  Only writes to the directory take place immediately. In all other
                        ;  cases, the data will be moved from the DMA address into the disk
                        ;  buffer, and only written out when circumstances force the
                        ;  transfer. The number of physical disk operations can therefore
                        ;  be reduced considerably.
                        ;
                        WRITE:
FC15 3AF9FB                 LDA     Deblocking$Required     ;Check if deblocking is required
FC18 B7                     ORA     A                       ;(flag set in SELDSK call)
FC19 CA4DFD                 JZ      Write$No$Deblock

FC1C AF                     XRA     A                       ;Indicate that a write operation
FC1D 32F8FB                 STA     Read$Operation          ;  is required (i.e. NOT a read)
FC20 79                     MOV     A,C                     ;Save the BDOS write type
FC21 32E3FB                 STA     Write$Type
FC24 FE02                   CPI     Write$Unallocated       ;Check if the first write to an
                                                            ;  unallocated allocation block
FC26 C237FC                 JNZ     Check$Unallocated$Block ;No, check if in the middle of
                                                            ;  writing to an unallocated block
                                                            ;Yes, first write to unallocated
                                                            ;  allocation block -- initialize
                                                            ;  variables associated with
                                                            ;  unallocated writes.
FC29 3E10                   MVI     A,Allocation$Block$Size/128     ;Get number of 128-byte
                                                                    ;  sectors and
FC2B 32F5FB                 STA     Unallocated$Record$Count        ;  set up a count.
                                                                    ;
FC2E 21EAFB                 LXI     H,Selected$Dk$Trk$Sec           ;Copy disk, track, and sector
FC31 11F1FB                 LXI     D,Unallocated$Dk$Trk$Sec        ;  into unallocated variables
FC34 CD35FD                 CALL    Move$Dk$Trk$Sec
                        ;
                        ;  Check if this is not the first write to an unallocated
                        ;  allocation block -- if it is, the unallocated record count
                        ;  has just been set to the number of 128-byte sectors in the
                        ;  allocation block.
                        ;
                        Check$Unallocated$Block:
FC37 3AF5FB                 LDA     Unallocated$Record$Count
FC3A B7                     ORA     A
FC3B CA66FC                 JZ      Request$Preread         ;No, this is a write to an
                                                            ;  allocated block
                                                            ;Yes, this is a write to an
                                                            ;  unallocated block
FC3E 3D                     DCR     A                       ;Count down on number of 128-byte sectors
                                                            ;  left unwritten to in allocation block
FC3F 32F5FB                 STA     Unallocated$Record$Count        ;  and store back new value.

FC42 21EAFB                 LXI     H,Selected$Dk$Trk$Sec   ;Check if the selected disk, track,
FC45 11F1FB                 LXI     D,Unallocated$Dk$Trk$Sec;  and sector are the same as for
FC48 CD29FD                 CALL    Compare$Dk$Trk$Sec      ;  those in the unallocated block.
FC4B C266FC                 JNZ     Request$Preread         ;No, a preread is required
                                                            ;Yes, no preread is needed.
                                                            ;Now is a convenient time to
                                                            ;  update the current sector and see
                                                            ;  if the track also needs updating.
                        ;
                                                            ;By design, Compare$Dk$Trk$Sec
                                                            ; returns with
                                                            ; DE -> Unallocated$Sector
FC4E EB                     XCHG                            ; HL -> Unallocated$Sector
FC4F 34                     INR     M                       ;Update Unallocated$Sector
FC50 7E                     MOV     A,M                     ;Check if sector now > maximum
FC51 FE48                   CPI     CPM$Sec$Per$Track       ; on a track
FC53 DA5FFC                 JC      No$Track$Change         ;No (A < M)
                                                            ;Yes,
```

Figure 6-4. (Continued)

```
FC56 3600        MVI    M,0                    ;Reset sector to 0
FC58 2AF2FB      LHLD   Unallocated$Track      ;Increase track by 1
FC5B 23          INX    H
FC5C 22F2FB      SHLD   Unallocated$Track
                ;
                No$Track$Change:
                                               ;Indicate to later code that
                                               ;  no preread is needed.
FC5F AF          XRA    A
FC60 32F7FB      STA    Must$Preread$Sector    ;Must$Preread$Sector=0
FC63 C36EFC      JMP    Perform$Read$Write
                ;
                Request$Preread:
FC66 AF          XRA    A                      ;Indicate that this is not a write
FC67 32F5FB      STA    Unallocated$Record$Count       ;  into an unallocated block.
FC6A 3C          INR    A
FC6B 32F7FB      STA    Must$Preread$Sector    ;Indicate that a preread of the
                                               ;  physical sector is required.
                ;
                ;
                Perform$Read$Write:            ;Common code to execute both reads and
                                               ;  writes of 128-byte sectors.
FC6E AF          XRA    A                      ;Assume that no disk errors will
FC6F 32F6FB      STA    Disk$Error$Flag        ;  occur

FC72 3AEDFB      LDA    Selected$Sector        ;Convert selected 128-byte sector
FC75 1F          RAR                           ;  into physical sector by dividing by 4
FC76 1F          RAR
FC77 E63F        ANI    3FH                    ;Remove any unwanted bits
FC79 32EEFB      STA    Selected$Physical$Sector
                                               ;
FC7C 21E8FB      LXI    H,Data$In$Disk$Buffer  ;Check if disk buffer already has
FC7F 7E          MOV    A,M                    ;  data in it.
FC80 3601        MVI    M,1                    ;(Unconditionally indicate that
                                               ; the buffer now has data in it)
FC82 B7          ORA    A                      ;Did it indeed have data in it?
FC83 CAA3FC      JZ     Read$Sector$into$Buffer ;No, proceed to read a physical
                                               ;  sector into the buffer.
                                               ;
                                   ;The buffer does have a physical sector
                                   ;  in it.
                                   ;  Note: The disk, track, and PHYSICAL
                                   ;  sector in the buffer need to be
                                   ;  checked, hence the use of the
                                   ;  Compare$Dk$Trk subroutine.
                                   ;
FC86 11E4FB      LXI    D,In$Buffer$Dk$Trk$Sec ;Check if sector in buffer is the
FC89 21EAFB      LXI    H,Selected$Dk$Trk$Sec  ;  same as that selected earlier
FC8C CD24FD      CALL   Compare$Dk$Trk         ;Compare ONLY disk and track
FC8F C29CFC      JNZ    Sector$Not$In$Buffer   ;No, it must be read in

FC92 3AE7FB      LDA    In$Buffer$Sector       ;Get physical sector in buffer
FC95 21EEFB      LXI    H,Selected$Physical$Sector
FC98 BE          CMP    M                      ;Check if correct physical sector
FC99 CAB1FC      JZ     Sector$In$Buffer       ;Yes, it is already in memory
                ;
                Sector$Not$In$Buffer:
                                               ;No, it will have to be read in
                                               ;  over current contents of buffer
FC9C 3AE9FB      LDA    Must$Write$Buffer      ;Check if buffer has data in that
FC9F B7          ORA    A                      ;  must be written out first
FCA0 C495FD      CNZ    Write$Physical         ;Yes, write it out
                ;
                Read$Sector$into$Buffer:
FCA3 CD11FD      CALL   Set$In$Buffer$Dk$Trk$Sec       ;Set in buffer variables from
                                               ; selected disk, track, and sector
                                               ;  to reflect which sector is in the
                                               ;  buffer now
FCA6 3AF7FB      LDA    Must$Preread$Sector    ;In practice, the sector need only
FCA9 B7          ORA    A                      ;  be physically read in if a preread
                                               ;  is required
FCAA C49AFD      CNZ    Read$Physical          ;Yes, preread the sector
FCAD AF          XRA    A                      ;Reset the flag to reflect buffer
FCAE 32E9FB      STA    Must$Write$Buffer      ;  contents.
                ;
                Sector$In$Buffer:          ;Selected sector on correct track and
```

Figure 6-4. (Continued)

```
                                                 ;  disk is already in the buffer.
                                                 ;Convert the selected CP/M (128-byte)
                                                 ;  sector into a relative address down
                                                 ;  the buffer.
FCB1 3AEDFB         LDA     Selected$Sector      ;Get selected sector number
FCB4 E603           ANI     Sector$Mask          ;Mask off only the least significant bits
FCB6 6F             MOV     L,A                  ;Multiply by 128 by shifting 16-bit value
FCB7 2600           MVI     H,0                  ;  left 7 bits
FCB9 29             DAD     H                    ;* 2
FCBA 29             DAD     H                    ;* 4
FCBB 29             DAD     H                    ;* 8
FCBC 29             DAD     H                    ;* 16
FCBD 29             DAD     H                    ;* 32
FCBE 29             DAD     H                    ;* 64
FCBF 29             DAD     H                    ;* 128
               ;
FCC0 1133F6         LXI     D,Disk$Buffer        ;Get base address of disk buffer
FCC3 19             DAD     D                    ;Add on sector number * 128
                                                 ;HL -> 128-byte sector number start
                                                 ;  address in disk buffer
FCC4 EB             XCHG                         ;DE -> sector in disk buffer
FCC5 2A63FB         LHLD    DMA$Address          ;Get DMA address set in SETDMA call
FCC8 EB             XCHG                         ;Assume a read operation, so
                                                 ;  DE -> DMA address
                                                 ;  HL -> sector in disk buffer
FCC9 0E10           MVI     C,128/8              ;Because of the faster method used
                                                 ;  to move data in and out of the
                                                 ;  disk buffer, (eight bytes moved per
                                                 ;  loop iteration) the count need only
                                                 ;  be 1/8th of normal.
                                                 ;At this point -
                                                 ;       C = loop count
                                                 ;       DE -> DMA address
                                                 ;       HL -> sector in disk buffer
FCCB 3AF8FB         LDA     Read$Operation       ;Determine whether data is to be moved
FCCE B7             ORA     A                    ;  out of the buffer (read) or into the
FCCF C2D7FC         JNZ     Buffer$Move          ;  buffer (write)
                                                 ;Writing into buffer
                                                         ;(A must be 0 get here)
FCD2 3C             INR     A                            ;Set flag to force a write
FCD3 32E9FB         STA     Must$Write$Buffer            ;  of the disk buffer later on.
FCD6 EB             XCHG                                 ;Make DE -> sector in disk buffer
                                                         ;      HL -> DMA address
               ;
               ;
               Buffer$Move:                      ;The folowing move loop moves eight bytes
                                                 ;  at a time from (HL) to (DE), C contains
                                                 ;  the loop count.
FCD7 7E             MOV     A,M                  ;Get byte from source
FCD8 12             STAX    D                    ;Put into destination
FCD9 13             INX     D                    ;Update pointers
FCDA 23             INX     H
FCDB 7E             MOV     A,M                  ;Get byte from source
FCDC 12             STAX    D                    ;Put into destination
FCDD 13             INX     D                    ;Update pointers
FCDE 23             INX     H
FCDF 7E             MOV     A,M                  ;Get byte from source
FCE0 12             STAX    D                    ;Put into destination
FCE1 13             INX     D                    ;Update pointers
FCE2 23             INX     H
FCE3 7E             MOV     A,M                  ;Get byte from source
FCE4 12             STAX    D                    ;Put into destination
FCE5 13             INX     D                    ;Update pointers
FCE6 23             INX     H
FCE7 7E             MOV     A,M                  ;Get byte from source
FCE8 12             STAX    D                    ;Put into destination
FCE9 13             INX     D                    ;Update pointers
FCEA 23             INX     H
FCEB 7E             MOV     A,M                  ;Get byte from source
FCEC 12             STAX    D                    ;Put into destination
FCED 13             INX     D                    ;Update pointers
FCEE 23             INX     H
FCEF 7E             MOV     A,M                  ;Get byte from source
FCF0 12             STAX    D                    ;Put into destination
FCF1 13             INX     D                    ;Update pointers
```

Figure 6-4. (Continued)

```
FCF2 23         INX     H
FCF3 7E         MOV     A,M             ;Get byte from source
FCF4 12         STAX    D               ;Put into destination
FCF5 13         INX     D               ;Update pointers
FCF6 23         INX     H

FCF7 0D         DCR     C               ;Count down on loop counter
FCF8 C2D7FC     JNZ     Buffer$Move     ;Repeat until CP/M sector moved
                                        ;
FCFB 3AE3FB     LDA     Write$Type      ;If write to directory, write out
FCFE FE01       CPI     Write$Directory ;  buffer immediately
FD00 3AF6FB     LDA     Disk$Error$Flag ;Get error flag in case delayed write or read
FD03 C0         RNZ                     ;Return if delayed write or read
                                        ;
FD04 B7         ORA     A               ;Check if any disk errors have occurred
FD05 C0         RNZ                     ;Yes, abandon attempt to write to directory
                                        ;
FD06 AF         XRA     A               ;Clear flag that indicates buffer must be
FD07 32E9FB     STA     Must$Write$Buffer       ;  written out
FD0A CD95FD     CALL    Write$Physical  ;Write buffer out to physical sector
FD0D 3AF6FB     LDA     Disk$Error$Flag ;Return error flag to caller
FD10 C9         RET
              ;
              ;
              Set$In$Buffer$Dk$Trk$Sec:         ;Indicate selected disk, track, and
                                                ;  sector now residing in buffer
FD11 3AEAFB     LDA     Selected$Disk
FD14 32E4FB     STA     In$Buffer$Disk

FD17 2AEBFB     LHLD    Selected$Track
FD1A 22E5FB     SHLD    In$Buffer$Track

FD1D 3AEEFB     LDA     Selected$Physical$Sector
FD20 32E7FB     STA     In$Buffer$Sector

FD23 C9         RET
              ;
              Compare$Dk$Trk:           ;Compares just the disk and track
                                        ;  pointed to by DE and HL
FD24 0E03       MVI     C,3             ;Disk (1), track (2)
FD26 C32BFD     JMP     Compare$Dk$Trk$Sec$Loop ;Use common code

              Compare$Dk$Trk$Sec:       ;Compares the disk, track, and sector
                                        ;  variables pointed to by DE and HL
FD29 0E04       MVI     C,4             ;Disk (1), track (2), and sector (1)
              Compare$Dk$Trk$Sec$Loop:
FD2B 1A         LDAX    D               ;Get comparitor
FD2C BE         CMP     M               ;Compare with comparand
FD2D C0         RNZ                     ;Abandon comparison if inequality found
FD2E 13         INX     D               ;Update comparitor pointer
FD2F 23         INX     H               ;Update comparand pointer
FD30 0D         DCR     C               ;Count down on loop count
FD31 C8         RZ                      ;Return (with zero flag set)
FD32 C32BFD     JMP     Compare$Dk$Trk$Sec$Loop
              ;
              ;
              Move$Dk$Trk$Sec:          ;Moves the disk, track, and sector
                                        ;  variables pointed at by HL to
                                        ;  those pointed at by DE
FD35 0E04       MVI     C,4             ;Disk (1), track (2), and sector (1)
              Move$Dk$Trk$Sec$Loop:
FD37 7E         MOV     A,M             ;Get source byte
FD38 12         STAX    D               ;Store in destination
FD39 13         INX     D               ;Update pointers
FD3A 23         INX     H
FD3B 0D         DCR     C               ;Count down on byte count
FD3C C8         RZ                      ;Return if all bytes moved
FD3D C337FD     JMP     Move$Dk$Trk$Sec$Loop
              ;
              ;
              ;
              ;   There are two "smart" disk controllers on this system, one
              ;   for the 8" floppy diskette drives, and one for the 5 1/4"
              ;   mini-diskette drives.
              ;
              ;   The controllers are "hard-wired" to monitor certain locations
```

Figure 6-4. (Continued)

```
                ;  in memory to detect when they are to perform some disk
                ;  operation. The 8" controller monitors location 0040H, and
                ;  the 5 1/4" controller monitors location 0045H. These are
                ;  called their disk control bytes. If the most significant
                ;  bit of a disk control byte is set, the controller will
                ;  look at the word following the respective control bytes.
                ;  This word must contain the address of a valid disk control
                ;  table that specifies the exact disk operation to be performed.
                ;
                ;  Once the operation has been completed, the controller resets
                ;  its disk control byte to 00H. This indicates completion
                ;  to the disk driver code.
                ;
                ;  The controller also sets a return code in a disk status block --
                ;  both controllers use the SAME location for this; 0043H.
                ;  If the first byte of this status block is less than 80H, then
                ;  a disk error has occurred. For this simple BIOS, no further details
                ;  of the status settings are relevant. Note that the disk controller
                ;  has built-in retry logic -- reads and writes are attempted ten
                ;  times before the controller returns an error.
                ;
                ;  The disk control table layout is shown below. Note that the
                ;  controllers have the capability for control tables to be
                ;  chained together so that a sequence of disk operations can
                ;  be initiated. In this BIOS this feature is not used. However,
                ;  the controller requires that the chain pointers in the
                ;  disk control tables be pointed back to the main control bytes
                ;  in order to indicate the end of the chain.
                ;
0040 =          Disk$Control$8          EQU     40H     ;8" control byte
0041 =          Command$Block$8         EQU     41H     ;Control table pointer
                ;
0043 =          Disk$Status$Block       EQU     43H     ;8" AND 5 1/4" status block
                ;
0045 =          Disk$Control$5          EQU     45H     ;5 1/4" control byte
0046 =          Command$Block$5         EQU     46H     ;Control table pointer
                ;
                ;
                ;  Floppy Disk Control Tables
                ;
FD40 00         Floppy$Command:         DB      0       ;Command
0001 =          Floppy$Read$Code        EQU     01H
0002 =          Floppy$Write$Code       EQU     02H
FD41 00         Floppy$Unit:            DB      0       ;Unit (drive) number = 0 or 1
FD42 00         Floppy$Head:            DB      0       ;Head number = 0 or 1
FD43 00         Floppy$Track:           DB      0       ;Track number
FD44 00         Floppy$Sector:          DB      0       ;Sector number
FD45 0000       Floppy$Byte$Count:      DW      0       ;Number of bytes to read/write
FD47 0000       Floppy$DMA$Address:     DW      0       ;Transfer address
FD49 0000       Floppy$Next$Status$Block: DW    0       ;Pointer to next status block
                                                        ;  if commands are chained.
FD4B 0000       Floppy$Next$Control$Location: DW 0      ;Pointer to next control byte
                                                        ;  if commands are chained.
                ;
                ;
                ;
                Write$No$Deblock:                       ;Write contents of disk buffer to
                                                        ;  correct sector.
FD4D 3E02          MVI     A,Floppy$Write$Code          ;Get write function code
FD4F C354FD        JMP     Common$No$Deblock            ;Go to common code
                Read$No$Deblock:                        ;Read previously selected sector
                                                        ;  into disk buffer.
FD52 3E01          MVI     A,Floppy$Read$Code           ;Get read function code
                Common$No$Deblock:
FD54 3240FD        STA     Floppy$Command  ;Set command function code
                                           ;Set up nondeblocked command table
FD57 218000        LXI     H,128           ;Bytes per sector
FD5A 2245FD        SHLD    Floppy$Byte$Count
FD5D AF            XRA     A               ;8" floppy only has head 0
FD5E 3242FD        STA     Floppy$Head
                                           ;
FD61 3AEAFB        LDA     Selected$Disk   ;8" Floppy controller only has information
                                           ;  on units 0 and 1 so Selected$Disk must
                                           ;  be converted
FD64 E601          ANI     01H             ;Turn into 0 or 1
FD66 3241FD        STA     Floppy$Unit     ;Set unit number
```

Figure 6-4. (Continued)

```
                                                ;
FD69 3AEBFB         LDA     Selected$Track
FD6C 3243FD         STA     Floppy$Track    ;Set track number
                                            ;
FD6F 3AEDFB         LDA     Selected$Sector
FD72 3244FD         STA     Floppy$Sector   ;Set sector number
                                            ;
FD75 2A63FB         LHLD    DMA$Address     ;Transfer directly between DMA address
FD78 2247FD         SHLD    Floppy$DMA$Address      ;and 8" controller.
                                            ;
                                            ;The disk controller can accept chained
                                            ;  disk control tables, but in this case,
                                            ;  they are not used, so the "Next" pointers
                                            ;  must be pointed back at the initial
                                            ;  control bytes in the base page.
FD7B 214300         LXI     H,Disk$Status$Block             ;Point next status back at
FD7E 2249FD         SHLD    Floppy$Next$Status$Block        ;  main status block
                                                            ;
FD81 214000         LXI     H,Disk$Control$8                ;Point next control byte
FD84 224BFD         SHLD    Floppy$Next$Control$Location    ;  back at main control byte
                                                            ;
FD87 2140FD         LXI     H,Floppy$Command                ;Point controller at control table
FD8A 224100         SHLD    Command$Block$8
                                                            ;
FD8D 214000         LXI     H,Disk$Control$8                ;Activate controller to perform
FD90 3680           MVI     M,80H                           ;  operation.
FD92 C3F7FD         JMP     Wait$For$Disk$Complete

                ;
                ;

                Write$Physical:                         ;Write contents of disk buffer to
                                                        ;  correct sector.
FD95 3E02           MVI     A,Floppy$Write$Code         ;Get write function code
FD97 C39CFD         JMP     Common$Physical             ;Go to common code
                Read$Physical:                          ;Read previously selected sector
                                                        ;  into disk buffer.
FD9A 3E01           MVI     A,Floppy$Read$Code          ;Get read function code
                ;
                Common$Physical:
FD9C 3240FD         STA     Floppy$Command              ;Set command table

                ;
FD9F 3AFAFB         LDA     Disk$Type                   ;Get disk type (set in SELDSK)
FDA2 FE01           CPI     Floppy$5                    ;Confirm it is a 5 1/4" Floppy
FDA4 CAADFD         JZ      Correct$Disk$Type           ;Yes
FDA7 3E01           MVI     A,1                         ;No, indicate disk error
FDA9 32F6FB         STA     Disk$Error$Flag
FDAC C9             RET
                Correct$Disk$Type:                      ;Set up disk control table
                                                        ;
FDAD 3AE4FB         LDA     In$Buffer$Disk              ;Convert disk number to 0 or 1
FDB0 E601           ANI     1                           ;  for disk controller
FDB2 3241FD         STA     Floppy$Unit

FDB5 2AE5FB         LHLD    In$Buffer$Track             ;Set up track number
FDB8 7D             MOV     A,L                         ;Note: This is single byte value
FDB9 3243FD         STA     Floppy$Track                ; for the controller.
                                                        ;
                                                        ;The sector must be converted into a
                                                        ;  head number and sector number.
                                                        ;  Sectors 0 - 8 are head 0, 9 - 17
                                                        ;  are head 1
FDBC 0600           MVI     B,0                         ;Assume head 0
FDBE 3AE7FB         LDA     In$Buffer$Sector            ;Get physical sector number
FDC1 4F             MOV     C,A                         ;Save copy in case it is head 0
FDC2 FE09           CPI     9                           ;Check if < 9
FDC4 DACBFD         JC      Head$0                      ;Yes it is < 9
FDC7 D609           SUI     9                           ;No, modify sector number back
                                                        ;  in the 0 - 8 range.
FDC9 4F             MOV     C,A                         ;Put sector in B
FDCA 04             INR     B                           ;Set to head 1
                Head$0:
FDCB 78             MOV     A,B                         ;Set head number
FDCC 3242FD         STA     Floppy$Head
FDCF 79             MOV     A,C                         ;Set sector number
```

Figure 6-4. (Continued)

```
FDD0 3C           INR     A                         ;  (physical sectors start at 1)
FDD1 3244FD       STA     Floppy$Sector
                                                    ;
FDD4 210002       LXI     H,Physical$Sector$Size    ;Set byte count
FDD7 2245FD       SHLD    Floppy$Byte$Count
                                                    ;
FDDA 2133F6       LXI     H,Disk$Buffer             ;Set transfer address to be
FDDD 2247FD       SHLD    Floppy$DMA$Address        ;  disk buffer
                                                    ;
                                                    ;As only one control table is in
                                                    ;  use, close the status and busy
                                                    ;  chain pointers back to the
                                                    ;  main control bytes.
FDE0 214300       LXI     H,Disk$Status$Block
FDE3 2249FD       SHLD    Floppy$Next$Status$Block
FDE6 214500       LXI     H,Disk$Control$5
FDE9 224BFD       SHLD    Floppy$Next$Control$Location

FDEC 2140FD       LXI     H,Floppy$Command          ;Set up command block pointer
FDEF 224600       SHLD    Command$Block$5

FDF2 214500       LXI     H,Disk$Control$5          ;Activate 5 1/4" disk controller
FDF5 3680         MVI     M,80H
                ;
                Wait$For$Disk$Complete:             ;Wait until Disk Status Block indicates
                                                    ;  operation complete, then check
                                                    ;  if any errors occurred.
                                                    ;On entry HL -> disk control byte
FDF7 7E           MOV     A,M                       ;Get control byte
FDF8 B7           ORA     A
FDF9 C2F7FD       JNZ     Wait$For$Disk$Complete    ;Operation still not yet done
                                                    ;
FDFC 3A4300       LDA     Disk$Status$Block         ;Complete -- now check status
FDFF FE80         CPI     80H                       ;Check if any errors occurred
FE01 DA09FE       JC      Disk$Error                ;Yes
FE04 AF           XRA     A                         ;No
FE05 32F6FB       STA     Disk$Error$Flag           ;Clear error flag
FE08 C9           RET
                Disk$Error:
FE09 3E01         MVI     A,1                       ;Set disk-error flag nonzero
FE0B 32F6FB       STA     Disk$Error$Flag
FE0E C9           RET
                ;
                ;
                ;
                ;  Disk control table images for warm boot
                ;
                Boot$Control$Part$1:
FE0F 01           DB      1                         ;Read function
FE10 00           DB      0                         ;Unit (drive) number
FE11 00           DB      0                         ;Head number
FE12 00           DB      0                         ;Track number
FE13 02           DB      2                         ;Starting sector number
FE14 0010         DW      8*512                     ;Number of bytes to read
FE16 00E0         DW      CCP$Entry                 ;Read into this address
FE18 4300         DW      Disk$Status$Block         ;Pointer to next status block
FE1A 4500         DW      Disk$Control$5            ;Pointer to next control table
                Boot$Control$Part2:
FE1C 01           DB      1                         ;Read function
FE1D 00           DB      0                         ;Unit (drive) number
FE1E 01           DB      1                         ;Head number
FE1F 00           DB      0                         ;Track number
FE20 01           DB      1                         ;Starting sector number
FE21 0006         DW      3*512                     ;Number of bytes to read
FE23 00F0         DW      CCP$Entry + (8*512)       ;Read into this address
FE25 4300         DW      Disk$Status$Block         ;Pointer to next status block
FE27 4500         DW      Disk$Control$5            ;Pointer to next control table

                ;
                ;
                ;
                WBOOT:                ;Warm boot entry
                                      ;On warm boot, the CCP and BDOS must be reloaded
                                      ;  into memory. In this BIOS, only the 5 1/4"
                                      ;  diskettes will be used. Therefore this code
```

Figure 6-4. (Continued)

```
                                 ;  is hardware specific to the controller. Two
                                 ;  prefabricated control tables are used.
FE29 318000      LXI    SP,80H
FE2C 110FFE      LXI    D,Boot$Control$Part1      ;Execute first read of warm boot
FE2F CD3BFE      CALL   Warm$Boot$Read            ;Load drive 0, track 0,
                                                  ;  head 0, sectors 2 to 8
FE32 111CFE      LXI    D,Boot$Control$Part2      ;Execute second read
FE35 CD3BFE      CALL   Warm$Boot$Read            ;Load drive 0, track 0,
                                                  ;  head 1, sectors 1 - 3
FE38 C340F8      JMP    Enter$CPM                 ;Set up base page and enter CCP
              ;
              Warm$Boot$Read:                     ;On entry, DE -> control table image
                                                  ;This control table is moved into
                                                  ;  the main disk control table and
                                                  ;  then the controller activated.
FE3B 2140FD      LXI    H,Floppy$Command          ;HL -> actual control table
FE3E 224600      SHLD   Command$Block$5           ;Tell the controller its address
                                                  ;Move the control table image
                                                  ;  into the control table itself
FE41 0E0D        MVI    C,13                      ;Set byte count
              Warm$Boot$Move:
FE43 1A          LDAX   D                         ;Get image byte
FE44 77          MOV    M,A                       ;Store into actual control table
FE45 23          INX    H                         ;Update pointers
FE46 13          INX    D
FE47 0D          DCR    C                         ;Count down on byte count
FE48 C243FE      JNZ    Warm$Boot$Move            ;Continue until all bytes moved

FE4B 214500      LXI    H,Disk$Control$5          ;Activate controller
FE4E 3680        MVI    M,80H
              Wait$For$Boot$Complete:
FE50 7E          MOV    A,M                       ;Get status byte
FE51 B7          ORA    A                         ;Check if complete
FE52 C250FE      JNZ    Wait$For$Boot$Complete    ;No
                                                  ;Yes, check for errors
FE55 3A4300      LDA    Disk$Status$Block
FE58 FE80        CPI    80H
FE5A DA5EFE      JC     Warm$Boot$Error           ;Yes, an error occurred
FE5D C9          RET
              ;
              Warm$Boot$Error:
FE5E 2167FE      LXI    H,Warm$Boot$Error$Message
FE61 CD33F8      CALL   Display$Message
FE64 C329FE      JMP    WBOOT                     ;Restart warm boot
              ;
              Warm$Boot$Error$Message:
FE67 0D0A576172  DB     CR,LF,'Warm Boot Error - retrying...',CR,LF,0
              ;
              ;
FE89             END    ;Of simple BIOS listing
```

Figure 6-4. (Continued)

7

Building a New CP/M System

This chapter describes how to build a version of CP/M with your own BIOS built into it. It also shows you how to put CP/M onto a floppy disk and how to write a bootstrap loader to bring CP/M into memory.

The manufacturer of your computer system plays a significant role in building a new CP/M system. Several of CP/M's utility programs may be modified by manufacturers to adapt them to individual computer systems. Unfortunately, not all manufacturers customize these programs. You should therefore invest some time in studying the documentation provided with your system to see what and how much customizing may have already been done. You should also assemble and print out listings of all assembly language source files from your CP/M release diskette.

It is impossible to predict the details of customization and special procedures that the manufacturer may have installed on your particular system. Therefore, this chapter describes first the overall mechanism of building a CP/M system, and second the details of building a CP/M system around the example BIOS shown in the previous chapter as Figure 6-4.

The Major Steps

Building a new CP/M system consists of the following major steps:

- Create a new or modified BIOS with the appropriate device drivers in it. Assemble this so that it will execute at the top end of memory (by using an *origin* statement (ORG) to set the location counter).
- Create new versions of the CCP and BDOS with all addresses in the instructions changed so that they will be correctly located in memory just below the new BIOS. Digital Research provides a special utility called MOVCPM to do this.
- Create or modify a CP/M bootstrap loader that will be loaded by the firmware that executes when you first switch on your computer (or press the RESET button). Normally, the CP/M bootstrap loader executes in the low-address end of memory. The exact address and the details of any hardware initialization that it must perform will depend entirely on your particular computer system.
- Using Digital Research standard utility programs, bring the bootstrap loader, the CCP and BDOS, and the BIOS together in the low part of memory. Then write this new version of CP/M onto a disk in the appropriate places. Again, depending on the design of your computer system, you may be able to use the standard utility program, SYSGEN, to write the entire CP/M *image* onto disk. Otherwise you may have to write a special program to do this.

When CP/M is already running on your computer system and you want to add new features to the BIOS, all you need to do is change the BIOS and rebuild the system. The CCP and BDOS will need to be moved down in memory if the changes expand the BIOS significantly. If this happens, you will have to make minor changes in the bootstrap loader so that it reads the new CP/M image into memory at a lower address and transfers control to the correct location (the first instruction of the BIOS jump vector).

Building Your First System

The first time that you build CP/M, it is a good idea to make no changes to the BIOS at all. Simply reassemble the BIOS source code and proceed with the system build. Then, if the new system does not run, you know that it must be something in the procedure you used rather than any new features or modification to the BIOS source code. Changes in the BIOS could easily obscure any problems you have with the build procedure itself.

The Ingredients

To build CP/M, you will need the following files and utility programs:

- The assembly language source code for your BIOS. Check your CP/M release diskette for a file with a name like CBIOS.ASM (Customized Basic Input/Output System). Some manufacturers do not supply you with the source code for their BIOS; it may be sold separately or not released at all. If you cannot get hold of the source code, the only way that you can add new features to the BIOS is by writing the entire BIOS from scratch.
- The source code for the CP/M bootstrap loader. This too may be on the release diskette or available separately from your computer's manufacturer.
- The Digital Research assembler, which converts source code into machine language in hexadecimal form. This program, called ASM.COM, will be on your CP/M release diskette. Equivalent assemblers, such as Digital Research's macro-assemblers MAC and RMAC or Microsoft's M80, can also be used.
- The Digital Research utility called MOVCPM, which prepares a memory image of the CCP and BDOS with all addresses adjusted to the right values.
- The Digital Research debugging utility, called DDT (Dynamic Debugging Tool), or the more enhanced version for the Z80 CPU chip, ZSID (Z80 Symbolic Interactive Debugger). DDT is used to read in the various program files and piece together a memory image of the CP/M system.
- The Digital Research utility program SYSGEN. This writes the composite memory image of the bootstrap, CCP, BDOS, and BIOS onto the disk. SYSGEN was designed to work on floppy disk systems. If your computer uses a hard disk, you may have a program with a name like PUTCPM or WRITECPM that performs the same function.

The Ultimate Goal

In Figure 6-4, lines 0044 to 0065, you can see the equates that define the base addresses for the CCP, the BDOS, and the BIOS. Figure 7-1 shows how the top of memory will look when this version of CP/M has been loaded into memory.

Life would be simple if you could build this image in memory at the addresses shown and write the image out to disk. Building this image, however, would probably overwrite the version of CP/M that you were operating since it too lives at the top of memory. Therefore, the goal is to create a replica of this image lower down in memory, but with all the instruction addresses set to *execute* at the addresses shown in Figure 7-1.

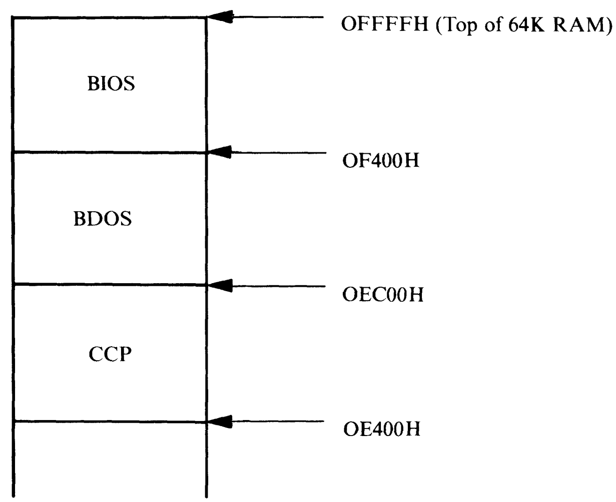


Figure 7-1. Memory layout of CP/M

Using SYSGEN to Write CP/M to Disk

The SYSGEN utility writes a memory image onto a specified logical disk. It can use a memory image that you arrange to be in memory before you invoke SYSGEN, or you can direct SYSGEN to read in a disk file that contains the image. You can also use SYSGEN to transport an existing CP/M system from one diskette to another by directing it to load the CP/M image from one diskette into memory and then to write that image out to another diskette.

Check the documentation supplied by your computer's manufacturer to make sure that you can use SYSGEN on your system. SYSGEN, as released by Digital Research, is constructed to run on 8-inch, single-sided, single-density diskettes. If your system does not use these standard diskettes, SYSGEN must be customized to your disk system.

When SYSGEN loads a CP/M image into memory, it will place the bootstrap, CCP, BDOS, and BIOS at the predetermined addresses shown in Figure 7-2, regardless of where this CP/M originated.

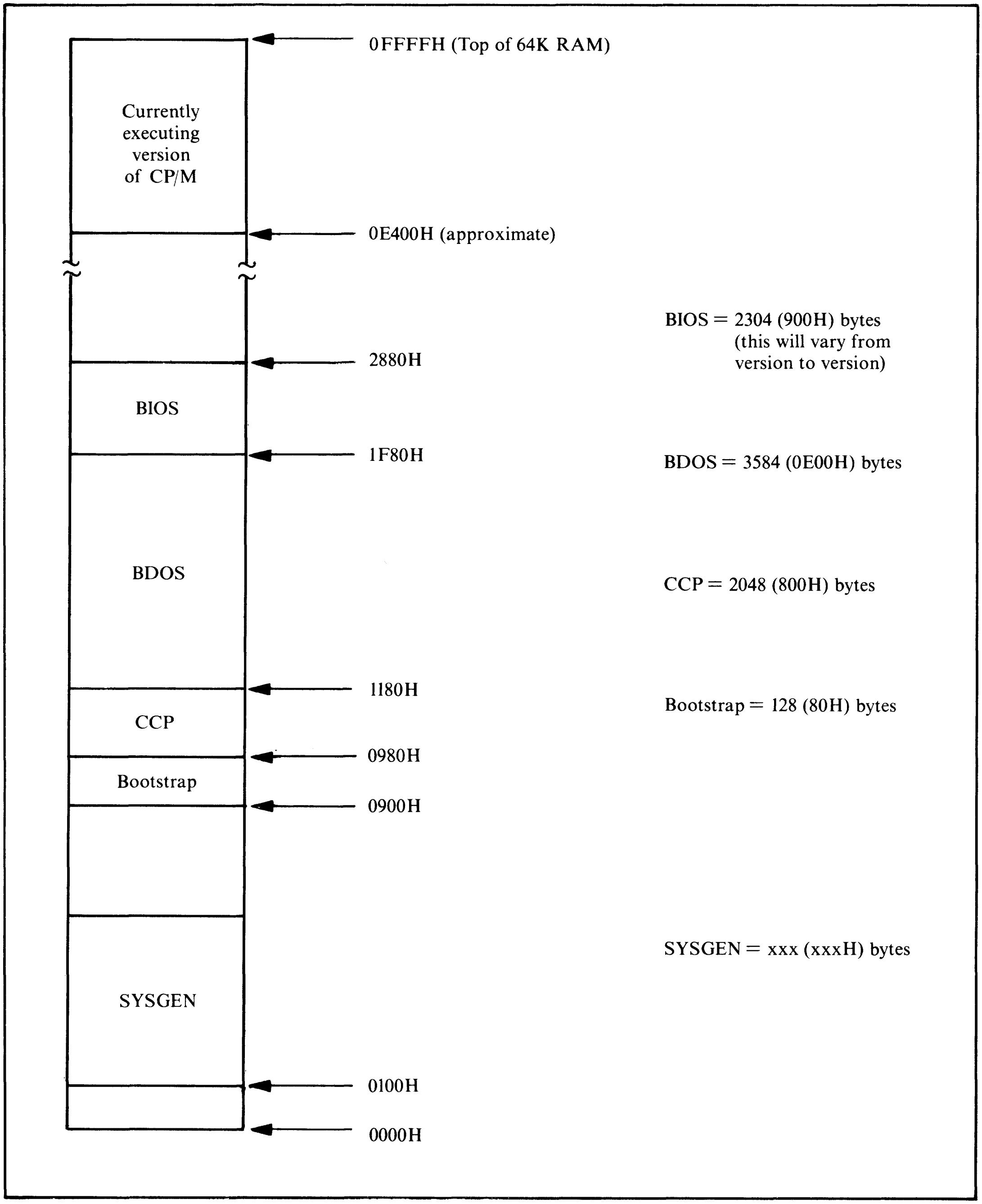


Figure 7-2. SYSGEN's memory layout

You can see that the *relative* arrangement between the components has not changed; the whole image has simply been moved down in memory well below the currently executing version of CP/M. The bootstrap has been added to the picture just beneath the CCP.

The SYSGEN utility writes this image onto a floppy diskette starting at sector 1 of track 0 and continuing to sector 26 on track 1. Refer back to Figure 2-2 to see the layout of CP/M on a standard 8-inch, single-sided, single-density diskette.

If you request SYSGEN to read the memory image from a file (which you do by calling SYSGEN with the file name on the same line as the SYSGEN call), then SYSGEN presumes that you have previously created the correct memory image and saved it (with the SAVE command). SYSGEN then skips over the first 16 sectors of the file so as to avoid overwriting itself.

Here is an example of how to use SYSGEN to move the CP/M image from one diskette to another:

```
A>SYSGEN<CR>
SYSGEN VER 2.0
SOURCE DRIVE NAME (OR RETURN TO SKIP) A
SOURCE ON A:, THEN TYPE RETURN <cr>
FUNCTION COMPLETE
DESTINATION DRIVE NAME (OR RETURN TO REBOOT) B
DESTINATION ON B: THEN TYPE RETURN <cr>
FUNCTION COMPLETE
DESTINATION DRIVE NAME (OR RETURN TO REBOOT) <cr>
A>_
```

As you can see, SYSGEN gives you the choice of specifying the source drive name or typing CARRIAGE RETURN. If you enter a CARRIAGE RETURN, SYSGEN assumes that the CP/M image is already in memory. Note that you need to call up SYSGEN only once to write out the same CP/M image to more than one disk.

A larger than standard BIOS can cause difficulties in using SYSGEN. The standard SYSGEN format only allows for six 128-byte sectors to contain the BIOS, so if your BIOS is larger than 768 (300H) bytes, it will be a problem. The CP/M image will not fit on the first two tracks of a standard 8-inch diskette.

Nowadays it is rare to find an 8-inch floppy diskette system where you must load CP/M from a single-sided, single-density diskette. Most systems now use double-sided or double-density diskettes as the normal format, but can switch to single-sided, single-density diskettes to interchange information with other computer systems.

Because there is no "standard" format for 8-inch, double-sided and double-density diskettes, you probably won't be able to read diskettes written on systems of a different make or model. Therefore, you need only be concerned about using a disk layout that will keep your disks compatible with other machines that are exactly the same as yours.

This is also true if you have 5 1/4-inch diskettes. There is no industry standard for these either, so your main consideration is to place the file directory in the same place as it will be on diskettes written by other users of your model of computer. You must also be sure to use the same sector skewing. Otherwise, you will get a garbled version whenever you try to read files originating on other systems.

With the higher capacity diskettes, you can reserve more space to hold the CP/M image on the diskette. For example, in the case of the BIOS shown in Figure 6-4, the CP/M image is written to a 5 1/4-inch, double-sided, double-density diskette using 512-byte sectors. Figure 7-3 shows the layout of this diskette. Note that the bootstrap loader is placed in a 512-byte sector all by itself. Doing so makes the bootstrap code and warm boot code in the BIOS much simpler.

The memory image must be altered to reflect the fact that the bootstrap now occupies an entire 512-byte sector. Rather than change all of the addresses, the bootstrap is loaded into memory 384 (180H) bytes lower, so that it ends at the same address as before. Figure 7-4 shows the revised memory image.

Writing a PUTCPM Utility

Because the example system uses 5 1/4-inch floppy diskettes with 512-byte sectors, the standard version of SYSGEN cannot be used to write the CP/M image onto a diskette. You will have to use a functional replacement provided by your computer's manufacturer or develop a small utility program to do the job.

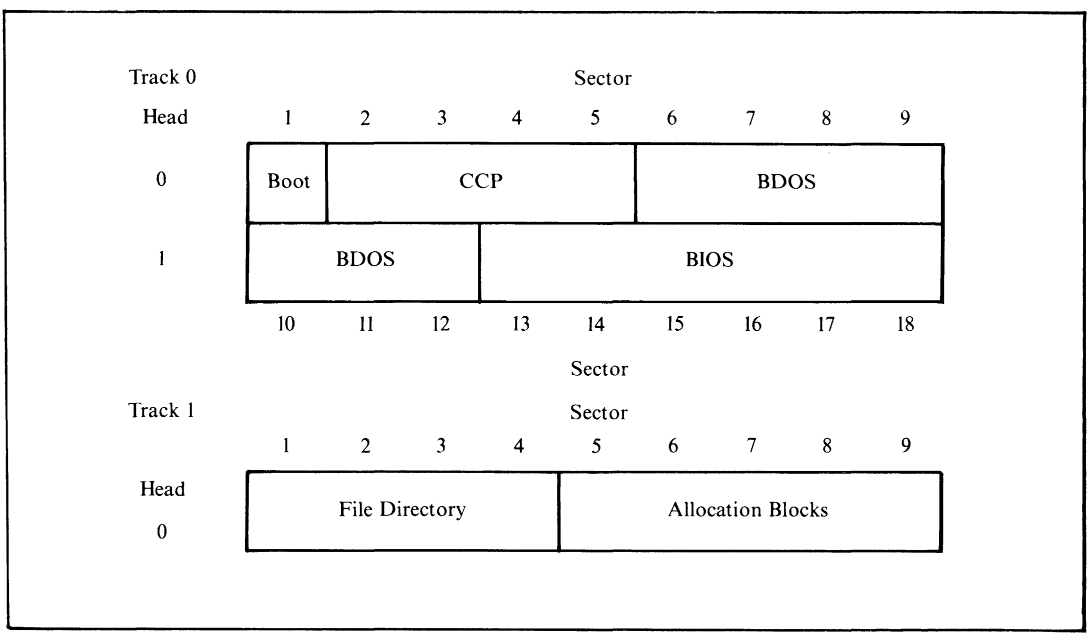


Figure 7-3. Disk layout for example BIOS on 5 1/4-inch diskettes

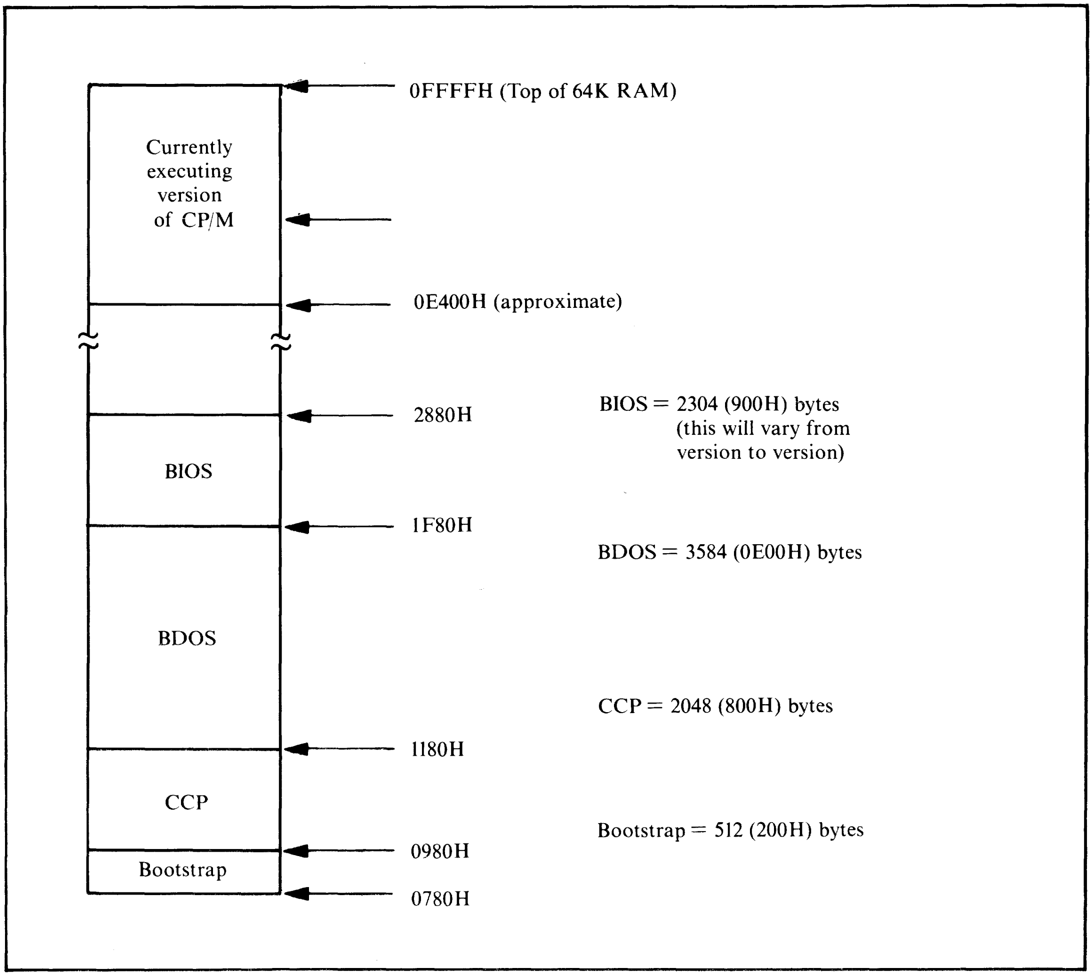


Figure 7-4. Addresses for example BIOS image

Figure 7-5 shows an example of such a program. It is written in a general-purpose way, so that you may be able to use it for your system by changing the equates at the front of the program to reflect the specifics of your disk drives.

Note that there are two problems to be solved. First, the area of the disk on which the CP/M image resides cannot be accessed by the BDOS, as it is outside the file system area on the disk. Second, it is rare to write the CP/M image onto the disk with any kind of sector skewing; to do so would slow down the loading process. In any case, skewing would be redundant, since the loader is doing no processing other than reading the disk and can therefore read the disk without skewing.

```
                ;       This program writes out the CP/M cold boot loader,
                ;       CCP, BDOS, and BIOS to a floppy diskette. It runs
                ;       under CP/M as a normal transient program.
                ;
3130 =          Version         EQU     '01'    ;Equates used in the sign-on
                                                ; message
3730 =          Month           EQU     '07'
3432 =          Day             EQU     '24'
3238 =          Year            EQU     '82'
                ;
                ;
                ;       The actual PUTCPMF5.COM program consists of this code,
                ;       plus the BOOTF5.HEX, CCP, BDOS, and BIOS.
                ;
                ;       When this program executes, the memory image should
                ;       look like this:
                ;
                ;             Component       Base Address
                ;               BIOS             1F80H
                ;               BDOS             1180H
                ;               CCP              0980H
                ;               BOOTF5           0780H
                ;
                ;       The components are produced as follows:
                ;
                ;               BIOS.HEX         By assembling source code
                ;               BDOS )           From a CPMnn.COM file output
                ;               CCP  )             by MOVCPM and SAVEd on disk
                ;               BOOTF5.HEX       By assembling source code
                ;
                ;       The components are pieced together using DDT with the
                ;       following commands:
                ;
                ;               DDT CPMnn.COM
                ;               IPUTCPMF5.HEX
                ;               R                       (Reads in this program)
                ;               IBOOTF5.HEX
                ;               R680                    (Reads in BOOT at 0780H)
                ;               IBIOS.HEX
                ;               R2980                   (Reads in BIOS at 1F80H)
                ;               G0                      (Exit from DDT)
                ;               SAVE 40 PUTCPMF5.COM    (Create final .COM file)
                ;
                ;       The actual layout of the diskette is as follows:
                ;
                ; Track 0                   Sector
                ;          1     2     3     4     5     6     7     8     9
                ; Head    +-----+-----+-----+-----+-----+-----+-----+-----+-----+
                ;  0      |Boot |<======== CCP ========>|<=======  BDOS ========|
                ;         +-----+-----+-----+-----+-----+-----+-----+-----+-----+
                ;  1      |====== BDOS ====>|<============= BIOS  ===========>|
                ;         +-----+-----+-----+-----+-----+-----+-----+-----+-----+
                ;          10    11    12    13    14    15    16    17    18
                ;                           Sector
                ;
                ;       Equates for defining memory size and the base address and
                ;       length of the system components
                ;
0040 =          Memory$Size     EQU     64      ;Number of Kbytes of RAM
                ;
                ;       The BIOS Length must match that declared in the BIOS.
                ;
0900 =          BIOS$Length     EQU     0900H
                ;
0200 =          Boot$Length     EQU     512
0800 =          CCP$Length      EQU     0800H   ;Constant
0E00 =          BDOS$Length     EQU     0E00H   ;Constant
                ;
1F00 =          Length$In$Bytes EQU     CCP$Length + BDOS$Length + BIOS$Length
                ;
0780 =          Start$Image     EQU     980H - Boot$Length      ;Address of CP/M image
2100 =          Length$Image    EQU     Length$In$Bytes + Boot$Length
                ;
                ;
```

Figure 7-5. Example PUTCPM

```
                ;       Disk characteristics
                ;
                ;       These equates describe the physical characteristics of
                ;       the floppy diskette so that the program can move from
                ;       one sector to the next, updating the track and resetting
                ;       the sector when necessary.
                ;
0001 =          First$Sector$on$Track   EQU     1
0012 =          Last$Sector$on$Track    EQU     18
0009 =          Last$Sector$on$Head$0   EQU     9
0200 =          Sector$Size             EQU     512
                ;
                ;
                ;       Controller characteristics
                ;
                ;       On this computer system, the floppy disk controller can write
                ;       multiple sectors in a single command. However, in order
                ;       to produce a more general example it is shown only reading one
                ;       sector at a time.
                ;
0001 =          Sectors$Per$Write       EQU     1
                ;
                ;
                ;       Cold boot characteristics
                ;
0000 =          Start$Track             EQU     0       ;Initial values for CP/M image
0001 =          Start$Sector            EQU     1       ;= " =
0011 =          Sectors$To$Write        EQU     (Length$Image + Sector$Size - 1) / Sector$Size
                ;

                ;
0009 =          B$PRINTS        EQU     9               ;Print string terminated by $
0005 =          BDOS            EQU     5               ;BDOS entry point
                ;
                ;

0100                    ORG     100H
                Put$CPM:
0100 C33F01             JMP     Main$Code               ;Enter main code body
                                                        ;For reasons of clarity, the main
                                                        ;  data structures are shown before the
                                                        ;  executable code.
000D =          CR      EQU     0DH                     ;Carriage return
000A =          LF      EQU     0AH                     ;Line feed
                ;
                Signon$Message:
0103 0D0A507574         DB      CR,LF,'Put CP/M on Diskette'
0119 0D0A               DB      CR,LF
011B 5665727369         DB      'Version '
0123 3031               DW      Version
0125 20                 DB      ' '
0126 3037               DW      Month
0128 2F                 DB      '/'
0129 3234               DW      Day
012B 2F                 DB      '/'
012C 3832               DW      Year
012E 0D0A24             DB      CR,LF,'$'

                ;
                ;        Disk control tables
                ;
0045 =          Disk$Control$5  EQU     45H     ;5 1/4" control byte
0046 =          Command$Block$5 EQU     46H     ;Control table pointer
0043 =          Disk$Status     EQU     43H     ;Completion status
                ;
                ;
                ;       The command table track and DMA$Address can also be used
                ;       as working storage and updated as the load process
                ;       continues. The sector in the command table cannot be
                ;       used directly as the disk controller requires it to be
                ;       the sector number on the specified head (1 -- 9) rather
                ;       than the sector number on track. Hence a separate variable
                ;       must be used.
                ;
```

Figure 7-5. (Continued)

```
0131 01         Sector:         DB      Start$Sector
                ;
0132 02         Command$Table:  DB      02H             ;Command -- Write
0133 00         Unit:           DB      0               ;Unit (drive) number = 0 or 1
0134 00         Head:           DB      0               ;Head number = 0 or 1
0135 00         Track:          DB      Start$Track     ;Used as working variable
0136 00         Sector$on$head: DB      0               ;Converted by low-level driver
0137 0002       Byte$Count:     DW      Sector$Size * Sectors$Per$Write
0139 8007       DMA$Address:    DW      Start$Image
013B 4300       Next$Status:    DW      Disk$Status     ;Pointer to next status block
                                                        ;  if commands are chained
013D 4500       Next$Control:   DW      Disk$Control$5  ;Pointer to next control byte
                                                        ;  if commands are chained

                Main$Code:
013F 310001             LXI     SP,Put$CPM      ;Stack grows down below code

0142 110301             LXI     D,Signon$Message        ;Sign on
0145 0E09               MVI     C,B$PRINTS              ;Print string until $
0147 CD0500             CALL    BDOS

014A 213201             LXI     H,Command$Table         ;Point the disk controller at

014D 224600             SHLD    Command$Block$5         ;  the command block

0150 0E11               MVI     C,Sectors$To$Write      ;Set sector count
                Write$Loop:
0152 CD7C01             CALL    Put$CPM$Write           ;Write data onto diskette
0155 0D                 DCR     C                       ;Downdate sector count
0156 CA0000             JZ      0                       ;Warm boot

0159 213101             LXI     H,Sector                ;Update sector number
015C 3E01               MVI     A,Sectors$Per$Write     ;  by adding on number of sectors
015E 86                 ADD     M                       ;  by controller
015F 77                 MOV     M,A                     ;Save result
0160 3E13               MVI     A,Last$Sector$On$Track + 1      ;Check if at end of track
0162 BE                 CMP     M
0163 C26F01             JNZ     Not$End$Track

0166 3601               MVI     M,First$Sector$On$Track ;Yes, reset to beginning
0168 2A3501             LHLD    Track                   ;Update track number
016B 23                 INX     H
016C 223501             SHLD    Track

                Not$End$Track:
016F 2A3901             LHLD    DMA$Address             ;Update DMA address
0172 110002             LXI     D,Sector$Size * Sectors$Per$Write
0175 19                 DAD     D
0176 223901             SHLD    DMA$Address
0179 C35201             JMP     Write$Loop              ;Write next block
                ;
                Put$CPM$Write:                  ;At this point, the description of the
                                                ;  operation required is in the variables
                                                ;  contained in the command table, along
                                                ;  with the sector variable.

017C C5                 PUSH    B                       ;Save sector count in C

                ;------ Change this routine to match the disk controller in use ------

017D 0600               MVI     B,0                     ;Assume head 0
017F 3A3101             LDA     Sector                  ;Get requested sector
0182 4F                 MOV     C,A                     ;Take a copy of it
0183 FE0A               CPI     Last$Sector$on$Head$0+1 ;Check if on head 1
0185 DA8C01             JC      Head$0                  ;No
0188 D609               SUI     Last$Sector$on$Head$0   ;Bias down for head 1
018A 4F                 MOV     C,A                     ;Save copy
018B 04                 INR     B                       ;Set head 1
                Head$0:
018C 78                 MOV     A,B                     ;Get head
018D 323401             STA     Head
0190 79                 MOV     A,C                     ;Get sector
0191 323601             STA     Sector$On$Head
```

Figure 7-5. (Continued)

```
0194 214500              LXI     H,Disk$Control$5        ;Activate controller
0197 3680                MVI     M,80H

                 Wait$For$Boot$Complete:
0199 7E                  MOV     A,M                     ;Get status byte
019A B7                  ORA     A                       ;Check if complete
019B C29901              JNZ     Wait$For$Boot$Complete  ;No
                                                         ;Yes, check for errors
019E 3A4300              LDA     Disk$Status
01A1 FE80                CPI     80H
01A3 DAA801              JC      Put$CPM$Error           ;Yes, an error occurred

                 ;------ End of physical write routine ------

01A6 C1                  POP     B                       ;Recover sector count in C
01A7 C9                  RET
                 ;
                 Put$CPM$Error:
01A8 11B301              LXI     D,Put$CPM$Error$Message
01AB 0E09                MVI     C,B$PRINTS              ;Print string until $
01AD CD0500              CALL    BDOS                    ;Output error message
01B0 C33F01              JMP     Main$Code               ;Restart the loader
                 ;

                 Put$CPM$Error$Message:
01B3 0D0A457272          DB      CR,LF,'Error in writing CP/M - retrying...',CR,LF,'$'
01DB                     END     Put$CPM
```

Figure 7-5. (Continued)

Using DDT to Build the CP/M Memory Image

DDT, the Digital Research debug program, is used to read files of type ".COM" and ".HEX" into memory. Understanding the internal structure of these file types is important, both to understand what DDT can do and to understand how the MOVCPM utility can effectively change a machine code file so that it can be executed at a new address in memory.

".COM" File Structure

A COM file is a memory image. It is a replica of the bit patterns that are to be created when the file is loaded into memory. COM files are normally designed to load at location 100H upwards. No internal structure to the file requires this, however, so if you know what the contents of a COM file are, there is nothing to preclude you from loading it into memory starting at some address other than 100H.

As you may recall from the description of the CCP in Chapter 4, the SAVE command built into the CCP allows you to create a COM file by specifying the number of 256-byte "pages" of memory and the name of the file. The CCP will write out an exact image of memory from location 100H up.

".HEX" File Structure

HEX files are output by the assembler. They contain an ASCII character representation of hexadecimal values. For example, the contents of a single byte of memory with the binary value 10101111 would be represented by two ASCII characters, A F, in a HEX file.

The HEX file has a higher level structure than just a series of ASCII characters however. Each line of ASCII characters is terminated by CARRIAGE RETURN/LINE FEED. The overall structure is shown in Figure 7-6.

The most important aspect of a HEX file is that each line contains the address at which the data bytes are loaded. Each line is processed independently, so the load addresses of succeeding lines need not be in order.

DDT can read in a HEX file at an address different from the address where the code must be in order to execute. For example, you can read in the HEX file of the BIOS at the correct place for the memory image (shown in Figure 7-4). There are two ways of using DDT to read in a COM or HEX file. You can specify the name of the file on the same command line with DDT. For example:

```
A>DDT B:XYZ.HEX<cr>          <- Call up DDT with file name
DDT VERS 2.0                 <- DDT signs on
NEXT  PC
0180 0100                    <- ... and displays next free byte
                                 and entry point address
                             <- ... and prompts for a commmand
```

The advantage of this method of loading a file is that you can specify which logical disk is to be searched for the file. The second way of using DDT is to load DDT first, and then, when it has given its prompt, specify the file name and request that DDT load it like this:

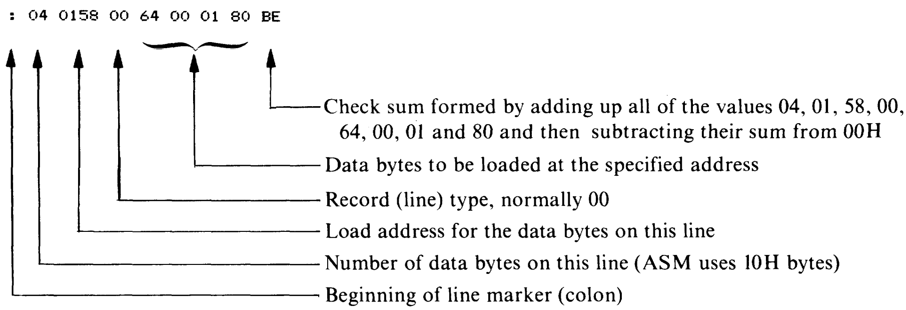
```
-Ifilename.typ<cr>           <- Enter the file name and type
-R<cr>                       <- Read in the file
```

The "I" command initializes the default file control block in the base page (at location 005CH) with the file name and type; it does *not* set up the logical disk. If you need to do this, you must set the first byte of the default FCB manually like this:

```
-Ifilename.typ<cr>           <- Specify file name
-S5C<cr>                     <- "S"et location 5C
005C 00 02<cr>               <- Was 00, you enter 02<cr>
005D 41 .<cr>                <- Enter "." to terminate
-R<cr>                       <- Read in the file
```

Location 005CH should be set to 01H for Drive A, 02H for B, and so on.

The "R" command will read in HEX files to the *execution* addresses specified in each line of the HEX file, so be careful—if you forget to put an ORG (origin)



NOTE: HEX files do not have embedded blank characters; the example above is shown with gaps between individual fields only for clarity.

Figure 7-6. Example line from HEX file

statement at the front of the assembly language source code, reading in the resultant HEX file will overwrite location 0000H on up, destroying the contents of the base page. Similarly, if you were trying to read in the HEX file for a BIOS, there is an excellent chance that you will overwrite the currently executing CP/M system.

DDT reacts to the file type you enter as part of the file name. For file types other than .HEX, DDT loads the file starting at location 0100H on up.

The "R" command can also be used to read files into memory at different addresses. You do this by typing a hexadecimal number immediately after the R, with no intervening punctuation. For HEX files, the number that you enter is added to the address in each line of the HEX file and the sum is used as the address into which the data bytes are loaded. The data bytes themselves are not changed, just the load address.

For COM files, the number that you enter is added to 0100H and the sum is used as the starting address for loading the file.

The sum is performed as 16-bit, unsigned arithmetic with any carry ignored, so you can load a BIOS HEX file into low memory by using the "R" command with what is called an "offset value."

If a HEX file has been assembled to execute at address "exec," and you need to use DDT to read in this file to address "load," you need to solve the following equation:

offset = load − exec.

DDT's "H" command performs hexadecimal arithmetic. It calculates and displays the sum of and difference between two hexadecimal values. For example, the BIOS in Figure 6-4 has been assembled to *execute* at location 0F600H, but needs to be *loaded* into memory at location 1F80H. Here is how to compute the correct offset for the "R" command:

```
-H1F80,F600<cr>          <- Use the H command
1580,2980                <- Sum, difference
```

Thus, to read in the BIOS HEX file called FIG6-4.HEX at location 1F80H, you would enter the following commands to DDT:

```
-IFIG6-4.HEX<cr>         <- Specify file name and type
-R2980<cr>               <- Load at 0F600H + 2980H (= 1F80H)
```

In this way, using DDT, you can read in the HEX files for both the BIOS and the bootstrap loader.

The CP/M Bootstrap Loader

The bootstrap loader is brought into memory by PROM-based firmware in the computer system. It loads in the CCP, BDOS, and BIOS and then transfers control to the cold boot entry point in the BIOS—the first jump instruction in the BIOS jump vector.

The bootstrap loader is a stand-alone program; it cannot make use of any CP/M functions because no part of CP/M is in memory when the bootstrap loader is needed. The firmware in the PROM that loaded the bootstrap may contain some subroutines that can be used by the bootstrap, but this will vary from system to system.

Figure 7-7 shows the bootstrap code for the example BIOS (from Figure 6-4). This code has been written in a general way, so that you can adapt it to your system. The disk controller on the example system can in fact read in multiple sectors from the disk, but for generality the code shown reads in only one sector at a time. This considerably increases the time it takes to load CP/M, but does make the bootstrap loader more general.

Note that almost the first thing that the bootstrap does is to output to the console a sign-on message. Not only does this confirm the version number, but it shows that the bootstrap has been successfully loaded.

The PROM-based code has been designed to load the CP/M bootstrap into location 100H, allowing the code to be debugged as though it were a normal transient program, albeit with minor changes to the address at which it loads the CP/M image from disk. Clearly, this feature is not very helpful if CP/M is being brought up for the first time on a computer system. It helps a great deal, however, if you need to modify the bootstrap or add the capability to boot your system from a new type of disk drive.

```
                        ;       Example CP/M cold bootstrap loader
                        ;
                        ;       This program is written out to track 0, head 0, sector 1
                        ;       by the PUTCPMF5 program.
                        ;       It is loaded into memory at location 100H on up by the
                        ;       PROM-based bootstrap mechanism that gets control of the
                        ;       CPU on power up or system reset.
                        ;
3130 =                  Version         EQU     '01'    ;Equates used in the sign-on message
3730 =                  Month           EQU     '07'
3432 =                  Day             EQU     '24'
3238 =                  Year            EQU     '82'
                        ;
0000 =                  Debug           EQU     0       ;Set nonzero to debug as normal
                                                        ;  transient program
                        ;
                        ;       The actual layout of the diskette is as follows :
                        ;
                        ; Track 0                     Sector
                        ;         1     2     3     4     5     6     7     8     9
                        ; Head  +-----+-----+-----+-----+-----+-----+-----+-----+-----+
                        ;  0    |Boot |<======== CCP ========>|<=======  BDOS ========|
                        ;       +-----+-----+-----+-----+-----+-----+-----+-----+-----+
                        ;  1    |====== BDOS ====>|<=============  BIOS  ============>|
                        ;       +-----+-----+-----+-----+-----+-----+-----+-----+-----+
                        ;         10    11    12    13    14    15    16    17    18
                        ;                             Sector
                        ;
                        ;       Equates for defining memory size and the base address and
                        ;       length of the system components.
                        ;
0040 =                  Memory$Size     EQU     64      ;Number of Kbytes of RAM
                        ;
                        ;       The BIOS Length must match that declared in the BIOS.
                        ;
0900 =                  BIOS$Length     EQU     0900H
                        ;
0800 =                  CCP$Length      EQU     0800H   ;Constant
0E00 =                  BDOS$Length     EQU     0E00H   ;Constant
                        ;
0008 =                  Length$In$K     EQU     ((CCP$Length + BDOS$Length + BIOS$Length) / 1024) + 1
1F00 =                  Length$In$Bytes EQU     CCP$Length + BDOS$Length + BIOS$Length
                        ;
                                IF      NOT Debug
E000 =                  CCP$Entry       EQU     (Memory$Size - Length$In$K) * 1024
                                ENDIF
                                IF      Debug
                        CCP$Entry       EQU     3980H   ;Read into a lower address.
                                                        ;This address is chosen to be above
                                                        ;  the area into which DDT initially loads
                                                        ;  and the 980H makes the addresses similar
                                                        ;  to the SYSGEN values so that the memory
                                                        ;  image can be checked with DDT.
                                ENDIF

E806 =                  BDOS$Entry      EQU     CCP$Entry + CCP$Length + 6
F600 =                  BIOS$Entry      EQU     CCP$Entry + CCP$Length + BDOS$Length
                        ;
                        ;
                        ;       Disk characteristics
                        ;
                        ;       These equates describe the physical characteristics of
                        ;       the floppy diskette so that the program can move from
                        ;       one sector to the next, updating the track and resetting
                        ;       the sector when necessary.
                        ;
0001 =                  First$Sector$on$Track   EQU     1
0012 =                  Last$Sector$on$Track    EQU     18
0009 =                  Last$Sector$on$Head$0   EQU     9
0200 =                  Sector$Size             EQU     512
                        ;
                        ;
                        ;       Controller characteristics
                        ;
```

Figure 7-7. Example CP/M cold bootstrap loader

```
                ;       On this computer system, the floppy disk controller can read
                ;       multiple sectors in a single command. However, in order to
                ;       produce a more general example it is shown only reading one
                ;       sector at a time.
                ;
0001 =          Sectors$Per$Read        EQU     1
                ;
                ;
                ;       Cold boot characteristics
                ;
0000 =          Start$Track             EQU     0       ;Initial values for CP/M image
0002 =          Start$Sector            EQU     2       ;= " =
0010 =          Sectors$To$Read         EQU     (Length$In$Bytes + Sector$Size - 1) / Sector$Size
                ;
                ;
                ;

0100                    ORG     100H
                Cold$Boot$Loader:
0100 C34001             JMP     Main$Code       ;Enter main code body
                                                ;For reasons of clarity, the main
                                                ;  data structures are shown before the
                                                ;  executable code.
000D =          CR      EQU     0DH             ;Carriage return
000A =          LF      EQU     0AH             ;Line feed
                ;
                Signon$Message:
0103 0D0A43502F         DB      CR,LF,'CP/M Bootstrap Loader'
                        IF      Debug
                        DB      ' (Debug)'
                        ENDIF
011A 0D0A               DB      CR,LF
011C 5665727369         DB      'Version '
0124 3031               DW      Version
0126 20                 DB      ' '
0127 3037               DW      Month
0129 2F                 DB      '/'
012A 3234               DW      Day
012C 2F                 DB      '/'
012D 3832               DW      Year
012F 0D0A00             DB      CR,LF,0

                ;
                ;        Disk Control Tables
                ;
0045 =          Disk$Control$5  EQU     45H     ;5 1/4" control byte
0046 =          Command$Block$5 EQU     46H     ;Control table pointer
0043 =          Disk$Status     EQU     43H     ;Completion status
                ;
                ;
                ;       The command table track and DMA$Address can also be used
                ;       as working storage and updated as the load process
                ;       continues. The sector in the command table cannot be
                ;       used directly as the disk controller requires it to be
                ;       the sector number on the specified head (1 -- 9) rather
                ;       than the sector number on track. Hence a separate variable
                ;       must be used.
                ;
0132 02         Sector:         DB      Start$Sector
                ;
0133 01         Command$Table:  DB      01H     ;Command -- read
0134 00         Unit:           DB      0       ;Unit (drive) number = 0 or 1
0135 00         Head:           DB      0       ;Head number = 0 or 1
0136 00         Track:          DB      Start$Track     ;Used as working variable
0137 00         Sector$on$head: DB      0       ;Converted by low-level driver
0138 0002       Byte$Count:     DW      Sector$Size * Sectors$Per$Read
013A 00E0       DMA$Address:    DW      CCP$Entry
013C 4300       Next$Status:    DW      Disk$Status     ;Pointer to next status block
                                                        ;  if commands are chained.
013E 4500       Next$Control:   DW      Disk$Control$5  ;Pointer to next control byte
                                                        ;  if commands are chained.

                Main$Code:
0140 310001             LXI     SP,Cold$Boot$Loader     ;Stack grows down below code
```

Figure 7-7. (Continued)

```
0143 210301            LXI     H,Signon$Message         ;Sign on
0146 CDD901            CALL    Display$Message

0149 213301            LXI     H,Command$Table          ;Point the disk controller at
014C 224600            SHLD    Command$Block$5          ;  the command block

014F 0E10              MVI     C,Sectors$To$Read        ;Set sector count
               Load$Loop:
0151 CD7B01            CALL    Cold$Boot$Read           ;Read data into memory
0154 0D                DCR     C                        ;Downdate sector count

                       IF      NOT Debug
0155 CA00F6            JZ      BIOS$Entry               ;Enter BIOS when load done
                       ENDIF
                       IF      Debug
                       JZ      0                        ;Warm boot
                       ENDIF

0158 213201            LXI     H,Sector                 ;Update sector number
015B 3E01              MVI     A,Sectors$Per$Read       ;  by adding on number of sectors
015D 86                ADD     M                        ;  by controller
015E 77                MOV     M,A                      ;Save result
015F 3E13              MVI     A,Last$Sector$On$Track + 1       ;Check if at end of track
0161 BE                CMP     M
0162 C26E01            JNZ     Not$End$Track

0165 3601              MVI     M,First$Sector$On$Track ;Yes, reset to beginning
0167 2A3601            LHLD    Track                    ;Update track number
016A 23                INX     H
016B 223601            SHLD    Track

               Not$End$Track:
016E 2A3A01            LHLD    DMA$Address              ;Update DMA Address
0171 110002            LXI     D,Sector$Size * Sectors$Per$Read
0174 19                DAD     D
0175 223A01            SHLD    DMA$Address
0178 C35101            JMP     Load$Loop                ;Read next block
               ;
               Cold$Boot$Read:                  ;At this point, the description of the
                                                ;  operation required is in the variables
                                                ;  contained in the command table, along
                                                ;  with the sector variable.

017B C5                PUSH    B                        ;Save sector count in C

               ;------ Change this routine to match the disk controller in use ------

017C 0600              MVI     B,0                      ;Assume head 0
017E 3A3201            LDA     Sector                   ;Get requested sector
0181 4F                MOV     C,A                      ;Take a copy of it
0182 FE0A              CPI     Last$Sector$on$Head$0+1  ;Check if on head 1
0184 DA8B01            JC      Head$0                   ;No
0187 D609              SUI     Last$Sector$on$Head$0    ;Bias down for head 1
0189 4F                MOV     C,A                      ;Save copy
018A 04                INR     B                        ;Set head 1
               Head$0:
018B 78                MOV     A,B                      ;Get head
018C 323501            STA     Head
018F 79                MOV     A,C                      ;Get sector
0190 323701            STA     Sector$On$Head

0193 214500            LXI     H,Disk$Control$5         ;Activate controller
0196 3680              MVI     M,80H

               Wait$For$Boot$Complete:
0198 7E                MOV     A,M                      ;Get status byte
0199 B7                ORA     A                        ;Check if complete
019A C29801            JNZ     Wait$For$Boot$Complete   ;No
                                                        ;Yes, check for errors
019D 3A4300            LDA     Disk$Status
01A0 FE80              CPI     80H
01A2 DAA701            JC      Cold$Boot$Error          ;Yes, an error occurred

               ;------ End of physical read routine ------
```

Figure 7-7. (Continued)

```
01A5 C1              POP    B                       ;Recover sector count in C
01A6 C9              RET
                ;
                Cold$Boot$Error:
01A7 21B001          LXI    H,Cold$Boot$Error$Message
01AA CDD901          CALL   Display$Message         ;Output error message
01AD C34001          JMP    Main$Code               ;Restart the loader
                ;
                Cold$Boot$Error$Message:
01B0 0D0A426F6F      DB     CR,LF,'Bootstrap Loader Error - retrying...',CR,LF,0
                ;
                ;    Equates for Terminal Output
                ;
0001 =          Terminal$Status$Port    EQU    01H
0002 =          Terminal$Data$Port      EQU    02H
                ;
0001 =          Terminal$Output$Ready   EQU    0000$0001B
                ;
                ;
                Display$Message:        ;Displays the specified message on the console.
                                        ;On entry, HL points to a stream of bytes to be
                                        ;output. A 00H-byte terminates the message.
01D9 7E              MOV    A,M               ;Get next message byte
01DA B7              ORA    A                 ;Check if terminator
01DB C8              RZ                       ;Yes, return to caller
01DC 4F              MOV    C,A               ;Prepare for output

                Output$Not$Ready:
01DD DB01            IN     Terminal$Status$Port    ;Check if ready for output
01DF E601            ANI    Terminal$Output$Ready
01E1 CADD01          JZ     Output$Not$Ready        ;No, wait
01E4 79              MOV    A,C                     ;Get data character
01E5 D302            OUT    Terminal$Data$Port      ;Output to screen

01E7 23              INX    H                 ;Move to next byte of message
01E8 C3D901          JMP    Display$Message ;Loop until complete message output

                                              ;The PROM-based bootstrap loader checks
                                              ;  to see that the characters "CP/M"
                                              ;  are on the diskette bootstrap sector
                                              ;  before it transfers control to it.
02E0                 ORG    2E0H
02E0 43502F4D        DB     'CP/M'
02E4                 END    Cold$Boot$Loader
```

Figure 7-7. (Continued)

In this case, the bootstrap code must be loaded at location 0780H, not the normal 0980H, because the bootstrap takes a complete 512-byte sector (200H). The same principle applies in determining the offset value to be used with DDT's "R" command to read the bootstrap HEX file, namely:

offset = load address − execution address.

In this case, the values are the following:

0680H = 0780H − 0100H

Using MOVCPM to Relocate the CCP and BDOS

MOVCPM builds a CP/M memory image at the correct locations for SYSGEN, but with the instructions modified to execute at a specific address. Inside MOVCPM is not only a complete replica of CP/M, but also enough information to tell MOVCPM which bytes of which instructions need be changed whenever the execution address of the image needs to be moved.

MOVCPM, as released from Digital Research, contains the bootstrap and BIOS for an Intel MDS-800 computer along with the generic CCP and BDOS. Unless you have an MDS-800, all you use is the CCP and BDOS. Some manufacturers have customized MOVCPM to include the correct bootstrap and BIOS for their own computers; consult their documentation to see if this applies to your computer system.

When you invoke MOVCPM, you have the following options:

- MOVCPM<cr>
  MOVCPM will relocate its built-in copy of CP/M to the top of available memory and will then transfer control to this new image of CP/M. Unless your manufacturer has included the correct BIOS into MOVCPM, using this option will cause an immediate system crash.
- MOVCPM nn<cr>
  This is similar to the option above, except that MOVCPM assumes that *nn*K bytes of memory are available and will relocate the CP/M image to the top of that before transferring control. Again, this will crash the system unless the correct BIOS has been installed into MOVCPM.
- MOVCPM * *<cr>
  MOVCPM will adjust all of the internal addresses inside the CP/M image so that the image could execute at the top of available memory, but instead of actually putting this image at the top of memory, MOVCPM will leave it in low memory at the correct place for SYSGEN to write it onto a disk. The SAVE command could also preserve the image on a disk.
- MOVCPM *nn* *<cr>
  MOVCPM proceeds as above for the "* *" option except that the CP/M image is modified to execute at the top of *nn*K.

MOVCPM has a fundamental problem. The *nn* value indicates that the top of available memory is computed, assuming that your BIOS is small—less that 890 (380H) bytes. If your BIOS is larger (as is the case with the example in Figure 6-4), then you will have to reduce the value of "*nn*" artificially.

Figure 7-8 shows the relationship between the size of the BIOS and the "*nn*" value to use with MOVCPM. It also shows, for different lengths of BIOS, the BIOS base address, the offset value to be used in DDT to read in the BIOS to location 1F80H (preparatory to using SYSGEN or PUTCPM to write it out), and also the base addresses for the CCP and the BDOS. The base address of the BDOS indicates how much memory is available for loading transient programs, as the CCP can be overwritten if necessary.

The numbers in Figure 7-8 are based on the assumption that you have 64K of memory in your computer system. If this is not the case, then proceed as follows:

1. Convert the amount of memory in your system to hex. Remember that 1K is 1024 bytes.
2. Determine the length of your BIOS in hex.
3. Locate the line in Figure 7-8 that shows a BIOS length equal to or greater than the length of your BIOS.
4. Using the "H" command in DDT, compute the BIOS Base Address using the formula:

   Memory in system − BIOS length from Figure 7-8
5. Find the line in Figure 7-8 that shows the same BIOS Base Address as the result of the computation above. Use this line to derive the other relevant numbers.

It is helpful to use DDT to examine a CP/M image in memory to check that all of the components are correctly placed, and, in the case of the CCP and BDOS, correctly relocated.

Figure 7-9 shows an example console dialog in which DDT is used first to examine the memory image produced by MOVCPM and second to examine the image built into the PUTCPMF utility shown in Figure 7-5.

BIOS Length	BIOS Base	DDT Offset	MOVCPM 'nn'	CCP Base	BDOS Base
600	FA00	2580	64	E400	EC00
A00	F600	2980	63	E000	E800
E00	F200	2D80	62	DC00	E400
1200	EE00	3180	61	D800	E000
1600	EA00	3580	60	D400	DC00
1A00	E600	3980	59	D000	D800
1E00	E200	3D80	58	CC00	D400
2200	DE00	4180	57	C800	D000
2600	DA00	4580	56	C400	CC00
2A00	D600	4980	55	C000	C800
2E00	D200	4D80	54	BC00	C400
3200	CE00	5180	53	B800	C000
3600	CA00	5580	52	B400	BC00
3A00	C600	5980	51	B000	B800
3E00	C200	5D80	50	AC00	B400
4200	BE00	6180	49	A800	B000
4600	BA00	6580	48	A400	AC00
4A00	B600	6980	47	A000	A800
4E00	B200	6D80	46	9C00	A400
5200	AE00	7180	45	9800	A000
5600	AA00	7580	44	9400	9C00
5A00	A600	7980	43	9000	9800
5E00	A200	7D80	42	8C00	9400
6200	9E00	8180	41	8800	9000
6600	9A00	8580	40	8400	8C00
6A00	9600	8980	39	8000	8800

Apart from the MOVCPM 'nn' value all other values are in hexadecimal

Figure 7-8. CP/M addresses for different BIOS lengths

```
                      Call up MOVCPM requesting a '63K' system
                      and the image to be left in memory.
A>Movcpm 63 *<cr>
CONSTRUCTING 63k CP/M vers 2.2
READY FOR "SYSGEN" OR
"SAVE 34 CPM63.COM"

                      Save the image from location 100H up. By
                      convention, the file name is CPMnn.COM, so
                      in this case it will be CPM63.COM
A>Save 34 cpm63.com<cr>

                      Call up DDT and request that it read in
                      CPM63.COM
A>ddt cpm63.com<cr>
DDT VERS 2.2
NEXT  PC
2300 0100

                      Display memory to show the first few bytes of
                      the CCP. Note the two JMP (C3H) instructions,
                      followed by 7FH, 00H, 20H's, and the Digital
                      Research Copyright notice. These identify the
                      code as being the CCP. Note that the first
                      JMP instruction is to 35CH into the CCP -- you
                      can therefore infer the base address of the
                      CCP. In this case the JMP is to locat;on E35C,
                      therefore this version of the CCP has been
                      configured to execute based at E000H.
-d980,9cf<cr>
0980 C3 5C E3 C3 58 E3 7F 00 20 20 20 20 20 20 20 20 .\..X...
0990 20 20 20 20 20 20 20 20 43 4F 50 59 52 49 47 48         COPYRIGH
09A0 54 20 28 43 29 20 31 39 37 39 2C 20 44 49 47 49 T (C) 1979, DIGI
09B0 54 41 4C 20 52 45 53 45 41 52 43 48 20 20 00 00 TAL RESEARCH  ..
09C0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................


                      Display the first few bytes of the BDOS. Note
                      the JMP instruction at 1186. This is the
                      instruction to which control is transferred
                      by the JMP in location 5.
-d1180,118F<cr>
1180 00 16 00 00 09 85 C3 11 E8 99 E8 A5 E8 AB E8 B1 ................

                      Displaying further up in the BDOS identifies
                      it unambiguously -- there are some ASCII error
                      messages.
-d1230,126f<cr>
1230 E8 21 DC E8 CD E5 E8 C3 00 00 42 64 6F 73 20 45 .!........Bdos E
1240 72 72 20 4F 6E 20 20 3A 20 24 42 61 64 20 53 65 rr On  : $Bad Se
1250 63 74 6F 72 24 53 65 6C 65 63 74 24 46 69 6C 65 ctor$Select$File
1260 20 52 2F 4F 24 E5 CD C9 E9 3A 42 EB C6 41 32 C6  R/O$....:B..A2.

                      Display the first few bytes of the BIOS.
                      Notice the BIOS JMP vector -- the series of C3H
                      instructions. Normally the first instruction
                      in the vector can be used to infer the base
                      address of the BIOS; in this case it is
                      F600H. But there is no rule that says that
                      the cold boot code must be close to the BIOS
                      JMP vector -- so this is only a rough guide.
-d1f80<cr>
1F80 C3 B3 F6 C3 C3 F6 C3 61 F7 C3 64 F7 C3 6A F7 C3 .......a..d..j..
1F90 6D F7 C3 72 F7 C3 75 F7 C3 78 F7 C3 7D F7 C3 A7 m..r..u..x..}...
1FA0 F7 C3 AC F7 C3 BB F7 C3 C1 F7 C3 CA F7 C3 70 F7 ..............p.
1FB0 C3 B1 F7 82 F6 00 00 00 00 00 00 6E F8 73 F6 0D ............n.s..
1FC0 F9 EE F8 82 F6 00 00 00 00 00 00 6E F8 73 F6 3C ............n.s.<
1FD0 F9 1D F9 82 F6 00 00 00 00 00 00 6E F8 73 F6 6B ............n.s.k
1FE0 F9 4C F9 82 F6 00 00 00 00 00 00 6E F8 73 F6 9A .L..........n.s..
1FF0 F9 7B F9 1A 00 03 07 00 F2 00 3F 00 C0 00 10 00 .{.........?.....
2000 02 00 01 07 0D 13 19 05 0B 11 17 03 09 0F 15 02 ................
2010 08 0E 14 1A 06 0C 12 18 04 0A 10 16 0D 0A 0A 36 ...............6
2020 33 6B 20 43 50 2F 4D 20 76 65 72 73 20 32 2E 32 3k CP/M vers 2.2
2030 0D 0A 00 31 00 01 21 9C F6 CD D3 F7 AF 32 04 00 ...1..!......2..
-^C
```

Figure 7-9. Using DDT to check CP/M images

```
                          In contrast, load DDT and request that it
                          load the PUTCPMF5.COM program.
A>ddt putcpmf5.com<cr>
DDT VERS 2.2
NEXT  PC
2900 0100

                          Display the special bootstrap loader that
                          starts at location 0780H (compared to the
                          MDS-800 bootstrap which is at 0980H). Note
                          the sign-on message.
-d780,7af<cr>
0780 C3 40 01 0D 0A 43 50 2F 4D 20 42 6F 6F 74 73 74 .@...CP/M Bootst
0790 72 61 70 20 4C 6F 61 64 65 72 0D 0A 56 65 72 73 rap Loader..Vers
07A0 69 6F 6E 20 30 31 20 30 37 2F 32 34 2F 38 32 0D ion 01 07/24/82.

                          Confirm that the CCP is loaded in the correct
                          place. Check the address of the first JMP
                          instruction (0E35CH).
-d980,9bf<cr>
0980 C3 5C E3 C3 58 E3 7F 00 20 20 20 20 20 20 20 20 .\..X...
0990 20 20 20 20 20 20 20 20 43 4F 50 59 52 49 47 48         COPYRIGH
09A0 54 20 28 43 29 20 31 39 37 39 2C 20 44 49 47 49 T (C) 1979, DIGI
09B0 54 41 4C 20 52 45 53 45 41 52 43 48 20 20 00 00 TAL RESEARCH  ..

                          Confirm that the BDOS is also in place.
-d1180,118f<cr>
1180 00 16 00 00 09 85 C3 11 E8 99 E8 A5 E8 AB E8 B1 ................

                          Confirm that the BIOS has been loaded in the
                          correct place. Check the first JMP to get
                          some idea of the BIOS base address. Note the
                          sign-on message.
-d1f80<cr>
1F80 C3 F9 F6 C3 0C FE C3 62 F8 C3 78 F8 C3 86 F8 C3 .......b..x.....
1F90 A4 F8 C3 B4 F8 C3 C5 F8 C3 B6 FB C3 0E FB C3 3B ...............;
1FA0 FB C3 41 FB C3 48 FB C3 DE FB C3 F8 FB C3 94 F8 ..A..H..........
1FB0 C3 B0 FB ED 06 00 00 00 42 6E 25 DF 01 B6 DE 02 ........Bn%.....
1FC0 38 00 00 43 50 2F 4D 20 32 2E 32 2E 30 30 20 30 8..CP/M 2.2.00 0
1FD0 37 2F 31 35 2F 38 32 0D 0A 0A 53 69 6D 70 6C 65 7/15/82...Simple
1FE0 20 42 49 4F 53 0D 0A 0A 44 69 73 6B 20 43 6F 6E  BIOS...Disk Con
1FF0 66 69 67 75 72 61 74 69 6F 6E 20 3A 0D 0A 0A 20 figuration :...
2000 20 20 20 20 41 3A 20 30 2E 33 35 20 4D 62 79 74     A: 0.35 Mbyt
2010 65 20 35 22 20 46 6C 6F 70 70 79 0D 0A 20 20 20 e 5" Floppy..
2020 20 20 42 3A 20 30 2E 33 35 20 4D 62 79 74 65 20   B: 0.35 Mbyte
2030 35 22 20 46 6C 6F 70 70 79 0D 0A 0A 20 20 20 20 5" Floppy...
-^C
A>_
```

Figure 7-9. Using DDT to check CP/M images (continued)

Putting it all Together

Figure 7-10 shows an annotated console dialog for the complete generation of a new CP/M system. Note that the following file names appear in the dialog:

```
BIOS1.ASM      Figure 6-4.
PUTCPMF5.ASM   Figure 7-5.
BOOTF5.ASM     Figure 7-7.
```

Assemble the CP/M Bootstrap Loader, with the source code and HEX file on drive C:, no listing output.

```
C>asm bootf5.ccz<cr>
CP/M ASSEMBLER - VER 2.0
02E4
004H USE FACTOR
END OF ASSEMBLY
```

Assemble the PUTCPMF5 program (that writes CP/M onto the disk), with the source code and HEX file on drive C:, no listing output.

```
C>asm putcpmf5.ccz<cr>
CP/M ASSEMBLER - VER 2.0
01DB
003H USE FACTOR
END OF ASSEMBLY
```

Assemble the BIOS with the source code and HEX file on drive C:, no listing output.

```
C>asm bios1.ccz<cr>
CP/M ASSEMBLER - VER 2.0
FE6C
011H USE FACTOR
END OF ASSEMBLY
```

Start piecing the CP/M image together. Load DDT and ask it to read in the file previously SAVEd after a MOVCPM 63 *.

```
C>ddt cpm63.com<cr>
DDT VERS 2.2
NEXT  PC
2300 0100
```

Indicate the file name of PUTCPMF5.HEX, and read in without any offset (i.e. it will load at 100H because of the ORG 100H it contains). -iputcpmf5.hex<cr>

```
-r<cr>
NEXT  PC
2300 0100
```

Indicate the file name of BOOTF5.HEX and read in with an offset of 680H to make it load at 780H on up (it contains ORG 100H too).

```
-ibootf5.hex<cr>
-r680<cr>
NEXT  PC
2300 0100
```

Indicate the file name of the BIOS HEX file, and read it in with an offset of 2980 such that it will load at 1F80H (it contains an ORG 0F600H).

```
-ibios1.hex<cr>
-r2980<cr>
NEXT  PC
27EC 0000
```

Exit from DDT by going to location 0000H and executing a warm boot.

```
-g0<cr>
```

Save the complete CP/M image on disk. Saving 40 256-byte pages from location 100H to 2900H.

```
C>save 40 putcpmf5.com<cr>
```

Figure 7-10. Console dialog for system build

	Load and execute the PUTCPMF5 program.
C>putcpmf5<cr>	
	PUTCPMF5 signs on
Put CP/M on Diskette Version 01 07/24/82	
	and writes the CP/M image to disk.
C>	

Figure 7-10. Console dialog for system build (continued)

8

Writing An Enhanced BIOS

This chapter describes ways in which you can enhance your BIOS to make CP/M easier to use, faster, and more versatile.

Get a standard BIOS working on your computer system, and then install the additional features. Although you can write an enhanced BIOS from the outset, it will take considerably longer to get it functioning correctly.

A complete listing of an enhanced BIOS is included at the end of this chapter. It is quite large: approximately 4500 lines of source code, with extensive comments and long variable names to make it more understandable.

The sections that follow describe the main concepts embodied in the enhanced BIOS listing.

BIOS Enhancements

BIOS enhancements fall into two classes: those that add new capabilities and those that extend existing features.

Some enhancements are normally accompanied by utility programs that allow you to select the enhancement option from the console. For example, when the BIOS is enhanced to include a *real time clock,* you need a utility program to set the clock to the correct time. Other enhancements will not require supporting utilities. For example, if the disk drivers are improved to read and write data faster, the enhancement is "transparent." As a user, you are aware of the results of the enhancement but not of the enhancement itself.

Viewed at its simplest, the BIOS deals with two broad classes of input/output:

Character input/output
This includes the console, auxiliary, and list devices.

Disk input/output
This can accommodate several types of floppy and hard disks.

Enhancements in these areas do not fundamentally change the way that the BDOS and CCP interact with these devices. Instead, enhancements improve the way in which the *device drivers* deal with the devices. They can improve the speed of manipulating data, the way of handling external devices, or the user's control over the behavior of the system.

The example enhanced BIOS has capabilities not found in standard CP/M systems. These can be grouped in several main categories:

Character input/output
This area probably benefits most from enhancement. This is partly because such a wide range of peripheral devices needs to be supported and partly because this is the most visible area of interaction between you and your computer. Any improvements here will therefore be immediate and obvious to you as a user.

Error handling
CP/M's error handling is, at best, startling in its simplicity. Enhanced error handling gives you more information about the nature of the failure, and then gives you the options of retrying the operation, ignoring the error, or aborting the program. This topic is covered in detail in Chapter 9.

System date and time
This is the ability to maintain a time-of-day clock and the current date. It allows your programs to set and access the date and time. In addition, your system can react to the passing of time, and you can move certain operations into the time domain. For example, you can set upper limits on the

number of seconds, or milliseconds, that each operation should take, and arrange for emergency action if the operation takes too long.

Logical-to-physical device assignment
CP/M's logical-to-physical device assignment is primitive. With enhancements, you can use any character input/output device as the system console, and output data to several devices at the same time.

Disk input/output
CP/M only knows about the 128-byte sector. Even with the deblocking routines shown in Figure 6-4, overall disk performance can be slow. Performance can be improved dramatically by "track buffering" (in which entire tracks are read and written at one time) or by using a *memory disk* (that is, using large areas of RAM as though they were a disk). These have a cost, though, in increased memory requirements.

Public files
CP/M's user number system needs improvements to function well in conjunction with large hard disks.

Preserving User-Settable Options

A by-product of adding features to the BIOS is that many of these features have options that you can alter, either from the console using a utility program or from within one of your programs.

Each of these options, once set according to your preferences, or to the requirements of your hardware, do not normally change from day to day. Therefore, the BIOS should be designed so that options set by the user can be "frozen" or preserved on the disk by using a utility program, FREEZE. All of the variables recording these options are gathered into a single area and then this area is written out to the disk.

This area is called the *configuration block*. In practice, there are two configuration blocks: one short term and the other long term. The short term block is not preservable — you can set options within it, but they cannot be preserved after you switch your computer off. The system date, for example, is normally set each time you turn your computer on, and therefore is kept in the short term block. The baud rate for your printer, on the other hand, is kept in the long term block so that it can be saved permanently.

An extra BIOS entry point, CBGetAddress, has been built into the enhanced BIOS so that utility programs can locate variables in both configuration blocks. For example, when a utility needs to know where the date is kept in memory, it calls CBGetAddress using a code number (specific for date) in a register. CBGetAddress returns the address of the date in memory. If a new version of the BIOS is produced with the date in a different location, CBGetAddress will still hand the correct, although different, address back to the utility program.

Two other variables that CBGetAddress can access pertain to the configuration block itself. One is the relative address of the start of the long term configuration block. The other is the length of the long term block. These are used by the FREEZE utility when it needs to preserve the long term block on a disk. FREEZE must (1) read in the sectors containing the long term block from the CP/M BIOS image on the reserved area of the disk, (2) copy the current RAM-resident version of the long term block over the disk image version, and then (3) write the sectors back onto the disk.

Figure 8-1 shows how the long term block appears on disk and in memory. The

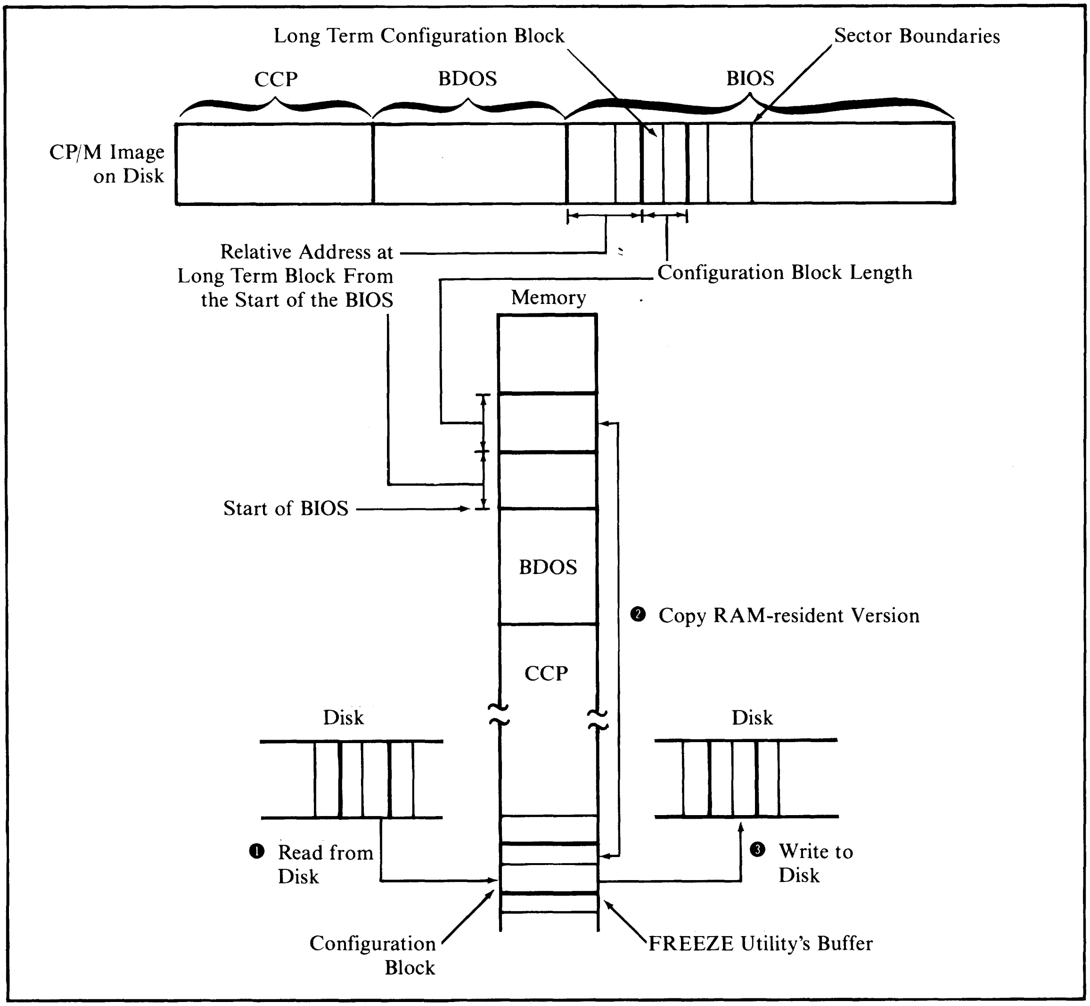


Figure 8-1. Saving the long term configuration block

size of the CCP and BDOS do not change, even if the BIOS does. Therefore, the sector containing the start of the BIOS will not change. The formula (using decimal numbers)

BIOS Start Sector + INT(Relative LTB Address / 128)

then gives the start sector number to be read in. The number of sectors to read is calculated as follows:

(Long Term Block Length + 127) / 128

The relative address and length can be used to locate the long term block in the BIOS executing in RAM.

Character Input/Output

The character I/O drivers shown in the example BIOS, Figure 8-10, have been enhanced to have the following features:

- A single set of driver subroutines controlling all character devices
- Preservation of option settings
- Flexible redirection of input/output between logical and physical devices
- Interrupt-driven input drivers, to get user "type-ahead" capability
- Support of several different protocols to avoid loss of data during high-speed output to printers or other operations
- Forced input of characters into the console input stream, allowing automatic commands at system start-up
- Conversion of terminal function keys into useful character strings
- Ability to recognize "escape sequences" output to the console and to take special action as a result
- Ability to read the current time and date as though they were typed on the console
- "Timeout" signaling when the printer is busy for too long.

Each of these features is discussed in the following sections, as an introduction to the actual code example.

Single Set of Driver Subroutines

In the following examples, only a single set of subroutines is used to process the input and output for all of the physical devices in the system.

This is made possible by grouping all of the individual device's characteristics into a table called the *device table*. For example, in order to get a character from the current console device, the address of its device table will be handed over to the subroutines. These in turn will use the appropriate values from the device table when they need to access a port number or any unique attribute of that device.

In our example, the drivers assume that all of the physical devices use serial input/output. To support a device with parallel input/output, you would need to extend the device table to include a field that would enable the drivers to detect whether they were operating on a serial or parallel device. You would probably also have to add different device initialization and input/output routines more suited to the problems of dealing with a parallel port.

The device table structure consists of a series of equate (EQU) instructions. These define the relative offset of each field in the table. Each definition is expressed by referencing the *preceding* field so that you can insert additional fields without revising the definitions for all the other fields.

Individual instances of device tables are then defined as a series of define byte (DB) and define word (DW) lines. The drivers are given the base address of the device table whenever they need to do something with a device. By adding the base address to the relative address (defined by the equate), the drivers can determine the actual address in memory that contains the required value. The detailed contents of the device table are described later in this chapter.

Permanent Setting of Options

About the only options that need preserving in the long term configuration block are the values used to initialize the hardware chips. Other options can be set during automatic execution of the command file when CP/M is first loaded.

Redirection of Input/Output Between Devices

As you recall, the BDOS only "knows about" the *logical* devices console, reader, punch, and list. Using the IOBYTE at location 0003H in conjunction with the STAT utility, you can redirect the BDOS to assign the logical devices to specific physical devices. However, the redirection provided by CP/M is rather primitive. It permits only four physical devices per logical device. Input and output of a logical device must always come from the same physical device. Output data can only be sent to a single destination, or (using the CONTROL-P toggle) to the console and the list device.

The system in Figure 8-10 supports up to 16 physical devices. Any one of these devices can act as the console, reader, punch, or list device. Input can come from any single device. Output can be sent to any or all of the devices. Each logical device's input and output are separate—that is, console input can come from physical device X while the output can be sent to physical devices Y and Z.

Device redirection can be done dynamically, either from within a program or by using a system utility program. For example, if you have some special input device, your program can momentarily switch over to reading input from this device as though it were the console, and then revert back to reading data from the "real" console.

This redirection scheme is achieved by defining a 16-bit word, called the *redirection word,* in the long term configuration block for each of the following logical devices:

- Console input
- Console output
- Auxiliary (reader/punch) input
- Auxiliary (reader/punch) output
- List input (printers need to send data, too)
- List output.

Each bit in a given redirection word is assigned to a physical device. For input, the drivers use the device corresponding to the first 1 bit that they find in the redirection word. For output, the drivers send the character to be output to all of the devices for which the corresponding bit is set.

The example code does not select a different driver for each bit set — it selects a specific device table and then hands over the base address of this table to the common driver used for all character operations.

Interrupt-Driven Input Drivers

With a standard CP/M BIOS, character data is read from the hardware chips only when control is transferred to the CONIN or READER subroutines. If this character data arrives faster than the BIOS can handle, data overrun occurs and incoming characters are lost.

By using interrupts, the hardware can transfer control to the appropriate interrupt service routine whenever an incoming character arrives. This routine reads the data character and places it into a buffer area to wait for the next CONIN or READER call, which will get the character from the buffer and feed it into the incoming data stream.

User programs and the CCP are "unaware" of this process, perceiving only that data characters are available. However, users will become aware of the process; they will be able to enter data characters from the keyboard before the program is ready for them. This gives the technique its other name — "type-ahead." Although this technique does not alter the speed of execution of any programs running under CP/M, it does create the illusion of greater speed, since pauses while a program accepts data vanish completely. The user can enter data at a rate convenient to the tasks or thoughts at hand, without regard to the rate at which the program can accept that data.

The example contains the code necessary to handle arriving characters under interrupt control. In order to be of general applicability, the code assumes a "flat" interrupt structure: that is, all character input interrupts cause control to be transferred to the same address in memory. The address is determined by the actual hardware interrupt architecture.

The simplest interrupt schemes use the restart (RST) instructions built into the 8080 CPU chip. In the RST scheme, the external hardware interrupts what the CPU chip is doing and forces one of the eight RST instructions into the processor. Each RST instruction causes the processor to execute what is, in effect, a CALL instruction to a predetermined address in memory.

In more complicated systems, a specific interrupt controller chip (such as the Intel 8259A) will be used. In addition to providing very sophisticated (and complicated) prioritization of interrupts, the interrupt controller can transfer control to a *different* address depending on which physical device causes the interrupt. It does this by forcing the CPU to execute a CALL instruction to a different address for each device.

In both architectures, it is the responsibility of the BIOS writer to initialize all the hardware chips so that an interrupt occurs under the correct circumstances. The BIOS writer also must plant instructions at the correct places in memory to receive control from an RST instruction or from the fake CALL instruction emitted by the interrupt controller.

Some hardware requires that the interrupt service subroutine inform it as soon as the interrupt has been serviced and the character has been input. The example drivers provide for this.

This section deals with using interrupts for the *input* drivers, not the output drivers. All of today's microcomputers can output data much faster than external peripherals can handle. After the first few minutes of output, the computer will fill any reasonably sized buffer—and from this point there is no advantage in having a buffered output system. The computer still must slow down to the peripheral's data rate for each character, although now it is waiting to put the character in the output buffer rather than out to the peripheral.

One exception to this is where you have a large amount of "spare" memory and a "slow" printer (which most of them are). Increasing numbers of systems have more than 64K of RAM. The 8080 or Z80 can't address more than this, but a "bank switched" memory system can switch blocks of memory in and out of that 64K address space.

Using this trick, you can access memory "unknown" to CP/M, store some characters in it, switch back to the normal 64K memory, and return control to the caller of the BIOS output routine. When the physical device is ready to accept another output data character from the CPU, it will generate an interrupt. The interrupt service routine then will access the "secret" buffer, output the characters to the device, and switch back to the normal memory.

For example, if you have a printer that prints at 80 characters per second and

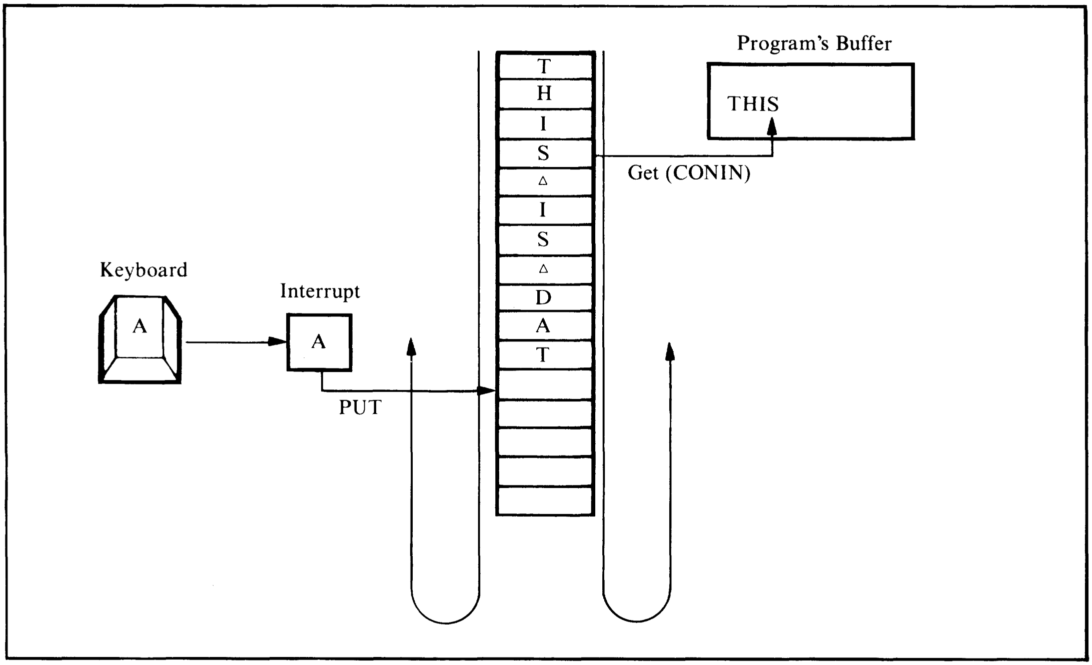


Figure 8-2. Circular buffer type-ahead

you can afford to use 64K of bank switched memory, you can squirrel away 13 minutes of printing—or even more if you design a scheme to compress blanks, storing them in the hidden buffer as a special control sequence.

From the point of view of software, interrupt-driven input drivers are divided into two major groups: the interrupt service routine that reads the characters and stacks them in a buffer, and the non-interrupt routines that get the characters from the buffer and handle the other BIOS functions such as returning console status.

The input character buffer serves as a transfer mechanism between the two groups of subroutines, although the device table also plays an important role.

The example code uses a circular buffer, as shown in Figure 8-2.

The drivers start putting data into the beginning of the buffer. When the last character in the buffer has been reached, the drivers reset to the beginning of the buffer and start over. This, of course, assumes that the non-interrupt drivers have been getting data from the front of the buffer, thus creating space for additional incoming data.

Each device table contains the address of the input buffer, a "put" pointer (for the interrupt service routine), and a "get" pointer (for the non-interrupt service routine). It also contains two character counts: the total number of characters and the number of control characters in the input buffer. You can see how the put and get pointers operate asynchronously. The put pointer is used every time an incoming character generates an interrupt. The get pointer is used for each CONIN call.

The get and put pointers are only single-byte values and are more accurately described as "relative offsets." That is, they contain a value which, when converted to a word and added to the base address of the buffer, will point directly to the appropriate position inside the buffer.

By making the buffer a binary number of characters long—32 characters, for example—a programming trick can be used to make the buffer appear circular. The device tables contain a mask value formed from the buffer's length minus one (length − 1). Whenever the get or put pointers are incremented by one (to "point" to the next character position), the updated value is ANDed with this (length − 1) mask. In this example, if the get value goes from 31 (the relative address of the last character in the buffer) to 32 (which would be "off the end"), the masking operation will reset it to zero (the relative address of the first character of the buffer). This avoids having to compare pointers to know when to reset them.

It is also simpler to use a count of the number of characters in the buffer, rather than comparing the get and put pointers, to distinguish between an empty and a full buffer. To support different serial protocols, the driver must be able to react when the buffer is within five characters of being full and when it drops below half empty. Both of these conditions are much easier to detect using a simple count that is incremented as a character is put into the buffer and decremented as a character is retrieved from the buffer.

The count of control characters is used to deal with a class of programs that incessantly "gobble" characters, thereby rendering any type-ahead useless. An example is Microsoft's BASIC interpreter. When it is interpreting a program, you can enter a CONTROL-C from the keyboard and the interpreter will come to an orderly stop. It does this by constantly making calls to CONST (console status). If it ever detects an incoming character, it makes a call to CONIN to input the character. A character that is not CONTROL-C is discarded without further ado. Thus, any characters that are input are consumed, destroying the effect of type-ahead.

To deal with this problem, the CONST routine shown in the example can be told to "lie" about the console's status. In this mode, CONST will only indicate that characters are waiting in the input buffer if a control character is received. It uses the control character count to determine whether there are control characters in the buffer; this count is incremented by the interrupt service routine when it detects one, and decremented by the CONIN routine when it gets a control character from the buffer.

Protocol Support

In this context, a protocol is a scheme to avoid loss of data that would otherwise occur if a device sent data faster than the receiving device could handle it. For example, protocols are used to prevent the CPU sending data out to a printer faster than the printer can print the characters and move the paper. The drivers also support input protocols, indicating to a transmitting device when the input buffer gets close to being full.

Two basic methods are used to implement protocols. The first uses the control lines found in the normal RS-232C serial interface cables. For data being output by the computer, the data terminal ready (DTR) signal is used, and for incoming data, the request to send (RTS) signal. These signals conform to the electrical standards for the RS-232C interface; they are considered true when they are at some positive voltage between +3 and +12 volts, and false when they are between −3 and −12 volts.

The second method uses ASCII control characters instead of control signals. Two separate protocols are supported by this method. One uses the ASCII characters XON and XOFF. Before the sending device (the computer or some peripheral device) sends a data character, it checks to see if an XOFF character has been received. If so, the sender will wait for an XON character. The receiving device will only send an XON when it is ready to receive more data.

The second protocol uses the characters ETX (end of transmission) and ACK (acknowledge). This method is normally used only when transmitting data from the computer to a buffered printer. A message length (usually half the printer's buffer size) is defined. When this number of characters has been output, the computer will send an ETX character. No further output will occur until the computer receives an ACK character from the printer.

The example drivers support the DTR high-to-send, the XON/XOFF, and the ETX/ACK protocols for output data. For input, they support RTS high-to-receive and XON/XOFF.

The input protocols are invoked when the input buffer gets within five characters of being full. Then the drivers output an XOFF character or lower the RTS signal voltage, or do both. Only when the input buffer has been emptied to 50% capacity will the drivers send XON or raise the RTS line, or both.

As an emergency measure, if the input buffer becomes completely full, notwithstanding protocols, the drivers will output a predetermined character (defined in the device table) each time they discard an incoming character. This is normally the ASCII BEL (bell) character. When you type too far ahead, the terminal will start beeping to tell you that data is being dropped.

Forced Input into the Console Stream

All application languages provide a means of reading data from the console keyboard. This makes the console input stream a useful gateway to the system. A simple enhancement to the CONIN/CONST routines makes it easy to "fool" the system into acting as if data had been input from the keyboard when in fact the data is coming in from a character string in memory.

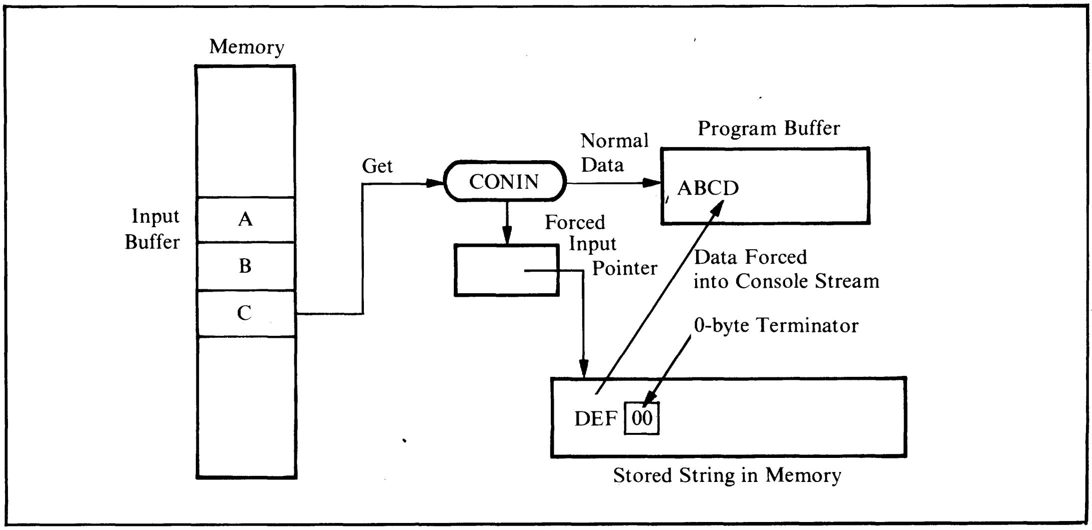


Figure 8-3. CONIN uses forced input data if pointer points to nonzero byte

In the enhanced BIOS, both CONIN and CONST are extended to check a pointer in the long term configuration block, as shown in Figure 8-3.

If this pointer is pointing at a nonzero byte, then that byte is returned as though it had come from the console keyboard. The forced input pointer is then moved up one byte in memory. The process of forcing input continues until a zero byte is encountered.

Forced input serves several purposes. It can be used to force a command or commands into the system when the system first starts up. In conjunction with a utility program, it can allow the user to enter several CP/M commands on a single command line, injecting the characters as each of the commands is executed. It also makes possible the features described in the next two sections.

Support of Terminal Function Keys

Many terminals on the market today have special function keys on their keyboards. When you press one of these keys, the terminal will emit several characters, the first of which is normally the ASCII ESC (escape) character. The remaining one or two characters identify the specific function key that was pressed.

For these function keys to be of any practical use, an applications program must detect the incoming escape sequence and take appropriate action. The problem is that not all terminal manufacturers support the ANSI standard escape sequences.

The example drivers avoid this problem by providing a general-purpose method, shown in Figure 8-4, of detecting escape sequences and of substituting a user-defined character string that is injected into the console input stream as though it had been entered from the keyboard.

This scheme permits function keys to be used very flexibly, even for off-the-shelf programs that have not been designed specifically to accept function key input.

There is, however, one stumbling block. When an ESCAPE character is received, the progam must detect whether this is the start of a function key sequence or the user pressing the ESCAPE key on the terminal's keyboard. In the former case, the

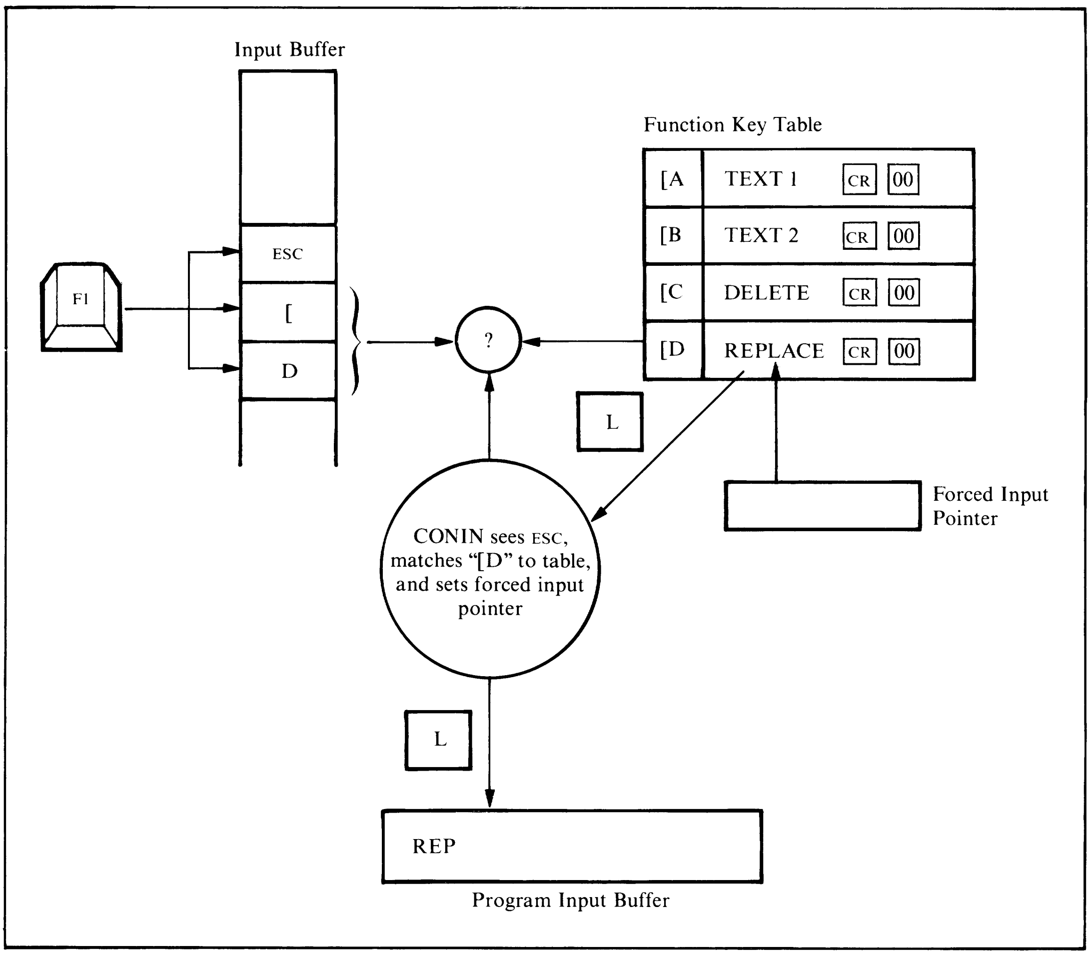


Figure 8-4. CONIN decodes terminal function keys

driver must wait to determine whether a function key string must be substituted for the escape sequence. In the latter case, the driver must input the ESCAPE character as it would other incoming data characters.

This recognition can only be done by moving into the time domain. When the CONIN routine (the non-interrupt routine) gets an ESCAPE character from the input buffer, it delays for approximately 90 milliseconds, enough time for a terminal-generated character sequence to arrive. CONIN then checks the input buffer to see if it contains at least two characters. If it does, the driver checks for a match in a function key table in the long term configuration block. If the characters match a defined function key, then the string associated with the function key will be injected into the console stream by pointing the forced input pointer at it. If the characters do not match anything in the function key table, then the ESCAPE and subsequent characters are handed over as normal data characters.

If after the 90-millisecond delay no further characters have arrived, the ESCAPE character is handed over as a normal character, on the basis that it must have been a manually entered ESCAPE character rather than part of a terminal-generated sequence.

The example drivers show the necessary code and tables for function keys that emit three characters. You could modify them easily for two-character sequences, or, if you are fortunate enough to have a keyboard that uses all eight bits of a byte, to recognize single incoming characters.

Processing Output Escape Sequences

The output side of the console driver, the CONOUT routine, can also be enhanced to recognize escape sequences. It uses a vectored JMP instruction to keep track of the current state of affairs. The CONOUT driver gets an address from the vector and transfers control to it. Normally this vector is set to direct control to the output byte routine. However, if an ESCAPE character is detected in the output stream, the vector is changed to transfer control to a routine that will recognize the character following the ESCAPE. If recognition does not occur, the driver will output an ESCAPE followed by the character that arrived after it.

If the second character is recognized, then the driver can transfer control to the correct escape-sequence processor. This processor can then take whatever action is appropriate. It must also make sure that when all processing is finished, the console output vector is set to process normal output characters again.

This technique is described in more practical detail in the next section, where it is used to preset and read the date and time. You can easily extend the recognition tables in the long term configuration block to perform any special processing that you need, ranging from altering the I/O redirection words to changing any other variable in the system or programming special hardware in your computer.

Be careful not to embed any pure binary values in the sequence of characters going out to the CONOUT routine. If you attempt to send a value of 09H (the TAB character) out via the BDOS, it will gratuitously expand the tab out to some number of blanks. If you need to send out a bit pattern, such as the I/O redirection word, split it up into a series of 7-bit long values. Then send it out with each byte having the most significant bit set to 1. A value of 09H will then become 89H, preventing the BDOS from expanding it to blanks.

Reading Date and Time From Console

For the moment, set aside the question of how the date and time get into the system. Since the date and time are stored in the short term configuration block (there being no need to save them from one work session to the next), all that the BIOS needs to be able to do is recognize a request from an applications program to read either the date or the time and then set the forced input pointer to the appropriate string in memory. Both the date and time strings are terminated by a LINE FEED followed by a 00 byte.

This sequence of events is shown in Figure 8-5.

You can see that the characters "ESC d" output to CONOUT cause it to point the forced input pointer at the date in memory. Subsequent calls to CONIN bring the characters in the date into the program as though they were being entered on the keyboard.

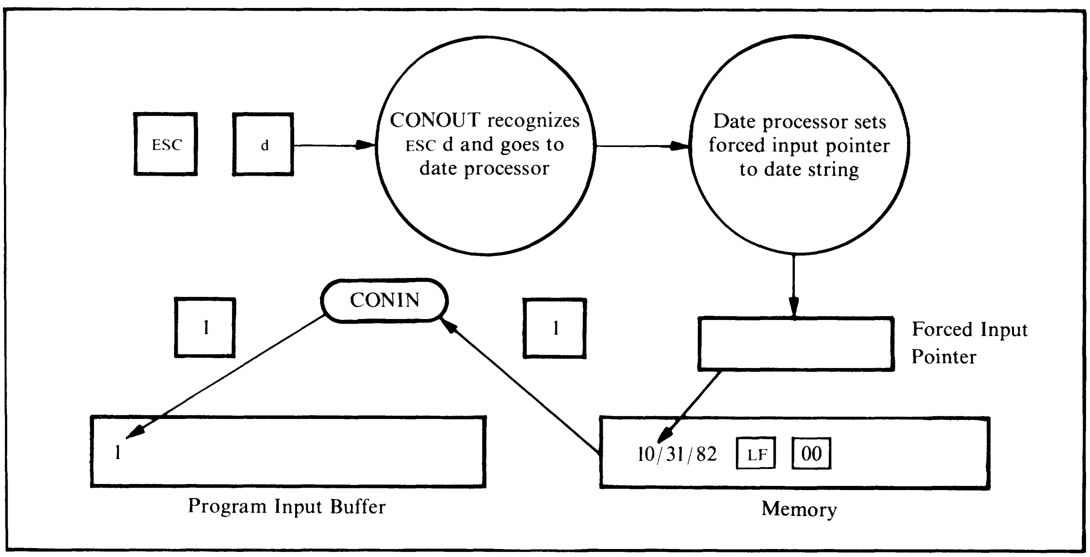


Figure 8-5. Escape sequences sent to CONOUT allow the date to be read by CONIN

"Watchdog" Timeout on Printer

There is no provision in CP/M to deal with a hardware device that for one reason or another is permanently unavailable. Unless special steps are taken in the drivers, the system will screech to a halt in a loop, reading status and testing for the peripheral to be ready.

The example enhancement code shows a scheme, using a real time clock, that can detect when a device such as a printer fails to come ready for more than 30 seconds. On detecting this situation, the code outputs a message to all of the console devices that are not also being used as printers. This type of output is needed to avoid "deadly embraces" where a printer not being ready generates a message that cannot be output because the printer is not ready.

The code that performs the timing function is known as a *watchdog timer*. Each time the real time clock "ticks," the interrupt service routine checks the watchdog count. If the count is nonzero, it is decremented. If the watchdog timer reaches zero, exceeding the time allowed, the drivers will display a message on the console indicating that the printer has been busy for too long. The user then has the option of making the printer ready and trying again to output data, ignoring the error and carrying on, or aborting the program by doing a BDOS System Reset (function 0).

Although sending an error message to the console sounds simple, it is complicated if console output is directed to the offending printer itself. The drivers attempt to solve this problem by sending the message only to those devices being used as consoles and *not* as printers. If all consoles are being used as printer devices as well, the driver will send the message to device 0—normally the main console.

Keeping Time and Date

CP/M does not have provision for keeping the current time and date in the system. The example enhancement shows how to keep the time of day and the current date in the short term configuration block by using escape sequences output to the console (1) to set them to the correct values and (2) to "read" them from the console input stream.

The example presupposes that the system has a hardware chip that can be programmed to generate an interrupt every 1/60th of a second (16.666 milliseconds). This provides a divide-down counter to measure seconds elapsed. Of course, if your computer has a *true* real time clock that you can read and get the current time in hours, minutes, and seconds, your code will be very simple. You still will need to have the clock generate a periodic interrupt, however, in order to use the watchdog feature for timing printer and disk operations.

Actual time is kept as ASCII characters, using another ASCII control table to determine when "carry and reset to zero" should occur. By changing two bytes in this table, the time can be kept in 12- or 24-hour format.

The date is simply stored as a string. The example code does not attempt to make sure that the date is valid, nor to update when midnight rolls around. This could be done easily by the BIOS — but it would take a fairly large amount of code.

Watchdog Timer

Having a periodic source of interrupts also opens the door to building in an emergency or watchdog timer. This is nothing more than a 16-bit counter. Each time the real time clock interrupts, or ticks, the interrupt service routine checks the watchdog count. If it is already at zero, nothing more happens — the watchdog is not in use. If it is nonzero, the routine decrements the count by one. If this results in a zero value, the interrupt service routine CALLs a predetermined address. This will be the address of some emergency interrupt service routine that can then take special action, such as investigating the cause of the timeout.

The watchdog routine has a non-interrupt-level subroutine associated with it. Calling this set watchdog subroutine provides a means of setting the count to a predetermined number of real time clock "ticks" and setting the address to which control should be transferred if the count reaches zero.

Having called the set watchdog subroutine, the driver can then sit in a status loop, with interrupts enabled, waiting for some event to occur. If the event happens before the watchdog count hits zero, the driver must call the set watchdog routine again to set the count back to zero, thereby disabling the watchdog mechanism.

The watchdog timer can be used to detect printers that are busy for too long or disk drives that take too long to complete an action either because of a hardware failure or because the user has not loaded the disk into the drive.

Data Structures

As already stated, each character I/O device has its own device table that describes all of its unique characteristics.

The other major data structure is the configuration blocks — both short and long term.

This section describes each field in these data structures.

Device Table

Figure 8-6 shows the contents of a device table. More correctly, it shows a series of equates that define the offsets of each field in the device table. The drivers are given the base address of a specific device table. They then access each field by adding the required offset to this base address.

The first part of the device table is devoted to the physical aspect of the device, defining which port numbers are to be used to communicate with it. The drivers need to know several different port numbers since each one is used for a particular

```
                ;       The drivers use a device table for each
                ;       physical device they service. The equates that follow
                ;       are used to access the various fields within the
                ;       device table.
                ;
                ;                               Port numbers and status bits
0000 =          DT$Status$Port          EQU     0               ;Device status port number
0001 =          DT$Data$Port            EQU     DT$Status$Port+1
                                                                ;Device data port number
0002 =          DT$Output$Ready         EQU     DT$DataPort+1
                                                                ;Output ready status mask
0003 =          DT$Input$Ready          EQU     DT$Output$Ready+1
                                                                ;Input ready status mask
0004 =          DT$DTR$Ready            EQU     DT$Input$Ready+1
                                                                ;DTR ready to send mask
0005 =          DT$Reset$Int$Port       EQU     DT$DTR$Ready+1
                                                                ;Port number used to reset an
                                                                ;  interrupt
0006 =          DT$Reset$Int$Value      EQU     DT$Reset$Int$Port+1
                                                                ;Value output to reset interrupt
0007 =          DT$Detect$Error$Port    EQU     DT$Reset$Int$Value+1
                                                                ;Port number for error detect
0008 =          DT$Detect$Error$Value   EQU     DT$Detect$Error$Port+1
                                                                ;Mask for detecting error (parity etc.)
0009 =          DT$Reset$Error$Port     EQU     DT$Detect$Error$Value+1
                                                                ;Output to port to reset error
000A =          DT$Reset$Error$Value    EQU     DT$Reset$Error$Port+1
                                                                ;Value to output to reset error
000B =          DT$RTS$Control$Port     EQU     DT$Reset$Error$Value+1
                                                                ;Control port for lowering RTS
000C =          DT$Drop$RTS$Value       EQU     DT$RTS$Control$Port+1
                                                                ;Value, when output, to drop RTS
000D =          DT$Raise$RTS$Value      EQU     DT$Drop$RTS$Value+1
                                                                ;Value, when output, to raise RTS
                ;
                ;               Device logical status (incl. protocols)
000E =          DT$Status               EQU     DT$Raise$RTS$Value+1
                                                                ;Status bits
0001 =          DT$Output$Suspend       EQU     0000$0001B      ;Output suspended pending
                                                                ;  protocol action
0002 =          DT$Input$Suspend        EQU     0000$0010B      ;Input suspended until
                                                                ;  buffer empties
0004 =          DT$Output$DTR           EQU     0000$0100B      ;Output uses DTR-high-to-send
0008 =          DT$Output$Xon           EQU     0000$1000B      ;Output uses Xon/Xoff
0010 =          DT$Output$Etx           EQU     0001$0000B      ;Output uses Etx/Ack
0020 =          DT$Output$Timeout       EQU     0010$0000B      ;Output uses Timeout
0040 =          DT$Input$RTS            EQU     0100$0000B      ;Input uses RTS-high-to-receive
0080 =          DT$Input$Xon            EQU     1000$0000B      ;Input uses Xon/Xoff
                ;
000F =          DT$Status$2             EQU     DT$Status+1     ;Secondary status byte
0001 =          DT$Fake$Typeahead       EQU     0000$0001B      ;Requests Input$Status to
                                                                ;  return "Data Ready" when
                                                                ;  control characters are in
                                                                ;  input buffer
                ;
0010 =          DT$Etx$Count            EQU     DT$Status$2+1
                                                                ;No. of chars. sent in Etx protocol
0012 =          DT$Etx$Message$Length   EQU     DT$Etx$Count+2
                                                                ;Specified message length
                ;
                ;                               Input buffer values
0014 =          DT$Buffer$Base          EQU     DT$Etx$Message$Length+2
                                                                ;Address of input buffer
0016 =          DT$Put$Offset           EQU     DT$Buffer$Base+2
                                                                ;Offset for putting chars. into buffer
0017 =          DT$Get$Offset           EQU     DT$Put$Offset+1
                                                                ;Offset for getting chars. from buffer
0018 =          DT$Buffer$Length$Mask   EQU     DT$Get$Offset+1
                                                                ;Length of buffer - 1
                                                                ;Note: Buffer length must always be
                                                                ;  a binary number; e.g. 32, 64, or 128,
                                                                ;This mask then becomes:
                                                                ;  32 ->  31 (0001$1111B)
                                                                ;  64 ->  63 (0011$1111B)
                                                                ; 128 -> 127 (0111$1111B)
```

Figure 8-6. Device table equates

```
                                                        ;After the get/put offset has been
                                                        ; incremented it is ANDed with the mask
                                                        ; to reset it to zero when the end of
                                                        ; the buffer has been reached.
0019 =         DT$Character$Count       EQU     DT$Buffer$Length$Mask+1
                                                        ;Count of the number of characters
                                                        ;  currently in the buffer
001A =         DT$Stop$Input$Count      EQU     DT$Character$Count+1
                                                        ;Stop input when the count reaches
                                                        ;  this value
001B =         DT$Resume$Input$Count    EQU     DT$Stop$Input$Count+1
                                                        ;Resume input when the count reaches
                                                        ;  this value
001C =         DT$Control$Count         EQU     DT$Resume$Input$Count+1
                                                        ;Count of the number of control
                                                        ;  characters in the buffer
001D =         DT$Function$Delay        EQU     DT$Control$Count+1
                                                        ;Number of clock ticks to delay to
                                                        ;  allow all characters after function
                                                        ;  key lead-in to arrive
001E =         DT$Initialize$Stream     EQU     DT$Function$Delay+1
                                                        ;Address of byte stream necessary to
                                                        ;  initialize this device
```

Figure 8-6. Device table equates (continued)

function. Depending upon your hardware, each port number could be different; however, with standard Intel or Zilog chips, you will often find that the same port number is used for several functions. The drivers also need to know what bit patterns to expect when they read some ports and what values to output to ports in order to obtain particular results.

The layout of the device table and the manner in which the equates are declared are designed to make it easy for you to change the contents of the table to meet your own special requirements. The fields in this first section of the device table are discussed in the sections that follow.

DT$Status$Port The driver reads this port to determine whether the hardware chip has incoming data ready to be input to the computer or whether the chip is capable of accepting another data character for output to the physical device.

DT$Data$Port The driver reads from this port to access the next data character from the physical device. The driver also writes to this port to output the next data character to the device.

If your computer hardware requires that the input data port be a different number from the output data port, you will have to alter the coding in the device table equates as well as make the necessary changes in the input and output subroutines in the body of the code.

DT$Output$Ready This is the bit mask that the driver will AND with the current device status (obtained by reading the DT$Status$Port) to see whether the device is ready to accept another output character. It assumes that the device is ready if the result of the AND instruction is nonzero. You may have to change some JNZ (jump nonzero) instructions to JZ (jump zero) instructions if your hardware device uses inverted logic, with bits in the status byte set to 0 to indicate that the device can accept another character for output.

Note that this status check relates only to the output chip — it is completely separate from the question of whether the peripheral itself is ready to accept data.

DT$Input$Ready This is the bit mask that the driver will AND with the current device status to see if there is an incoming data character. The drivers again presume that if the result of the AND is nonzero, then an incoming data character is waiting to be read from the data port. You will need to make changes similar to those for the output subroutines described in the previous section if your hardware uses inverted logic (0 bit means incoming data).

DTDTRReady DTR stands for *data terminal ready*. It refers to one of the control lines connected from the actual peripheral device to the I/O chip (via several other integrated circuits). The drivers, as an option, will only output data to the device when the DTR signal is at a positive voltage. If the peripheral, in order to stop the flow of data characters being output to it, lowers the DTR signal to a negative voltage, the drivers will wait. Once DTR goes positive again, the drivers will resume sending data. Many hard-copy devices use this scheme to give themselves a chance to print out data received from the computer. They may have to lower DTR for several seconds, while they perform paper movement, for example.

The value in this field is a bit mask that the drivers use on the device status to determine the state of the data-terminal-ready control signal.

DT$Reset$Int$Port Since the input side of the drivers uses interrupts, when an incoming character is ready to be input by the CPU, the hardware generates an interrupt signal, and control is transferred to the interrupt service routine. This routine "services" the interrupt by reading the incoming data character, saving it in memory, and then transferring control back to whatever was being executed when the interrupt occurred.

The more complicated interrupt controller chips (such as the Intel 8259A) must be told as soon as a given interrupt has been serviced so that they can permit servicing of any lower priority interrupts that may be waiting.

This field contains the port number that will be used to "reset" the interrupt, or more correctly, to indicate the end of the previous interrupt's servicing.

DT$Reset$Int$Value This is the value that will be output to the DT$Reset$Int$Port to tell the hardware that the previous interrupt service has been completed.

DT$Detect$Error$Port Before the driver attempts to read any incoming data from the DT$Data$Port, it checks to see if any hardware errors have occurred. It does so by reading status from this port.

DT$Detect$Error$Value The status byte that is input from the DT$Detect$Error$Port is ANDed with this value. If the result is nonzero, the driver assumes that an error has occurred.

DT$Reset$Error$Port If an error has occurred, the driver outputs an error reset value to this port number.

DT$Reset$Error$Value This is the value that will be output to the DT$Reset$Error$Port to reset an error.

DTRTSControl$Port The drivers use this port number to control the request-to-send line if the RTS protocol option is selected.

DT$Drop$RTS$Value This value is output to the RTS control port to lower the RTS line so that some external device will stop sending data to the computer.

DT$Raise$RTS$Value This value is output to raise the RTS line so that the external device will resume sending data to the computer.

DT$Status This is the first of two status bytes. It contains bit flags that are set to a 1 bit to indicate the following conditions:

DT$Output$Suspend
Because of protocol, the device is currently suspended from receiving any further output characters.

DT$Input$Suspend
Because of protocol, the device has been requested not to send any more input characters.

DT$Output$DTR
The driver will maintain DTR-high-to-send protocol for output data.

DT$Output$Xon
The driver will maintain XON/XOFF protocol for output data.

DT$Output$Etx
The driver will maintain ETX/ACK protocol for output data.

DT$Input$RTS
The driver will maintain RTS-high-to-receive protocol for input data.

DT$Input$Xon
The driver will maintain XON/XOFF protocol for input data.

DT$Status$2 This is another status byte, also with the following bit flag:

DT$Fake$Typeahead
CONST will "lie" about the availability of incoming console characters. It

will only indicate that data is waiting if there are control characters other than CARRIAGE RETURN, LINE FEED, or TAB in the input buffer.

DTEtxCount This value is only used for ETX/ACK protocol. It is a count of the number of characters sent in the current message. When this count reaches the defined message length, then the driver will send an ETX character and suspend any further output.

DTEtxMessage$Length This value is the defined message length for the ETX/ACK protocol. It is used to reset the DTEtxCount.

DT$Buffer$Base This is the address of the first byte of the device's input buffer.

DTPutOffset This *byte* contains the relative offset indicating where the next incoming character is to be "put" in the input buffer. This byte must then be converted into a word value and added to the DT$Buffer$Base address to get the absolute memory location.

DTGetOffset This byte contains the relative offset indicating where the next character is to be "got" in the input buffer.

DT$Buffer$Length$Mask This byte contains the length of the buffer minus one. The length of the buffer must always be a binary number (8, 16, 32, 64...). Therefore, one less than the length forms a mask value. Both the get and put offsets, after being incremented, are masked with this value. When the offset reaches the end of the buffer, this masking operation will "automatically" reset the offset to zero.

DT$Character$Count This is a count of the total number of characters in the buffer. It is incremented by the interrupt service routine each time a character is placed in the buffer, and decremented by the CONIN routine each time it gets a character from the buffer.

CONST uses this value to determine whether any characters are available for input.

DT$Stop$Input$Count When the interrupt service routines detect that the DT$Character$Count is equal to this value (normally buffer length minus five), the drivers will invoke the selected input protocol, lowering RTS or sending XOFF, to shut off the incoming data stream.

DT$Resume$Input$Count When the CONIN routine detects that the DT$Character$-Count has become equal to this value, the drivers will again invoke the selected input protocol, either raising RTS or sending XON to resume receiving input data.

DT$Control$Count This is a count of the number of control characters in the input buffer. CARRIAGE RETURN, LINE FEED, and TAB characters are not included in this count.

It is incremented by the interrupt service routine and decremented by CONIN. CONST uses the count when the DT$Fake$Typeahead mode is active; it will only indicate that characters are waiting in the input buffer if the control count is nonzero.

DT$Function$Delay This is the number of clock ticks that should be allowed to elapse after the first character of an incoming escape sequence has been detected. It allows time for the remaining characters in the escape sequence to arrive, assuming that these are being emitted by a terminal at maximum baud rate. Normally, this will correspond to a delay of approximately 90 milliseconds.

DT$Initialize$Stream This is the address of the first byte of a string. This string has the following format:

```
DB  ppH        Port number
DB  nnH        Number of bytes to be output
DB  vvH,vvH... Initialization bytes to be output to the specified port number
```

This sequence can be repeated as many times as is necessary, with a "port" number of 00H acting as a terminator.

Disk Input/Output

The example drivers show three main disk I/O enhancements:

- Full track buffering
- Using memory as an ultra-fast disk
- Improved error handling.

Full Track Buffering

The 5 1/4″ diskettes used in the example system are double-sided. Each side has a separate read/write head in the disk drive. The disk controller is fast enough that, if so commanded, it can read in a complete track's worth of data from one side of the diskette in a single revolution of the diskette.

The drivers have been modified to do just this. The main disk buffer has been dramatically enlarged to accommodate nine 512-byte sectors.

In the earlier standard BIOS, CP/M was configured for tracks of 18 512-byte sectors. The data from each head on a given track was laid "end-to-end" to create the illusion of a single surface with twice as much data on it. For track buffering, performance would be reduced if each read required two revolutions of the diskette, and so in this BIOS the tables and the low-level driver logic have been changed. Each surface is separated, with even numbered tracks on head 0, odd on head 1.

The track number given to the low-level drivers serves two purposes. The least significant bit identifies the head number. When the track number is shifted one bit right, the result is the *physical* track number to which the head assembly must be positioned.

The deblocking algorithm has also been modified by deleting references to sectors. The code is now concerned only with whether the correct disk and track are in the buffer. If this is true, the correct sector must, by definition, be in the buffer.

The deblocking code no longer takes any note when the BDOS indicates that it is writing to an unallocated allocation block—knowledge it used to bypass a sector preread in the standard BIOS. The track size in this enhanced BIOS is much larger than an allocation block, and so the question is meaningless; the whole track must be preread to write just a single sector.

This enhancement really excels when the BDOS is doing directory operations, which always involve a series of sequential reads. The entire directory can be brought into memory, updated, and written back in just two disk revolutions.

One point to watch out for is what is known as "deferred writes." Imagine a program instructed to write on a sector on track 20. The drivers will read in track 20, copy the contents of the designated sector into the track buffer, and return to the program *without* actually writing the data to the disk. The program could "write" to all of the sectors on this track without any actual disk writes. During all this time, this data would exist only in memory and not on the disk drive, so if a power failure occurred, several thousand bytes of data would be lost. Writing to the directory is an exception. The drivers always physically write to the disk when the BDOS indicates that it is writing to a directory sector.

In reality, the increased risk is small. Most programs are constantly reading and writing files, so that the track buffer will be written out frequently in order to read in another track. When programs end, they close output files. This in turn triggers directory writes that force data tracks onto the disk.

If high security is a requirement for your computer, you could extend the watchdog routine to include another separate timer. You could preset this timer for, say, a ten-second delay each time you write into the track buffer but do not write the buffer to the disk. When the count expires, it would set a flag that could be tested by all of the BIOS entry points. If set, they would initiate a write of the track buffer to the disk.

Using Memory as an Ultra-Fast Disk

As you can see from the preceding section, increased performance tends to go hand in hand with increased memory requirements. This is certainly true with a "memory disk," commonly called a RAM-disk or M-disk. In fact, to have an M-disk with reasonable storage capacity, your computer must have at least 128K bytes of additional memory.

Since the 8080 or Z80 can only address 64K of memory at one time, to get access to any of this additional memory, some part of your computer's "normal" memory must be removed from the 64K address space and the additional memory must be switched in. This is known as bank-switched memory.

Figure 8-7 shows the memory organization that is supported by the example M-disk drivers.

You can see that the system has a total of 256K bytes of RAM, organized with the top 16K, from 64K down to 48K, being "common"—that is, switched into the address space all the time. The lower 48K can be selected from five banks, numbered 0 to 4. Bank 0 is switched in for normal CP/M operations.

The M-disk parameter blocks describe a disk with eight "tracks," numbered 0 to 7. The least significant bit of the track number determines whether the base address of the track will be 0000H or 6000H. Shifting the track number right one bit gives the bank number. Each track consists of 192 sectors. To get the relative address of a sector within its "track," shift the sector number eight bits left, thus multiplying it by 128.

The M-disk is referenced by logical disk M:. A few special-case instructions are required to return the special M-disk parameter header in SELDSK.

One problem, fortunately easily solved, is that the user's DMA address coexists in the address space with the M-disk image itself. There is no direct way to move data between bank 0 and any other bank. The M-disk uses an intermediary buffer in common memory (above 48K), moving data into this, switching banks, and then moving the data down again. Figure 8-8 shows an example of this sequence, as used when reading from the M-disk.

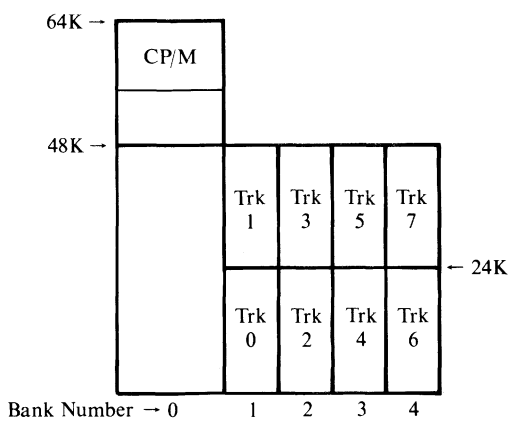


Figure 8-7. Memory organization for M-disk

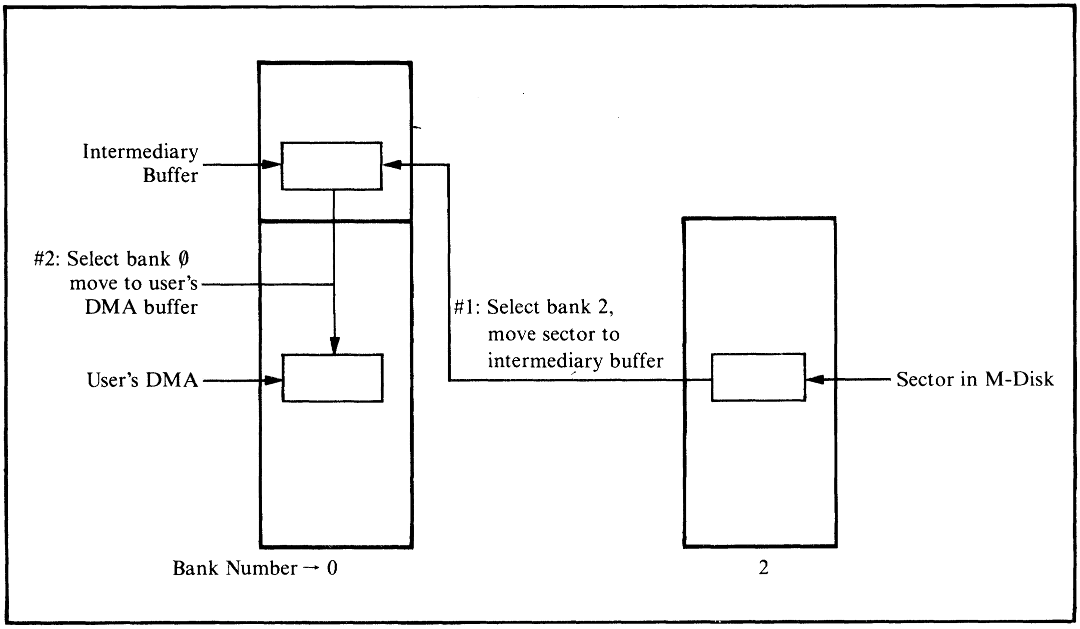


Figure 8-8. Reading a sector from the M-disk image

During cold boot initialization, the M-disk driver checks the very first directory entry (in bank 1) to see if it matches a dummy entry for a file called "M$Disk." If this entry is present, the M-disk is assumed to contain valid information. If the entry is absent, the initialization code makes this special directory entry and fills the remainder of the directory with 0E5H, making it appear empty. The dummy entry makes it appear that the "M$Disk" file is in user 15, marked System status and Read-Only—all of which are designed to prevent its accidental erasure.

Custom Patches to CP/M

Two features shown in the enhanced BIOS, one in the CCP and one in the BDOS, require changes to CP/M itself. These features are implemented by modifying the CCP and BDOS to transfer control to the BIOS at specific points, execute a few instructions in the BIOS, and then return to CP/M. The patches could be made by modifying the MOVCPM program to install the changes permanently. The changed version of MOVCPM, however, *must* be used with a specific version of the BIOS. Therefore, patching CP/M "on the fly" ensures that there will be no mismatch between the BIOS and the rest of CP/M.

Both of these patches were produced with the assistance of Digital Research.

User 0 Files Made Public

The first change permits files created in user area 0 to be accessible from all other user numbers. This feature comes into its own only with hard disk systems. On a hard disk, user numbers can partition the disk, but the frequently used utilities must then be duplicated in each user area. Allowing files in user area 0 to be public means that these files will be accessible from all the other user numbers. Hence the files need not be copied into each user area.

The public files feature alters the way that the BDOS performs the Search Next function, allowing access to files declared in user area 0 even when the current user number is not 0. However, the feature is a double-edged sword—user 0 files can be accidentally erased or damaged as well as accessed. Therefore, user 0 files should be declared as System status and Read-Only to protect them. As an additional precaution, public files can be turned off by a control flag in the long term configuration block. This flag is set to an initial state that disables public files.

Modified User Prompt

This modification makes the CCP display the current user number as well as the default disk. For example,

```
3B>
```

indicates that you are currently in user number 3, with disk B: as the default. In addition, if you have enabled public files, the prompt is preceded by the letter "P" to serve as a reminder:

```
P3B>
```

An Enhanced BIOS

The remainder of this chapter consists of the assembly language source code for the enhanced BIOS described here. It is rather a daunting listing, but will be well worth your study. The copious commentary has been written to make this study easier, and emphasis has been placed on explaining *why* as well as *what* things are done.

As with the standard BIOS, each line is numbered so that you can use the functional index in Figure 8-9 to find areas of interest in the listing. Note that the line numbers are not contiguous. They jump several hundred at the start of each major section or subroutine. This facilitates minor changes in the listing without revision of the functional index. The full listing is given in Figure 8-10.

Start Line	Functional Component or Routine
00001	Introductory Comments and Equates
00200	BIOS Jump Table with Additional Private Entries
00400	Long Term Configuration Block
00800	Interrupt Vector
00900	Device Port Numbers and Other Equates
01100	Display$Message Subroutine
01200	Enter$CPM Setup
01300	Device Table Equates
01500	Device Table Declarations
01700	General Device Initialization
01800	Specific Device Initialization
02000	Output Byte Stream
02100	CONST Routine
02200	CONIN Routine with Function Key Processing
02500	Console Output
02700	CONOUT Routine with Escape Sequence Processing
02900	AUXIST—Auxiliary Input Status Routine
03000	AUXOST—Auxiliary Output Status Routine
03100	AUXIN—Auxiliary Input Routine
03200	AUXOUT—Auxiliary Output Routine
03300	LISTST—List Status Routine
03400	LIST—List Output Routine
03500	Request User Choice—Request Action After Error
03600	Output Error Message
03656	Get Composite Status from Selected Output Devices
03800	Multiple Output of Byte to All Output Devices
04000	Check Output Device Logically (Protocol) Ready
04200	Process ETX/ACK Protocol
04400	Select Device Table from I/O Redirection Bit Map
04600	Get Input Character from Input Buffer
04800	Introductory Comments for Interrupt-Driven Drivers
04900	Character Interrupt Service Routine
05000	Service Device—Puts Character into Input Buffer
05300	Get Address of Character in Input Buffer
05400	Check if Control Character (not CR, LF, TAB)
05500	Output Data Byte
05700	Input Status Routine
05900	Set Watchdog Timer Routine
06000	Real Time Clock Interrupt Service Routine
06200	Shift HL Right One Bit Routine
06300	Introductory Comments for High-Level Disk Drivers
06400	Disk Parameter Headers
06600	Disk Parameter Blocks
06800	SELDSK—Select Disk Routine
07000	SETTRK—Set Track Routine
07100	SETSEC—Set Sector Routine

Figure 8-9. Functional index for listing in Figure 8-10

FIGURE 8-9. Functional index for listing in Figure 8-10 (continued)

```
               00001  ;        This is a skeletal example of an enhanced BIOS.
               00010  ;        It includes fragments of the standard BIOS
               00011  ;        shown as Figure 6-4 in outline, so as to
               00012  ;        avoid cluttering up the enhancements with the
               00013  ;        supporting substructure. Many of the original
               00014  ;        comment blocks have been abbreviated or deleted
               00015  ;        entirely.
               00016  ;
               00017  ;< --  NOTE:   The line numbers at the left are included
               00018  ;               to allow reference to the code from the text.
               00019  ;               There are deliberate discontinuities in the
               00020  ;               numbers to allow space for expansion.
               00021  ;
3030 =         00022  VERSION       EQU    '00'    ;Equates used in the sign-on message
3230 =         00023  MONTH         EQU    '02'
3632 =         00024  DAY           EQU    '26'
3338 =         00025  YEAR          EQU    '83'
               00026  ;
               00027  ;*******************************************************************
               00028  ;*                                                                 *
               00029  ;*      This BIOS is for a computer system with the following      *
               00030  ;*      hardware configuration :                                   *
               00031  ;*                                                                 *
               00032  ;*             -- 8080 CPU                                         *
               00033  ;*             -- 64K bytes of RAM                                 *
               00034  ;*             -- 3 serial I/O ports (using signetics 2651) for:   *
               00035  ;*                console, communications and list                 *
               00036  ;*             -- Two 5 1/4" mini floppy, double-sided, double-    *
               00037  ;*                density drives. These drives use 512-byte sectors. *
               00038  ;*                These are used as logical disks A: and B:.       *
               00039  ;*                Full track buffering is supported.               *
```

Figure 8-10. Enhanced BIOS listing

```
                ;*              -- Two 8" standard diskette drives (128-byte sectors)   *
                ;*                 These are used as logical disks C: and D:.           *
                ;*              -- A memory-based disk (M-disk) is supported.           *
                ;*                                                                      *
                ;*                 Two intelligent disk controllers are used, one for   *
                ;*                  each diskette type. These controllers access memory *
                ;*                  directly, both to read the details of the           *
                ;*                  operations they are to perform and also to read     *
                ;*                  and write data from and to the diskettes.           *
                ;*                                                                      *
                ;*                                                                      *
                ;************************************************************************


                ;       Equates for characters in the ASCII character set
                ;
0011 =          XON     EQU     11H     ;Reenables transmission of data
0013 =          XOFF    EQU     13H     ;Disables transmission of data
0003 =          ETX     EQU     03H     ;End of transmission
0006 =          ACK     EQU     06H     ;Acknowledge
000D =          CR      'EQU    0DH     ;Carriage return
000A =          LF      EQU     0AH     ;Line feed
0009 =          TAB     EQU     09H     ;Horizontal tab
0007 =          BELL    EQU     07H     ;Sound terminal's bell
                ;
                ;
                ;       Equates for defining memory size and the base address and
                ;       length of the system components
                ;
0040 =          Memory$Size     EQU     64      ;Number of Kbytes of RAM
                ;
                ;       The BIOS length must be determined by inspection.
                ;       Comment out the ORG BIOS$Entry line below by changing the first
                ;       character to a semicolon (this will make the assembler start
                ;       the BIOS at location 0). Then assemble the BIOS and round up to
                ;       the nearest 100H the address displayed on the console at the end
                ;       of the assembly.
                ;
2500 =          BIOS$Length     EQU     2500H   ;<-- Revised to an approximate value
                                                ;     to reflect enhancements
                ;
0800 =          CCP$Length      EQU     0800H   ;Constant
0E00 =          BDOS$Length     EQU     0E00H   ;Constant
                ;
000F =          Overall$Length  EQU     (CCP$Length + BDOS$Length + BIOS$Length + 1023) / 1024
                ;
C400 =          CCP$Entry       EQU     (Memory$Size - Overall$Length) * 1024
CC06 =          BDOS$Entry      EQU     CCP$Entry + CCP$Length + 6
DA00 =          BIOS$Entry      EQU     CCP$Entry + CCP$Length + BDOS$Length
                ;
0005 =          BDOS            EQU     0005H   ;BDOS entry point (used for making
                                                ;  system reset requests)
                ;
                ;#
                ;       ORG     BIOS$Entry      ;Assemble code at BIOS address
                ;
                ;       BIOS jump vector
                ;
0000 C31311             JMP     BOOT    ;Cold boot -- entered from CP/M bootstrap loader
                Warm$Boot$Entry:        ;  Labelled so that the initialization code can
                                        ;  put the warm boot entry address in location
                                        ;  0001H and 0002H of the base page
0003 C3750E             JMP     WBOOT   ;Warm boot -- entered by jumping to location 0000H
                                        ;  Reloads the CCP, which could have been
                                        ;  overwritten by previous program in transient
                                        ;  program area
0006 C32D03             JMP     CONST   ;Console status -- returns A = 0FFH if there is a
                                        ;  console keyboard character waiting
0009 C33A03             JMP     CONIN   ;Console input -- returns the next console keyboard
                                        ;  character in A
000C C3D703             JMP     CONOUT  ;Console output -- outputs the character in C to
                                        ;  the console device
000F C3F504             JMP     LIST    ;List output -- outputs the character in C to the
                                        ;  list device
0012 C3CE04             JMP     AUXOUT  ;Auxiliary output -- outputs the character in C to the
                                        ;  logical auxiliary device
```

Figure 8-10. (Continued)

```
0015 C3A104             JMP     AUXIN   ;Auxiliary input -- returns the next input character from
                                        ;  the logical auxiliary device in A
0018 C3160A             JMP     HOME    ;Homes the currently selected disk to track 0
001B C36309             JMP     SELDSK  ;Selects the disk drive specified in register C and
                                        ;  returns the address of the disk parameter header
001E C39B09             JMP     SETTRK  ;Sets the track for the next read or write operation
                                        ;  from the BC register pair
0021 C3A109             JMP     SETSEC  ;Sets the sector for the next read or write operation
                                        ;  from the A register
0024 C3A809             JMP     SETDMA  ;Sets the direct memory address (disk read/write)
                                        ;  address for the next read or write operation
                                        ;  from the DE register pair
0027 C3370A             JMP     READ    ;Reads the previously specified track and sector from
                                        ;  the selected disk into the DMA address
002A C34B0A             JMP     WRITE   ;Writes the previously specified track and sector onto
                                        ;  the selected disk from the DMA address
002D C3D704             JMP     LISTST  ;Returns A = 0FFH if the list device(s) are
                                        ;  logically ready to accept another output byte
0030 C3100A             JMP     SECTRAN ;Translates a logical sector into a physical one
                ;
                ;       Additional "private" BIOS entry points
                ;
0033 C38F04             JMP     AUXIST  ;Returns A = 0FFH if there is input data for
                                        ;  the logical auxiliary device
0036 C39B04             JMP     AUXOST  ;Returns A = 0FFH if the auxiliary device(s) are
                                        ;  logically ready to accept another output byte
0039 C3FA02             JMP     Specific$CIO$Initialization
                                        ;Initializes character device whose device
                                        ;  number is in register A on entry
003C C36D08             JMP     Set$Watchdog
                                        ;Sets up watchdog timer to CALL address specified
                                        ;  in HL, after BC clock ticks have elapsed
003F C33C0F             JMP     CB$Get$Address
                                        ;Configuration block get address
                                        ;  Returns address in HL of data element whose
                                        ;  code number is specified in C
                ;
                ;#
                ;       Long term configuration block
                ;
                Long$Term$CB:
                ;
                ;
                ;       Public files (files in user 0 accessible from all
                ;       other user numbers) enabled when this flag is set
                ;       nonzero.
                ;
0042 00         CB$Public$Files:        DB      0       ;Default is OFF
                ;
                ;
                ;       The forced input pointer is initialized to point to the
                ;       following string of characters. These are injected into
                ;       the console input stream on system start-up.
                ;
0043 5355424D49 CB$Startup:             DB      'SUBMIT STARTUP',LF,0,0,0,0,0,0
                ;
                ;       Logical to physical device redirection
                ;
                ;               Each logical device has a 16-bit word associated
                ;               with it. Each bit in the word is assigned to a
                ;               specific physical device. For input, only one bit
                ;               can be set -- input will be read from the
                ;               corresponding physical device. Output can be
                ;               directed to several devices, so more than one
                ;               bit can be set.
                ;
                ;               The following equates are used to indicate
                ;               specific physical devices.
                ;
                ;                       1111 11             )
                ;                       5432 1098 7654 3210 )<- Device number
0001 =          Device$0        EQU     0000$0000$0000$0001B
0002 =          Device$1        EQU     0000$0000$0000$0010B
0004 =          Device$2        EQU     0000$0000$0000$0100B
                ;
                ;               The following words are tested by the logical
                ;               device drivers to transfer control to
```

Figure 8-10. (Continued)

```
                ;               the appropriate physical device drivers
                ;
0058 0100       CB$Console$Input:       DW      Device$0
005A 0100       CB$Console$Output:      DW      Device$0
                ;
005C 0200       CB$Auxiliary$Input:     DW      Device$1
005E 0200       CB$Auxiliary$Output:    DW      Device$1
                ;
0060 0400       CB$List$Input:          DW      Device$2
0062 0400       CB$List$Output:         DW      Device$2
                ;
                ;               The table below relates specific bits in the
                ;               redirection words above to specific device
                ;               tables used by the physical drivers
                ;
                CB$Device$Table$Addresses:
0064 8E02               DW      DT$0
0066 AE02               DW      DT$1
0068 CE02               DW      DT$2
006A 000000000          DW      0,0,0,0,0,0,0,0,0,0,0,0,0       ;Unassigned
                ;
                ;
                ;       Device initialization byte streams
                ;
                ;       These initialization streams are output during the device
                ;       initialization phase, or on request whenever the baud rate
                ;       needs to be changed. They are defined in the long term
                ;       configuration block so as to "freeze" their contents from one
                ;       system startup until the next.
                ;
                ;       The address of each stream is contained in each device table.
                ;
                ;       The stream format is:
                ;
                ;               DB      xx              ;Port number (00H terminates)
                ;               DB      nn              ;Number of bytes to output to port
                ;               DB      vv,vv,vv..      ;Values to be output
                ;
                D0$Initialize$Stream:           ;Example data for an 8251A chip
0084 ED                 DB      0EDH            ;Port number for 8251A
0085 06                 DB      6               ;Number of bytes
0086 000000             DB      0,0,0           ;Dummy bytes to get chip ready
0089 42                 DB      0100$0010B      ;Reset and raise DTR
008A 6E                 DB      01$10$11$10B    ;1 stop, no parity, 8 bits/char,
                                                ;  divide down of 16
008B 25                 DB      0010$0101B      ;RTS high, enable Tx/Rx
                                                ;Example data for an 8253 chip
008C DF                 DB      0DFH            ;Port number for 8253 mode
008D 01                 DB      1               ;Number of bytes to output
008E 76                 DB      01$11$011$0B    ;Select:
                                                ;       Counter 1
                                                ;       Load LS byte first
                                                ;       Mode 3, binary count
008F DE                 DB      0DEH            ;Port number for counter
0090 02                 DB      2               ;Number of bytes to output
                D0$Baud$Rate$Constant:          ;Label used by utilities
0091 0700               DW      0007H           ;9600 Baud (based on 16x divider)
0093 00                 DB      0               ;Port number of 00 terminates stream

                D1$Initialize$Stream:           ;Example data for an 8251A chip
0094 DD                 DB      0DDH            ;Port number for 8251A
0095 06                 DB      6               ;Number of bytes
0096 000000             DB      0,0,0           ;Dummy bytes to get chip ready
0099 42                 DB      0100$0010B      ;Reset and raise DTR
009A 6E                 DB      01$10$11$10B    ;1 stop, no parity, 8 bits/char,
                                                ;  divide down of 16
009B 25                 DB      0010$0101B      ;RTS high, enable Tx/Rx

                                                ;Example data for an 8253 chip
009C DF                 DB      0DFH            ;Port number for 8253 mode
009D 01                 DB      1               ;Number of bytes to output
009E B6                 DB      10$11$011$0B    ;Select:
                                                ;       Counter 2
                                                ;       Load LS byte first
                                                ;       Mode 3, binary count
009F DE                 DB      0DEH            ;Port number for counter
00A0 02                 DB      2               ;Number of bytes to output
```

Figure 8-10. (Continued)

```
                D1$Baud$Rate$Constant:
00A1 3800               DW      0038H                   ;1200 baud (based on 16x divider)
00A3 00                 DB      0                       ;Port number of 00 terminates stream

                D2$Initialize$Stream:           ;Example data for an 8251A chip
00A4 DD                 DB      0DDH                    ;Port number for 8251A
00A5 06                 DB      6                       ;Number of bytes
00A6 000000             DB      0,0,0                   ;Dummy bytes to get chip ready
00A9 42                 DB      0100$0010B              ;Reset and raise DTR
00AA 6E                 DB      01$10$11$10B            ;1 stop, no parity, 8 bits/char,
                                                        ;  divide down of 16
00AB 25                 DB      0010$0101B              ;RTS high, enable Tx/Rx

                                                ;Example data for an 8253 chip
00AC DF                 DB      0DFH                    ;Port number for 8253 mode
00AD 01                 DB      1                       ;Number of bytes to output
00AE F6                 DB      11$11$011$0B            ;Select:
                                                        ;       Counter 3
                                                        ;       Load LS byte first
                                                        ;       Mode 3, binary count
00AF DE                 DB      0DEH                    ;Port number for counter
00B0 02                 DB      2                       ;Number of bytes to output
                D2$Baud$Rate$Constant:
00B1 3800               DW      0038H                   ;1200 baud (based on 16x divider)
00B3 00                 DB      0                       ;Port number of 00 terminates stream

                ;
                ;       This following table is used to determine the maximum
                ;       value for each character position in the ASCII time
                ;       value above (except the ":"). Note -- this table is
                ;       in the long term configuration block so that the clock
                ;       can be set "permanently" to either 12 or 24 hour format.
                ;
                ;       NOTE: The table is processed backwards -- to correspond
                ;       with the ASCII time.
                ;       Each character represents the value for the corresponding
                ;       character in the ASCII time at which a carry-and-reset-to-zero
                ;       should occur.
                ;
00B4 00                 DB      0               ;"Terminator"
                CB$12$24$Clock:
00B5 3334               DB      '34'            ;Change to '23' for a 12-hour clock
00B7 FF                 DB      0FFH            ;"Skip" character
00B8 363A               DB      '6:'            ;Maximum minutes are 59
00BA FF                 DB      0FFH            ;"Skip" character
00BB 363A               DB      '6:'            ;Maximum seconds are 59
                Update$Time$End:                ;Used when updating the time
                ;
                ;
                ;       Variables for the real time clock and watchdog
                ;       timer
                ;
00BD 3C         RTC$Ticks$per$Second    DB      60      ;Number of real time clock
                                                        ;  ticks per elapsed second
00BE 3C         RTC$Tick$Count          DB      60      ;Residual count before next
                                                        ;  second will elapse
00BF 0000       RTC$Watchdog$Count      DW      0       ;Watchdog timer tick count
                                                        ;(0 = no watchdog timer set)
00C1 0000       RTC$Watchdog$Address    DW      0       ;Address to which control
                                                        ;  will be transferred if the
                                                        ;  watchdog count hits 0


                ;
                ;       Function key table
                ;
                ;       This table consists of a series of entries, each one having the
                ;       following structure:
                ;
                ;               DB      Second character of sequence emitted by
                ;                       terminal's function key
                ;       (       DB      Third character of sequence -- NOTE: this       )
                ;       (               field will not be present if the source code    )
                ;       (               has been configured to accept only two characters )
                ;       (               in function key sequences.                      )
                ;       (               NOTE: Adjust the equates for:                   )
                ;       (                       Function$Key$Length                     )
                ;       (                       Three$Character$Function                )
```

Figure 8-10. (Continued)

```
                ;
                ;               DB      A character string to be forced into the console
                ;                       input stream when the corresponding function key
                ;                       is pressed. The last byte of this string must be
                ;                       00H to terminate the forced input.
                ;
001B =          Function$Key$Lead       EQU     1BH     ;Signals function key sequence
0003 =          Function$Key$Length     EQU     3       ;Number of characters in function
                                                        ;  key input sequence (NOTE: this
                                                        ;  can only be 3 or 2 characters).

                ;
                                                        ;The logic associated with function
                                                        ;  key recognition is made easier with
                                                        ;  the following equate
0001 =          Three$Character$Function        EQU     Function$Key$Length - 2
                                                ;Three$Character$Function will be TRUE if the
                                                ;  function keys emit a three character
                                                ;  sequence, FALSE if they emit a two character
                                                ;  sequence.

                ;       Each entry in the table must be the same length, as defined by:
                ;
0013 =          CB$Function$Key$Entry$Size      EQU     16 + 1 + Function$Key$Length - 1
                ;                                       ^    ^                         ^
                ;                                       |    |                         |
                ;               Maximum length of substitute |         Lead character is not
                ;               string                       |         in table entry
                ;                                            For the terminating 00H
                ;
                ;       The last entry in the table is marked by a 00-byte.
                ;
                ;       The example values shown below are for a VT-100 terminal.
                ;
                CB$Function$Key$Table:
                ;                          123456789.1234  5  6 7 <- Use to check length
00C3 4F5046756E         DB      'O','P','Function Key 1',LF,0,0
00D6 4F5146756E         DB      'O','Q','Function Key 2',LF,0,0
00E9 4F5246756E         DB      'O','R','Function Key 3',LF,0,0
00FC 4F5346756E         DB      'O','S','Function Key 4',LF,0,0
                ;
                ;                          123456789.1
010F 5B41557020         DB      '[','A','Up Arrow',LF,0,0,0,0,0,0,0,0
0122 5B42446F77         DB      '[','B','Down Arrow',LF,0,0,0,0,0,0
0135 5B43526967         DB      '[','C','Right Arrow',LF,0,0,0,0,0
0148 5B444C6566         DB      '[','D','Left Arrow',LF,0,0,0,0,0,0

015B 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0   ;Spare entries
016E 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
0181 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
0194 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
01A7 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
01BA 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
01CD 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
01E0 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
01F3 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
0206 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0

0219 FFFF               DB      0FFH,0FFH       ;Terminator for utility that preprograms
                                                ;  function key sequence
                ;
                ;
                ;       Console output escape sequence control table
                ;
                ;       This table is referenced after a Function$Key$Lead character
                ;       has been detected in the CONOUT routine. The next character
                ;       to be output to the console is compared to the first byte
                ;       in each 3-byte table entry. If a match is found, then control
                ;       is transferred to the address following the byte that matched.
                ;
                CONOUT$Escape$Table:
021B 74                 DB      't'             ;Read current time
021C 4804               DW      CONOUT$Time
021E 64                 DB      'd'             ;Read current date
021F 4104               DW      CONOUT$Date
0221 75                 DB      'u'             ;Set current time
0222 5D04               DW      CONOUT$Set$Time
```

Figure 8-10. (Continued)

```
0224 65                    DB      'e'             ;Set current date
0225 4E04                  DW      CONOUT$Set$Date

0227 00                    DB      0               ;Terminator
                   ;
                   Long$Term$CB$End:
                   ;
                   ;#
                   ;
                   ;       Interrupt vector
                   ;
                   ;       Control is transferred here by the programmable interrupt
                   ;       controller -- an Intel 8259A.
                   ;
                   ;       NOTE: The interrupt controller chip requires that the
                   ;               interrupt vector table start on a paragraph
                   ;               boundary. This is achieved by the following ORG line
0240                       ORG     ($ AND OFFEOH) + 20H
                   Interrupt$Vector:
                                           ;Interrupt number
0240 C37808                JMP     RTC$Interrupt           ;0 -- clock
0243 00                    DB      0                       ;Skip a byte
0244 C3E806                JMP     Character$Interrupt     ;1 -- character I/O
0247 00                    DB      0
0248 C3D80E                JMP     Ghost$Interrupt         ;2 -- not used
024B 00                    DB      0
024C C3D80E                JMP     Ghost$Interrupt         ;3 -- not used
024F 00                    DB      0
0250 C3D80E                JMP     Ghost$Interrupt         ;4 -- not used
0253 00                    DB      0
0254 C3D80E                JMP     Ghost$Interrupt         ;5 -- not used
0257 00                    DB      0
0258 C3D80E                JMP     Ghost$Interrupt         ;6 -- not used
025B 00                    DB      0
025C C3D80E                JMP     Ghost$Interrupt         ;7 -- not used
                   ;
                   ;#

                   ;       Device port numbers and other equates
                   ;
0080 =             CIO$Base$Port   EQU     80H             ;Base port number

0080 =             D0$Base$Port    EQU     CIO$Base$Port           ;Device 0
0080 =             D0$Data$Port    EQU     D0$Base$Port
0081 =             D0$Status$Port  EQU     D0$Base$Port + 1
0082 =             D0$Mode$Port    EQU     D0$Base$Port + 2
0083 =             D0$Command$Port EQU     D0$Base$Port + 3
                   ;

0084 =             D1$Base$Port    EQU     CIO$Base$Port + 4       ;Device 1
0084 =             D1$Data$Port    EQU     D1$Base$Port
0085 =             D1$Status$Port  EQU     D1$Base$Port + 1
0086 =             D1$Mode$Port    EQU     D1$Base$Port + 2
0087 =             D1$Command$Port EQU     D1$Base$Port + 3

0088 =             D2$Base$Port    EQU     CIO$Base$Port + 8       ;Device 2
0088 =             D2$Data$Port    EQU     D2$Base$Port
0089 =             D2$Status$Port  EQU     D2$Base$Port + 1
008A =             D2$Mode$Port    EQU     D2$Base$Port + 2
008B =             D2$Command$Port EQU     D2$Base$Port + 3

004E =             D$Mode$Value$1  EQU     01$00$11$10B
                                                   ;1 stop bit, no parity
                                                   ;8 bits, Async. 16x rate
003C =             D$Mode$Value$2  EQU     00$11$1100B
                                                   ;Tx/Rx on internal clock
                                                   ;9600 baud
0027 =             D$Command$Value EQU     00$100111B
                                                   ;Normal mode
                                                   ;Enable Tx/Rx
                                                   ;RTS and DTR active
0038 =             D$Error         EQU     0011$1000B
0037 =             D$Error$Reset   EQU     00$110111B
                                                   ;Same as command value plus error reset
0001 =             D$Output$Ready  EQU     0000$0001B
0002 =             D$Input$Ready   EQU     0000$0010B
0080 =             D$DTR$High      EQU     1000$0000B      ;Note: this is actually the
```

Figure 8-10. (Continued)

```
                                                                ;  data-set-ready pin
                                                                ;  on the chip. It is connected
                                                                ;  to the DTR pin on the cable
0027 =          D$Raise$RTS       EQU     00$1$00111B          ;Raise RTS, Tx/Rx enable
0007 =          D$Drop$RTS        EQU     00$0$00111B          ;Drop RTS, Tx/Rx enable
                ;
                ;
                ;       Interrupt controller ports (Intel 8259A)
                ;
                ;       Note : these equates are placed here so that they
                ;               follow the  definition of the interrupt vector
                ;               and thus avoid 'P' (phase) errors in ASM.
                ;
00D9 =          IC$OCW1$Port      EQU     0D9H    ;Operational control word 1
00D8 =          IC$OCW2$Port      EQU     0D8H    ;Operational control word 2
00D8 =          IC$OCW3$Port      EQU     0D8H    ;Operational control word 3
00D8 =          IC$ICW1$Port      EQU     0D8H    ;Initialization control word 1
00D9 =          IC$ICW2$Port      EQU     0D9H    ;Initialization control word 2
                ;
0020 =          IC$EOI            EQU     20H     ;Nonspecific end of interrupt
                ;
0056 =          IC$ICW1           EQU     (Interrupt$Vector AND 1110$0000B) + 000$10110B
                                                  ;Sets the A7 - A5 bits of the interrupt
                                                  ;  vector address plus:
                                                  ;       Edge triggered
                                                  ;       4-byte interval
                                                  ;       Single 8259 in system
                                                  ;       No ICW4 needed
0002 =          IC$ICW2           EQU     Interrupt$Vector SHR 8
                                                  ;Address bits A15 - A8 of the interrupt
                                                  ;  vector address. Note the interrupt
                                                  ;  vector is the first structure in
                                                  ;  the long term configuration block

00FC =          IC$OCW1           EQU     1111$1100B      ;Interrupt mask
                                                  ;Interrupt 0 (clock) enabled
                                                  ;Interrupt 1 (character input) enabled
                ;
                ;#
                ;
                ;
                Display$Message:          ;Displays the specified message on the console.
                                          ;On entry, HL points to a stream of bytes to be
                                          ;output. A 00H-byte terminates the message.
025F 7E                 MOV     A,M               ;Get next message byte
0260 B7                 ORA     A                 ;Check if terminator
0261 C8                 RZ                        ;Yes, return to caller
0262 4F                 MOV     C,A               ;Prepare for output
0263 E5                 PUSH    H                 ;Save message pointer
0264 CDD703             CALL    CONOUT            ;Go to main console output routine
0267 E1                 POP     H                 ;Recover message pointer
0268 23                 INX     H                 ;Move to next byte of message
0269 C35F02             JMP     Display$Message ;Loop until complete message output
                ;
                ;#
                ;
                Enter$CPM:      ;This routine is entered either from the cold or warm
                                ; boot code. It sets up the JMP instructions in the
                                ; base page, and also sets the high-level disk driver's
                                ; input/output address (the DMA address).
                ;
026C 3EC3               MVI     A,JMP             ;Get machine code for JMP
026E 320000             STA     0000H             ;Set up JMP at location 0000H
0271 320500             STA     0005H             ; and at location 0005H
                ;
0274 210300             LXI     H,Warm$Boot$Entry       ;Get BIOS vector address
0277 220100             SHLD    0001H             ;Put address at location 0001H

027A 2106CC             LXI     H,BDOS$Entry      ;Get BDOS entry point address
027D 220600             SHLD    6                 ;Put address at location 0005H
                ;
0280 018000             LXI     B,80H             ;Set disk I/O address to default
0283 CDA809             CALL    SETDMA            ;Use normal BIOS routine
                ;
0286 FB                 EI                        ;Ensure interrupts are enabled
0287 3A0400             LDA     Default$Disk      ;Handover current default disk to
028A 4F                 MOV     C,A               ; console command processor
```

Figure 8-10. (Continued)

```
028B C300C4             JMP     CCP$Entry       ;Transfer to CCP
                ;
                ;#
                ;
                ;       Device table equates
                ;       The drivers use a device table for each
                ;       physical device they service. The equates that follow
                ;       are used to access the various fields within the
                ;       device table.
                ;
                ;                       Port numbers and status bits
0000 =          DT$Status$Port          EQU     0       ;Device status port number
0001 =          DT$Data$Port            EQU     DT$Status$Port+1
                                                        ;Device data port number
0002 =          DT$Output$Ready         EQU     DT$DataPort+1
                                                        ;Output ready status mask
0003 =          DT$Input$Ready          EQU     DT$Output$Ready+1
                                                        ;Input ready status mask
0004 =          DT$DTR$Ready            EQU     DT$Input$Ready+1
                                                        ;DTR ready to send mask
0005 =          DT$Reset$Int$Port       EQU     DT$DTR$Ready+1
                                                        ;Port number used to reset an
                                                        ;  interrupt
0006 =          DT$Reset$Int$Value      EQU     DT$Reset$Int$Port+1
                                                        ;Value output to reset interrupt
0007 =          DT$Detect$Error$Port    EQU     DT$Reset$Int$Value+1
                                                        ;Port number for detecting error
0008 =          DT$Detect$Error$Value   EQU     DT$Detect$Error$Port+1
                                                        ;Mask for detecting error (parity etc.)
0009 =          DT$Reset$Error$Port     EQU     DT$Detect$Error$Value+1
                                                        ;Output to port to reset error
000A =          DT$Reset$Error$Value    EQU     DT$Reset$Error$Port+1
                                                        ;Value to output to reset error
000B =          DT$RTS$Control$Port     EQU     DT$Reset$Error$Value+1
                                                        ;Control port for lowering RTS
000C =          DT$Drop$RTS$Value       EQU     DT$RTS$Control$Port+1
                                                        ;Value, when output, to drop RTS
000D =          DT$Raise$RTS$Value      EQU     DT$Drop$RTS$Value+1
                                                        ;Value, when output, to raise RTS
                ;
                ;               Device logical status (incl. protocols)
000E =          DT$Status               EQU     DT$Raise$RTS$Value+1
                                                                ;Status bits
0001 =          DT$Output$Suspend       EQU     0000$0001B      ;Output suspended pending
                                                                ;  protocol action
0002 =          DT$Input$Suspend        EQU     0000$0010B      ;Input suspended until
                                                                ;  buffer empties
0004 =          DT$Output$DTR           EQU     0000$0100B      ;Output uses DTR-high-to-send
0008 =          DT$Output$Xon           EQU     0000$1000B      ;Output uses XON/XOFF
0010 =          DT$Output$Etx           EQU     0001$0000B      ;Output uses ETX/ACK
0020 =          DT$Output$Timeout       EQU     0010$0000B      ;Output uses timeout
0040 =          DT$Input$RTS            EQU     0100$0000B      ;Input uses RTS-high-to-receive
0080 =          DT$Input$Xon            EQU     1000$0000B      ;Input uses XON/XOFF
                ;
000F =          DT$Status$2             EQU     DT$Status+1     ;Secondary status byte
0001 =          DT$Fake$Typeahead       EQU     0000$0001B      ;Requests Input$Status to
                                                                ; return "Data Ready" when
                                                                ; control dharacters are in
                                                                ; input buffer
                ;
0010 =          DT$Etx$Count            EQU     DT$Status$2+1
                                                        ;No. of chars. sent in Etx protocol
0012 =          DT$Etx$Message$Length   EQU     DT$Etx$Count+2
                                                        ;Specified message length
                ;
                ;                       Input buffer values
0014 =          DT$Buffer$Base          EQU     DT$Etx$Message$Length+2
                                                        ;Address of Input buffer
0016 =          DT$Put$Offset           EQU     DT$Buffer$Base+2
                                                        ;Offset for putting chars. into buffer
0017 =          DT$Get$Offset           EQU     DT$Put$Offset+1
                                                        ;Offset for getting chars. from buffer
0018 =          DT$Buffer$Length$Mask   EQU     DT$Get$Offset+1
                                                        ;Length of buffer - 1
                                                        ;Note: Buffer length must always be
                                                        ;  a binary number; e.g. 32, 64 or 128
```

Figure 8-10. (Continued)

```
                                                          ;This mask then becomes:
                                                          ;  32 ->  31 (0001$1111B)
                                                          ;  64 ->  63 (0011$1111B)
                                                          ; 128 -> 127 (0111$1111B)
                                                          ;After the get/put offset has been
                                                          ;  incremented, it is ANDed with the mask
                                                          ;  to reset it to zero when the end of
                                                          ;  the buffer has been reached
0019 =          DT$Character$Count      EQU     DT$Buffer$Length$Mask+1
                                                          ;Count of the number of characters
                                                          ;  currently in the buffer
001A =          DT$Stop$Input$Count     EQU     DT$Character$Count+1
                                                          ;Stop input when the count reaches
                                                          ;  this value
001B =          DT$Resume$Input$Count   EQU     DT$Stop$Input$Count+1
                                                          ;Resume input when the count reaches
                                                          ;  this value
001C =          DT$Control$Count        EQU     DT$Resume$Input$Count+1
                                                          ;Count of the number of control
                                                          ;  characters in the buffer
001D =          DT$Function$Delay       EQU     DT$Control$Count+1
                                                          ;Number of clock ticks to delay to
                                                          ;  allow all characters after function
                                                          ;  key lead-in to arrive
001E =          DT$Initialize$Stream    EQU     DT$Function$Delay+1
                                                          ;Address of byte stream necessary to
                                                          ;  initialize this device

                ;#
                ;
                ;       Device tables
                ;
                DT$0:
028E 81                 DB      D0$Status$Port  ;Status port (8251A chip)
028F 80                 DB      D0$Data$Port    ;Data port
0290 01                 DB      D$Output$Ready  ;Output data ready
0291 02                 DB      D$Input$Ready   ;Input data ready
0292 80                 DB      D$DTR$High      ;DTR ready to send
0293 D8                 DB      IC$OCW2$Port    ;Reset interrupt port (00H is an unused port)
0294 20                 DB      IC$EOI          ;Reset interrupt value (nonspecific EOI)
0295 81                 DB      D0$Status$Port  ;Detect error port
0296 38                 DB      D$Error         ;Mask: framing, overrun, parity errors
0297 83                 DB      D0$Command$Port ;Reset error port
0298 37                 DB      D$Error$Reset   ;Reset error: RTS high, reset, Tx/Rx enable
0299 83                 DB      D0$Command$Port ;Drop/raise RTS port
029A 07                 DB      D$Drop$RTS      ;Drop RTS Value (keep Tx & Rx enabled)
029B 27                 DB      D$Raise$RTS     ;Raise RTS value (keep Tx & Rx enabled)
029C C0                 DB      DT$Input$Xon + DT$Input$RTS     ;Protocol and status
029D 00                 DB      0               ;Status #2
029E 0004               DW      1024            ;Etx/Ack message count
02A0 0004               DW      1024            ;Etx/Ack message length
02A2 2422               DW      D0$Buffer       ;Input buffer
02A4 00                 DB      0               ;Put offset into buffer
02A5 00                 DB      0               ;Get offset into buffer
02A6 1F                 DB      D0$Buffer$Length -1 ;Buffer length mask
02A7 00                 DB      0               ;Count of characters in buffer
02A8 1B                 DB      D0$Buffer$Length - 5 ;Stop input when count hits this value
02A9 10                 DB      D0$Buffer$Length / 2 ;Resume input when count hits this value
02AA 00                 DB      0               ;Count of control characters in buffer
02AB 06                 DB      6               ;Number of 16.66ms ticks to allow function
                                                ;  key sequence to arrive (approx. 90ms)
02AC 8400               DW      D0$Initialize$Stream    ;Address of initialization stream
                ;
                DT$1:
02AE 85                 DB      D1$Status$Port  ;Status port (8251A chip)
02AF 84                 DB      D1$Data$Port    ;Data port
02B0 01                 DB      D$Output$Ready  ;Output data ready
02B1 02                 DB      D$Input$Ready   ;Input data ready
02B2 80                 DB      D$DTR$High      ;DTR ready to send
02B3 D8                 DB      IC$OCW2$Port        ;Reset interrupt port (00H is an unused port)
02B4 20                 DB      IC$EOI          ;Reset interrupt value (nonspecific EOI)
02B5 85                 DB      D1$Status$Port  ;Detect error port
02B6 38                 DB      D$Error         ;Mask: framing, overrun, parity errors
02B7 87                 DB      D1$Command$Port ;Reset error port
02B8 37                 DB      D$Error$Reset   ;Reset error: RTS high, reset, Tx/Rx enable
02B9 87                 DB      D1$Command$Port ;Drop/raise RTS port
02BA 07                 DB      D$Drop$RTS      ;Drop RTS value (keep Tx & Rx enabled)
```

Figure 8-10. (Continued)

```
02BB 27                    DB      D$Raise$RTS     ;Raise RTS value (keep Tx & Rx enabled)
02BC C0                    DB      DT$Input$Xon + DT$Input$RTS     ;Protocol and status
02BD 00                    DB      0               ;Status #2
02BE 0004                  DW      1024            ;Etx/Ack message count
02C0 0004                  DW      1024            ;Etx/Ack message length
02C2 4422                  DW      D1$Buffer       ;Input buffer
02C4 00                    DB      0               ;Put offset into buffer
02C5 00                    DB      0               ;Get offset into buffer
02C6 1F                    DB      D1$Buffer$Length -1 ;Buffer length mask
02C7 00                    DB      0               ;Count of characters in buffer
02C8 1B                    DB      D1$Buffer$Length - 5 ;Stop input when count hits this value
02C9 10                    DB      D1$Buffer$Length / 2 ;Resume input when count hits this value
02CA 00                    DB      0               ;Count of control characters in buffer
02CB 06                    DB      6               ;Number of 16.66ms ticks to allow function
                                                   ;  key sequence to arrive (approx. 90ms)
02CC 9400                  DW      D1$Initialize$Stream    ;Address of initialization stream
                   ;
                   ;
                   DT$2:
02CE 89                    DB      D2$Status$Port  ;Status port (8251A chip)
02CF 88                    DB      D2$Data$Port    ;Data port
02D0 01                    DB      D$Output$Ready  ;Output data ready
02D1 02                    DB      D$Input$Ready   ;Input data ready
02D2 80                    DB      D$DTR$High      ;DTR ready to send
02D3 D8                    DB      IC$OCW2$Port    ;Reset interrupt port (00H is an unused port)
02D4 20                    DB      IC$EOI          ;Reset interrupt value (nonspecific EOI)
02D5 89                    DB      D2$Status$Port  ;Detect error port
02D6 38                    DB      D$Error         ;Mask: framing, overrun, parity errors
02D7 8B                    DB      D2$Command$Port ;Reset error port
02D8 37                    DB      D$Error$Reset   ;Reset error: RTS high, reset, Tx/Rx enable
02D9 8B                    DB      D2$Command$Port ;Drop/raise RTS port
02DA 07                    DB      D$Drop$RTS      ;Drop RTS value (keep Tx & Rx enabled)
02DB 27                    DB      D$Raise$RTS     ;Raise RTS value (keep Tx & Rx enabled)
02DC C0                    DB      DT$Input$Xon + DT$Input$RTS     ;Protocol and status
02DD 00                    DB      0               ;Status #2
02DE 0004                  DW      1024            ;Etx/Ack message count
02E0 0004                  DW      1024            ;Etx/Ack message length
02E2 6422                  DW      D2$Buffer       ;Input buffer
02E4 00                    DB      0               ;Put offset into buffer
02E5 00                    DB      0               ;Get offset into buffer
02E6 1F                    DB      D2$Buffer$Length -1 ;Buffer length mask
02E7 00                    DB      0               ;Count of characters in buffer
02E8 1B                    DB      D2$Buffer$Length - 5 ;Stop input when count hits this value
02E9 10                    DB      D2$Buffer$Length / 2 ;Resume input when count hits this value
02EA 00                    DB      0               ;Count of control characters in buffer
02EB 06                    DB      6               ;Number of 16.66ms ticks to allow function
                                                   ;  Key sequence to arrive (approx. 90ms)
02EC A400                  DW      D2$Initialize$Stream    ;Address of initialization stream
                   ;
                   ;#
                   ;       General character I/O device initialization
                   ;
                   ;       This routine will be called from the main CP/M
                   ;       initialization code.
                   ;
                   ;       It makes repeated calls to the specific character I/O
                   ;       device initialization routine.
                   ;
                   General$CIO$Initialization:
02EE AF                    XRA     A               ;Set device number (used to access the
                                                   ;  table of device table addresses in the
                                                   ;  configuration block)
02EF 4F                    MOV     C,A             ;Match to externally CALLable interface
                   GCI$Next$Device:
02F0 CDFA02                CALL    Specific$CIO$Initialization     ;Initialize the device
02F3 3C                    INR     A               ;Move to next device
02F4 FE10                  CPI     16              ;Check if all possible devices (0 - 15)
02F6 C8                    RZ                      ;  have been initialized
02F7 C3F002                JMP     GCI$Next$Device
                   ;
                   ;#
                   ;
                   ;       Specific character I/O initialization
                   ;
                   ;       This routine outputs the specified byte values to the specified
                   ;       ports as controlled by the initialization streams in the
                   ;       configuration block. Each device table contains a pointer to
```

Figure 8-10. (Continued)

```
                ;       these streams. The device table itself is selected according
                ;       to the device NUMBER -- this is an entry parameter for this
                ;       routine.
                ;       This routine will be called either from the general device
                ;       initialization routine above, or directly by a BIOS call from
                ;       a system utility executing in the TPA.
                ;
                ;       Entry parameters
                ;
                ;               C = device number
                ;
                ;       Exit parameters
                ;
                ;               A = Device number (preserved)
                ;
                ;=========================
                Specific$CIO$Initialization:            ;<=== BIOS entry point (private)
                ;=========================
02FA 79                 MOV     A,C             ;Get device number
02FB F5                 PUSH    PSW             ;Preserve device number
02FC 87                 ADD     A               ;Make device number into word pointer
02FD 4F                 MOV     C,A
02FE 0600               MVI     B,0             ;Make into a word
0300 216400             LXI     H,CB$Device$Table$Addresses     ;Get table base
0303 09                 DAD     B               ;HL -> device table address
0304 5E                 MOV     E,M             ;Get LS byte
0305 23                 INX     H
0306 56                 MOV     D,M             ;Get MS byte: DE -> device table

0307 7A                 MOV     A,D             ;Check if device table address = 0
0308 B3                 ORA     E
0309 CA1703             JZ      SCI$Exit        ;Yes, device table nonexistent

030C 211E00             LXI     H,DT$Initialize$Stream
030F 19                 DAD     D               ;HL -> initialization stream address
0310 5E                 MOV     E,M             ;Get LS byte
0311 23                 INX     H
0312 56                 MOV     D,M             ;Get MS byte
0313 EB                 XCHG                    ;HL -> initialization stream itself
0314 CD1903             CALL    Output$Byte$Stream      ;Output byte stream to various
                                                        ;  ports
                ;
                SCI$Exit:
0317 F1                 POP     PSW             ;Recover user's device number in C
0318 C9                 RET
                ;
                ;#
                ;       Output byte stream
                ;
                ;       This routine outputs initialization bytes to port
                ;       numbers. The byte stream has the following format:
                ;
                ;               DB      ppH     Port number
                ;               DB      nn      Number of bytes to output
                ;               DB      vvH,vvH...   Bytes to be output
                ;               :
                ;               :       Repeated
                ;               :
                ;               DB      00H     Port number of 0 terminates
                ;
                ;       Entry parameters
                ;
                ;               HL -> Byte stream
                ;
                Output$Byte$Stream:
                OBS$Loop:
0319 7E                 MOV     A,M             ;Get port number
031A B7                 ORA     A               ;Check if 00H (terminator)
031B C8                 RZ                      ;Exit if at end of stream
031C 322503             STA     OBS$Port        ;Store in port number below
031F 23                 INX     H               ;HL -> count of bytes
0320 4E                 MOV     C,M             ;Get count
0321 23                 INX     H               ;HL -> first initialization byte
                ;
                OBS$Next$Byte:
0322 7E                 MOV     A,M             ;Get next byte
0323 23                 INX     H               ;HL -> next data byte (or port number)
```

Figure 8-10. (Continued)

```
0324 D3                  DB      OUT
                 OBS$Port:
0325 00                  DB      0               ;<- Set up in instruction above
0326 0D                  DCR     C               ;Count down on byte counter
0327 C22203              JNZ     OBS$Next$Byte   ;Output next data byte
032A C31903              JMP     OBS$Loop        ;Go back for next port number
                 ;
                 ;#
                 ;       CONST - Console status
                 ;
                 ;       This routine checks both the forced input pointer and
                 ;       the character count for the appropriate input buffer.
                 ;       The A register is set to indicate whether or not there
                 ;       is data waiting.
                 ;
                 ;       Entry parameters: none.
                 ;
                 ;       Exit parameters
                 ;
                 ;               A = 000H if there is no data waiting
                 ;               A = 0FFH if there is data waiting
                 ;
                 ;=========================
                 CONST:                                  ;<=== BIOS entry point (standard)
                 ;=========================
032D 2A5800              LHLD    CB$Console$Input        ;Get redirection word
0330 116400              LXI     D,CB$Device$Table$Addresses
0333 CD6F06              CALL    Select$Device$Table     ;Get device table address
0336 C34708              JMP     Get$Input$Status        ;Get status from input device
                                                         ;  and return to caller
                 ;#
                 ;
                 ;       CONIN -- console input
                 ;
                 ;       This routine returns the next character for the console input
                 ;       stream. Depending on the circumstances, this can be a character
                 ;       from the console input buffer, or from a previously stored
                 ;       string of characters to be "forced" into the input stream, for
                 ;       the automatic execution of system initialization routines.
                 ;       The "forced input" can come from any previously stored character
                 ;       string in memory. It is used to inject the current time and date
                 ;       or a string associated with a function key into the console
                 ;       stream. On system startup, a string of "SUBMIT STARTUP" is
                 ;       forced into the console input stream to provide a mechanism.
                 ;
                 ;       Normal ("unforced") input comes from whichever physical device
                 ;       is specified in the console input redirection word (see the
                 ;       configuration block).
                 ;
0339 00          CONIN$Delay$Elapsed:    DB      0       ;Flag used during function key
                                                         ;  processing to indicate that
                                                         ;  a predetermined delay has
                                                         ;  elapsed
                 ;
                 ;=========================
                 CONIN:                                  ;<=== BIOS entry point (standard)
                 ;=========================
033A 2A8D0F              LHLD    CB$Forced$Input         ;Get the forced input pointer
033D 7E                  MOV     A,M                     ;Get the next character of input
033E B7                  ORA     A                       ;Check if a null
033F CA4703              JZ      CONIN$No$FI             ;Yes, no forced input
0342 23                  INX     H                       ;Yes, update the pointer
0343 228D0F              SHLD    CB$Forced$Input         ;  and store it back
0346 C9                  RET
                 ;
                 CONIN$No$FI                             ;No forced input
0347 2A5800              LHLD    CB$Console$Input        ;Get redirection word
034A 116400              LXI     D,CB$Device$Table$Addresses
034D CD6F06              CALL    Select$Device$Table     ;Get device table address
0350 CD9106              CALL    Get$Input$Character     ;Get next character from input device

                                                 ;Function key processing
0353 FE1B                CPI     Function$Key$Lead       ;Check if first character of function
                                                         ;  key sequence (normally escape)
0355 C0                  RNZ                             ;Return to BIOS caller if not
0356 F5                  PUSH    PSW                     ;Save lead in character
```

Figure 8-10. (Continued)

```
0357 211D00             LXI     H,DT$Function$Delay     ;Get delay time constant for
                                                        ;  delay while waiting for subsequent
                                                        ;  characters of function key sequence
                                                        ;  to arrive
035A 19                 DAD     D
035B 4E                 MOV     C,M                     ;Get delay value
035C 0600               MVI     B,0                     ;Make into word value
035E AF                 XRA     A                       ;Indicate timer not yet out of time
035F 323903             STA     CONIN$Delay$Elapsed
0362 217B03             LXI     H,CONIN$Set$Delay$Elapsed ;Address to resume at after delay
0365 CD6D08             CALL    Set$Watchdog            ;Sets up delay based on real time
                                                        ;  clock such that control will be
                                                        ;  transferred to specified address
                                                        ;  after time interval has elapsed
                CONIN$Wait$for$Delay:                   ;Wait here until delay has elapsed
0368 3A3903             LDA     CONIN$Delay$Elapsed     ;Check flag set by watchdog routine
036B B7                 ORA     A
036C CA6803             JZ      CONIN$Wait$for$Delay

                CONIN$Check$for$Function:
036F 211900             LXI     H,DT$Character$Count    ;Now check if the remaining characters
                                                        ;  of the sequence have been input
0372 19                 DAD     D
0373 7E                 MOV     A,M                     ;Get count of characters in buffer
0374 FE02               CPI     Function$Key$Length - 1
0376 D28103             JNC     CONIN$Check$Function    ;Enough characters in buffer for
                                                        ;  possible function key sequence
0379 F1                 POP     PSW                     ;Insufficient characters in buffer
                                                        ;  to be a function key, so return
                                                        ;  to caller with lead character
037A C9                 RET

                ;
                ;       The following routine is called by the watchdog routine
                ;       when the specified delay has elapsed.
                ;
                CONIN$Set$Delay$Elapsed:
037B 3EFF               MVI     A,0FFH                  ;Indicate watchdog timer out of time
037D 323903             STA     CONIN$Delay$Elapsed
0380 C9                 RET                             ;Return to watchdog routine
                ;
                ;
                CONIN$Check$Function:
0381 211700             LXI     H,DT$Get$Offset         ;Save the current "get pointer"
0384 19                 DAD     D                       ;  in the buffer
0385 7E                 MOV     A,M                     ;Get the pointer
0386 F5                 PUSH    PSW                     ;Save pointer on the stack

0387 211700             LXI     H,DT$Get$Offset         ;Check the second (and possibly third)
038A CDF007             CALL    Get$Address$in$Buffer   ;  character in the sequence
038D 46                 MOV     B,M                     ;Get the second character

                        IF      Three$Character$Function
038E C5                 PUSH    B                       ;Save for later use
038F 211700             LXI     H,DT$Get$Offset         ;Retrieve the third character
0392 CDF007             CALL    Get$Address$in$Buffer
0395 C1                 POP     B                       ;Recover second character
0396 4E                 MOV     C,M                     ;Now BC = Char 2, Char 3
                        ENDIF

0397 D5                 PUSH    D                       ;Save device table pointer
0398 21B000             LXI     H,CB$Function$Key$Table - CB$Function$Key$Entry$Size
                                                        ;Get pointer to function key table
                                                        ;  in configuration block
039B 111300             LXI     D,CB$Function$Key$Entry$Size    ;Get entry size ready for loop
                CONIN$Next$Function:
039E 19                 DAD     D                       ;Move to next (or first) entry
039F 7E                 MOV     A,M                     ;Get second character of sequence
03A0 B7                 ORA     A                       ;Check if end of function key table
03A1 CAC203             JZ      CONIN$Not$Function      ;Yes -- it is not a function key
03A4 B8                 CMP     B                       ;Compare second characters
03A5 C29E03             JNZ     CONIN$Next$Function     ;No match, so try next entry in table

                        IF      Three$Character$Function
03A8 23                 INX     H                       ;HL -> third character
03A9 7E                 MOV     A,M                     ;Get third character of sequence
03AA 2B                 DCX     H                       ;Simplify logic for 2 & 3 char. seq.
```

Figure 8-10. (Continued)

```
03AB B9                 CMP     C                       ;Compare third characters
03AC C29E03             JNZ     CONIN$Next$Function     ;No match, so try next entry in table
03AF 23                 INX     H                       ;When match found, compensate for
                                                        ;  extra decrement
                        ENDIF

03B0 23                 INX     H                       ;HL -> first character of substitute
                                                        ;  string of characters (00-byte term.)
03B1 228D0F             SHLD    CB$Forced$Input         ;Make the CONIN routine inject the
                                                        ;  substitute string into the input
                                                        ;  stream

                                                        ;Now that a function sequence has been
                                                        ;  identified, the stack must be
                                                        ;  balanced prior to return
03B4 D1                 POP     D                       ;Get the device table pointer
03B5 F1                 POP     PSW                     ;Dump the "get" offset value
03B6 F1                 POP     PSW                     ;Dump the function sequence lead char.

03B7 211900             LXI     H,DT$Character$Count    ;Downdate the character count
03BA 19                 DAD     D                       ;  to reflect the characters removed
                                                        ;  from the buffer
03BB 7E                 MOV     A,M                     ;Get the count
03BC D602               SUI     Function$Key$Length -1  ; (the lead character has already
03BE 77                 MOV     M,A                     ;  been deducted)
03BF C33A03             JMP     CONIN                   ;Return to CONIN processing to get
                                                        ;  the forced input characters
                CONIN$Not$Function:
                                                ;Attempts to recognize a function key sequence
                                                ;  have failed. The "get" offset pointer must be
                                                ;  restored to its previous value so that
                                                ;  the character(s) presumed to be part of
                                                ;  the function sequence are not lost.

03C2 D1                 POP     D                       ;Recover device table pointer
03C3 F1                 POP     PSW                     ;Recover previous "get" offset
03C4 211700             LXI     H,DT$Get$Offset
03C7 19                 DAD     D                       ;HL -> "get" offset in table
03C8 77                 MOV     M,A                     ;Reset "get" offset as it was after
                                                        ;  the lead character was detected
03C9 F1                 POP     PSW                     ;Recover lead character
03CA C9                 RET                             ;Return the lead character to the user
                ;
                ;#
                ;       Console output
                ;
                ;       This routine outputs data characters to the console device(s).
                ;       It also "traps" escape sequences being output to the console,
                ;       triggering specific actions according to the sequences.
                ;       A primitive "state-machine" is used to step through escape
                ;       sequence recognition.
                ;       In addition to outputting the next character to all of the
                ;       devices currently selected in the console output redirection word,
                ;       it checks to see that output to the selected device has not been
                ;       suspended by XON/XOFF protocol, and that DTR is high if
                ;       it should be.
                ;       Once the character has been output, if ETX/ACK protocol is in use,
                ;       and the specified length of message has been output, an Etx
                ;       character is output and the device is flagged as being suspended.
                ;
                ;   '   Entry parameters
                ;
                ;               C = character to be output
                ;
                ;       CONOUT storage variables
                ;
03CB 00         CONOUT$Character:       DB      0       ;Save area for character to be output

03CC DB03       CONOUT$Processor:       DW      CONOUT$Normal
                                                        ;This is the address of the piece of
                                                        ;  code that will process the next
                                                        ;  character. The default case is
                                                        ;  CONOUT$Normal
03CE 0000       CONOUT$String$Pointer:  DW      0       ;This points to a string (normally
                                                        ;  in the configuration block) that
                                                        ;  is being preset by characters from
                                                        ;  the console output stream
```

Figure 8-10. (Continued)

```
03D0 00          CONOUT$String$Length:    DB      0          ;This contains the maximum number of
                                                             ; characters to be preset into a
                                                             ; from the console output stream

                 ;
                 ;       *** WARNING ***
                 ;       The output error message routine shares the code in this
                 ;       subroutine. On entry here, the data byte to be output
      -          ;       will be on the stack, and the DE registers set up correctly.
                 ;
                 ;
                 CONOUT$OEM$Entry:
03D1 32CB03              STA     CONOUT$Character            ;Save data byte
03D4 C3E803              JMP     CONOUT$Entry2               ;HL already has special bit map
                 ;
                 ;====================
                 CONOUT:                              ;<=== BIOS entry point (standard)
                 ;====================
03D7 2ACC03              LHLD    CONOUT$Processor            ;Get address of processor to handle
                                                             ; the next character to be output
                                                             ;(Default is CONOUT$Normal)
03DA E9                  PCHL                                ;Transfer control to the processor
                 ;
                 ;
                 CONOUT$Normal:                              ;Normal processor for console output
03DB 79                  MOV     A,C                         ;Check if possible start of escape
03DC FE1B                CPI     Function$Key$Lead           ; sequence
03DE CA1204              JZ      CONOUT$Escape$Found         ;Perhaps
                 CONOUT$Forced:
03E1 79                  MOV     A,C                         ;Forced output entry point
03E2 32CB03              STA     CONOUT$Character            ;Not escape sequence -- Save data byte

03E5 2A5A00              LHLD    CB$Console$Output           ;Get console redirection word
                 ;
                 CONOUT$Entry2:                       ;<=== output error message entry point
                 ;
03E8 116400              LXI     D,CB$Device$Table$Addresses     ;Addresses of dev. tables
03EB D5                  PUSH    D                               ;Put onto stack ready for loop
03EC E5                  PUSH    H

                 CONOUT$Next$Device:
03ED E1                  POP     H                           ;Recover redirection bit map
03EE D1                  POP     D                           ;Recover device table addresses pointer
03EF CD6F06              CALL    Select$Device$Table         ;Get device table in DE
03F2 B7                  ORA     A                           ;Check if a device has been
                                                             ; selected (i.e. bit map not all zero)
03F3 CA0D04              JZ      CONOUT$Exit                 ;No, exit
03F6 C5                  PUSH    B       ;Yes - B..          ;Save redirection bit map
03F7 E5                  PUSH    H                           ;Save device table addresses pointer
                 CONOUT$Wait:
03F8 CD0F06              CALL    Check$Output$Ready          ;Check if device not suspended and
                                                             ; (if appropriate) DTR is high
03FB CAF803              JZ      CONOUT$Wait                 ;No, wait

03FE F3                  DI                                  ;Interrupts off to avoid
                                                             ; involuntary re-entrance
03FF 3ACB03              LDA     CONOUT$Character            ;Recover the data byte
0402 4F                  MOV     C,A                         ;Ready for output
0403 CD2608              CALL    Output$Data$Byte            ;Output the data byte
0406 FB                  EI

0407 CD3A06              CALL    Process$Etx$Protocol        ;Deal with Etx/Ack protocol
040A C3ED03              JMP     CONOUT$Next$Device          ;Loop back for next device

                 CONOUT$Exit:
040D 3ACB03              LDA     CONOUT$Character            ;Recover data character
0410 79                  MOV     A,C                         ;CP/M "convention"
0411 C9                  RET
                 ;
                 CONOUT$Escape$Found:                        ;Possible escape sequence
0412 211904              LXI     H,CONOUT$Process$Escape ;Vector processing of next character
                 CONOUT$Set$Processor:
0415 22CC03              SHLD    CONOUT$Processor            ;Set vector address
0418 C9                  RET                                 ;Return to BIOS caller
                 ;#
                 ;
                 ;       Console output: escape sequence processing
```

Figure 8-10. (Continued)

```
                ;
                CONOUT$Process$Escape:                          ;Control arrives here with character
                                                                ;  after escape in C
0419 211B02             LXI     H,CONOUT$Escape$Table           ;Get base of recognition table
                CONOUT$Next$Entry:
041C 7E                 MOV     A,M                             ;Check if at end of table
041D B7                 ORA     A
041E CA2B04             JZ      CONOUT$No$Match                 ;Yes, no match found
0421 B9                 CMP     C                               ;Compare to data character
0422 CA3B04             JZ      CONOUT$Match                    ;They match
0425 23                 INX     H                               ;Move to next entry in table
0426 23                 INX     H
0427 23                 INX     H
0428 C31C04             JMP     CONOUT$Next$Entry               ;Go back and check again
                ;
                CONOUT$No$Match:                                ;No match found, so original
                                                                ;  escape and following character
                                                                ;  must be output
042B C5                 PUSH    B                               ;Save character after escape
042C 0E1B               MVI     C,Function$Key$Lead             ;Get escape character
042E CDE103             CALL    CONOUT$Forced                   ;Output to console devices
0431 C1                 POP     B                               ;Get character after escape
0432 CDE103             CALL    CONOUT$Forced                   ;Output it, too
                ;
                CONOUT$Set$Normal:
0435 21DB03             LXI     H,CONOUT$Normal                 ;Set vector back to normal
0438 C31504             JMP     CONOUT$Set$Processor            ;  for subsequent characters
                ;

                CONOUT$Match:
043B 23                 INX     H                               ;HL -> LS byte of subprocessor
043C 5E                 MOV     E,M                             ;Get LS byte
043D 23                 INX     H
043E 56                 MOV     D,M                             ;Get MS byte
043F EB                 XCHG                                    ;HL -> subprocessor
0440 E9                 PCHL                                    ;Goto subprocessor
                ;
                CONOUT$Date:                            ;Subprocessor to inject current date
                                                        ;  into console input stream (using
                                                        ;  forced input)
0441 218F0F             LXI     H,Date
                CONOUT$Set$Forced$Input:
0444 228D0F             SHLD    CB$Forced$Input
0447 C9                 RET                             ;Return to BIOS' caller
                ;
                CONOUT$Time:                            ;Subprocessor to inject time into
                                                        ;  console input stream
0448 21990F             LXI     H,Time$In$ASCII
044B C34404             JMP     CONOUT$Set$Forced$Input
                ;
                CONOUT$Set$Date:                        ;Subprocessor to set the date by taking
                                                        ;  the next 8 characters of console output
                                                        ;  and storing them in the date string
044E 21A30F             LXI     H,Time$Date$Flags               ;Set flag to indicate that the
0451 3E02               MVI     A,Date$Set                      ;  date has been set by program
0453 B6                 ORA     M
0454 77                 MOV     M,A
0455 3E08               MVI     A,8                             ;Set character count
0457 218F0F             LXI     H,Date                          ;Set address
045A C36C04             JMP     CONOUT$Set$String$Pointer
                ;
                ;
                CONOUT$Set$Time:                        ;Subprocessor to set the time by taking
                                                        ;  the next 8 characters of console output
                                                        ;  and storing them in the time string
045D 21A30F             LXI     H,Time$Date$Flags               ;Set flag to indicate that the
0460 3E01               MVI     A,Time$Set                      ;  time has been set by program
0462 B6                 ORA     M
0463 77                 MOV     M,A
0464 3E08               MVI     A,8                             ;Set character count
0466 21990F             LXI     H,Time$in$ASCII                 ;Set address
0469 C36C04             JMP     CONOUT$Set$String$Pointer
                ;
                CONOUT$Set$String$Pointer:                      ;HL -> string, A = count
046C 32D003             STA     CONOUT$String$Length            ;Save count
046F 22CE03             SHLD    CONOUT$String$Pointer           ;Save address
0472 217804             LXI     H,CONOUT$Process$String ;Vector further output
```

Figure 8-10. (Continued)

```
0475 C31504            JMP     CONOUT$Set$Processor
                ;
                CONOUT$Process$String:          ;Control arrives here for each character
                                                ;  in the string in register C. The
                                                ;  characters are stacked into the
                                                ;  receiving string until either a 00-byte
                                                ;  is encountered or the specified number
                                                ;  of characters is stacked.
0478 2ACE03            LHLD    CONOUT$String$Pointer   ;Get current address for stacking chars
047B 79                MOV     A,C                     ;Check if current character is 00H
047C B7                ORA     A
047D CA3504            JZ      CONOUT$Set$Normal       ;Revert to normal processing
0480 77                MOV     M,A                     ;Otherwise, stack character
0481 23                INX     H                       ;Update pointer
0482 3600              MVI     M,00H                   ;Stack fail-safe terminator
0484 22CE03            SHLD    CONOUT$String$Pointer   ;Save updated pointer
0487 21D003            LXI     H,CONOUT$String$Length  ;Downdate count
048A 35                DCR     M
048B CA3504            JZ      CONOUT$Set$Normal       ;Revert to normal processing
                                                       ;  if count hits 0
048E C9                RET                             ;Return with output vectored back
                                                       ; to CONOUT$Process$String
                ;
                ;#
                ;
                ;       Auxiliary input status
                ;
                ;       This routine checks the character count in the
                ;       appropriate input buffer.
                ;       The A register is set to indicate whether or not
                ;       data is waiting.
                ;
                ;       Entry parameters: none.
                ;
                ;       Exit parameters
                ;
                ;               A = 000H if there is no data waiting
                ;               A = 0FFH if there is data waiting
                ;
                ;========================
                AUXIST:                         ;<=== BIOS entry point (Private)
                ;========================
048F 2A5C00            LHLD    CB$Auxiliary$Input              ;Get redirection word
0492 116400            LXI     D,CB$Device$Table$Addresses     ;  and table pointer
0495 CD6F06            CALL    Select$Device$Table     ;Get device table address
0498 C34708            JMP     Get$Input$Status        ;Get status from input device
                                                       ;  and return to caller
                ;
                ;#
                ;
                ;       Auxiliary output status
                ;
                ;       This routine sets the A register to indicate whether the
                ;       Auxiliary device(s) is/are ready to accept output data.
                ;       As more than one device can be used for auxiliary output, this
                ;       routine returns a Boolean AND of all of their statuses.
                ;
                ;       Entry parameters: none
                ;
                ;       Exit parameters
                ;
                ;               A = 000H if one or more list devices are not ready
                ;               A = 0FFH if all list devices are ready
                ;
                ;
                ;=====================
                AUXOST:                         ;<=== BIOS entry point (Private)
                ;=====================
049B 2A5E00            LHLD    CB$Auxiliary$Output             ;Get list redirection word
049E C37905            JMP     Get$Composite$Status
                ;
                ;#
                ;
                ;       Auxiliary input (replacement for READER)
                ;
                ;       This routine returns the next input character from the
```

Figure 8-10. (Continued)

```
                 ;       appropriate logical auxiliary device.
                 ;
                 ;       Entry parameters: none.
                 ;
                 ;       Exit parameters
                 ;
                 ;               A = data character
                 ;
                 ;========================
                 AUXIN:                          ;<=== BIOS entry point (standard)
                 ;========================
04A1 2A5C00              LHLD    CB$Auxiliary$Input              ;Get redirection word
04A4 116400              LXI     D,CB$Device$Table$Addresses     ;  and table pointer
04A7 CD6F06              CALL    Select$Device$Table     ;Get device table address
04AA C39106              JMP     Get$Input$Character     ;Get next input character
                                                         ;  and return to caller
                 ;
                 ;#
                 ;       Auxiliary output (replaces PUNCH)
                 ;
                 ;       This routine outputs a data byte to the auxiliary device(s).
                 ;       It is similar to CONOUT except that it uses the watchdog
                 ;       timer to detect if a device stays busy for more than
                 ;       30 seconds at a time. It outputs a message to the console
                 ;       if this happens.
                 ;
                 ;       Entry parameters
                 ;
                 ;               C = data byte
                 ;
04AD 0D0A074175  AUXOUT$Busy$Message:    DB      CR,LF,7,'Auxiliary device not Ready?',CR,LF,0
                 ;
                 ;====================
                 AUXOUT:                         ;<=== BIOS entry point (standard)
                 ;====================
04CE 2A5E00              LHLD    CB$Auxiliary$Output     ;Get aux. redirection word
04D1 11AD04              LXI     D,AUXOUT$Busy$Message   ;Message to be output if time
                                                         ;  runs out
04D4 C3A205              JMP     Multiple$Output$Byte
                 ;
                 ;#
                 ;
                 ;       List status
                 ;
                 ;       This routine sets the A register to indicate whether the
                 ;       List Device(s) is/are ready to accept output data.
                 ;       As more than one device can be used for list output, this
                 ;       routine returns a Boolean AND of all of their statuses.
                 ;
                 ;       Entry parameters: none
                 ;
                 ;       Exit parameters
                 ;
                 ;               A = 000H if one or more list devices are not ready
                 ;               A = 0FFH if all list devices are ready
                 ;
                 ;
                 ;=====================
                 LISTST:                                 ;<=== BIOS entry point (standard)
                 ;=====================
04D7 2A6200              LHLD    CB$List$Output          ;Get list redirection word
04DA C37905              JMP     Get$Composite$Status
                 ;
                 ;#
                 ;       List output
                 ;
                 ;       This routine outputs a data byte to the list device.
                 ;       It is similar to CONOUT except that it uses the watchdog
                 ;       timer to detect if the printer stays busy for more
                 ;       than 30 seconds at a time. It outputs a message to the console
                 ;       if this happens.
                 ;
                 ;       Entry parameters
                 ;
                 ;               C = data byte
                 ;
```

Figure 8-10. (Continued)

```
04DD 0D0A075072  LIST$Busy$Message:      DB      CR,LF,7,'Printer not Ready?',CR,LF,0
                 ;
                 ;====================
                 LIST:                           ;<=== BIOS entry point (standard)
                 ;====================
04F5 2A6200              LHLD    CB$List$Output          ;Get list redirection word
04F8 11DD04              LXI     D,LIST$Busy$Message     ;Message to be output if time
                                                         ;  runs out
04FB C3A205              JMP     Multiple$Output$Byte
                 ;
                 ;#
                 ;       Request user choice
                 ;
                 ;       This routine displays an error message, requesting
                 ;       a choice of:
                 ;
                 ;               R -- Retry the operation that caused the error
                 ;               I -- Ignore the error and attempt to continue
                 ;               A -- Abort the program and return to CP/M
                 ;
                 ;       This routine accepts a character from the console,
                 ;       converts it to uppercase and returns to the caller
                 ;       with the response in the A register.
                 ;
                 RUC$Message:
04FE 0D0A                        DB      CR,LF
0500 2020202020                  DB      '     Enter R - Retry, I - Ignore, A - Abort : ',0
                 ;
                 ;
                 Request$User$Choice:
052F CD2D03              CALL    CONST                   ;Gobble up any type-ahead
0532 CA3B05              JZ      RUC$Buffer$Empty
0535 CD3A03              CALL    CONIN
0538 C32F05              JMP     Request$User$Choice

                 RUC$Buffer$Empty:
053B 21FE04              LXI     H,RUC$Message           ;Display prompt
053E CD5305              CALL    Output$Error$Message

0541 CD3A03              CALL    CONIN                   ;Get console character
0544 CD3B0E              CALL    A$To$Upper              ;Make uppercase for comparisons
0547 32B00D              STA     Disk$Action$Confirm     ;Save in confirmatory message
054A F5                  PUSH    PSW                     ;Save for later

054B 21B00D              LXI     H,Disk$Action$Confirm
054E CD5305              CALL    Output$Error$Message

0551 F1                  POP     PSW                     ;Recover action code
0552 C9                  RET
                 ;
                 ;#
                 ;
                 ;       Output error message
                 ;
                 ;       This routine outputs an error message to all the currently
                 ;       selected console devices except those being used to receive
                 ;       LIST output as well. This is to avoid "deadly embrace" situations
                 ;       where the printer's being busy for too long causes an error message
                 ;       to be output -- and console output is being directed to the
                 ;       printer as well.
                 ;
                 ;       This subroutine makes use of most of the CONOUT subroutine.
                 ;       For memory economy it enters CONOUT using a private
                 ;       entry point.
                 ;
                 ;       Entry parameters
                 ;
                 ;               HL -> 00-byte terminated error message
                 ;
                 Output$Error$Message:
0553 E5                  PUSH    H                       ;Save message address
0554 2A5A00              LHLD    CB$Console$Output       ;Get console redirection bit map
0557 EB                  XCHG
0558 2A6200              LHLD    CB$List$Output          ;Get list redirection bit map
                                                         ;HL = list, DE = console
                                                         ;Now set to 0 all bits in the console
```

Figure 8-10. (Continued)

```
                                                         ;  bit map that are set to 1 in the
                                                         ;  list bit map
055B 7C                     MOV     A,H                  ;Get MS byte of list
055C 2F                     CMA                          ;Invert
055D A2                     ANA     D                    ;Preserve only bits with 0's
055E 67                     MOV     H,A                  ;Save result
055F 7D                     MOV     A,L                  ;Repeat for LS byte of list
0560 2F                     CMA
0561 A3                     ANA     E
0562 6F                     MOV     L,A                  ;HL now has only pure console
                                                         ;  devices
0563 B4                     ORA     H                    ;Ensure that at least one device
0564 CA6A05                 JZ      OEM$Device$Present   ;  is selected
0567 210100                 LXI     H,0001H              ;Otherwise use default of device 0
                    OEM$Device$Present:
                    OEM$Next$Character:
056A D1                     POP     D                    ;Recover message address into DE
056B 1A                     LDAX    D                    ;Get next byte of message
056C 13                     INX     D                    ;Update message pointer
056D B7                     ORA     A                    ;Check if end of message
056E C8                     RZ                           ;Yes, exit
056F D5                     PUSH    D                    ;Save message address for later
0570 E5                     PUSH    H                    ;Save special bit map
                                                         ;Data character is in A
0571 CDD103                 CALL    CONOUT$OEM$Entry     ;Enter shared code
0574 E1                     POP     H                    ;Recover special bit map
0575 C36A05                 JMP     OEM$Next$Character
                    ;
                    ;
                    ;
                    ;       Get composite status
                    ;
                    ;       This routine sets the A register to indicate whether the
                    ;       output device(s) is/are ready to accept output data.
                    ;       As more than one device can be used for output, this
                    ;       routine returns a Boolean AND of all of their statuses.
                    ;
                    ;       Entry parameters
                    ;
                    ;               HL = I/O redirection bit map for output device(s)
                    ;
                    ;       Exit parameters
                    ;
                    ;               A = 000H if one or more list devices are not ready
                    ;               A = 0FFH if all list devices are ready
                    ;
0578 00             GCS$Status:     DB      0       ;Composite status of all devices
                    ;
                    ;
                    Get$Composite$Status:
0579 3EFF                   MVI     A,0FFH               ;Assume all devices are ready
057B 327805                 STA     GCS$Status           ;Preset composite status byte

057E 116400                 LXI     D,CB$Device$Table$Addresses    ;Addresses of dev. tables
0581 D5                     PUSH    D                    ;Put onto stack ready for loop
0582 E5                     PUSH    H                    ;Save bit map
                    GCS$Next$Device:
0583 E1                     POP     H                    ;Recover redirection bit map
0584 D1                     POP     D                    ;Recover device table addresses pointer
0585 CD6F06                 CALL    Select$Device$Table  ;Get device table in DE
0588 B7                     ORA     A                    ;Check if a device has been
                                                         ;  selected (i.e. bit map not all zero)
0589 CA9905                 JZ      GCS$Exit             ;No, exit
058C C5                     PUSH    B       ;Yes - B..   ;Save redirection bit map
058D E5                     PUSH    H                    ;Save device table addresses pointer
058E CD0F06                 CALL    Check$Output$Ready   ;Check if device ready
0591 217805                 LXI     H,GCS$Status         ;AND together with previous devices
0594 A6                     ANA     M                    ;  status
0595 77                     MOV     M,A                  ;Save composite status

0596 C38305                 JMP     GCS$Next$Device      ;Loop back for next device
                    ;
                    GCS$Exit:
0599 3A7805                 LDA     GCS$Status           ;Return with composite status
059C B7                     ORA     A
059D C9                     RET
```

Figure 8-10. (Continued)

```
              ;
              ;#
              ;
              ;       Multiple output byte
              ;
              ;       This routine outputs a data byte to the all of the
              ;       devices specified in the I/O redirection word.
              ;       It is similar to CONOUT except that it uses the watchdog
              ;       timer to detect if any of the devices stays busy for more
              ;       than 30 seconds at a time. It outputs a message to the console
              ;       if this happens.
              ;
              ;       Entry parameters
              ;
              ;               HL = I/O redirection bit map
              ;               DE -> Message to be output if time runs out
              ;               C = data byte
              ;
0708 =        MOB$Maximum$Busy        EQU     1800    ;Number of clock ticks (each at
                                                      ;  16.666 milliseconds) for  which the
                                                      ;  device might be busy
059E 00       MOB$Character:          DB      0       ;Character to be output
059F 0000     MOB$Busy$Message:       DW      0       ;Address of message to be
                                                      ;  output if time runs out
05A1 00       MOB$Need$Message:       DB      0       ;Flag used to detect that the
                                                      ;  watchdog timer timed out
              ;
              Multiple$Output$Byte:
05A2 79               MOV     A,C                     ;Get data byte
05A3 320807           STA     MOB$Maximum$Busy        ;Save copy
05A6 EB               XCHG                            ;HL -> timeout message
05A7 229F05           SHLD    MOB$Busy$Message        ;Save for later use
05AA EB               XCHG                            ;HL = bit map again

05AB 116400           LXI     D,CB$Device$Table$Addresses     ;Addresses of dev. tables
05AE D5               PUSH    D                       ;Save on stack ready for loop
05AF E5               PUSH    H                       ;Save I/O redirection bit map
              MOB$Next$Device:
05B0 E1               POP     H                       ;Recover redirection bit map
05B1 D1               POP     D                       ;Recover device table addresses pointer
05B2 CD6F06           CALL    Select$Device$Table     ;Get device table in DE
05B5 B7               ORA     A                       ;Check if any device selected
05B6 CAEC05           JZ      MOB$Exit

05B9 C5               PUSH    B       ;<- Yes : B     ;Save device table addresses pointer
05BA E5               PUSH    H                       ;Save redirection bit map
              ;
              MOB$Start$Watchdog:
05BB AF               XRA     A                               ;Reset message needed flag
05BC 32A105           STA     MOB$Need$Message
05BF 010807           LXI     B,MOB$Maximum$Busy              ;Time delay
05C2 210906           LXI     H,MOB$Not$Ready                 ;Address to go to
05C5 CD6D08           CALL    Set$Watchdog                    ;Start timer

              MOB$Wait:
05C8 3AA105           LDA     MOB$Need$Message                ;Check if watchdog timed out
05CB B7               ORA     A
05CC C2EE05           JNZ     MOB$Output$Message              ;Yes, output warning message
05CF CD0F06           CALL    Check$Output$Ready              ;Check if device ready
05D2 CAC805           JZ      MOB$Wait                        ;No, wait
              ;
05D5 F3               DI                                      ;Interrupts off to avoid
                                                              ;  involuntary reentrance
05D6 010000           LXI     B,0                             ;Turn off watchdog
05D9 CD6D08           CALL    Set$Watchdog                    ; (HL setting is irrelevant)

05DC 3A9E05           LDA     MOB$Character                   ;Get data byte
05DF 4F               MOV     C,A
05E0 CD2608           CALL    Output$Data$Byte                ;Output the data byte
05E3 FB               EI
05E4 CD3A06           CALL    Process$Etx$Protocol            ;Deal with ETX/ACK protocol
05E7 C3B005           JMP     MOB$Next$Device
              ;
              MOB$Ignore$Exit:                                ;Ignore timeout error
05EA E1               POP     H                               ;Balance the stack
05EB D1               POP     D
```

Figure 8-10. (Continued)

```
                ;
                MOB$Exit:
05EC 79                 MOV     A,C                             ;CP/M "convention"
05ED C9                 RET
                ;
                MOB$Output$Message:
05EE 2A9F05             LHLD    MOB$Busy$Message                ;Display warning message
05F1 CD5305             CALL    Output$Error$Message            ;  on selected console devices
                MOB$Request$Choice:
05F4 CD2F05             CALL    Request$User$Choice             ;Display message and get
                                                                ;  action character
05F7 FE52               CPI     'R'                             ;Retry
05F9 CABB05             JZ      MOB$Start$Watchdog              ;Restart watchdog and try again
05FC FE49               CPI     'I'                             ;Ignore
05FE CAEA05             JZ      MOB$Ignore$Exit
0601 FE41               CPI     'A'                             ;Abort
0603 CA360E             JZ      System$Reset                    ;  Give BDOS function 0
0606 C3F405             JMP     MOB$Request$Choice
                ;
                MOB$Not$Ready:                          ;Watchdog timer routine will call this
                                                        ;  routine if the device is busy
                                                        ;  for more than approximately 30 seconds
                                                        ;Note: This is an interrupt service routine
0609 3EFF               MVI     A,0FFH                  ;Set request to output message
060B 32A105             STA     MOB$Need$Message
060E C9                 RET                             ;Return to the watchdog routine
                ;
                ;#
                ;       Check output ready
                ;
                ;       This routine checks to see if the specified device is ready
                ;       to receive output data.
                ;       It does so by checking to see if the device has been suspended
                ;       for protocol reasons and if DTR is low.
                ;
                ;       NOTE: This routine does NOT check if the USART itself is ready.
                ;             This test is done in the output data byte routine itself.
                ;
                ;       Entry parameters
                ;
                ;               DE -> device table
                ;
                ;       Exit parameters
                ;
                ;               A = 000H (Zero-flag set)   : Device not ready
                ;               A = 0FFH (Zero-flag clear) : Device ready
                ;
                Check$Output$Ready:
060F 210E00             LXI     H,DT$Status             ;Get device status
0612 19                 DAD     D                       ;HL -> status byte
0613 7E                 MOV     A,M                     ;Get status byte
0614 47                 MOV     B,A                     ;Take a copy of the status byte
0615 E601               ANI     DT$Output$Suspend       ;Check if output is suspended
0617 C23806             JNZ     COR$Not$Ready           ;Yes, indicate not ready

061A 3E04               MVI     A,DT$Output$DTR         ;Check if DTR must be high to send
061C A0                 ANA     B                       ;Mask with device status from table
061D CA3406             JZ      COR$Ready               ;No, device is logically ready

0620 210000             LXI     H,DT$Status$Port        ;Set up to read device status
0623 19                 DAD     D
0624 7E                 MOV     A,M                     ;Get status port number
0625 322906             STA     COR$Status$Port         ;Set up instruction below

0628 DB                 DB      IN
                COR$Status$Port:
0629 00                 DB      0               ;<-- Set up by instruction above
062A 4F                 MOV     C,A                     ;Save hardware status

062B 210400             LXI     H,DT$DTR$Ready          ;Yes, set up to check chip status
062E 19                 DAD     D                       ;  to see if DTR is high
062F 7E                 MOV     A,M                     ;Get DTR high status mask
0630 A1                 ANA     C                       ;Test chip status
0631 CA3806             JZ      COR$Not$Ready           ;DTR low, indicate not ready
                ;
                COR$Ready:
```

Figure 8-10. (Continued)

```
0634 3EFF              MVI     A,0FFH                  ;Indicate device ready for output
0636 B7                ORA     A
0637 C9                RET
               ;
               COR$Not$Ready:                          ;Indicate device not ready for output
0638 AF                XRA     A
0639 C9                RET
               ;
               ;#
               ;
               ;       Process ETX/ACK protocol
               ;
               ;       This routine maintains ETX/ACK protocol.
               ;       After a specified number of data characters have been output
               ;       to the device, an ETX character is output and the device
               ;       put into output suspended state. Only when an incoming
               ;       ACK character is received (under interrupt control) will
               ;       output be resumed to the device.
               ;
               ;       Entry parameters
               ;
               ;               DE -> device table
               ;
               ;       Exit parameters
               ;
               ;               Message count downdated (and reset if necessary)
               ;
               Process$Etx$Protocol:
063A 210E00            LXI     H,DT$Status             ;Check if ETX/ACK protocol enabled
063D 19                DAD     D
063E 7E                MOV     A,M
063F E610              ANI     DT$Output$Etx
0641 C8                RZ                              ;No, so return immediately
0642 211000            LXI     H,DT$Etx$Count          ;Yes, so downdate count
0645 19                DAD     D
0646 E5                PUSH    H                       ;Save address of count for later
0647 4E                MOV     C,M                     ;Get LS byte
0648 23                INX     H
0649 46                MOV     B,M                     ;Get MS byte
064A 0B                DCX     B
064B 78                MOV     A,B
064C B1                ORA     C                       ;Check if count now zero
064D C25706            JNZ     PEP$Save$Count          ;No
0650 211200            LXI     H,DT$Etx$Message$Length ;Yes, reset to message length
0653 19                DAD     D
0654 4E                MOV     C,M                     ;Get LS byte
0655 23                INX     H
0656 46                MOV     B,M                     ;Get MS byte
               PEP$Save$Count:
0657 E1                POP     H                       ;Recover address of count
0658 71                MOV     M,C                     ;Save count back in table
0659 23                INX     H
065A 70                MOV     M,B
               ;
065B B7                ORA     A                       ;Reestablish whether count hit 0
065C C0                RNZ                             ;No, no further processing required
065D 0E03              MVI     C,ETX                   ;Yes, send ETX to device
065F F3                DI                              ;Avoids involuntary reentrance
0660 CD2608            CALL    Output$Data$Byte
0663 FB                EI
0664 210E00            LXI     H,DT$Status             ;Flag device as output suspended
0667 19                DAD     D
0668 F3                DI                              ;Avoid interaction with interrupts
0669 7E                MOV     A,M                     ;Get status byte
066A F601              ORI     DT$Output$Suspend       ;Set bit
066C 77                MOV     M,A                     ;Save back in table
066D FB                EI
066E C9                RET
               ;
               ;#
               ;
               ;       Select device table
               ;
               ;       This routine scans a 16-bit word, and depending on which is the
               ;       first 1-bit set, selects the corresponding device table address.
               ;
```

Figure 8-10. (Continued)

```
               ;       Entry parameters
               ;
               ;               HL = Bit map
               ;               DE -> Table of device table addresses
               ;                       The first address in the list is called
               ;                       if the least significant bit of the bit map is
               ;                       nonzero, and so on.
               ;
               ;       Exit parameters
               ;
               ;               BC -> Current entry in device table addresses
               ;               DE = Selected device table address
               ;               HL = Shifted bit map
               ;                    Nonzero if a 1-bit was found
               ;                    Zero if bit map now entirely 0000
               ;
               ;       Note: If HL is 0000H on input, then the first entry in the
               ;       device table addresses will be returned in DE.
               ;
               Select$Device$Table:
066F 7C                MOV     A,H             ;Get most significant byte of bit map
0670 B5                ORA     L               ;Check if HL completely 0
0671 C8                RZ                      ;Return indicating no more bits set
0672 7D                MOV     A,L             ;Check if the LS bit is nonzero
0673 E601              ANI     1
0675 C28006            JNZ     SDT$Bit$Set     ;Yes, return corresponding address
0678 13                INX     D               ;No, update table pointer
0679 13                INX     D
067A CDDB08            CALL    SHLR            ;Shift HL right one bit
067D C36F06            JMP     Select$Device$Table     ;Check next bit
               SDT$Bit$Set:
0680 E5                PUSH    H               ;Save shifted bit map
0681 42                MOV     B,D             ;Take copy of table pointer
0682 4B                MOV     C,E
0683 EB                XCHG                    ;HL -> address in table
0684 5E                MOV     E,M
0685 23                INX     H
0686 56                MOV     D,M             ;DE -> selected device table
                                               ;Set up registers for another
                                               ;  entry
0687 E1                POP     H               ;Recover shifted bit map
0688 CDDB08            CALL    SHLR            ;Shift bit map right one bit
068B 03                INX     B               ;Update DT address table pointer to
068C 03                INX     B               ;  entry
068D 3E01              MVI     A,1             ;Indicate that a one bit was found
068F B7                ORA     A               ;  and registers are set up correctly
0690 C9                RET
               ;
               ;#
               ;
               ;       Get input character
               ;
               ;       This routine gets the next input character from the device
               ;       specified in the device table handed over as an input
               ;       parameter.
               ;
               Get$Input$Character:
0691 211900            LXI     H,DT$Character$Count    ;Check if any characters have
0694 19                DAD     D                       ;  been stored in the buffer
               GIC$Wait:
0695 FB                EI                              ;Ensure that incoming chars. will
                                                       ;  be detected
0696 7E                MOV     A,M                     ;Get character count
0697 B7                ORA     A
0698 CA9506            JZ      GIC$Wait                ;No characters, so wait
069B 35                DCR     M                       ;Down date character count for
                                                       ;  the character about to be
                                                       ;  removed from the buffer
069C 211700            LXI     H,DT$Get$Offset         ;Use the get offset to access
069F CDF007            CALL    Get$Address$in$Buffer   ;Returns HL -> character
                                                       ;  and with get offset updated
06A2 7E                MOV     A,M                     ;Get the actual data character
06A3 F5                PUSH    PSW                     ;Save until later

06A4 211900            LXI     H,DT$Character$Count    ;Check downdated count of chars. in
06A7 19                DAD     D                       ;  buffer, checking if input should be
```

Figure 8-10. (Continued)

```
0702 11CE02             LXI     D,DT$2          ;Device 2
0705 CD1607             CALL    Service$Device

0708 3E20               MVI     A,IC$EOI        ;Tell the interrupt controller chip
070A D3D8               OUT     IC$OCW2$Port    ;  that the interrupt has been serviced
070C D1                 POP     D               ;Restore registers
070D C1                 POP     B
070E F1                 POP     PSW
070F 2A8422             LHLD    PI$User$Stack   ;Switch back to user's stack
0712 F9                 SPHL
0713 E1                 POP     H
0714 FB                 EI                      ;Re-enable interrupts in the CPU
0715 C9                 RET                     ;Resume pre-interrupt processing
                ;
                ;#
                ;
                ;       Service device
                ;
                ;       This routine performs the device interrupt servicing,
                ;       checking to see if the device described in the specified
                ;       device table (address in DE) is actually interrupting,
                ;       and if so, inputs the character. Depending on which data character
                ;       is input, this routine will either stack it in the input buffer
                ;       (shutting off the input stream if the buffer is nearly full),
                ;       or will suspend or resume the output to the device.
                ;
                ;       Entry parameters
                ;
                ;               DE -> device table
                ;
                Service$Device:
0716 210000             LXI     H,DT$Status$Port        ;Check if this device is really
0719 19                 DAD     D                       ;  interrupting
071A 7E                 MOV     A,M                     ;Get status port number
071B 321F07             STA     SD$Status$Port          ;Store in instruction below

071E DB                 DB      IN                      ;Input status
                SD$Status$Port:
071F 00                 DB      0       ;<-- Set up by instruction above
                ;
0720 210300             LXI     H,DT$Input$Ready        ;Check if status indicates data ready
0723 19                 DAD     D
0724 A6                 ANA     M                       ;Mask with input ready value
0725 C8                 RZ                              ;No, return to interrupt service
                                                        ;Check if any errors have occurred
0726 210700             LXI     H,DT$Detect$Error$Port  ;Set up to read error status
0729 19                 DAD     D                       ;  interrupting
072A 7E                 MOV     A,M                     ;Get status port number
072B 322F07             STA     SD$Error$Port           ;Store in instruction below

072E DB                 DB      IN                      ;Input error status
                SD$Error$Port:
072F 00                 DB      0       ;<-- Set up by instruction above
                ;
0730 210800             LXI     H,DT$Detect$Error$Value ;Mask with error bit(s)
0733 19                 DAD     D
0734 A6                 ANA     M
0735 CA4707             JZ      SD$No$Error             ;No bit(s) set
0738 210900             LXI     H,DT$Reset$Error$Port   ;Set up to reset error
073B 19                 DAD     D
073C 7E                 MOV     A,M                     ;Get reset port number
073D 324607             STA     SD$Reset$Error$Port     ;Store in instruction below
0740 210A00             LXI     H,DT$Reset$Error$Value
0743 19                 DAD     D
0744 7E                 MOV     A,M                     ;Get reset interrupt value

0745 D3                 DB      OUT
                SD$Reset$Error$Port:
0746 00                 DB      0       ;<-- Set up in instruction above

                SD$No$Error:
0747 210100             LXI     H,DT$Data$Port          ;Input the data character (this may
074A 19                 DAD     D                       ;  be garbled if an error occurred)
074B 7E                 MOV     A,M                     ;Get data port number
074C 325007             STA     SD$Data$Port            ;Store in instruction below
```

Figure 8-10. (Continued)

```
074F DB                  DB      IN                      ;Input data character
                 SD$Data$Port:
0750 00                  DB      0       ;<-- Set up by instruction above

0751 47                  MOV     B,A                     ;Take copy of data character above
0752 210E00              LXI     H,DT$Status             ;Check if either XON or ETX protocols
0755 19                  DAD     D                       ;  is currently active
0756 7E                  MOV     A,M                     ;Get protocol byte
0757 E618                ANI     DT$Output$Xon + DT$Output$Etx
0759 CA8107              JZ      SD$No$Protocol          ;Neither is active
075C E608                ANI     DT$Output$Xon           ;Check if XON/XOFF is active
075E C26E07              JNZ     SD$Check$if$Xon         ;Yes, check if XON char. input
                                                         ;No, assume ETX/ACK active
0761 3E06                MVI     A,ACK                   ;Check if input character is ACK
0763 B8                  CMP     B
0764 C28107              JNZ     SD$No$Protocol          ;No, process character as data
                 SD$Output$Desuspend:                    ;Yes, device now ready
                                                         ;  to accept more data, so indicate
                                                         ;  output to device can resume
                                                         ;The noninterrupt driven output
                                                         ;  routine checks the suspend bit
0767 7E                  MOV     A,M                     ;Get status/protocol byte again
0768 E6FE                ANI     OFFH AND NOT DT$Output$Suspend ;Preserve all bits BUT suspend
076A 77                  MOV     M,A                     ;Save back with suspend = 0
076B C3D907              JMP     SD$Exit                 ;Exit to interrupt service without
                                                         ;  saving data character
                 ;
                 SD$Check$if$Xon:                        ;XON/XOFF protocol active, so
                                                         ;  if XOFF received, suspend output
                                                         ;  if XON received, resume output
                                                         ;The noninterrupt driven output
                                                         ;  routine checks the suspend bit
076E 3E11                MVI     A,XON                   ;Check if XON character input
0770 B8                  CMP     B
0771 CA6707              JZ      SD$Output$Desuspend     ;Yes, enable output to device
0774 3E13                MVI     A,XOFF                  ;Check if XOFF character input
0776 B8                  CMP     B
0777 C28107              JNZ     SD$No$Protocol          ;No, process character as data
                 SD$Output$Suspend:                      ;Device needs pause in output of
                                                         ;  data, so indicate output suspended
077A 7E                  MOV     A,M                     ;Get status/protocol byte again
077B F601                ORI     DT$Output$Suspend       ;Set suspend bit to 1
077D 77                  MOV     M,A                     ;Save back in device table
077E C3D907              JMP     SD$Exit                 ;Exit to interrupt service without
                                                         ;  saving the input character
                 ;
                 SD$No$Protocol:
0781 211800              LXI     H,DT$Buffer$Length$Mask ;Check if there is still space
0784 19                  DAD     D                       ;  in the input buffer
0785 7E                  MOV     A,M                     ;Get length - 1
0786 3C                  INR     A                       ;Update to actual length
0787 211900              LXI     H,DT$Character$Count    ;Get current count of characters
078A 19                  DAD     D                       ;  in buffer
078B BE                  CMP     M                       ;Check if count = length
078C CAEB07              JZ      SD$Buffer$Full          ;Yes, output bell character
078F C5                  PUSH    B                       ;Save data character
0790 211600              LXI     H,DT$Put$Offset         ;Compute address of character in
                                                         ;  input buffer
0793 CDF007              CALL    Get$Address$In$Buffer   ;HL -> character position
0796 C1                  POP     B                       ;Recover input character
0797 70                  MOV     M,B                     ;Save character in input buffer
                                                 ;Update number of characters in input
                                                 ;  buffer, checking if input should
                                                 ;  be temporarily halted
0798 211900              LXI     H,DT$Character$Count
079B 19                  DAD     D
079C 34                  INR     M                       ;Update character count
079D 7E                  MOV     A,M                     ;Get updated count
079E 211A00              LXI     H,DT$Stop$Input$Count   ;Check if current count matches
07A1 19                  DAD     D                       ;  buffer-full threshold
07A2 BE                  CMP     M
07A3 C2CE07              JNZ     SD$Check$Control        ;Not at threshold, check if control
                                                         ;  character input
07A6 210E00              LXI     H,DT$Status             ;At threshold, check which means
07A9 19                  DAD     D                       ;  for pausing input are to be used
```

Figure 8-10. (Continued)

```
07AA 7E                    MOV    A,M                      ;Get status/protocol byte
07AB F602                  ORI    DT$Input$Suspend         ;Indicate input is suspended
07AD 77                    MOV    M,A                      ;Save updated status in table
07AE F5                    PUSH   PSW                      ;Save for later use
07AF E640                  ANI    DT$Input$RTS             ;Check if clear to send to be dropped
07B1 CAC307                JZ     SD$Check$Input$Xon       ;No
07B4 210B00                LXI    H,DT$RTS$Control$Port    ;Yes, get control port number
07B7 19                    DAD    D
07B8 7E                    MOV    A,M
07B9 32C207                STA    SD$Drop$RTS$Port                 ;Store in instruction below
07BC 210C00                LXI    H,DT$Drop$RTS$Value
07BF 19                    DAD    D
07C0 7E                    MOV    A,M                      ;Get value needed to drop RTS

07C1 D3                    DB     OUT
             SD$Drop$RTS$Port:
07C2 00                    DB     0               ;<- Set up in instruction above
                                                           ;Drop into input XON test
             SD$Check$Input$Xon:                  ;Check if XON/XOFF protocol being used
                                                  ;  to temporarily suspend input
07C3 F1                    POP    PSW                      ;Recover status/protocol byte
07C4 E680                  ANI    DT$Input$Xon             ;Check if XON bit set
07C6 CACE07                JZ     SD$Check$Control         ;No, see if control char. input
07C9 0E13                  MVI    C,XOFF                   ;Yes, output XOFF character
07CB CD2608                CALL   Output$Data$Byte         ;Output data byte
             ;
             SD$Check$Control:                    ;Check if control character (other than
                                                  ;  CR, LF, or TAB) input, and update
                                                  ;  count of control characters in buffer
07CE CD0808                CALL   Check$Control$Char       ;Check if control character
07D1 CAD907                JZ     SD$Exit                  ;No,it is not a control character
07D4 211C00                LXI    H,DT$Control$Count
07D7 19                    DAD    D
07D8 34                    INR    M                        ;Update count of control chars.
             ;
             SD$Exit:                             ;Reset hardware interrupt system
07D9 210500                LXI    H,DT$Reset$Int$Port
07DC 19                    DAD    D
07DD 7E                    MOV    A,M                      ;Get reset port number
07DE B7                    ORA    A                        ;Check if port specified
                                                           ;  (assumes it will always be NZ)
07DF C8                    RZ                              ;Bypass reset if no port specified
07E0 32E907                STA    SD$Reset$Int$Port        ;Store in instruction below
07E3 210600                LXI    H,DT$Reset$Int$Value
07E6 19                    DAD    D
07E7 7E                    MOV    A,M                      ;Get reset interrupt value

07E8 D3                    DB     OUT
             SD$Reset$Int$Port:
07E9 00                    DB     0       ;<-- Set up in instruction above
07EA C9                    RET                             ;Return to interrupt service routine
             ;
             SD$Buffer$Full:                               ;Input buffer completely full
07EB 0E07                  MVI    C,BELL                   ;Send bell character as desperate
07ED C32608                JMP    Output$Data$Byte         ;  measure. Note JMP return to
                                                           ;  caller will be done by subroutine
             ;
             ;#
             ;
             ;       Get address in buffer
             ;
             ;       This routine computes the address of the next character to
             ;       access in a device buffer.
             ;
             ;       Entry parameters
             ;
             ;               DE -> appropriate device table
             ;               HL =  offset in the device table of either the
             ;                       Get$Offset or the Put$Offset
             ;
             ;       Exit parameters
             ;
             ;               DE unchanged
             ;               HL -> address in character buffer
             ;
             Get$Address$In$Buffer:
```

Figure 8-10. (Continued)

```
07F0 19                 DAD     D                       ;HL -> get/put offset in dev. table
07F1 E5                 PUSH    H                       ;Preserve pointer to table
07F2 4E                 MOV     C,M                     ;Get offset value
07F3 0600               MVI     B,0                     ;Make into word value
                                                ;Update offset value, resetting to
                                                ;  0 at end of buffer
07F5 79                 MOV     A,C                     ;Get copy of offset
07F6 3C                 INR     A                       ;Update to next position
07F7 211800             LXI     H,DT$Buffer$Length$Mask
07FA 19                 DAD     D
07FB A6                 ANA     M                       ;Mask LS bits with length - 1
07FC E1                 POP     H                       ;Recover pointer to offset in table
07FD 77                 MOV     M,A                     ;Save new value (set to 0 if nec.)
07FE 211400             LXI     H,DT$Buffer$Base        ;Get base address of input buffer
0801 19                 DAD     D                       ;HL -> address of buffer in table
0802 7E                 MOV     A,M                     ;Get LS byte of address
0803 23                 INX     H                       ;HL -> MS byte of address
0804 66                 MOV     H,M                     ;H = MS byte
0805 6F                 MOV     L,A                     ;L = LS byte
0806 09                 DAD     B                       ;Add on offset to base
0807 C9                 RET

                ;
                ;#
                ;
                ;       Check control character
                ;
                ;       This routine checks the character in A to see if it is a
                ;       control character other than CR, LF, or TAB. The result is
                ;       returned in the Z-flag.
                ;
                ;       Entry parameters
                ;
                ;               A = character to be checked
                ;
                ;       Exit parameters
                ;
                ;               Zero status if A does not contain a control character
                ;                       or if it is CR, LF, or TAB
                ;
                ;               Nonzero if A contains a control character other than
                ;                       CR, LF, or TAB.
                Check$Control$Char:
0808 3E1F               MVI     A,' ' - 1               ;Space is first noncontrol char.
080A B8                 CMP     B
080B DA2408             JC      CCC$No                  ;Not a control character
080E 3E0D               MVI     A,CR                    ;Check if carriage return
0810 B8                 CMP     B
0811 CA2408             JZ      CCC$No                  ;Not really a control character
0814 3E0A               MVI     A,LF                    ;Check if LF
0816 B8                 CMP     B
0817 CA2408             JZ      CCC$No                  ;Not really a control character
081A 3E09               MVI     A,TAB                   ;Check if horizontal tab
081C B8                 CMP     B
081D CA2408             JZ      CCC$No                  ;Not really a control character
0820 3E01               MVI     A,1                     ;Indicate a control character
0822 B7                 ORA     A
0823 C9                 RET
                CCC$No:                                 ;Indicate A does not contain
0824 AF                 XRA     A                       ;  a control character
0825 C9                 RET
                ;
                ;#
                ;
                ;       Output data byte
                ;
                ;       This is a simple polled output routine that outputs a single
                ;       character (in register C on entry) to the device specified in
                ;       the device table.
                ;       Preferably, this routine would have been re-entrant; however
                ;       it does have to store the port numbers. Therefore, to use it
                ;       from code executed with interrupts enabled, the instruction
                ;       sequence must be:
                ;
                ;               DI              ;Interrupts off
                ;               CALL    Output$Data$Byte
```

Figure 8-10. (Continued)

```
                ;               EI              ;Interrupts on
                ;
                ;       Failure to do this may cause involuntary re-entrance.
                ;
                ;       Entry parameters
                ;
                ;               C = character to be output
                ;               DE -> device table
                ;
                Output$Data$Byte:
0826 C5                 PUSH    B                       ;Save registers
0827 210200             LXI     H,DT$Output$Ready       ;Get output ready status mask
082A 19                 DAD     D
082B 46                 MOV     B,M
082C 210000             LXI     H,DT$Status$Port        ;Get status port number
082F 19                 DAD     D
0830 7E                 MOV     A,M
0831 323508             STA     ODB$Status$Port         ;Store in instruction below
                ODB$Wait$until$Ready:

0834 DB                 DB      IN                      ;Read status
                ODB$Status$Port:
0835 00                 DB      0       ;<--  Set up in instruction above

0836 A0                 ANA     B                       ;Check if ready for output
0837 CA3408             JZ      ODB$Wait$until$Ready    ;No
083A 210100             LXI     H,DT$Data$Port          ;Get data port
083D 19                 DAD     D
083E 7E                 MOV     A,M
083F 324408             STA     ODB$Data$Port           ;Store in instruction below
0842 79                 MOV     A,C                     ;Get character to output

0843 D3                 DB      OUT
                ODB$Data$Port:
0844 00                 DB      0       ;<--  Set up in instruction above

0845 C1                 POP     B                       ;Restore registers
0846 C9                 RET
                ;
                ;#
                ;
                ;
                ;       Input status routine
                ;
                ;       This routine returns a value in the A register indicating whether
                ;       one or more data characters is/are waiting in the input buffer.
                ;       Some products, such as Microsoft BASIC, defeat normal type-ahead
                ;       by constantly "gobbling" characters in order to see if an incoming
                ;       Control-S, -Q or -C has been received. In order to preserve
                ;       type-ahead under these circumstances, the input status return
                ;       can, as an option selected by the user, return "data waiting" only
                ;       if the input buffer contains a Control-S, -Q or -C. This fools
                ;       Microsoft BASIC into allowing type-ahead.
                ;
                ;       Entry parameters
                ;
                ;               DE -> device table
                ;
                ;       Exit parameters
                ;
                ;               A = 000H if no characters are waiting in the input
                ;                       buffer
                ;
                ;
                Get$Input$Status:
0847 210F00             LXI     H,DT$Status$2           ;Check if fake mode enabled
084A 19                 DAD     D                       ;HL -> status byte in table
084B 7E                 MOV     A,M                     ;Get status byte
084C E601               ANI     DT$Fake$Typeahead       ;Isolate status bit
084E CA5B08             JZ      GIS$True$Status         ;Fake mode disabled
                                                        ;
                                                        ;Fake mode -- only indicates data
                                                        ;ready if control chars. in buffer
0851 211C00             LXI     H,DT$Control$Count      ;Check if any control characters
0854 19                 DAD     D                       ;  in the input buffer
0855 AF                 XRA     A                       ;Cheap 0
```

Figure 8-10. (Continued)

```
0856 B6                  ORA     M                       ;Set flags according to count
0857 C8                  RZ                              ;Return indicating zero
              GIS$Data$Ready:
0858 AF                  XRA     A                       ;Cheap 0
0859 3D                  DCR     A                       ;Set A = 0FFH and flags NZ
085A C9                  RET                             ;Return to caller
              ;
              GIS$True$Status:                          ;
                                                        ;True status, based on any characters
                                                        ;ready in input buffer
085B 2A8D0F              LHLD    CB$Forced$Input         ;Check if any forced input waiting
085E 7E                  MOV     A,M                     ;Get next character of forced input
085F B7                  ORA     A                       ;Check if nonzero
0860 C25808              JNZ     GIS$Data$Ready          ;Yes, indicate data waiting

0863 211900              LXI     H,DT$Character$Count    ;Check if any characters
0866 19                  DAD     D                       ;  in buffer
0867 7E                  MOV     A,M                     ;Get character count
0868 B7                  ORA     A
0869 C8                  RZ                              ;Empty buffer, A = 0, Z-set
086A C35808              JMP     GIS$Data$Ready
              ;
              ;
              ;#
              ;
              ;       Real time clock processing
              ;
              ;       Control is transferred to the RTC$Interrupt routine each time
              ;       the real time clock ticks. The tick count is downdated to see
              ;       if a complete second has elapsed. If so, the ASCII time in
              ;       the configuration block is updated.
              ;
              ;       With each tick, the watchdog count is downdated to see if control
              ;       must be "forced" to a previously specified address on return
              ;       from the RTC interrupt. The watchdog timer can be used to pull
              ;       control out of what would otherwise be an infinite loop, such
              ;       as waiting for the printer to come ready.
              ;
              ;
              ;       Set watchdog
              ;
              ;       This is a noninterrupt level subroutine that simply sets the
              ;       watchdog count and address
              ;
              ;       Entry parameters
              ;
              ;               BC = number of clock ticks before watchdog should
              ;                       "time out"
              ;               HL = address to which control will be transferred when
              ;                       watchdog times out
              ;
              Set$Watchdog:
086D F3                  DI                              ;Avoid interference from interrupts
086E 22C100              SHLD    RTC$Watchdog$Address    ;Set address
0871 60                  MOV     H,B
0872 69                  MOV     L,C
0873 22BF00              SHLD    RTC$Watchdog$Count      ;Set count
0876 FB                  EI
0877 C9                  RET
              ;
              ;
              ;#      /
              ;
                                              ;Control is received here each time the
                                              ;  real time clock ticks
              RTC$Interrupt:
0878 F5                  PUSH    PSW                     ;Save other registers
0879 228622              SHLD    PI$User$HL              ;Switch to local stack
087C 210000              LXI     H,0
087F 39                  DAD     SP                      ;Get user's stack
0880 228422              SHLD    PI$User$Stack           ;Save it
0883 31B022              LXI     SP,PI$Stack             ;Switch to local stack
0886 C5                  PUSH    B
0887 D5                  PUSH    D

0888 21BE00              LXI     H,RTC$Tick$Count        ;Downdate tick count
```

Figure 8-10. (Continued)

```
088B 35                 DCR     M
088C C2B008             JNZ     RTC$Check$Watchdog      ;Is not at 0 yet
                                                        ;One second has elapsed so
088F 3ABD00             LDA     RTC$Ticks$per$Second    ;  reset to original value
0892 77                 MOV     M,A
                                                        ;Update ASCII real time clock
0893 11A10F             LXI     D,Time$in$ASCII$End     ;DE -> 1 character after ASCII time
0896 21BD00             LXI     H,Update$Time$End       ;HL -> 1 character after control table
                RTC$Update$Digit:
0899 1B                 DCX     D                       ;Downdate pointer to time in ASCII
089A 2B                 DCX     H                       ;Downdate pointer to control table
089B 7E                 MOV     A,M                     ;Get next control character
089C B7                 ORA     A                       ;Check if end of table and therefore
089D CAB008             JZ      RTC$Clock$Updated       ;  all digits of clock updated
08A0 FA9908             JM      RTC$Update$Digit        ;Skip over ":" in ASCII time
08A3 1A                 LDAX    D                       ;Get next ASCII time digit
08A4 3C                 INR     A                       ;Update it
08A5 12                 STAX    D                       ;  and store it back
08A6 BE                 CMP     M                       ;Compare to maximum value
08A7 C2B008             JNZ     RTC$Clock$Updated       ;No carry needed so update complete
08AA 3E30               MVI     A,'0'                   ;Reset digit to ASCII 0
08AC 12                 STAX    D                       ;  and store back in ASCII time
08AD C39908             JMP     RTC$Update$Digit        ;Go back for next digit
                ;
                RTC$Clock$Updated:
                RTC$Check$Watchdog:
08B0 2ABF00             LHLD    RTC$Watchdog$Count      ;Get current watchdog count
08B3 2B                 DCX     H                       ;Downdate it
08B4 7C                 MOV     A,H                     ;Check if it is now 0FFFFH
08B5 B7                 ORA     A
08B6 FACB08             JM      RTC$Dog$Not$Set         ;It must have been 0 beforehand
08B9 B5                 ORA     L                       ;Check if it is now 0
08BA C2C808             JNZ     RTC$Dog$NZ              ;No, it is not out of time

                                                        ;Watchdog time elapsed, so "call"
                                                        ;  appropriate routine
08BD 21C508             LXI     H,RTC$Watchdog$Return   ;Set up return address
08C0 E5                 PUSH    H                       ;  ready for return
08C1 2AC100             LHLD    RTC$Watchdog$Address    ;Transfer control as though by CALL
08C4 E9                 PCHL
                RTC$Watchdog$Return:                    ;Control will come back here from
                                                        ;  the user's watchdog routine
08C5 C3CB08             JMP     RTC$Dog$Not$Set         ;Behave as though watchdog not active

                RTC$Dog$NZ:
08C8 22BF00             SHLD    RTC$Watchdog$Count      ;Save downdated count
                RTC$Dog$Not$Set:                        ;  (Leaves count unchanged)
08CB 3E20               MVI     A,IC$EOI                ;Reset the interrupt controller chip
08CD D3D8               OUT     IC$OCW2$Port

08CF D1                 POP     D                       ;Restore registers from local stack
08D0 C1                 POP     B
08D1 2A8422             LHLD    PI$User$Stack           ;Switch back to user's stack
08D4 F9                 SPHL
08D5 2A8622             LHLD    PI$User$HL              ;Recover user's registers
08D8 F1                 POP     PSW
08D9 FB                 EI                              ;Re-enable interrupts
08DA C9                 RET
                ;
                ;#
                ;
                ;       Shift HL Right one bit
                ;
                SHLR:
08DB B7                 ORA     A               ;Clear carry
08DC 7C                 MOV     A,H             ;Get MS byte
08DD 1F                 RAR                     ;Bit 7 set from previous carry
                                                ;Bit 0 goes into carry
08DE 67                 MOV     H,A             ;Put shifted MS byte back
08DF 7D                 MOV     A,L             ;Get LS byte
08E0 1F                 RAR                     ;Bit 7 = bit 0 of MS byte
08E1 6F                 MOV     L,A             ;Put back into result
08E2 C9                 RET

                ;
                ;#
```

Figure 8-10. (Continued)

```
                ;       High level diskette drivers
                ;
                ;       These drivers perform the following functions:
                ;
                ;       SELDSK  Select a specified disk and return the address of
                ;               the appropriate disk parameter header
                ;       SETTRK  Set the track number for the next read or write
                ;       SETSEC  Set the sector number for the next read or write
                ;       SETDMA  Set the DMA (read/write) address for the next read or write
                ;       SECTRAN Translate a logical sector number into a physical
                ;       HOME    Set the track to 0 so that the next read or write will
                ;               be on Track 0
                ;
                ;       In addition, the high level drivers are responsible for making
                ;       the 5 1/4" floppy diskettes that use a 512-byte sector appear
                ;       to CP/M as though they used a 128-byte sector. They do this
                ;       by using blocking/deblocking code. This blocking/deblocking
                ;       code is described in more detail later in this listing,
                ;       just prior to the code itself.
                ;
                ;
                ;
                ;       Disk parameter tables
                ;
                ;       As discussed in Chapter 3, these describe the physical
                ;       characteristics of the disk drives. In this example BIOS,
                ;       there are two types of disk drives: standard single-sided,
                ;       single-density 8", and double-sided, double-density 5 1/4"
                ;       mini-diskettes.
                ;
                ;       The standard 8" diskettes do not need to use the blocking/
                ;       deblocking code, but the 5 1/4" drives do. Therefore an additional
                ;       byte has been prefixed onto the disk parameter block to
                ;       tell the disk drivers what each logical disk's physical
                ;       diskette type is, and whether or not it needs deblocking.
                ;
                ;
                ;       Disk definition tables
                ;
                ;       These consist of disk parameter headers, with one entry
                ;       per logical disk driver, and disk parameter blocks with
                ;       either one parameter block per logical disk, or the same
                ;       parameter block for several logical disks.
                ;
                ;#
                ;
                Disk$Parameter$Headers:                 ;Described in Chapter 3
                ;
                                        ;Logical disk A: (5 1/4" diskette)
08E3 AE09               DW      Floppy$5$Skewtable      ;5 1/4" skew table
08E5 000000000000       DW      0,0,0                   ;Reserved for CP/M
08EB B022               DW      Directory$Buffer
08ED 3409               DW      Floppy$5$Parameter$Block
08EF B023               DW      Disk$A$Workarea
08F1 1024               DW      Disk$A$Allocation$Vector
                ;
                                        ;Logical disk B: (5 1/4" diskette)
08F3 AE09               DW      Floppy$5$Skewtable      ;Shares same skew table as A:
08F5 000000000000       DW      0,0,0                   ;Reserved for CP/M
08FB B022               DW      Directory$Buffer        ;Shares same buffer as A:
08FD 3409               DW      Floppy$5$Parameter$Block        ;Same DPB as A:
08FF D023               DW      Disk$B$Workarea                 ;Private work area
0901 2624               DW      Disk$B$Allocation$Vector        ;Private allocation vector
                ;
                                        ;Logical disk C: (8" floppy)
0903 F609               DW      Floppy$8$Skewtable      ;8" skew table
0905 000000000000       DW      0,0,0                   ;Reserved for CP/M
090B B022               DW      Directory$Buffer        ;Shares same buffer as A:
090D 4409               DW      Floppy$8$Parameter$Block
090F F023               DW      Disk$C$Workarea                 ;Private work area
0911 3C24               DW      Disk$C$Allocation$Vector        ;Private allocation vector
                ;
                                        ;Logical disk D: (8" floppy)
0913 AE09               DW      Floppy$5$Skewtable      ;Shares same skew table as A:
0915 000000000000       DW      0,0,0                   ;Reserved for CP/M
091B B022               DW      Directory$Buffer        ;Shares same buffer as A:
```

Figure 8-10. (Continued)

```
091D 4409                  DW      Floppy$8$Parameter$Block        ;Same DPB as C:
091F 0024                  DW      Disk$D$Workarea                 ;Private work area
0921 5B24                  DW      Disk$D$Allocation$Vector        ;Private allocation vector

                                   ;Logical disk M: (memory disk)
                   M$Disk$DPH:
0923 0000                  DW      0                             ;No skew required
0925 000000000000          DW      0,0,0                         ;Reserved for CP/M
092B B022                  DW      Directory$Buffer
092D 5409                  DW      M$Disk$Parameter$Block
092F 0000                  DW      0                             ;Disk cannot be changed, therefore
                                                                 ;  no work area is required
0931 7A24                  DW      M$Disk$Allocation$Vector
                   ;
                   ;
                   ;       Equates for disk parameter block
                   ;
                   ;       Disk Types
                   ;
0001 =             Floppy$5        EQU     1           ;5 1/4" mini floppy
0002 =             Floppy$8        EQU     2           ;8" floppy (SS SD)
0003 =             M$Disk          EQU     3           ;Memory disk
                   ;
                   ;       Blocking/deblocking indicator
                   ;
0080 =             Need$Deblocking EQU     1000$0000B      ;Sector size > 128 bytes
                   ;
                   ;#
                   ;
                   ;       Disk parameter blocks
                   ;
                   ;       5 1/4" mini floppy
                   ;
                                                       ;Extra byte prefixed to indicate
                                                       ;  disk type and blocking required
0933 81                    DB      Floppy$5 + Need$Deblocking
                                                       ;The parameter block has been amended
                                                       ;  to reflect the new layout of one
                                                       ;  track per diskette side, rather
                                                       ;  than viewing one track as both
                                                       ;  sides on a given head position.
                                                       ;It has also been adjusted to reflect
                                                       ;  one "new" track more being used for
                                                       ;  the CP/M image, with the resulting
                                                       ;  change in the number of allocation
                                                       ;  blocks and the number of reserved
                                                       ;  tracks.
                   Floppy$5$Parameter$Block:
0934 2400                  DW      36                  ;128-byte sectors per track
0936 04                    DB      4                   ;Block shift
0937 0F                    DB      15                  ;Block mask
0938 01                    DB      1                   ;Extent mask
0939 AB00                  DW      171                 ;Maximum allocation block number
093B 7F00                  DW      127                 ;Number of directory entries - 1
093D C0                    DB      1100$0000B          ;Bit map for reserving 1 alloc. block
093E 00                    DB      0000$0000B          ;  for file directory
093F 2000                  DW      32                  ;Disk-changed work area size
0941 0300                  DW      3                   ;Number of tracks before directory
                   ;
                   ;
                   ;       Standard 8" Floppy
                                                       ;Extra byte prefixed to DPB for
                                                       ; this version of the BIOS
0943 02                    DB      Floppy$8            ;Indicates disk type and the fact
                                                       ; that no deblocking is required
                   Floppy$8$Parameter$Block:
0944 1A00                  DW      26                  ;Sectors per track
0946 03                    DB      3                   ;Block shift
0947 07                    DB      7                   ;Block mask
0948 00                    DB      0                   ;Extent mask
0949 F200                  DW      242                 ;Maximum allocation block number
094B 3F00                  DW      63                  ;Number of directory entries - 1
094D C0                    DB      1100$0000B          ;Bit map for reserving 2 alloc. blocks
094E 00                    DB      0000$0000B          ;  for file directory
094F 1000                  DW      16                  ;Disk-changed work area size
0951 0200                  DW      2                   ;Number of tracks before directory
```

Figure 8-10. (Continued)

```
                 ;
                 ;        M$Disk
                 ;
                                                 ;The M$Disk presumes that 4 x 48K memory
                                                 ;  banks are available. The following
                                                 ;  table describes the disk as having
                                                 ;  8 tracks: two tracks per memory bank
                                                 ;  with each track having 192 128-byte
                                                 ;  sectors.
                                                 ;  The track number divided by 2 will be
                                                 ;  used to select the bank
0953 03                   DB      M$Disk         ;Type is M$Disk, no deblocking
                 M$Disk$Parameter$Block:
0954 C000                 DW      192            ;Sectors per "track". Each track is
                                                 ;  24K of memory
0956 03                   DB      3              ;Block shift (1024 byte allocation)
0957 07                   DB      7              ;Block mask
0958 00                   DB      0              ;Extent mask
0959 C000                 DW      192            ;Maximum allocation block number
095B 3F00                 DW      63             ;Number of directory entries -1
095D C0                   DB      1100$0000B     ;Bit map for reserving 2 allocation blocks
095E 00                   DB      0000$0000B     ;  for file directory
095F 0000                 DW      0              ;Disk cannot be changed, therefore no
                                                 ;  work area
0961 0000                 DW      0              ;No reserved tracks
                 ;
0004 =           Number$of$Logical$Disks         EQU     4
                 ;
                 ;#
                 ;
                 SELDSK:          ;Select disk in register C
                                         ;C = 0 for drive A, 1 for B, etc.
                                         ;Return the address of the appropriate
                                         ;  disk parameter header in HL, or 0000H
                                         ;  if the selected disk does not exist.
                                         ;
0963 210000               LXI     H,0            ;Assume an error
0966 79                   MOV     A,C            ;Check if requested disk valid

0967 FE0C                 CPI     'M' - 'A'      ;Check if memory disk
0969 CA9509               JZ      SELDSK$M$Disk  ;Yes

096C FE04                 CPI     Number$of$Logical$Disks
096E D0                   RNC                    ;Return if > maximum number of disks
                 ;
096F 322D0A               STA     Selected$Disk  ;Save selected disk number
                                                 ;Set up to return DPH address
0972 6F                   MOV     L,A            ;Make disk into word value
0973 2600                 MVI     H,0
                                                 ;Compute offset down disk parameter
                                                 ; header table by multiplying by
                                                 ; parameter header length (16 bytes)
0975 29                   DAD     H              ;*2
0976 29                   DAD     H              ;*4
0977 29                   DAD     H              ;*8
0978 29                   DAD     H              ;*16
0979 11E308               LXI     D,Disk$Parameter$Headers      ;Get base address
097C 19                   DAD     D              ;DE -> appropriate DPH
097D E5                   PUSH    H              ;Save DPH address
                 ;
                                                 ;Access disk parameter block to
                                                 ;  extract special prefix byte that
                                                 ;  identifies disk type and whether
                                                 ;  deblocking is required
                                                 ;
097E 110A00               LXI     D,10           ;Get DPB pointer offset in DPH
0981 19                   DAD     D              ;DE -> DPB address in DPH
0982 5E                   MOV     E,M            ;Get DPB address in DE
0983 23                   INX     H
0984 56                   MOV     D,M
0985 EB                   XCHG                   ;DE -> DPB

                 SELDSK$Set$Disk$Type:
0986 2B                   DCX     H              ;DE -> prefix byte
0987 7E                   MOV     A,M            ;Get prefix byte
0988 E60F                 ANI     0FH            ;Isolate disk type
```

Figure 8-10. (Continued)

```
098A 32360A              STA     Selected$Disk$Type      ;Save for use in low level driver
098D 7E                  MOV     A,M             ;Get another copy of prefix byte
098E E680                ANI     Need$Deblocking         ;Isolate deblocking flag
0990 32350A              STA     Selected$Disk$Deblock   ;Save for use in low level driver
0993 E1                  POP     H               ;Recover DPH pointer
0994 C9                  RET
                 ;
                 SELDSK$M$Disk:                          ;M$Disk selected
0995 212309              LXI     H,M$Disk$DPH            ;Return correct parameter header
0998 C38609              JMP     SELDSK$Set$Disk$Type    ;Resume normal processing
                 ;
                 ;#
                 ;
                 ;       Set logical track for next read or write
                 ;
                 SETTRK:
099B 60                  MOV     H,B             ;Selected track in BC on entry
099C 69                  MOV     L,C
099D 222E0A              SHLD    Selected$Track  ;Save for low level driver
09A0 C9                  RET
                 ;
                 ;#
                 ;
                 ;       Set logical sector for next read or write
                 ;
                 ;
                 SETSEC:                         ;Logical sector in C on entry
09A1 79                  MOV     A,C
09A2 32300A              STA     Selected$Sector ;Save for low level driver
09A5 C9                  RET
                 ;
                 ;#
                 ;
                 ;       Set disk DMA (Input/Output) address for next read or write
                 ;
09A6 0000        DMA$Address:    DW      0               ;DMA address
                 ;
                 SETDMA:                         ;Address in BC on entry
09A8 69                  MOV     L,C             ;Move to HL to save
09A9 60                  MOV     H,B
09AA 22A609              SHLD    DMA$Address     ;Save for low level driver
09AD C9                  RET
                 ;
                 ;#
                 ;
                 ;       Translate logical sector number to physical
                 ;
                 ;       Sector translation tables
                 ;       These tables are indexed using the logical sector number,
                 ;       and contain the corresponding physical sector number.
                 ;
                 Floppy$5$Skewtable:     ;Each physical sector contains four
                                         ;128-byte sectors.
                         ;       Physical 128b   Logical 128b       Physical 512-byte
09AE 00010203            DB      00,01,02,03     ;00,01,02,03            0  )
09B2 10111213            DB      16,17,18,19     ;04,05,06,07            4  )
09B6 20212223            DB      32,33,34,35     ;08,09,10,11            8  )
09BA 0C0D0E0F            DB      12,13,14,15     ;12,13,14,15            3  ) Head
09BE 1C1D1E1F            DB      28,29,30,31     ;16,17,18,19            7  )  0
09C2 08090A0B            DB      08,09,10,11     ;20,21,22,23            2  )
09C6 18191A1B            DB      24,25,26,27     ;24,25,26,27            6  )
09CA 04050607            DB      04,05,06,07     ;28,29,30,31            1  )
09CE 14151617            DB      20,21,22,23     ;32,33,34,35            5  )
                 ;
09D2 24252627            DB      36,37,38,39     ;36,37,38,39            0  ]
09D6 34353637            DB      52,53,54,55     ;40,41,42,43            4  ]
09DA 44454647            DB      68,69,70,71     ;44,45,46,47            8  ]
09DE 30313233            DB      48,49,50,51     ;48,49,50,51            3  ] Head
09E2 40414243            DB      64,65,66,67     ;52,53,54,55            7  ]  1
09E6 2C2D2E2F            DB      44,45,46,47     ;56,57,58,59            2  ]
09EA 3C3D3E3F            DB      60,61,62,63     ;60,61,62,63            6  ]
09EE 28292A2B            DB      40,41,42,43     ;64,65,66,67            1  ]
09F2 38393A3B            DB      56,57,58,59     ;68,69,70,71            5  ]
                 ;
                 ;
                 Floppy$8$Skewtable:     ;Standard 8" Driver
```

Figure 8-10. (Continued)

```
                         ;      01,02,03,04,05,06,07,08,09,10    Logical sectors
09F6 01070D1319          DB     01,07,13,19,25,05,11,17,23,03    ;Physical sectors
                         ;
                         ;      11,12,13,14,15,16,17,18,19,20    Logical sectors
0A00 090F150208          DB     09,15,21,02,08,14,20,26,06,12    ;Physical sectors
                         ;
                         ;      21,22,23,24,25,26        Logical sectors
0A0A 1218040A10          DB     18,24,04,10,16,22        ;Physical sectors
              ;#
              ;
              SECTRAN:                       ;Translate logical sector into physical
                                             ;On entry, BC = logical sector number
                                             ;          DE -> appropriate skew table
                                             ;
                                             ;on exit, HL = physical sector number
0A10 EB                  XCHG                        ;HL -> skew table base
0A11 09                  DAD    B                    ;Add on logical sector number
0A12 6E                  MOV    L,M                  ;Get physical sector number
0A13 2600                MVI    H,0                  ;Make into a 16-bit value
0A15 C9                  RET
              ;
              ;#
              ;
              ;
              HOME:                          ;Home the selected logical disk to track 0
                                             ;Before doing this, a check must be made to see
                                             ;  if the physical disk buffer has information in
                                             ;  it that must be written out. This is indicated by
                                             ;  a flag, Must$Write$Buffer, that is set in the
                                             ;  deblocking code.
                                             ;
0A16 3A2C0A              LDA    Must$Write$Buffer    ;Check if physical buffer must
0A19 B7                  ORA    A                    ;  be written to a disk
0A1A C2200A              JNZ    HOME$No$Write
0A1D 322B0A              STA    Data$In$Disk$Buffer  ;No, so indicate that buffer
                                                     ;  is now unoccupied
              HOME$No$Write:
0A20 0E00                MVI    C,0                  ;Set to track 0 (logically,
0A22 CD9B09              CALL   SETTRK               ;  no actual disk operation occurs)
0A25 C9                  RET

              ;
              ;#
              ;       Data written to or read from the mini-floppy drive is transferred
              ;       via a physical buffer that is one complete track in length,
              ;       9 * 512 bytes. It is declared at the end of the BIOS, and has
              ;       some small amount of initialization code "hidden" in it.
              ;
              ;       The blocking/deblocking code attempts to minimize the amount
              ;       of actual disk I/O by storing the disk and track
              ;       currently residing in the physical buffer.
              ;       If a read request occurs of a 128-byte CP/M "sector"
              ;       that already is in the physical buffer, no disk access occurs
              ;       If a write request occurs if and the 128-byte CP/M 'sector'
              ;       is already in the physical buffer, no disk access will occur,
              ;       UNLESS the BDOS indicates that it is writing to the directory.
              ;       Directory writes cause an immediate write to disk of the entire
              ;       track in the physical buffer.
              ;
              ;
0800 =        Allocation$Block$Size    EQU     2048
0009 =        Physical$Sec$Per$Track   EQU     9       ;Adjusted to reflect a "new"
                                                       ;  track is only one side of the
                                                       ;  disk
0200 =        Physical$Sector$Size     EQU     512     ;This is the actual sector size
                                                       ;  for the 5 1/4" mini-floppy diskettes
                                                       ;The 8" diskettes and memory disk
                                                       ;  use 128-byte sectors
                                                       ;Declare the physical disk buffer for the
                                                       ;  5 1/4" diskettes
0004 =        CPM$Sec$Per$Physical     EQU     Physical$Sector$Size/128
0024 =        CPM$Sec$Per$Track        EQU     CPM$Sec$Per$Physical*Physical$Sec$Per$Track
1200 =        Bytes$Per$Track          EQU     Physical$Sec$Per$Track*Physical$Sector$Size
0003 =        Sector$Mask              EQU     CPM$Sec$Per$Physical-1
0002 =        Sector$Bit$Shift         EQU     2               ;LOG2(CPM$Sec$Per$Physical)
```

Figure 8-10. (Continued)

```
                   ;
                                          ;These are the values handed over by the BDOS
                                          ;  when it calls the write operation.
                                          ;The allocated/unallocated indicates whether the
                                          ;  BDOS wishes to write to an unallocated allocation
                                          ;  block (it only indicates this for the first
                                          ;  128-byte sector write), or to an allocation block
                                          ;  that has already been allocated to a file.
                                          ;The BDOS also indicates if it wishes to write to
                                          ;  the file directory.
                                          ;
0000 =             Write$Allocated        EQU     0
0001 =             Write$Directory        EQU     1
0002 =             Write$Unallocated      EQU     2       ;<== ignored for track buffering
                   ;
0A26 00            Write$Type:            DB      0       ;Contains the type of write
                                                          ;  indicated by the BDOS
                   ;
                   ;
                   In$Buffer$Dk$Trk:              ;Variables for physical sector currently
                                                  ;  in Disk$Buffer in memory
0A27 00            In$Buffer$Disk:        DB      0       ;) These are moved and compared
0A28 0000          In$Buffer$Track:       DW      0       ;) as a group, so do not alter
                                                          ;  these lines
0A2A 00            In$Buffer$Disk$Type:   DB      0       ;Disk type for sector in buffer
                   ;
0A2B 00            Data$In$Disk$Buffer:   DB      0       ;When nonzero, the disk buffer has
                                                          ;  data from the disk in it
0A2C 00            Must$Write$Buffer:     DB      0       ;Nonzero when data has been written
                                                          ;  into Disk$Buffer but not yet
                                                          ;  written out to disk
                   ;
                   Selected$Dk$Trk:               ;Variables for selected disk, track and sector
                                                  ; (Selected by SELDSK, SETTRK and SETSEC)
0A2D 00            Selected$Disk:         DB      0       ;) These are moved and compared
0A2E 0000          Selected$Track:        DW      0       ;) as a group so do not alter order

0A30 00            Selected$Sector:       DB      0       ;Not part of group but needed here

0A31 00            Selected$Physical$Sector: DB   0       ;Selected physical sector derived
                                                          ;  from selected (CP/M) sector by
                                                          ;  shifting it right the number of
                                                          ;  bits specified by Sector$Bit$Shift

                   ;

0A32 00            Disk$Error$Flag:       DB      0       ;Nonzero to indicate an error
                                                          ;  that could not be recovered
                                                          ;  by the disk drivers. The BDOS
                                                          ;  will output a "Bad Sector" message
0A33 00            Disk$Hung$Flag:        DB      0       ;Nonzero if a watchdog timeout
                                                          ;  occurs
0258 =             Disk$Timer             EQU     600     ;Number of 16.66 ms clock ticks
                                                          ;  for a 10 second timeout
                   ;
                                                  ;Flags used inside the deblocking code

0A34 00            Read$Operation:        DB      0       ;Nonzero when a CP/M 128-byte
                                                          ;  sector is to be read
0A35 00            Selected$Disk$Deblock: DB      0       ;Nonzero when the selected disk
                                                          ;  needs deblocking (set in SELDSK)
0A36 00            Selected$Disk$Type:    DB      0       ;Indicates 8" or 5 1/4" floppy or
                                                          ;  M$Disk selected. (set in SELDSK)
                   ;
                   ;#
                   ;
                   ;       Read in the 128-byte CP/M sector specified by previous calls
                   ;       to Select Disk, Set Track and Sector. The sector will be read
                   ;       into the address specified in the previous Set DMA Address call.
                   ;
                   ;       If reading from a disk drive using sectors larger than 128 bytes,
                   ;       deblocking code will be used to "unpack" a 128-byte sector from
                   ;       the physical sector.
                   READ:
0A37 3A350A                LDA     Selected$Disk$Deblock  ;Check if deblocking needed
0A3A B7                    ORA     A                      ;  (flag was set in SELDSK call)
```

Figure 8-10. (Continued)

```
0A3B CA2F0B            JZ      Read$No$Deblock        ;No, use normal nondeblocked
                                               ;The deblocking algorithm used is such
                                               ; that a read operation can be viewed
                                               ; until the actual data transfer as though
                                               ; it was the first write to an unallocated
                                               ; allocation block
0A3E 3E01              MVI     A,1                    ;Indicate that a read actually
0A40 32340A            STA     Read$Operation         ; is to be performed
0A43 3E00              MVI     A,Write$Allocated      ;Fake deblocking code into believing
0A45 32260A            STA     Write$Type             ; that this is a write to an
                                                      ; allocated allocation block
0A48 C35C0A            JMP     Perform$Read$Write     ;Use common code to execute read
                ;
                ;#
                ;      Write a 128-byte sector from the current DMA address to
                ;      the previously selected disk, track and sector.
                ;
                ;      On arrival here, the BDOS will have set register C to indicate
                ;      whether this write operation is to an already allocated allocation
                ;      block (which means a preread of the sector may be needed), or
                ;      to the directory (in which case the data will be written to the
                ;      disk immediately).
                ;
                ;      Only writes to the directory take place immediately. In all other
                ;      cases, the data will be moved from the DMA address into the disk
                ;      buffer, and only be written out when circumstances force the
                ;      transfer. The number of physical disk operations can therefore
                ;      be reduced considerably.
                ;
                WRITE:
0A4B 3A350A            LDA     Selected$Disk$Deblock  ;Check if deblocking is required
0A4E B7                ORA     A                      ; (flag set in SELDSK call)
0A4F CA2A0B            JZ      Write$No$Deblock
0A52 AF                XRA     A                      ;Indicate that a write operation
0A53 32340A            STA     Read$Operation         ; is required (i.e NOT a read)
0A56 79                MOV     A,C                    ;Save the BDOS write type
0A57 E601              ANI     1                      ; but only distinguish between
                                                      ; write to allocated block or
0A59 32260A            STA     Write$Type             ; directory write
                ;
                ;
                ;#
                ;
                Perform$Read$Write:     ;Common code to execute both reads and
                                        ; writes of 128-byte sectors.
0A5C AF                XRA     A              ;Assume that no disk errors will
0A5D 32320A            STA     Disk$Error$Flag ; occur
0A60 3A300A            LDA     Selected$Sector ;Convert selected 128-byte sector
0A63 1F                RAR                     ; into physical sector by dividing by 4
0A64 1F                RAR
0A65 E63F              ANI     3FH             ;Remove any unwanted bits
0A67 32310A            STA     Selected$Physical$Sector
                                                      ;
0A6A 212B0A            LXI     H,Data$In$Disk$Buffer  ;Check if disk buffer already has
0A6D 7E                MOV     A,M                    ; data in it
0A6E 3601              MVI     M,1                    ;(Unconditionally indicate that
                                                      ; the buffer now has data in it)
0A70 B7                ORA     A                      ;Did it indeed have data in it?
0A71 CA870A            JZ      Read$Track$into$Buffer ;No, proceed to read a physical
                                                      ; track into the buffer
                                               ;
                                               ;The buffer does have a physical track
                                               ; in it. Check if it is the right one
                                               ;
0A74 11270A            LXI     D,In$Buffer$Dk$Trk     ;Check if track in buffer is the
0A77 212D0A            LXI     H,Selected$Dk$Trk      ; same as that selected earlier
0A7A CDE10A            CALL    Compare$Dk$Trk         ;Compare ONLY disk and track
0A7D CA910A            JZ      Track$In$Buffer        ;Yes, it is already in buffer
                                                      ;No, it will have to be read in
                                                      ; over current contents of buffer
0A80 3A2C0A            LDA     Must$Write$Buffer      ;Check if buffer has data in that
```

Figure 8-10. (Continued)

```
0A83 B7                 ORA     A                       ;  must be written out first
0A84 C4E50B             CNZ     Write$Physical          ;Yes, write it out
                ;
                Read$Track$into$Buffer:
0A87 CDCE0A             CALL    Set$In$Buffer$Dk$Trk    ;Set in buffer variables from
                                                        ;  selected disk, track
                                                        ;  to reflect which track is in the
                                                        ;  buffer now
0A8A CDEA0B             CALL    Read$Physical           ;Read the track into the buffer
0A8D AF                 XRA     A                       ;Reset the flag to reflect buffer
0A8E 322C0A             STA     Must$Write$Buffer       ;  contents
                ;
                Track$In$Buffer:                ;Selected track and
                                                ;  disk is already in the buffer
                                                ;Convert the selected CP/M (128-byte)
                                                ;  sector into a relative address down
                                                ;  the buffer
0A91 3A300A             LDA     Selected$Sector ;Get selected sector number
0A94 6F                 MOV     L,A             ;Multiply by 128 by shifting 16-bit value
0A95 2600               MVI     H,0             ;left 7 bits
0A97 29                 DAD     H               ;* 2
0A98 29                 DAD     H               ;* 4
0A99 29                 DAD     H               ;* 8
0A9A 29                 DAD     H               ;* 16
0A9B 29                 DAD     H               ;* 32
0A9C 29                 DAD     H               ;* 64
0A9D 29                 DAD     H               ;* 128
                ;
0A9E 11A40F             LXI     D,Disk$Buffer   ;Get base address of disk buffer
0AA1 19                 DAD     D               ;Add on sector number * 128
                                                ;HL -> 128-byte sector number start
                                                ;  address in disk buffer
0AA2 EB                 XCHG                    ;DE -> sector in disk buffer
0AA3 2AA609             LHLD    DMA$Address     ;Get DMA address set in SETDMA call
0AA6 EB                 XCHG                    ;Assume a read operation, so
                                                ;  DE -> DMA address
                                                ;  HL -> sector in disk buffer
0AA7 0E10               MVI     C,128/8         ;Because of the faster method used
                                                ;  to move data in and out of the
                                                ;  disk buffer, (eight bytes moved per
                                                ;  loop iteration) the count need only
                                                ;  be 1/8 of normal
                                                ;At this point,
                                                ;  C = loop count
                                                ;  DE -> DMA address
                                                ;  HL -> sector in disk buffer
0AA9 3A340A             LDA     Read$Operation  ;Determine whether data is to be moved
0AAC B7                 ORA     A               ;  out of the buffer (read) or into the
0AAD C2B50A             JNZ     Buffer$Move     ;  buffer (write)
                                                ;Writing into buffer
                                                        ;(A must be 0 get here)
0AB0 3C                 INR     A                       ;Set flag to force a write
0AB1 322C0A             STA     Must$Write$Buffer       ;  of the disk buffer later on.
0AB4 EB                 XCHG                            ;Make DE -> sector in disk buffer
                                                        ;     HL -> DMA address
                ;
                ;
                Buffer$Move:
0AB5 CDF80A             CALL    Move$8          ;Moves 8 bytes * C times from (HL)
                                                ;  to (DE)

                                                ;
0AB8 3A260A             LDA     Write$Type      ;If write to directory, write out
0ABB FE01               CPI     Write$Directory ;  buffer immediately
0ABD 3A320A             LDA     Disk$Error$Flag ;Get error flag in case delayed write or read
0AC0 C0                 RNZ                     ;Return if delayed write or read
                                                ;
0AC1 B7                 ORA     A               ;Check if any disk errors have occured
0AC2 C0                 RNZ                     ;Yes, abandon attempt to write to directory
                                                ;
0AC3 AF                 XRA     A               ;Clear flag that indicates buffer must be
0AC4 322C0A             STA     Must$Write$Buffer       ;  written out
0AC7 CDE50B             CALL    Write$Physical  ;Write buffer out to physical track
0ACA 3A320A             LDA     Disk$Error$Flag ;Return error flag to caller
0ACD C9                 RET
                ;
```

Figure 8-10. (Continued)

```
                 ;
                 ;
                 Set$In$Buffer$Dk$Trk:                        ;Indicate selected disk, track
                                                              ;  now residing in buffer
0ACE 3A2D0A              LDA     Selected$Disk
0AD1 32270A              STA     In$Buffer$Disk

0AD4 2A2E0A              LHLD    Selected$Track
0AD7 22280A              SHLD    In$Buffer$Track

0ADA 3A360A              LDA     Selected$Disk$Type           ;Also reflect disk type
0ADD 322A0A              STA     In$Buffer$Disk$Type

0AE0 C9                  RET
                 ;
                 ;
                 Compare$Dk$Trk:                      ;Compares just the disk and track
                                                      ;  pointed to by DE and HL
0AE1 0E03                MVI     C,3                  ;Disk (1), track (2)
                 Compare$Dk$Trk$Loop:
0AE3 1A                  LDAX    D                    ;Get comparitor
0AE4 BE                  CMP     M                    ;Compare with comparand
0AE5 C0                  RNZ                          ;Abandon comparison if inequality found
0AE6 13                  INX     D                    ;Update comparitor pointer
0AE7 23                  INX     H                    ;Update comparand pointer
0AE8 0D                  DCR     C                    ;Count down on loop count
0AE9 C8                  RZ                           ;Return (with zero flag set)
0AEA C3E30A              JMP     Compare$Dk$Trk$Loop
                 ;
                 ;
                 Move$Dk$Trk:                         ;Moves the disk, track
                                                      ;  variables pointed at by HL to
                                                      ;  those pointed at by DE
0AED 0E03                MVI     C,3                  ;Disk (1), Track (2)
                 Move$Dk$Trk$Loop:
0AEF 7E                  MOV     A,M                  ;Get source byte
0AF0 12                  STAX    D                    ;Store in destination
0AF1 13                  INX     D                    ;Update pointers
0AF2 23                  INX     H
0AF3 0D                  DCR     C                    ;Count down on byte count
0AF4 C8                  RZ                           ;Return if all bytes moved
0AF5 C3EF0A              JMP     Move$Dk$Trk$Loop
                 ;
                 ;#
                 ;
                 ;       Move eight bytes
                 ;
                 ;       This routine moves eight bytes in a block, C times, from
                 ;       (HL) to (DE). It uses "drop through" coding to speed
                 ;       up execution.
                 ;
                 ;       Entry Parameters
                 ;
                 ;               C = number of 8-byte blocks to move
                 ;               DE -> destination address
                 ;               HL -> source address
                 ;
                 Move$8:
0AF8 7E                  MOV     A,M                  ;Get byte from source
0AF9 12                  STAX    D                    ;Put into destination
0AFA 13                  INX     D                    ;Update pointers
0AFB 23                  INX     H
0AFC 7E                  MOV     A,M                  ;Get byte from source
0AFD 12                  STAX    D                    ;Put into destination
0AFE 13                  INX     D                    ;Update pointers
0AFF 23                  INX     H
0B00 7E                  MOV     A,M                  ;Get byte from source
0B01 12                  STAX    D                    ;Put into destination
0B02 13                  INX     D                    ;Update pointers
0B03 23                  INX     H
0B04 7E                  MOV     A,M                  ;Get byte from source
0B05 12                  STAX    D                    ;Put into destination
0B06 13                  INX     D                    ;Update pointers
0B07 23                  INX     H
0B08 7E                  MOV     A,M                  ;Get byte from source
0B09 12                  STAX    D                    ;Put into destination
```

Figure 8-10. (Continued)

```
0B0A 13                 INX     D               ;Update pointers
0B0B 23                 INX     H
0B0C 7E                 MOV     A,M             ;Get byte from source
0B0D 12                 STAX    D               ;Put into destination
0B0E 13                 INX     D               ;Update pointers
0B0F 23                 INX     H
0B10 7E                 MOV     A,M             ;Get byte from source
0B11 12                 STAX    D               ;Put into destination
0B12 13                 INX     D               ;Update pointers
0B13 23                 INX     H
0B14 7E                 MOV     A,M             ;Get byte from source
0B15 12                 STAX    D               ;Put into destination
0B16 13                 INX     D               ;Update pointers
0B17 23                 INX     H

0B18 0D                 DCR     C               ;Count down on loop counter
0B19 C2F80A             JNZ     Move$8          ;Repeat until done
0B1C C9                 RET

                ;
                ;#
                ;
                ;       Introduction to the disk controllers on this computer system
                ;
                ;       There are two "smart" disk controllers on this system, one
                ;       for the 8" floppy diskette drives, and one for the 5 1/4"
                ;       mini-diskette drives.
                ;
                ;       The controllers are "hard-wired" to monitor certain locations
                ;       in memory to detect when they are to perform some disk
                ;       operation. The 8" controller looks at location 0040H, and
                ;       the 5 1/4" controller looks at location 0045H. These are
                ;       called their disk control bytes. If the most significant
                ;       bit of a disk control byte is set, the controller will then
                ;       look at the word following the respective control bytes.
                ;       This word must contain the address of a valid disk control
                ;       table that specifies the exact disk operation to be performed.
                ;
                ;       Once the operation has been completed, the controller resets
                ;       its disk control byte to 00H, and this indicates completion
                ;       to the disk driver code.
                ;
                ;       The controller also sets a return code in a disk status block.
                ;       Both controllers use the same location (0043H) for this.
                ;       If the first byte of this status block is less than 80H, then
                ;       a disk error has occurred. For this simple BIOS, no further details
                ;       of the status settings are relevant. Note that the disk controller
                ;       has built-in retry logic; reads and writes are attempted ten
                ;       times before the controller returns an error.
                ;
                ;       The disk control table layout is shown below. Note that the
                ;       controllers have the capability for control tables to be
                ;       chained together so that a sequence of disk operations can
                ;       be initiated. In this BIOS this feature is not used. However,
                ;       the controller requires that the chain pointers in the
                ;       disk control tables be pointed back to the main control bytes
                ;       in order to indicate the end of the chain.
                ;
0040 =          Disk$Control$8          EQU     40H     ;8" control byte
0041 =          Command$Block$8         EQU     41H     ;Control table pointer
                ;
0043 =          Disk$Status$Block       EQU     43H     ;8" AND 5 1/4" status block
                ;
0045 =          Disk$Control$5          EQU     45H     ;5 1/4" control byte
0046 =          Command$Block$5         EQU     46H     ;Control table pointer
                ;
                ;
                ;       Floppy Disk Control Tables
                ;
0B1D 00         Floppy$Command:                 DB      0       ;Command
0001 =          Floppy$Read$Code                EQU     01H
0002 =          Floppy$Write$Code               EQU     02H
0B1E 00         Floppy$Unit:                    DB      0       ;Unit (drive) number = 0 or 1
0B1F 00         Floppy$Head:                    DB      0       ;Head number = 0 or 1
0B20 00         Floppy$Track:                   DB      0       ;Track number
0B21 00         Floppy$Sector:                  DB      0       ;Sector number
```

Figure 8-10. (Continued)

```
0B22 0000        Floppy$Byte$Count:              DW      0       ;Number of bytes to read/write
0B24 0000        Floppy$DMA$Address:             DW      0       ;Transfer address
0B26 0000        Floppy$Next$Status$Block:       DW      0       ;Pointer to next status block
                                                                 ;  if commands are chained.
0B28 0000        Floppy$Next$Control$Location:   DW      0       ;Pointer to next control byte
                                                                 ;  if commands are chained
                 ;
                 ;#
                 ;
                 ;
                 Write$No$Deblock:               ;Write contents of disk buffer to
                                                 ;  correct sector
0B2A 3E02                MVI     A,Floppy$Write$Code     ;Get write function code
0B2C C3310B              JMP     Common$No$Deblock       ;Go to common code
                 Read$No$Deblock:                ;Read previously selected sector
                                                 ;  into disk buffer.
0B2F 3E01                MVI     A,Floppy$Read$Code      ;Get read function code
                 Common$No$Deblock:
0B31 321D0B              STA     Floppy$Command  ;Set command function code
                                                 ;Set up nondeblocked command table

0B34 3A360A              LDA     Selected$Disk$Type      ;Check if memory disk operation
0B37 FE03                CPI     M$Disk
0B39 CA7A0B              JZ      M$Disk$Transfer ;Yes, it is M$Disk

                 No$Deblock$Retry:               ;Re-entry point to retry after error
0B3C 218000              LXI     H,128           ;Bytes per sector
0B3F 22220B              SHLD    Floppy$Byte$Count
0B42 AF                  XRA     A               ;8" floppy only has head 0
0B43 321F0B              STA     Floppy$Head
                                                 ;
0B46 3A2D0A              LDA     Selected$Disk   ;8" floppy controller only knows about
                                                 ;  units 0 and 1 so Selected$Disk must
                                                 ;  be converted
0B49 E601                ANI     01H             ;Turn into 0 or 1
0B4B 321E0B              STA     Floppy$Unit     ;Set unit number
                                                 ;
0B4E 3A2E0A              LDA     Selected$Track
0B51 32200B              STA     Floppy$Track    ;Set track number
                                                 ;
0B54 3A300A              LDA     Selected$Sector
0B57 32210B              STA     Floppy$Sector   ;Set sector number
                                                 ;
0B5A 2AA609              LHLD    DMA$Address     ;Transfer directly between DMA Address
0B5D 22240B              SHLD    Floppy$DMA$Address      ; and 8" controller.
                                                 ;
                                                 ;The disk controller can accept chained
                                                 ;  disk control tables, but in this case,
                                                 ;  they are not used, so the "Next" pointers
                                                 ;  must be pointed back at the initial
                                                 ;  control bytes in the base page.
0B60 214300              LXI     H,Disk$Status$Block             ;Point next status back at
0B63 22260B              SHLD    Floppy$Next$Status$Block        ;  main status block
                                                                 ;
0B66 214000              LXI     H,Disk$Control$8                ;Point next control byte
0B69 22280B              SHLD    Floppy$Next$Control$Location    ;  back at main control byte
                                                         ;
0B6C 211D0B              LXI     H,Floppy$Command        ;Point controller at control table
0B6F 224100              SHLD    Command$Block$8
                                                         ;
0B72 214000              LXI     H,Disk$Control$8        ;Activate controller to perform
0B75 3680                MVI     M,80H                   ;  operation
0B77 C33B0C              JMP     Wait$For$Disk$Complete

                 ;
                 ;#
                 ;       Memory disk driver
                 ;
                 ;       This routine must use an intermediary buffer, since the
                 ;       DMA address in bank ("track") 0 occupies the same
                 ;       place in the overall address space as the M$Disk itself.
                 ;       The M$Disk$Buffer is above the 48K mark, and therefore
                 ;       remains in the address space regardless of which bank/track
                 ;       is selected.
                 ;
                 ;
```

Figure 8-10. (Continued)

```
                ;       For writing, the 128-byte sector must be processed:
                ;
                ;               1. Move sector DMA$Address -> M$Disk$Buffer
                ;               2. Select correct track (+1 to get bank number)
                ;               3. Move sector M$Disk$Buffer -> M$Disk image
                ;               4. Select bank 0
                ;
                ;       For reading, the processing is:
                ;
                ;               1. Select correct track/bank
                ;               2. Move sector M$Disk image -> M$Disk$Buffer
                ;               3. Select Bank 0
                ;               4. Move sector M$Disk$Buffer -> DMA$Address
                ;
                ;       If there is any risk of any interrupt causing control
                ;       to be transferred to an address below 48K, interrupts must
                ;       be disabled when any bank other than 0 is selected.
                ;
                M$Disk$Transfer:
0B7A 3A300A             LDA     Selected$Sector ;Compute address in memory
0B7D 6F                 MOV     L,A             ;  by muliplying sector * 128
0B7E 2600               MVI     H,0
0B80 29                 DAD     H               ;* 2
0B81 29                 DAD     H               ;* 4
0B82 29                 DAD     H               ;* 8
0B83 29                 DAD     H               ;* 16
0B84 29                 DAD     H               ;* 32
0B85 29                 DAD     H               ;* 64
0B86 29                 DAD     H               ;* 128

0B87 3A2E0A             LDA     Selected$Track  ;Compute which half of bank sector
                                                ;  is in by using LS bit of track
0B8A 47                 MOV     B,A             ;Save copy for later
0B8B E601               ANI     1               ;Isolate lower/upper indicator
0B8D CA940B             JZ      M$Disk$Lower$Half

0B90 110060             LXI     D,(48 * 1024) / 2       ;Upper half,  so bias address
0B93 19                 DAD     D

                M$Disk$Lower$Half:                      ;HL -> sector in memory
0B94 78                 MOV     A,B             ;Recover selected track
0B95 1F                 RAR                     ;Divide by 2 to get bank number
0B96 3C                 INR     A               ;Bank 1 is first track
0B97 47                 MOV     B,A             ;Preserve for later use

0B98 3A1D0B             LDA     Floppy$Command  ;Check if reading or writing
0B9B FE02               CPI     Floppy$Write$Code
0B9D CABE0B             JZ      M$Disk$Write    ;Writing
                                                ;Reading

0BA0 CDDD0B             CALL    Select$Bank     ;Select correct memory bank
0BA3 113023             LXI     D,M$Disk$Buffer ;DE -> M$Disk$Buffer, HL -> M$Disk image
0BA6 0E10               MVI     C,128/8         ;Number of 8-byte blocks to move
0BA8 CDF80A             CALL    Move$8

0BAB 0600               MVI     B,0             ;Revert to normal memory bank
0BAD CDDD0B             CALL    Select$Bank

0BB0 2AA609             LHLD    DMA$Address     ;Get user's DMA address
0BB3 113023             LXI     D,M$Disk$Buffer
0BB6 EB                 XCHG                    ;DE -> User's DMA, HL -> M$Disk buffer
0BB7 0E10               MVI     C,128/8         ;Number of 8-byte blocks to move
0BB9 CDF80A             CALL    Move$8

0BBC AF                 XRA     A               ;Indicate no error
0BBD C9                 RET

                M$Disk$Write:                           ;Writing
0BBE E5                 PUSH    H               ;Save sector's address in M$Disk image
0BBF 2AA609             LHLD    DMA$Address     ;Move sector into M$Disk$Buffer
0BC2 113023             LXI     D,M$Disk$Buffer
0BC5 0E10               MVI     C,128/8         ;Number of 8-byte blocks to move
0BC7 CDF80A             CALL    Move$8          ;(Does not use B register)
                                                ;B = memory bank to select
0BCA CDDD0B             CALL    Select$Bank
```

Figure 8-10. (Continued)

```
0BCD D1                        POP     D               ;Recover sector's M$Disk image address
0BCE 213023                    LXI     H,M$Disk$Buffer
0BD1 0E10                      MVI     C,128/8
0BD3 CDF80A                    CALL    Move$8          ;Move into M$Disk image

0BD6 0600                      MVI     B,0             ;Select bank 0
0BD8 CDDD0B                    CALL    Select$Bank

0BDB AF                        XRA     A               ;Indicate no error
0BDC C9                        RET
                       ;
                       ;#
                       ;       Select bank
                       ;
                       ;       This routine switches in the required memory bank.
                       ;       Note that the hardware port that controls bank selection
                       ;       also has other bits in it. These are preserved across
                       ;       bank selections.
                       ;
                       ;       Entry parameter
                       ;
                       ;               B = bank number
                       ;
0040 =                 Bank$Control$Port       EQU     40H
00F8 =                 Bank$Mask               EQU     1111$1000B      ;To preserve other bits
                       ;
                       Select$Bank:
0BDD DB40                      IN      Bank$Control$Port       ;Get current setting in port
0BDF E6F8                      ANI     Bank$Mask               ;Preserve all other bits
0BE1 B0                        ORA     B                       ;Set bank code
0BE2 D340                      OUT     Bank$Control$Port       ;Select the bank
0BE4 C9                        RET
                       ;
                       ;#
                       ;

                       Write$Physical:                 ;Write contents of disk buffer to
                                                       ;  correct sector
0BE5 3E02                      MVI     A,Floppy$Write$Code     ;Get write function code
0BE7 C3EC0B                    JMP     Common$Physical ;Go to common code
                       Read$Physical:                  ;Read previously selected sector
                                                       ;  into disk buffer
0BEA 3E01                      MVI     A,Floppy$Read$Code      ;Get read function code
                       ;
                       Common$Physical:
0BEC 321D0B                    STA     Floppy$Command  ;Set command table

                       ;
                       Deblock$Retry:                  ;Re-entry point to retry after error
0BEF 3A2A0A                    LDA     In$Buffer$Disk$Type     ;Get disk type currently in buffer
0BF2 FE01                      CPI     Floppy$5                ;Confirm it is a 5 1/4" floppy
0BF4 CAFD0B                    JZ      Correct$Disk$Type       ;Yes
0BF7 3E01                      MVI     A,1                     ;No, indicate disk error
0BF9 32320A                    STA     Disk$Error$Flag
0BFC C9                        RET
                       Correct$Disk$Type:              ;Set up disk control table
                                                       ;
0BFD 3A270A                    LDA     In$Buffer$Disk  ;Convert disk number to 0 or 1
0C00 E601                      ANI     1               ;  for disk controller
0C02 321E0B                    STA     Floppy$Unit

0C05 2A280A                    LHLD    In$Buffer$Track ;Set up head and track number
0C08 7D                        MOV     A,L             ;Even numbered tracks will be on
0C09 E601                      ANI     1               ;  head 0, odd numbered on head 1
0C0B 321F0B                    STA     Floppy$Head     ;Set head number

0C0E 7D                        MOV     A,L             ;Note: this is single byte value
0C0F 1F                        RAR                     ; /2 for track (carry off from ANI above)
0C10 32200B                    STA     Floppy$Track

0C13 3E01                      MVI     A,1                     ;Start with sector 1 as a whole
0C15 32210B                    STA     Floppy$Sector           ;  track will be transferred
                                                               ;
0C18 210012                    LXI     H,Bytes$Per$Track       ;Set byte count for complete
0C1B 22220B                    SHLD    Floppy$Byte$Count       ;  track to be transferred
                                                               ;
```

Figure 8-10. (Continued)

```
0C1E 21A40F          LXI     H,Disk$Buffer           ;Set transfer address to be
0C21 22240B          SHLD    Floppy$DMA$Address      ;  disk buffer
                                                     ;
                                                     ;As only one control table is in
                                                     ;  use, close the status and busy
                                                     ;  chain pointers back to the
                                                     ;  main control bytes
0C24 214300          LXI     H,Disk$Status$Block
0C27 22260B          SHLD    Floppy$Next$Status$Block
0C2A 214500          LXI     H,Disk$Control$5
0C2D 22280B          SHLD    Floppy$Next$Control$Location
0C30 211D0B          LXI     H,Floppy$Command        ;Set up command block pointer
0C33 224600          SHLD    Command$Block$5
0C36 214500          LXI     H,Disk$Control$5        ;Activate 5 1/4" disk controller
0C39 3680            MVI     M,80H
                ;
                Wait$For$Disk$Complete:         ;Wait until disk status block indicates
                                                ;  operation has completed, then check
                                                ;  if any errors occurred.
                                                ;On entry HL -> disk control byte
0C3B AF              XRA     A                       ;Ensure hung flag clear
0C3C 32330A          STA     Disk$Hung$Flag
0C3F 21570C          LXI     H,Disk$Timed$Out        ;Set up watchdog timer
0C42 015802          LXI     B,Disk$Timer            ;Time delay
0C45 CD6D08          CALL    Set$Watchdog
                Disk$Wait$Loop:
0C48 7E              MOV     A,M                     ;Get control byte
0C49 B7              ORA     A
0C4A CA5D0C          JZ      Disk$Complete           ;Operation done
0C4D 3A330A          LDA     Disk$Hung$Flag          ;Also check if time expired
0C50 B7              ORA     A
0C51 C2B40D          JNZ     Disk$Error              ;Will be set to 40H
0C54 C3480C          JMP     Disk$Wait$Loop
                Disk$Timed$Out:                 ;Control arrives here from watchdog
                                                ;  routine itself -- so this is effectively
                                                ;  part of the interrupt service routine.
0C57 3E40            MVI     A,40H                   ;Set disk hung error code
0C59 32330A          STA     Disk$Hung$Flag          ;  into error flag to pull
                                                     ;  control out of loop
0C5C C9              RET                             ;Return to watchdog routine
                Disk$Complete:
0C5D 010000          LXI     B,0                     ;Reset watchdog timer
                                                     ;HL is irrelevant here
0C60 CD6D08          CALL    Set$Watchdog
0C63 3A4300          LDA     Disk$Status$Block       ;Complete, now check status
0C66 FE80            CPI     80H                     ;Check if any errors occurred
0C68 DAB40D          JC      Disk$Error              ;Yes
                ;
                Disk$Error$Ignore:
0C6B AF              XRA     A                       ;No
0C6C 32320A          STA     Disk$Error$Flag         ;Clear error flag
0C6F C9              RET
                ;
                ;#
                ;       Disk error message handling
                ;
                ;
                Disk$Error$Messages:            ;This table is scanned, comparing the
                                                ;  disk error status with those in the
                                                ;  table. Given a match, or even when
                                                ;  then end of the table is reached, the
                                                ;  address following the status value
                                                ;  points to the correct message text.
0C70 40              DB      40H
0C71 9D0C            DW      Disk$Msg$40
0C73 41              DB      41H
0C74 A20C            DW      Disk$Msg$41
```

Figure 8-10. (Continued)

```
0C76 42                DB      42H
0C77 AC0C              DW      Disk$Msg$42
0C79 21                DB      21H
0C7A BC0C              DW      Disk$Msg$21
0C7C 22                DB      22H
0C7D C10C              DW      Disk$Msg$22
0C7F 23                DB      23H
0C80 C80C              DW      Disk$Msg$23
0C82 24                DB      24H
0C83 DA0C              DW      Disk$Msg$24
0C85 25                DB      25H
0C86 E60C              DW      Disk$Msg$25
0C88 11                DB      11H
0C89 F90C              DW      Disk$Msg$11
0C8B 12                DB      12H
0C8C 070D              DW      Disk$Msg$12
0C8E 13                DB      13H
0C8F 140D              DW      Disk$Msg$13
0C91 14                DB      14H
0C92 220D              DW      Disk$Msg$14
0C94 15                DB      15H
0C95 310D              DW      Disk$Msg$15
0C97 16                DB      16H
0C98 3D0D              DW      Disk$Msg$16
0C9A 00                DB      0                       ;<== Terminator
0C9B 4D0D              DW      Disk$Msg$Unknown        ;Unmatched code
                  ;
0003 =            DEM$Entry$Size  EQU     3       ;Disk error message table entry size
                  ;
                  ;       Message texts
                  ;
0C9D 48756E6700   Disk$Msg$40:    DB      'Hung',0        ;Timeout message
0CA2 4E6F742052   Disk$Msg$41:    DB      'Not Ready',0
0CAC 5772697465   Disk$Msg$42:    DB      'Write Protected',0
0CBC 4461746100   Disk$Msg$21:    DB      'Data',0
0CC1 466F726D61   Disk$Msg$22:    DB      'Format',0
0CC8 4D69737369   Disk$Msg$23:    DB      'Missing Data Mark',0
0CDA 4275732054   Disk$Msg$24:    DB      'Bus Timeout',0
0CE6 436F6E7472   Disk$Msg$25:    DB      'Controller Timeout',0
0CF9 4472697665   Disk$Msg$11:    DB      'Drive Address',0
0D07 4865616420   Disk$Msg$12:    DB      'Head Address',0
0D14 547261636B   Disk$Msg$13:    DB      'Track Address',0
0D22 536563746F   Disk$Msg$14:    DB      'Sector Address',0
0D31 4275732041   Disk$Msg$15:    DB      'Bus Address',0
0D3D 496C6C6567   Disk$Msg$16:    DB      'Illegal Command',0
0D4D 556E6B6E6F   Disk$Msg$Unknown:       DB      'Unknown',0
                  ;
                  Disk$EM$1:                      ;Main disk error message -- part 1
0D55 070D0A                       DB      BELL,CR,LF
0D58 4469736B20                   DB      'Disk ',0
                  ;
                                                  ;Error text output next
                  ;
                  Disk$EM$2:                      ;Main disk error message -- part 2
0D5E 204572726F                   DB      ' Error ('
0D66 0000         Disk$EM$Status: DB      0,0     ;Status code in Hex.
0D68 290D0A2020                   DB      ')',CR,LF,'     Drive '
0D76 00           Disk$EM$Drive:  DB      0       ;Disk drive code, A,B...
0D77 2C20486561                   DB      ', Head '
0D7E 00           Disk$EM$Head:   DB      0       ;Head number
0D7F 2C20547261                   DB      ', Track '
0D87 0000         Disk$EM$Track:  DB      0,0     ;Track number
0D89 2C20536563                   DB      ', Sector '
0D92 0000         Disk$EM$Sector: DB      0,0     ;Sector number
0D94 2C204F7065                   DB      ', Operation - '
0DA2 00                           DB      0               ;Terminator
                  ;
0DA3 526561642E   Disk$EM$Read:   DB      'Read.',0       ;Operation names
0DA9 5772697465   Disk$EM$Write:  DB      'Write.',0
                  ;
                  ;
                  Disk$Action$Confirm:
0DB0 00                           DB      0       ;Set to character entered by user
0DB1 0D0A00                       DB      CR,LF,0
                  ;
                  ;       Disk error processor
```

Figure 8-10. (Continued)

```
                      ;
                      ;       This routine builds and outputs an error message.
                      ;       The user is then given the opportunity to:
                      ;
                      ;               R -- retry the operation that caused the error
                      ;               I -- ignore the error and attempt to continue
                      ;               A -- abort the program and return to CP/M.
                      ;
                      Disk$Error:
0DB4 F5                       PUSH    PSW             ;Preserve error code from controller
0DB5 21660D                   LXI     H,Disk$EM$Status        ;Convert code for message
0DB8 CD440E                   CALL    CAH                     ;Converts A to hex.

0DBB 3A270A                   LDA     In$Buffer$Disk          ;Convert disk id. for message
0DBE C641                     ADI     'A'                     ;Make into letter
0DC0 32760D                   STA     Disk$EM$Drive

0DC3 3A1F0B                   LDA     Floppy$Head             ;Convert head number
0DC6 C630                     ADI     '0'
0DC8 327E0D                   STA     Disk$EM$Head

0DCB 3A200B                   LDA     Floppy$Track            ;Convert track number
0DCE 21870D                   LXI     H,Disk$EM$Track
0DD1 CD440E                   CALL    CAH

0DD4 3A210B                   LDA     Floppy$Sector           ;Convert sector number
0DD7 21920D                   LXI     H,Disk$EM$Sector
0DDA CD440E                   CALL    CAH

0DDD 21550D                   LXI     H,Disk$EM$1             ;Output first part of message
0DE0 CD5305                   CALL    Output$Error$Message

0DE3 F1                       POP     PSW                     ;Recover error status code
0DE4 47                       MOV     B,A                     ;For comparisons
0DE5 216D0C                   LXI     H,Disk$Error$Messages - DEM$Entry$Size
                                                              ;HL -> table - one entry
0DE8 110300                   LXI     D,DEM$Entry$Size        ;Get entry size for loop below
                      Disk$Error$Next$Code:
0DEB 19                       DAD     D                       ;Move to next (or first) entry

0DEC 7E                       MOV     A,M                     ;Get code number from table
0DED B7                       ORA     A                       ;Check if end of table
0DEE CAF80D                   JZ      Disk$Error$Matched      ;Yes, pretend a match occurred
0DF1 B8                       CMP     B                       ;Compare to actual code
0DF2 CAF80D                   JZ      Disk$Error$Matched      ;Yes, exit from loop
0DF5 C3EB0D                   JMP     Disk$Error$Next$Code    ;Check next code
                      ;
                      Disk$Error$Matched:
0DF8 23                       INX     H                       ;HL -> address of text
0DF9 5E                       MOV     E,M                     ;Get address into DE
0DFA 23                       INX     H
0DFB 56                       MOV     D,M
0DFC EB                       XCHG                            ;HL -> text
0DFD CD5305                   CALL    Output$Error$Message    ;Display explanatory text

0E00 215E0D                   LXI     H,Disk$EM$2             ;Display second part of message
0E03 CD5305                   CALL    Output$Error$Message

0E06 21A30D                   LXI     H,Disk$EM$Read          ;Choose operation text
                                                              ;  (assume a read)
0E09 3A1D0B                   LDA     Floppy$Command          ;Get controller command
0E0C FE01                     CPI     Floppy$Read$Code
0E0E CA140E                   JZ      Disk$Error$Read         ;Yes
0E11 21A90D                   LXI     H,Disk$EM$Write         ;No, change address in HL
                      Disk$Error$Read:
0E14 CD5305                   CALL    Output$Error$Message    ;Display operation type
                      ;
                      Disk$Error$Request$Action:              ;Ask the user what to do next
0E17 CD2F05                   CALL    Request$User$Choice     ;Display prompt and wait for input
                                                              ;  Returns with A = uppercase char.
0E1A FE52                     CPI     'R'                     ;Retry?
0E1C CA2C0E                   JZ      Disk$Error$Retry
0E1F FE41                     CPI     'A'                     ;Abort
0E21 CA360E                   JZ      System$Reset
0E24 FE49                     CPI     'I'                     ;Ignore
0E26 CA6B0C                   JZ      Disk$Error$Ignore
```

Figure 8-10. (Continued)

```
0E29 C3170E           JMP     Disk$Error$Request$Action
                ;
                Disk$Error$Retry:                       ;The decision on where to return
                                                        ;  depends on whether the operation
                                                        ;  failed on a deblocked or
                                                        ;  nondeblocked drive.
0E2C 3A350A           LDA     Selected$Disk$Deblock
0E2F B7               ORA     A
0E30 C2EF0B           JNZ     Deblock$Retry
0E33 C33C0B           JMP     No$Deblock$Retry
                ;
                System$Reset:                           ;This is a radical approach, but
                                                        ;  it does cause CP/M to restart.
0E36 0E00             MVI     C,0                       ;System reset
0E38 CD0500           CALL    BDOS

                ;
                ;
                ;     A to upper
                ;
                ;     Converts the contents of the A register to an upper-
                ;     case letter if it is currently a lowercase letter.
                ;
                ;     Entry parameters
                ;
                ;             A = character to be converted
                ;
                ;     Exit parameters
                ;
                ;             A = converted character
                ;
                A$To$Upper:
0E3B FE61             CPI     'a'             ;Compare to lower limit
0E3D D8               RC                      ;No need to convert
0E3E FE7B             CPI     'z' + 1         ;Compare to upper limit
0E40 D0               RNC                     ;No need to convert
0E41 E65F             ANI     5FH             ;Convert to uppercase
0E43 C9               RET
                ;
                ;     Convert A register to hexadecimal
                ;
                ;     This subroutine converts the A register to hexadecimal.
                ;
                ;     Entry parameters
                ;
                ;             A = value to be converted and output
                ;             HL -> buffer area to receive two characters of output
                ;
                ;     Exit parameters
                ;
                ;             HL -> byte following last hex byte output
                ;
                CAH:
0E44 F5               PUSH    PSW             ;Take a copy of the value to be converted
0E45 0F               RRC                     ;Shift A right four places
0E46 0F               RRC
0E47 0F               RRC
0E48 0F               RRC
0E49 CD4D0E           CALL    CAH$Convert     ;Convert to ASCII
0E4C F1               POP     PSW             ;Get original value again
                                              ;Drop into subroutine, which converts
                                              ;  and returns to caller
                CAH$Convert:
0E4D E60F             ANI     0000$1111B      ;Isolate LS four bits
0E4F C630             ADI     '0'             ;Convert to ASCII
0E51 FE3A             CPI     '9' + 1         ;Compare to maximum
0E53 DA580E           JC      CAH$Numeric     ;No need to convert to A -> F
0E56 C607             ADI     7               ;Convert to a letter
                CAH$Numeric:
0E58 77               MOV     M,A             ;Save character
0E59 23               INX     H               ;Update character pointer
0E5A C9               RET

                ;
                ;
                ;#
```

Figure 8-10. (Continued)

```
                ;
                ;       Disk control table images for warm boot
                ;
                Boot$Control$Part$1:
0E5B 01                 DB      1                       ;Read function
0E5C 00                 DB      0                       ;Unit (drive) number
0E5D 00                 DB      0                       ;Head number
0E5E 00                 DB      0                       ;Track number
0E5F 02                 DB      2                       ;Starting sector number
0E60 0010               DW      8*512                   ;Number of bytes to read
0E62 00C4               DW      CCP$Entry               ;Read into this address
0E64 4300               DW      Disk$Status$Block       ;Pointer to next status block
0E66 4500               DW      Disk$Control$5          ;Pointer to next control table
                Boot$Control$Part2:
0E68 01                 DB      1                       ;Read function
0E69 00                 DB      0                       ;Unit (drive) number
0E6A 01                 DB      1                       ;Head number
0E6B 00                 DB      0                       ;Track number
0E6C 01                 DB      1                       ;Starting sector number
0E6D 0006               DW      3*512                   ;Number of bytes to read
0E6F 00D4               DW      CCP$Entry + (8*512)     ;Read into this address
0E71 4300               DW      Disk$Status$Block       ;Pointer to next status block
0E73 4500               DW      Disk$Control$5          ;Pointer to next control table

                ;
                ;
                ;#
                ;
                WBOOT:          ;Warm boot entry
                                        ;On warm boot, the CCP and BDOS must be reloaded
                                        ;  into memory. In this BIOS, only the 5 1/4"
                                        ;  diskettes will be used, therefore this code
                                        ;  is hardware specific to the controller. Two
                                        ;  prefabricated control tables are used.
0E75 318000             LXI     SP,80H
0E78 115B0E             LXI     D,Boot$Control$Part1    ;Execute first read of warm boot
0E7B CD8A0E             CALL    Warm$Boot$Read          ;Load drive 0, track 0,
                                                        ;  head 0, sectors 2 - 8
0E7E 11680E             LXI     D,Boot$Control$Part2    ;Execute second read
0E81 CD8A0E             CALL    Warm$Boot$Read          ;Load drive 0, track 0,
                                                        ;  head 1, sectors 1 - 3
0E84 CDDF0E             CALL    Patch$CPM               ;Make custom enhancements patches
0E87 C36C02             JMP     Enter$CPM               ;Set up base page and enter CCP
                ;
                Warm$Boot$Read:                 ;On entry, DE -> control table image
                                                ;This control table is moved into
                                                ;  the main disk control table and
                                                ;  then the controller activated.
0E8A 211D0B             LXI     H,Floppy$Command        ;HL -> actual control table
0E8D 224600             SHLD    Command$Block$5         ;Tell the controller its address
                                                        ;Move the control table image
                                                        ;  into the control table itself.
0E90 0E0D               MVI     C,13            ;Set byte count
                Warm$Boot$Move:
0E92 1A                 LDAX    D               ;Get image byte
0E93 77                 MOV     M,A             ;Store into actual control table
0E94 23                 INX     H               ;Update pointers
0E95 13                 INX     D
0E96 0D                 DCR     C               ;Count down on byte count
0E97 C2920E             JNZ     Warm$Boot$Move  ;Continue until all bytes moved

0E9A 214500             LXI     H,Disk$Control$5        ;Activate controller
0E9D 3680               MVI     M,80H
                Wait$For$Boot$Complete:
0E9F 7E                 MOV     A,M                     ;Get status byte
0EA0 B7                 ORA     A                       ;Check if complete
0EA1 C29F0E             JNZ     Wait$For$Boot$Complete  ;No
                                                        ;Yes, check for errors
0EA4 3A4300             LDA     Disk$Status$Block
0EA7 FE80               CPI     80H
0EA9 DAAD0E             JC      Warm$Boot$Error         ;Yes, an error occurred
0EAC C9                 RET
                ;
                Warm$Boot$Error:
0EAD 21B60E             LXI     H,Warm$Boot$Error$Message
0EB0 CD5F02             CALL    Display$Message
```

Figure 8-10. (Continued)

```
0EB3 C3750E             JMP     WBOOT                   ;Restart warm boot
                ;
                Warm$Boot$Error$Message:
0EB6 0D0A576172         DB      CR,LF,'Warm Boot Error - retrying...',CR,LF,0
                ;
                ;
                ;#
                ;
                Ghost$Interrupt:        ;Control will only arrive here under the most
                                        ;  unusual circumstances, as the interrupt
                                        ;  controller will have been programmed to
                                        ;  suppress unused interrupts.
                                        ;
0ED8 F5                 PUSH    PSW                     ;Save pre-interrupt registers
0ED9 3E20               MVI     A,IC$EOI                ;Indicate end of interrupt
0EDB D3D8               OUT     IC$OCW2$Port
0EDD F1                 POP     PSW
0EDE C9                 RET
                ;
                ;
                ;#
                ;
                ;       Patch CP/M
                ;
                ;       This routine makes some very special patches to the
                ;       CCP and BDOS in order to make some custom enhancements
                ;
                ;       Public files:
                ;               On large hard disk systems it is extremely useful
                ;               to partition the disk using the user number features.
                ;               However, it becomes wasteful of disk space because
                ;               multiple copies of common programs must be stored in
                ;               each user area. This patch makes User 0 public --
                ;               accessible from any other user area.
                ;               *** WARNING ***
                ;               Files in User 0 MUST be set to system and read/only
                ;               status to avoid their being accidentally damaged.
                ;               Because of the side effects associated with public
                ;               files, the patch can be turned on or off using
                ;               a flag in the long term configuration block.
                ;
                ;       User prompt:
                ;               When using CP/M's USER command and user numbers
                ;               in general, it is all too easy to become confused
                ;               and forget which user number you are "in." This
                ;               patch modifies the CCP to display a prompt which
                ;               shows not only the default disk id., but also the
                ;               current user number, and an indication of whether
                ;               public files are enabled:
                ;
                ;                               P3B>  or  3B>
                ;                               ^
                ;                               When public files are enabled.
                ;
                ;       Equates for public files
                ;
D35E =          PF$BDOS$Exit$Point      EQU     BDOS$Entry + 758H
D37C =          PF$BDOS$Char$Matches    EQU     BDOS$Entry + 776H
D361 =          PF$BDOS$Resume$Point    EQU     BDOS$Entry + 75BH
000D =          PF$BDOS$Unused$Bytes    EQU     13
                ;
                ;
                ;       Equates for user prompt
                ;
C788 =          UP$CCP$Exit$Point       EQU     CCP$Entry + 388H
C78B =          UP$CCP$Resume$Point     EQU     CCP$Entry + 38BH
C513 =          UP$CCP$Get$User         EQU     CCP$Entry + 113H
C5D0 =          UP$CCP$Get$Disk$Id      EQU     CCP$Entry + 1D0H
C48C =          UP$CCP$CONOUT           EQU     CCP$Entry + 8CH
                ;
                ;
                ;       Set up the intervention points
                ;
                Patch$CPM:
0EDF 3EC3               MVI     A,JMP           ;Set up opcode
0EE1 325ED3             STA     PF$BDOS$Exit$Point
```

Figure 8-10. (Continued)

```
0EE4 3288C7          STA     UP$CCP$Exit$Point
0EE7 21F40E          LXI     H,Public$Patch
0EEA 225FD3          SHLD    PF$BDOS$Exit$Point + 1
0EED 21110F          LXI     H,Prompt$Patch  ;Get address of intervening code
0EF0 2289C7          SHLD    UP$CCP$Exit$Point + 1
0EF3 C9              RET                     ;Return to enter CP/M
             ;
             ;
             ;
             Public$Patch:                   ;Control arrives here from the BDOS
                                             ;The BDOS is in the process of scanning
                                             ;  down the target file name in the
                                             ;  search next function
                                             ;  HL -> the name of the file searched for
                                             ;  DE -> directory entry
                                             ;  B = character count
0EF4 3A4200          LDA     CB$Public$Files ;Check if public files are to be enabled
0EF7 B7              ORA     A
0EF8 CA0B0F          JZ      No$Public$Files ;No
0EFB 78              MOV     A,B             ;Get character count
0EFC B7              ORA     A               ;Check if looking at first byte
                                             ;  (that contains the user number)
0EFD C20B0F          JNZ     No$Public$Files ;No, ignore this patch
0F00 1A              LDAX    D               ;Get user number from directory entry
0F01 FEE5            CPI     0E5H            ;Check if active directory entry
0F03 CA0B0F          JZ      No$Public$Files ;Yes, ignore this patch
0F06 7E              MOV     A,M             ;Get user number
0F07 B7              ORA     A               ;Check if User 0
0F08 CA7CD3          JZ      PF$BDOS$Char$Matches    ;Force character match
             No$Public$Files:                ;Replaced patched out code
0F0B 78              MOV     A,B                     ;Check if count indicates that
0F0C FE0D            CPI     PF$BDOS$Unused$Bytes    ;  registers are pointing at
                                                     ;  unused bytes field of FCB
0F0E C361D3          JMP     PF$BDOS$Resume$Point    ;Return to BDOS
             ;
             Prompt$Patch:                   ;Control arrives here from the CCP
                                             ;The CCP is just about to get the
                                             ;  drive id. when control gets here.
                                             ;The CCP's version of CONOUT is used
                                             ;  so that the CCP can keep track of
                                             ;  the cursor position.
0F11 3A4200          LDA     CB$Public$Files ;Check if public files are enabled
0F14 B7              ORA     A
0F15 CA1D0F          JZ      UP$Private$Files        ;No
0F18 3E50            MVI     A,'P'
0F1A CD8CC4          CALL    UP$CCP$CONOUT   ;Use CCP's CONOUT routine
             UP$Private$Files:
0F1D CD13C5          CALL    UP$CCP$Get$User ;Get current user number
0F20 FE0A            CPI     9 + 1           ;Check if one or two digits
0F22 D2300F          JNC     UP$2$Digits
0F25 C630            ADI     '0'             ;Convert to ASCII
             UP$1$Digit:
0F27 CD8CC4          CALL    UP$CCP$CONOUT   ;Output the character
0F2A CDD0C5          CALL    UP$CCP$Get$Disk$Id      ;Get disk identifier
0F2D C38BC7          JMP     UP$CCP$Resume$Point     ;Return to CCP
             ;
             UP$2$Digits:
0F30 C626            ADI     '0' - 10        ;Subtract 10 and convert to ASCII
0F32 F5              PUSH    PSW             ;Save converted second digit
0F33 3E31            MVI     A,'1'           ;Output leading '1'
0F35 CD8CC4          CALL    UP$CCP$CONOUT
0F38 F1              POP     PSW             ;Recover second digit
0F39 C3270F          JMP     UP$1$Digit      ;Output remainder of prompt and return to
                                             ;  the CCP
             ;
             ;#
```

Figure 8-10. (Continued)

```
                ;
                ;       Configuration block get address
                ;
                ;       This routine is called by utility programs running in the TPA.
                ;       Given a specific code number, it returns the address of a specific
                ;       object in the configuration block.
                ;
                ;       By using this routine, utility programs need not know the exact
                ;       layout of the configuration block.
                ;
                ;       Entry parameters
                ;
                ;               C = Object identity code (in effect, this is the
                ;                       subscript of the object's address in the
                ;                       table below)
                ;
                ;========================
                CB$Get$Address:                         ;<=== BIOS entry point (private)
                ;========================
0F3C F5                 PUSH    PSW                     ;Save user's registers
0F3D C5                 PUSH    B
0F3E D5                 PUSH    D

0F3F 69                 MOV     L,C                     ;Make code into a word
0F40 2600               MVI     H,0
0F42 29                 DAD     H                       ;Convert code into word offset
0F43 114F0F             LXI     D,CB$Object$Table       ;Get base address of table
0F46 19                 DAD     D                       ;HL -> object's address in table
0F47 5E                 MOV     E,M                     ;Get LS byte
0F48 23                 INX     H
0F49 56                 MOV     D,M                     ;Get MS byte
0F4A EB                 XCHG                            ;HL = address of object

0F4B D1                 POP     D                       ;Recover user's registers
0F4C C1                 POP     B
0F4D F1                 POP     PSW

0F4E C9                 RET
                ;
                ;#
                ;
                CB$Object$Table:
                                                ;       Code
                                                ;        vv
0F4F 8F0F               DW      Date                    ;01 date in ASCII
0F51 990F               DW      Time$In$ASCII           ;02 time in ASCII
0F53 A30F               DW      Time$Date$Flags         ;03 flags indicated if time/date set
0F55 8D0F               DW      CB$Forced$Input         ;04 forced input pointer
0F57 4300               DW      CB$Startup              ;05 system startup message
                                                        ;   Redirection words
0F59 5800               DW      CB$Console$Input        ;06
0F5B 5A00               DW      CB$Console$Output       ;07
0F5D 5C00               DW      CB$Auxiliary$Input      ;08
0F5F 5E00               DW      CB$Auxiliary$Output     ;09
0F61 6000               DW      CB$List$Input           ;10
0F63 6200               DW      CB$List$Output          ;11

0F65 6400               DW      CB$Device$Table$Addresses ;12
0F67 B500               DW      CB$12$24$Clock          ;13 Selects 12/24 hr. format clock
0F69 BD00               DW      RTC$Ticks$per$Second    ;14
0F6B BF00               DW      RTC$Watchdog$Count      ;15
0F6D C100               DW      RTC$Watchdog$Address    ;16
0F6F C300               DW      CB$Function$Key$Table   ;17
0F71 1B02               DW      CONOUT$Escape$Table     ;18

0F73 8400               DW      D0$Initialize$Stream    ;19
0F75 9100               DW      D0$Baud$Rate$Constant   ;20
0F77 9400               DW      D1$Initialize$Stream    ;21
0F79 A100               DW      D1$Baud$Rate$Constant   ;22
0F7B A400               DW      D2$Initialize$Stream    ;23
0F7D B100               DW      D2$Baud$Rate$Constant   ;24
0F7F 4002               DW      Interrupt$Vector        ;25
0F81 890F               DW      LTCB$Offset             ;26
0F83 8B0F               DW      LTCB$Length             ;27
0F85 4200               DW      CB$Public$Files         ;30
```

Figure 8-10. (Continued)

```
0F87 A421               DW      Multi$Command$Buffer    ;31
                ;
                ;#
                ;       The short term configuration block.
                ;
                ;       This contains variables that can be set once CP/M
                ;       has been initiated, but that are never preserved
                ;       from one loading of CP/M to the next. This part of
                ;       the configuration block form the last initialized bytes
                ;       in the BIOS.
                ;
                ;       The two values below are used by utility programs that
                ;       need to read in the long term configuration block from disk.
                ;       The BIOS starts on a 256-byte page boundary, and therefore
                ;       will always be on a 128-byte sector boundary in the reserved
                ;       area on the disk. A utility program can then, using the
                ;       CB$Get$Address Private BIOS call, determine how many 128-byte
                ;       sectors need to be read in by the formula:
                ;
                ;               (LCTB$Offset + LTCB$Length) / 128
                ;
                ;       The LTCB$Offset is the offset from the start of the BIOS to
                ;       where the first byte of the long term configuration block
                ;       starts. Using the offset and the length, the utility can
                ;       copy the RAM version of the LTCB over the disk image
                ;       that it has read from the disk, and then write the
                ;       updated LTCB back onto the disk.
                ;
0F89 BED9       LTCB$Offset:    DW      BIOS$Entry - Long$Term$CB
0F8B E601       LTCB$Length:    DW      Long$Term$CB$End - Long$Term$CB
                ;
                ;       Forced input pointer
                ;
                ;       If CONIN ever finds that this pointer is pointing to a nonzero
                ;       byte, then this byte will be injected into the console input
                ;       stream as though it had been typed on the console. The
                ;       pointer is then updated to the next byte in memory.

0F8D 4300       CB$Forced$Input:        DW      CB$Startup
                ;
                ;
                Date:                   ;Current system date
0F8F 31302F3137         DB      '10/17/82',LF   ;Unless otherwise set to the contrary
                                                ;  this is the release date of the system
                                                ;Normally, it will be set by the DATE utility
0F98 00                 DB      0               ;00-byte terminator

                Time$in$ASCII:          ;Current system time
0F99 3030       HH:     DB      '00'            ;Hours
0F9B 3A                 DB      ':'
0F9C 3030       MM:     DB      '00'            ;Minutes
0F9E 3A                 DB      ':'
0F9F 3030       SS:     DB      '00'            ;Seconds
                Time$in$ASCII$End:              ;Used when updating the time
0FA1 0A                 DB      LF
0FA2 00                 DB      0               ;00-byte terminator
                ;
                ;
                Time$Date$Flags:        ;This byte contains two flags that are  used
                                        ;  to indicate whether the time and/or date
                                        ;  have been set either programmatically or
                                        ;  by using the TIME and DATE utilities. These
                                        ;  flags can be tested by utility programs that
                                        ;  need to have the correct time and date set.
0FA3 00                 DB      0
0001 =          Time$Set        EQU     0000$0001B
0002 =          Date$Set        EQU     0000$0010B

                ;
                ;#
                ;       Uninitialized buffer areas
                ;
                ;       With the exception of the main Disk$Buffer, which contains a few
                ;       bytes of code, all of the other uninitialized variables
                ;       occur here. This has the effect of reducing the number of
                ;       bytes that need be stored in the CP/M image on the disk,
```

Figure 8-10. (Continued)

```
                ;       since uninitialized areas do not need to be kept on the disk.
                ;
                ;
                ;#
                ;
                ;       The cold boot initialization code is only needed once.
                ;       It can be overwritten once it has been executed.
                ;       Therefore, it is "hidden" inside the main disk buffer.
                ;
                ;
0FA4            Disk$buffer:    DS      Physical$Sector$Size * Physical$Sec$Per$Track
                ;
                                                        ;Save the location counter
21A4 =          After$Disk$Buffer       EQU     $       ;$ = current value of location counter
                ;
0FA4                                    ORG     Disk$Buffer     ;Wind the location counter back
                ;
                Initialize$Stream:      ;This stream of data is used by the
                                        ;  Initialize subroutine. It has the following
                                        ;  format:
                                        ;
                                        ;       DB      Port number to be initialized
                                        ;       DB      Number of byte to be output
                                        ;       DB      xx,xx,xx,xx data to be output
                                        ;       :
                                        ;       :
                                        ;       DB      Port number of 00H terminates
                                        ;
                                        ;
                ;
                ;       Initialization stream declared here
0FA4 D8                 DB      IC$ICW1$Port    ;Program the 8259 interrupt controller
0FA5 01                 DB      1
0FA6 56                 DB      IC$ICW1

0FA7 D9                 DB      IC$ICW2$Port
0FA8 01                 DB      1
0FA9 02                 DB      IC$ICW2

0FAA D9                 DB      IC$OCW1$Port
0FAB 01                 DB      1
0FAC FC                 DB      IC$OCW1

0FAD 83                 DB      83H                     ;Program the 8253 clock generator
0FAE 01                 DB      1
0FAF 34                 DB      00$11$010$0B            ;Counter 0, periodic interrupt, mode 2

0FB0 80                 DB      80H                     ;RTC uses channel 0
0FB1 02                 DB      2
0FB2 0146               DW      17921                   ;19721 * 930 nanoseconds =
                                                        ;  16.666 milliseconds). 60 ticks/sec.
0FB4 00                 DB      0               ;Port number of 0 terminates
                ;
                ;
                Signon$Message:
0FB5 43502F4D20         DB      'CP/M 2.2.'
0FBE 3030               DW      VERSION         ;Current version number
0FC0 20                 DB      ' '
0FC1 3032               DW      MONTH           ;Current date
0FC3 2F                 DB      '/'
0FC4 3236               DW      DAY
0FC6 2F                 DB      '/'
0FC7 3833               DW      YEAR
0FC9 0D0A0A             DB      CR,LF,LF
0FCC 456E68616E         DB      'Enhanced BIOS',CR,LF,LF
0FDC 4469736B20         DB      'Disk Configuration :',CR,LF,LF
0FF3 2020202020         DB      '     A: 0.35 Mbyte 5" Floppy',CR,LF
1011 2020202020         DB      '     B: 0.35 Mbyte 5" Floppy',CR,LF,LF
1030 2020202020         DB      '     C: 0.24 Mbyte 8" Floppy',CR,LF
104E 2020202020         DB      '     D: 0.24 Mbyte 8" Floppy',CR,LF
106C 2020202020         DB      '     M: 0.19 Mbyte Memory Disk',CR,LF,LF
                ;
108D 00                 DB      0
                ;
                ;       Messages for M$Disk
                ;
```

Figure 8-10. (Continued)

```
                M$Disk$Setup$Message:
108E 2020202020         DB      '     M$Disk already contains valid information.',CR,LF,0
                M$Disk$Not$Setup$Message:
10C0 2020202020         DB      '     M$Disk has been initialized to empty state.',CR,LF,0
                ;
                M$Disk$Dir$Entry:                       ;Dummy directory entry used to determine
                                                        ;  if the M$Disk contains valid information
10F3 0F                 DB      15                      ;User 15
10F4 4D24446973         DB      'M$Disk  '
10FC A0A020             DB      ' '+80H,' '+80H,' '          ;System and read/only
10FF 00000000           DB      0,0,0,0
1103 0000000000         DB      0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0
                ;
0004 =          Default$Disk    EQU     0004H   ;Default disk in base page
                ;
                BOOT:                   ;Entered directly from the BIOS JMP Vector
                                        ;Control will be transferred here by the CP/M
                                        ;  bootstrap loader
                                        ;
                                                        ;Initialize system
                                                        ;This routine uses the Initialize$Stream
                                                        ;  declared above

1113 F3                 DI                              ;Disable interrupts to prevent any
                                                        ; side effects during initialization
1114 21A40F             LXI     H,Initialize$Stream     ;HL -> data stream
1117 CD1903             CALL    Output$Byte$Stream      ;Output it to the specified
                                                        ;  ports

111A CDEE02             CALL    General$CIO$Initialization ;Initialize character devices

111D 21B50F             LXI     H,Signon$Message        ;Display sign-on message on console
1120 CD5F02             CALL    Display$Message
                ;
1123 CDDF0E             CALL    Patch$CPM               ;Make necessary patches to CCP and BDOS
                                                        ;  for custom enhancements

                                                        ;Initialize M$Disk
                                                        ;If the M$Disk directory has the
                                                        ;  special reserved file name "M$disk"
                                                        ;  (with lowercase letters and marked
                                                        ;  SYS and R/O), then the M$Disk is
                                                        ;  assumed to contain valid data.
                                                        ;If the "M$Disk" file is absent, the
                                                        ;  M$Disk Directory entry is moved into
                                                        ;  the M$Disk image, and the remainder of
                                                        ;  the directory set to 0E5H.
1126 0601               MVI     B,1                     ;Select bank 1
1128 CDDD0B             CALL    Select$Bank             ;  which contains the M$Disk directory

                                                        ;Check if M$Disk directory entry present
112B 210000             LXI     H,0                     ;Start address for first directory
112E 11F310             LXI     D,M$Disk$Dir$Entry
1131 0E20               MVI     C,32                    ;Length to compare
                M$Disk$Test:
1133 1A                 LDAX    D                       ;Get byte from initialized variable
1134 BE                 CMP     M                       ;Compare with M$Disk image
1135 C24F11             JNZ     M$Disk$Not$Setup        ;Match fails
1138 13                 INX     D
1139 23                 INX     H
113A 0D                 DCR     C
113B CA4111             JZ      M$Disk$Setup            ;All bytes match
113E C33311             JMP     M$Disk$Test
                ;
                M$Disk$Setup:
1141 218E10             LXI     H,M$Disk$Setup$Message  ;Inform user
                ;
                M$Disk$Setup$Done:
1144 CD5F02             CALL    Display$Message

1147 AF                 XRA     A                       ;Set default disk drive to A:
1148 320400             STA     Default$Disk
114B FB                 EI                              ;Interrupts can now be enabled

114C C36C02             JMP     Enter$CPM               ;Go into CP/M
                ;
```

Figure 8-10. (Continued)

```
                 M$Disk$Not$Setup:
114F 110000              LXI     D,0             ;Move M$Disk directory entry into
1152 21F310              LXI     H,M$Disk$Dir$Entry       ;  M$Disk image
1155 0E04                MVI     C,32/8                   ;Number of 8-byte blocks to move
1157 CDF80A              CALL    Move$8
                 ;
                                                 ;DE -> next byte after M$Disk directory
                                                 ;  entry in image
115A 3EE5                MVI     A,0E5H                   ;Set up to do memory fill
115C 12                  STAX    D                        ;Store first byte in "source" area
115D 62                  MOV     H,D                      ;Set HL to DE +1
115E 6B                  MOV     L,E
115F 23                  INX     H
1160 0EFC                MVI     C,((2 * 1024) - 32) / 8 ;Two allocation blocks
                                                          ;  less 32 bytes for M$Disk entry
1162 CDF80A              CALL    Move$8                   ;Use Move$8 to do fill operation

1165 21C010              LXI     H,M$Disk$Not$Setup$Message
1168 C34411              JMP     M$Disk$Setup$Done        ;Output message and enter CP/M
                 ;
                 ;
116B 00                  DB      0               ;Dummy
                 Last$Initialized$Byte:          ;<== address of last initialized byte
                 ;
                 ;       End of cold boot initialization code
                 ;
21A4                     ORG     After$Disk$Buffer        ;Reset location counter
                 ;
21A4             Multi$Command$Buffer:   DS      128      ;This can be used to insert long
                                                          ;  command sequences into the
                                                          ;  console input stream by setting
                                                          ;  the forced input pointer here
                 ;
0020 =           D0$Buffer$Length        EQU     32       ;Must be binary number
2224             D0$Buffer:      DS      D0$Buffer$Length
                 ;
0020 =           D1$Buffer$Length        EQU     32       ;Must be binary number
2244             D1$Buffer:      DS      D1$Buffer$Length
                 ;
0020 =           D2$Buffer$Length        EQU     32       ;Must be binary number
2264             D2$Buffer:      DS      D2$Buffer$Length
                 ;
                 ;       Data areas for the character drivers
                 ;
2284             PI$User$Stack:  DS      2       ;Storage area for user's stack pointer
                                                 ;  when an interrupt occurs
2286             PI$User$HL:     DS      2       ;Save area for user's HL
2288                             DS      40      ;Stack area for use by interrupt service
                 PI$Stack:                       ;  routines to avoid overflowing the
                                                 ;  user's stack area
                 ;
22B0             Directory$Buffer:       DS      128      ;Disk directory buffer
                 ;
2330             M$Disk$Buffer:          DS      128      ;Intermediary buffer for
                                                          ;  M$Disk
                 ;
                 ;       Disk work areas
                 ;
                 ;       These are used by the BDOS to detect any unexpected
                 ;       change of diskettes. The BDOS will automatically set
                 ;       such a changed diskette to read-only status.
                 ;
23B0             Disk$A$Workarea:        DS      32       ; A:
23D0             Disk$B$Workarea:        DS      32       ; B:
23F0             Disk$C$Workarea:        DS      16       ; C:
2400             Disk$D$Workarea:        DS      16       ; D:
                 ;
                 ;
                 ;       Disk allocation vectors
                 ;
                 ;       These are used by the BDOS to maintain a bit map of
                 ;       which allocation blocks are used and which are free.
                 ;       One byte is used for eight allocation blocks, hence the
                 ;       expression of the form (allocation blocks/8)+1.
                 ;
2410             Disk$A$Allocation$Vector        DS      (174/8)+1       ; A:
```

Figure 8-10. (Continued)

```
2426    Disk$B$Allocation$Vector        DS      (174/8)+1       ; B:
        ;
243C    Disk$C$Allocation$Vector        DS      (242/8)+1       ; C:
245B    Disk$D$Allocation$Vector        DS      (242/8)+1       ; D:
        ;
247A    M$Disk$Allocation$Vector        DS      (192/8)+1       ; M$Disk

2493            END     ;of enhanced BIOS listing
```

Figure 8-10. (Continued)

Classes of Errors
BIOS Error-Handling Functions
Practical Error Handling
Character I/O Errors
Disk Errors
Improving Error Messages

9

Dealing with Hardware Errors

This chapter describes the enhancements you can make to improve CP/M's somewhat primitive error handling. It covers the general classes of errors that the BIOS may have to handle. It describes some of the underlying philosophical aspects of errors, how to detect them, and how to correct them or otherwise make the best of the situation.

At the end of the chapter are some example error-handling subroutines. Some of these have already been shown in the previous chapter as part of the enhanced BIOS (Figure 8-10); they are repeated here so that you can see them in isolation.

Classes of Errors

Basically, the user perceives only two classes of errors — those that are user-correctable and those that are not. There is a third, almost invisible class of errors—those that are recoverable by the hardware or software without the user's intervention.

The possible sources for hardware errors vary wildly from one computer system to another, since error detection is heavily dependent on the particular logic in the hardware. The BIOS can detect some hardware-related errors — mainly errors caused when something takes too long to happen, such as when a recalcitrant printer does not react in a specified length of time.

The BDOS has no built-in hardware detection code. It can detect *system* errors, such as an attempt to write to a disk file that is marked "Read-Only" in the file directory or attempts to access files that are not on the disk. These BDOS-detected errors, however, generally are unrelated to the well-being of the hardware. For example, a disk controller with a hardware problem could easily overwrite a sector of the directory, thereby deleting several files. This error would not show up until the user tried to use one of the now-departed files.

BIOS Error-Handling Functions

The error-handling code in the BIOS has to serve the following functions:

- Detection
- Analysis
- Indication
- Correction.

Error Detection

Clearly, before any later steps can be taken, an error must be detected. This can be done by the software alone or by the BIOS interacting with error-detecting logic in the hardware. In general, the only errors that the BIOS can detect unassisted are caused when certain operations take longer to complete than expected. Because the writer of the BIOS knows the operating environment of the specific peripherals in the system, the code can predict how long a particular operation should take and can signal an error when this time is exceeded. This would include such problems as printers that fail to react within a specified time period.

The BIOS can work in cooperation with the hardware to determine whether the hardware itself has detected an error. Armed with the hardware's specifications, the BIOS can input information on controller or device status to trigger error-detecting logic. How this should be done depends heavily on the peripheral devices in your computer system and the degree to which these devices have "smart" controllers capable of processing independently of the computer. Unfortunately, many manufacturers document the significance of individual status bits that indicate errors, but not combinations of errors, or what to do when a particular error occurs.

Error Analysis

Given that your BIOS has detected an error, it must first determine the class of error; that is, whether or not the error can be corrected by simply trying the operation again. Some errors appear at first to be correctable, but retrying the operation several times still fails to complete it. An example would be a check-sum error while reading a disk sector. If several attempts to read the sector all yield an error, then it becomes a "fatal" error. The code in your BIOS must be capable of initial classification and then subsequent reclassification if remedial action fails.

Other types of errors can be classified immediately as fatal errors—nothing can be done to save the situation. For example, if the floppy disk controller indicates that it cannot find a particular sector number on a diskette (due to an error in formatting), there is nothing that the BIOS can do other than inform the user of the problem and supply other helpful information.

Analysis of errors may require some basic research, such as inducing failures in the hardware and observing combinations of error indicators. For example, some printers (interfaced via a parallel port) indicate that they are "Out of Paper" or "Busy" when, in fact, they are switched off. The BIOS should detect this condition and tell the user to switch the printer on, not load more paper.

Error Indication

An incomplete or cryptic error message is infuriating. It is the functional equivalent of saying, "There has been an error. See if you can guess what went wrong!"

An error message, to be complete, should inform the recipient of the following:

- The fact that an error has occurred.
- Whether or not automatic recovery has been attempted and failed.
- The details of the error, if need be in technical terms to assist a hardware engineer.
- What possible choices the user has now.

To put these points into focus, consider the error message that can be output by CP/M after you have attempted to load a program by entering its name into the CCP. What you see on the console is the following dialog:

```
A>myprog<cr>
BAD LOAD
A>
```

All you know is that there has been an error, and you must guess what it is, even though the specific cause of the error was known to CP/M when it output the message. This error message is output by the CCP when it attempts to load a

".COM" file larger than the current transient program area. The message "BAD LOAD" is only understandable *after* you know what the error is. Even then, it does not tell you what went wrong, whether there is anything you can do about it, and how to go about doing it.

To be complete, this error message could say something like this:

```
A>myprog<cr>
    "MYPROG.COM" exceeds the available memory space by
    1,024 bytes, and therefore cannot be loaded under the
    current version of CP/M.
```

Notice how the message tells you what the problem is, and even quantifies it so that you can determine its severity (you need to get 1K more memory or reduce the program's size). It also tells you how you stand—you cannot load this program under the current version of CP/M, so retrying the operation is futile.

Not many systems programmers like to output messages like the example above. They argue that such a message is too long and too much work for something that does not happen often. Admittedly, the message *is* too long. It could be shortened to read

```
(131) Program 1,024 bytes too large to load.
```

This conveys the same information; the number in parentheses can serve as a reference to a manual where the full impact of the message should be described.

The major problem with the way error messages are designed is that they usually are written by programmers to be read by nontechnical lay users, and programmers are notoriously bad at guessing what nonexperts need to know.

Error indications you design should address the following issues, from the point of view of the user:

- The cause of the error
- The severity of the error
- The corrective action that has and can be taken.

Examine the error messages in the error processor for the example BIOS in Figure 8-10, from line 03600 onward. Although these are an improvement on the BDOS all-purpose

```
BDOS Error on A: Bad Sector
```

even these messages do not really meet all of the requirements of a good error message system.

Another often overlooked aspect of errors is that most hardware errors form a pattern. This pattern is normally only discernible to the trained eye of a hardware maintenance engineer. When these engineers are called to investigate a problem, they will quiz the user to determine whether a given failure is an isolated incident or part of an ongoing pattern. This is why an error message should contain additional technical details. For example, a disk error message should include the track and sector used in the operation that resulted in an error. Only with these details can the engineer piece together the context of a failure or group of failures.

Error Correction

Given that a lucid error message has been displayed on the console, the user is still confronted with the question: "Now what do I do?" Not only can this be difficult for the user to answer, but also the particular solution decided upon can be hard for the BIOS to execute.

Normally, there are three possible options in response to errors:

- Try the operation again
- Ignore the error and attempt to continue
- Abort the program causing the error and return to CP/M.

For some errors, retrying can be effective. For example, if you forget to put the printer on-line and get a "Printer Timeout" error message, it is easy to put the printer back on-line and ask the BIOS to try again to send data to the printer.

Seldom can you ignore an error and hope to get sensible results from the machine; many disk controllers do not even transfer data between themselves and the disk drive if an error has been detected. Only ignorant users, or brave ones in desperation, ignore errors.

Aborting the program causing the error is a drastic measure, although it does escape from what could otherwise be a "deadly embrace" situation. For example, if you misassign the printer to an inactive serial port and turn on printer echoing (with the CONTROL-P toggle), you will send the system into an endless series of "Printer Timeout" messages. If you abort the program, the error handler in the BIOS executes a System Reset function (function 0) in the BDOS, CP/M warm boots, and control is returned to the CCP. In the process, the printer toggle is reset and the circle is broken.

Practical Error Handling

This section discusses several errors, describing their causes and the way in which the BIOS and the user can handle them when they occur.

Character I/O Errors

At the BIOS level, most detectable errors related to character input or output will be found by the hardware chips.

Parity Error

Parity, in this context, refers to the number of bits set to 1 in an 8-bit character. The otherwise unused eighth bit in ASCII characters can be set to make this number always odd, or alternatively, always even. Your computer hardware can be programmed to count the number of 1 bits in each character and to generate an error if the number is odd (odd parity) or, alternatively, if it is even (even parity). If the hardware on the other end of the line is programmed to operate in the same mode, parity checking provides a primitive error-detection mechanism — you can tell that a character is bad, but not what it should have been.

CP/M does not provide a standard mechanism for reporting a parity error, so your only option is to reset the hardware and substitute an ASCII DEL (7FH; delete) character in the place of the erroneous character.

If your BIOS is operating in a highly specialized environment, you may need to count the number of such parity errors so that a utility program can report on the overall performance of the system.

Framing Error

When an 8-bit ASCII character is transmitted over a serial line, the eight bits are transmitted serially, one after the other. A *start* bit is transmitted first, followed by the data character and then a *stop* bit. If the hardware fails to find the stop and start bits in the correct positions, a *framing error* will occur. Again, the only option available to the BIOS is to reset the hardware chip and substitute an ASCII DEL.

Overrun Error

This error occurs when incoming data characters arrive faster than the program can handle them, so that the last characters overrun those being processed by the hardware chip. This error can normally be avoided by the use of serial line protocols, such as those in the example BIOS in Figure 8-10.

An *overrun error* implies that the protocol has broken down. As with the parity and framing errors, almost the only option is to reset the hardware and substitute a DEL character.

Printer Timeout Error

This is one of the few errors where the BIOS can sensibly attempt an error recovery. The error occurs when the BIOS tries to output a character to a serial printer and finds that the printer is not ready for more than, say, 30 seconds. The most common cause of this error is that the user forgets to put the printer on-line. Many printers require that they be off-line during a manual form feed, and users will often forget to push the on-line button afterward.

After a 30-second delay, the BIOS can send a message to the console device(s) informing the user of the error and asking the user to choose the appropriate course of action. Note that console output can be directed to more than one device.

Parallel Printers

Printers connected to your system by means of a parallel port can indicate their status to the computer much more easily than can serial printers. They can communicate such error states as "Out of Paper," "End of Ribbon," and "Off-line."

These single-error indicators can also be used in combination to indicate whether the printer cable is connected, or even whether the printer is receiving power. You need to experiment, deliberately putting the printer into these states and reading status in order to identify them. It is misleading to indicate to the inexperienced user that the printer is "Out of Paper" when the problem is that the data cable has inadvertently become disconnected.

However, each of these errors can be dealt with in the same way as the serial printer's timeout problem: display an error message and request the user's choice of action.

Example Printer Error Routine

Figure 9-1 shows an example of a program that handles printer errors. It consists of several subroutines, including

- The error detection classification and indication routine
- The error correction routine.

It uses other subroutines that are omitted from the figure to avoid obscuring the logic. These subroutines are listed in full in the example BIOS in Figure 8-10.

```
                ;       This example shows, in outline form, how to handle the
                ;       situation when a serial printer remains busy for too long.
                ;       It is intended that this generic example show how to
                ;       deal with this class of errors.
                ;
                ;       The example presupposes the existence of a clock interrupt
                ;       every 16.666 milliseconds (1/60th of a second), and that
                ;       control will be transferred to the Real Time Clock service
                ;       routine each time the clock "ticks".
                ;
                ;       Figure 8-10 shows a more complete example, installed in a real
                ;       BIOS.
                ;
                ;
0000 =          B$System$Reset          EQU     0       ;BDOS system reset function
0005 =          BDOS                    EQU     5       ;BDOS entry point
                ;
0000 00         Printer$Timeout$Flag:   DB      0       ;This flag is set by the interrupt
                                                        ;  service subroutine that is called
                                                        ;  when the watchdog timer subroutine
                                                        ;  count hits zero (after having
                                                        ;  counted down a 30-second delay)

0708 =          Printer$Delay$Count     EQU     1800    ;Given a clock period of 16.666 ms
                                                        ;  this represents a delay of 30 secs
```

Figure 9-1. Serial printer error handling

```
                  ;
000D =            CR                      EQU     0DH     ;Carriage return
000A =            LF                      EQU     0AH     ;Line feed
                  ;
                  Printer$Busy$Message:
0001 0D0A                 DB      CR,LF
0003 5072696E74           DB      'Printer has been busy for too long,',CR,LF
0028 436865636B           DB      'Check that it is on-line and ready.',CR,LF,0
                  ;
004E 00           Printer$Character:      DB      0       ;Save area for the data character
                                                          ;  to be output
                  ;
                  ;
                  LIST:                                   ;<=== Main BIOS entry point
                          ;.......                         ;<=== I/O redirection code occurs here

004F 79                   MOV     A,C                     ;Save the data character
0050 324E00               STA     Printer$Character

                  Printer$Retry:
0053 010807               LXI     B,Printer$Delay$Count   ;This is the count of the number
                                                          ;  of clock ticks before the watchdog
                                                          ;  subroutine call
0056 217E00               LXI     H,Printer$Timed$Out     ;  <== this address
0059 CDA300               CALL    Set$Watchdog            ;Sets the watchdog running

                  Printer$Wait:
005C CDA300               CALL    Get$Printer$Status      ;See if the printer is ready to
                                                          ;  accept a character for output
                                                          ;  This includes checking if the printer
                                                          ;  is "Busy" because the driver is
                                                          ;  waiting for XON, ACK, or DTR to
                                                          ;  come high
005F C26C00               JNZ     Printer$Ready           ;The printer is now ready

0062 3A0000               LDA     Printer$Timeout$Flag    ;Check if the watchdog timer has
                                                          ;  hit zero (if it does, the
                                                          ;  watchdog routine will call
                                                          ;  the Printer$Timed$Out code
                                                          ;  that sets this flag)
0065 B7                   ORA     A
0066 C28400               JNZ     Display$Busy$Message    ;Yes, so display message to
                                                          ;  indicate an error has occurred
0069 C35C00               JMP     Printer$Wait            ;Otherwise, check if printer is
                                                          ;  now not busy

                  Printer$Ready:                          ;The printer is now ready to output
                                                          ;  a character, but before doing so,
                                                          ;  the watchdog timer must be reset
006C F3                   DI                              ;Ensure no false timeout occurs
006D 010000               LXI     B,0                     ;This is done by setting the count
0070 CDA300               CALL    Set$Watchdog            ;  to zero
0073 FB                   EI

0074 3A4E00               LDA     Printer$Character       ;Get character to output
0077 11A300               LXI     D,Printer$Device$Table  ;DE -> device table for printer
007A CDA300               CALL    Output$Data$Byte        ;Output the character to the printer

007D C9                   RET                             ;Return to the BIOS's caller
                  ;
                  ;
                  Printer$Timed$Out:                      ;Control arrives here from the
                                                          ;  watchdog routine if the
                                                          ;  watchdog count ever hits zero
                                                          ;  This is an interrupt service
                                                          ;  routine
                                                          ;All registers have been saved
                                                          ;  before control arrives here
007E 3EFF                 MVI     A,0FFH                  ;Set printer timeout flag
0080 320000               STA     Printer$Timeout$Flag
0083 C9                   RET                             ;Return back to the watchdog

                                                          ;Interrupt service routine
```

Figure 9-1. (Continued)

```
                ;
                Display$Busy$Message:                   ;Printer has been busy for
                                                        ;  30 seconds or more
0084 AF                 XRA     A                       ;Reset timeout flag
0085 320000             STA     Printer$Timeout$Flag

0088 210100             LXI     H,Printer$Busy$Message  ;Output error message
008B CDA300             CALL    Output$Error$Message

008E CDA300             CALL    Request$User$Choice     ;Displays a Retry, Abort, Ignore?
                                                        ;  prompt, accepts a character from
                                                        ;  the keyboard, and returns with the
                                                        ;  character, converted to upper
                                                        ;  case in the A register
0091 FE52               CPI     'R'                     ;Check if Retry

0093 CA5300             JZ      Printer$Retry
0096 FE41               CPI     'A'                     ;Check if Abort
0098 CA9E00             JZ      Printer$Abort
009B FE49               CPI     'I'                     ;Check if Ignore
009D C8                 RZ
                ;
                Printer$Abort:
009E 0E00               MVI     C,B$System$Reset        ;Issue system reset
00A0 C30500             JMP     BDOS                    ;No need to give call as
                                                        ;  control will not be returned
                ;
                ;
                ;       Dummy subroutines
                ;       These are shown in full in Figure 8-10. The line numbers in
                ;       Figure 8-10 are shown in the comment field below
                ;
                Printer$Device$Table:           ;Line 01300 (example layout)
                Request$User$Choice:            ;Line 03400
                Output$Error$Message:           ;Line 03500
                Get$Printer$Status:             ;Line 03900 (similar code)
                Output$Data$Byte:               ;Line 05400 (similar code)
                Set$Watchdog:                   ;Line 05800
```

Figure 9-1. Serial printer error handling (continued)

Disk Errors

Disks are much more complicated than character I/O devices. Errors are possible in the electronics and in the disk medium itself. Most of the errors concerned with electronics need only be reported in enough detail to give a maintenance engineer information about the problem. This kind of error is rarely correctable by retrying the operation. In contrast, media errors often can be remedied by retrying the operation or by special error processing software built into the BIOS. This chapter discusses this class of errors.

Media errors occur when the BIOS tries to read a sector from the disk and the hardware detects a check-sum failure in the data. This is known as a *cyclical redundancy check* (CRC) error. Some disk controllers execute a read-after-write check, so a CRC error can also occur during an attempt to write a sector to the disk.

With floppy diskettes, the disk driver should retry the operation at least ten times before reporting the error to the user. Then, because diskettes are inexpensive and replaceable, the user can choose to discard the diskette and continue with a new one.

With hard disks, the media cannot be exchanged. The only way of dealing with bad sectors is to replace them logically, substituting other sectors in their place.

There are two fundamentally different ways of doing this. Figure 9-2 shows the scheme known as sector sparing—substituting sectors on an outer track for a sector that is bad.

The advantage of this scheme is that it is dynamic. If a sector is found to be bad in a read-after-write check, even after several retries, then the data intended for the failing sector can be written to a spare sector. The failing sector's number is placed into a spare-sector directory on the disk. Thereafter, the disk drivers will be redirected to the spare sector every time an attempt is made to read or write the bad sector.

The disadvantage of this system is that the read/write heads on the disk must move out to the spare sector and then back to access the next sector. This can be a problem if you attempt to make a high-speed backup on a streaming tape drive (one that writes data to a tape in a single stream rather than in discrete blocks). The delay caused by reading the spare sector interrupts the data flow to the streaming tape drive.

You need a special utility program to manipulate the spare-sector directory, both to substitute for a failing sector manually and to attempt to rewrite a spare sector back onto the bad sector.

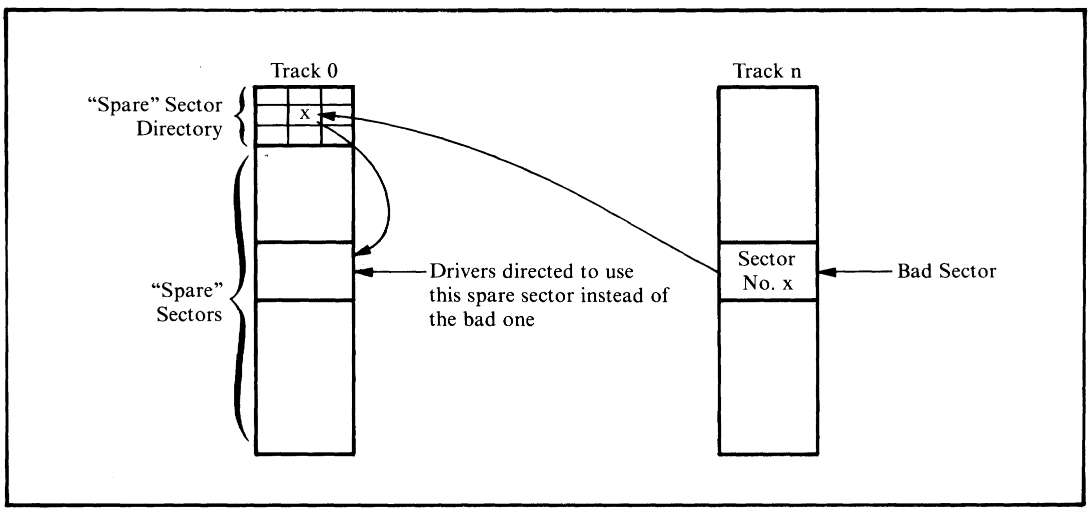


Figure 9-2. Sector sparing

Figure 9-3 shows another scheme for dealing with bad sectors. In this method, bad sectors are skipped rather than having sectors substituted for them.

The advantage of sector skipping is that the heads do not have to perform any long seeks. The failing sector is skipped, and the next sector is used in its place. Because of this, sector skipping can give much better performance. Data can be read off the disk fast enough to keep a streaming tape drive "fed" with data.

The disadvantage of sector skipping is that it does not lend itself to dynamic operation. The bad sector table is best built during formatting. Once data has been written to the disk, if a sector goes bad, all subsequent sectors on the disk must be "moved down one" to make space to skip the bad sector. On a large hard disk, this could take several minutes.

Example Bad Sector Management

Sector sparing and sector skipping use similar logic. Both require a spare-sector directory on each physical disk, containing the sector numbers of the bad sectors. This directory is read into memory during cold start initialization. Thereafter, all disk read and write operations refer to the memory-resident table to see if they are about to access a bad sector.

For sector sparing, if the sector about to be read or written is found in the spare directory, its position in the directory determines which spare sector should be read.

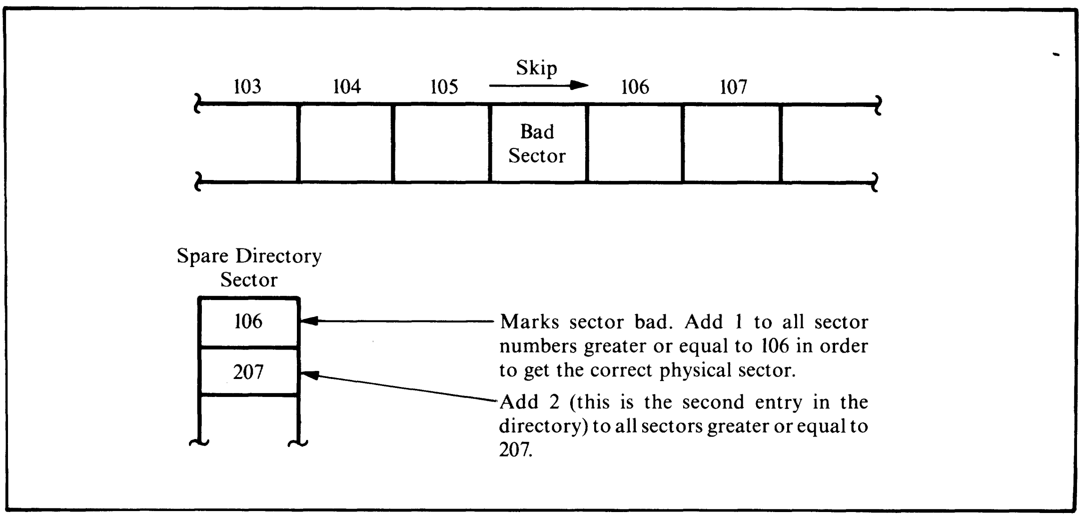


Figure 9-3. Sector skipping

In the case of sector skipping, every access to the disk makes the driver check the bad sector directory. The directory is used to tell how many bad sectors exist between the start of the disk and the failing bad sector. This number must be added to the requested track and sector to compensate for all the bad sectors.

The physical low-level drivers need four entry points:

- Read the specified sector without using bad sector management. This is used to read in the spare directory itself.
- Write the specified sector without using bad sector management. This is used to write the spare directory onto the disk, both to initialize it and to update it.
- Read and write the sector using bad sector management. These entry points are used for normal disk input/output.

Figure 9-4 shows the code necessary for both sector sparing and (using conditional code) sector skipping.

```
                ;       This example shows the modifications to be made in order
                ;       to implement bad sector management using sector sparing
                ;       and sector skipping.
                ;
0000 =          False           EQU     0
FFFF =          True            EQU     Not False
                ;
0000 =          Sector$Sparing  EQU     False
FFFF =          Sector$Skipping EQU     Not Sector$Sparing
                ;
                ;
                ;       Additional equates and definitions
                ;
                Spare$Directories:              ;Table of spare directory addresses
                                                        ;Note: The directories themselves
                                                        ;  are declared at the end of the
                                                        ;  BIOS
0000 D500               DW      Spare$Directory$0       ;Physical disk 0
0002 9701               DW      Spare$Directory$1       ;Physical disk 1
                ;
                Spare$Dir$In$Memory:            ;Flags used to indicate whether spare
0004 00                 DB      0               ;  directory for a given physical disk
0005 00                 DB      0               ;  has been loaded into memory. Set by SELDSK
                ;
                ;
0000 =          Spare$Track             EQU     0       ;Track containing spare directory
                                                        ;  sectors
0004 =          Spare$Sector            EQU     4       ;Sector containing directory
0005 =          First$Spare$Sector      EQU     Spare$Sector + 1
                ;
                ;       Variables set by SELDSK
                ;
                Selected$Spare$Directory:
0006 0000                               DW      0       ;Pointer to directory
0008 00         Selected$Disk:          DB      0       ;Logical disk number
0009 00         Disk$Type:              DB      0       ;Floppy/hard disks
000A 00         Deblocking$Required:    DB      0       ;Deblocking flag
000B 00         Selected$Physical$Disk: DB      0       ;Physical disk number
                ;
000C 0000       Disk$Track:     DW      0       ;) These variables are part of the command
000E 00         Disk$Sector:    DB      0       ;) block handed over to the disk controller
```

Figure 9-4. Bad sector management

```
                ;
8000 =          Maximum$Track           EQU     32768   ;Used as a terminator
0012 =          Sectors$Per$Track       EQU     18
0000 =          First$Sector$On$Track   EQU     0
                ;
                ;
                Disk$Parameter$Headers:
                        ;-------------------------
                        ;Standard DPH Declarations
                        ;-------------------------
                ;
                ;
                ;       Equates for disk parameter block
                ;
                ;       The special disk parameter byte that precedes each disk
                ;       parameter block, needs to be rearranged so that a
                ;       physical disk drive number can be added.
                ;
                ;       Disk types
                ;
                ;                               vvvv--- Physical disk number
0010 =          Floppy$5        EQU     0$001$0000B     ;5 1/4" mini floppy
0020 =          Floppy$8        EQU     0$010$0000B     ;8" floppy (SS SD)
0030 =          M$Disk          EQU     0$011$0000B     ;Memory disk
0040 =          H$Disk$10       EQU     0$100$0000B     ;Hard disk - 10 megabyte
                ;
0070 =          Disk$Type$$Mask         EQU     0$111$0000B     ;Masks to isolate values
000F =          Physical$Disk$Mask      EQU     0$000$1111B
                ;
                ;       Blocking/deblocking indicator
                ;
0080 =          Need$Deblocking EQU     1$000$0000B     ;Sector size > 128 bytes
                ;
                ;
                ;       Disk parameter blocks
                ;
                        ;----------------------------
                        ; Standard DPB's for A: and B:
                        ;----------------------------
                ;
                                                ;Logical disk C:
                                                ;Extra byte indicates disk type
                                                ; deblocking requirements and physical
                                                ; disk drive.
000F C0                 DB      H$Disk$10 + Need$Deblocking + 0 ; Physical drive 0
                Hard$5$Parameter$Block$C:
                        ;------------------------------
                        ;Standard format parameter block
                        ;------------------------------
                ;
                ;
0010 C0                 DB      H$Disk$10 + Need$Deblocking + 0 ; Physical drive 0
                Hard$5$Parameter$Block$D:
                        ;------------------------------
                        ;Standard format parameter block
                        ;------------------------------
                ;
                ;
0004 =          Number$of$Logical$Disks         EQU     4
                ;
                ;
                SELDSK:                 ;Select disk in register C
                                        ;C = 0 for drive A, 1 for B, etc.
                                        ;Return the address of the appropriate
                                        ; disk parameter header in HL, or 0000H
                                        ; if the selected disk does not exist.
                                        ;
0011 210000             LXI     H,0             ;Assume an error
0014 79                 MOV     A,C             ;Check if requested disk valid

0015 FE04               CPI     Number$of$Logical$Disks
0017 D0                 RNC                     ;Return if > maximum number of disks
```

Figure 9-4. (Continued)

```
0018 320800             STA     Selected$Disk   ;Save selected disk number
                                                ;Set up to return DPH address
001B 6F                 MOV     L,A             ;Make disk into word value
001C 2600               MVI     H,0
                                                ;Compute offset down disk parameter
                                                ; header table by multiplying by
                                                ; parameter header length (16 bytes)
001E 29                 DAD     H               ;*2
001F 29                 DAD     H               ;*4
0020 29                 DAD     H               ;*8
0021 29                 DAD     H               ;*16
0022 110F00             LXI     D,Disk$Parameter$Headers        ;Get base address
0025 19                 DAD     D               ;DE -> appropriate DPH
0026 E5                 PUSH    H               ;Save DPH address

                                                ;Access disk parameter block in order
                                                ; to extract special prefix byte that
                                                ; identifies disk type and whether
                                                ; deblocking is required
                                                ;
0027 110A00             LXI     D,10            ;Get DPB pointer offset in DPH
002A 19                 DAD     D               ;DE -> DPB address in DPH
002B 5E                 MOV     E,M             ;Get DPB address in DE
002C 23                 INX     H
002D 56                 MOV     D,M
002E EB                 XCHG                    ;DE -> DPB

                SELDSK$Set$Disk$Type:
002F 2B                 DCX     H               ;DE -> prefix byte
0030 7E                 MOV     A,M             ;Get prefix byte
0031 E670               ANI     Disk$Type$Mask  ;Isolate disk type
0033 320900             STA     Disk$Type       ;Save for use in low-level driver
0036 7E                 MOV     A,M             ;Get another copy of prefix byte
0037 E680               ANI     Need$Deblocking         ;Isolate deblocking flag
0039 320A00             STA     Deblocking$Required     ;Save for use in low-level driver

                                                ;Additional code to check if spare
                                                ;  directory for given disk has already
                                                ;  been read in.

003C 7E                 MOV     A,M             ;Get physical disk number
003D E60F               ANI     Physical$Disk$Mask
003F 320B00             STA     Selected$Physical$Disk  ;Save for low-level drivers

0042 5F                 MOV     E,A             ;Make into word
0043 1600               MVI     D,0
0045 210400             LXI     H,Spare$Dir$In$Memory   ;Make pointer into table
0048 19                 DAD     D

0049 7E                 MOV     A,M             ;Get flag
004A B7                 ORA     A
004B C27700             JNZ     Dir$In$Memory   ;Spare directory already in memory
004E 34                 INR     M               ;Set flag

004F 210000             LXI     H,Spare$Directories     ;Create pointer to spare
0052 19                 DAD     D                       ;  spare directory (added twice
0053 19                 DAD     D                       ;  as table has word entries)

                                                        ;HL -> word containing directory addr.
0054 5E                 MOV     E,M
0055 23                 INX     H
0056 56                 MOV     D,M                     ;Spare directory address in DE
0057 EB                 XCHG                            ;HL -> spare directory

0058 220600             SHLD    Selected$Spare$Directory ;Save for use in physical
                                                        ;  drivers later on

005B 110000             LXI     D,Spare$Track           ;Track containing spare directory
005E 3A0B00             LDA     Selected$Physical$Disk
0061 47                 MOV     B,A
0062 3E04               MVI     A,Spare$Sector          ;Sector containing spare directory
0064 0E18               MVI     C,Spare$Length/8        ;Number of bytes in spare directory / 8
0066 CDD500             CALL    Absolute$Read           ;Read in spare directory - without
                                                        ;  using bad sector management
```

Figure 9-4. (Continued)

```
0069 2A0600          LHLD    Selected$Spare$Directory ;Set end marker
006C 11C000          LXI     D,Spare$Length          ;  at back end of spare directory
006F 19              DAD     D
0070 110080          LXI     D,Maximum$Track         ;Use maximum track number
0073 73              MOV     M,E
0074 23              INX     H
0075 3602            MVI     M,D

             Dir$In$Memory:
0077 E1              POP     H               ;Recover DPH pointer
0078 C9              RET
             ;
             ;
             ;       In the low-level disk drivers, the following code must be
             ;       inserted just before the disk controller is activated to
             ;       execute a read or a write command.
             ;
0079 2A0C00          LHLD    Disk$Track              ;Get track number from disk
                                                     ;  controller command table
007C EB              XCHG                            ;DE = track
007D 2A0600          LHLD    Selected$Spare$Directory ;HL -> spare directory
0080 2B              DCX     H                       ;Back up one entry
0081 2B              DCX     H                       ;  (3 bytes)
0082 2B              DCX     H

0083 3A0E00          LDA     Disk$Sector             ;Get sector number
0086 4F              MOV     C,A                     ;Save for later

0087 06FF            MVI     B,0FFH                  ;Set counter (biased -1)

             Check$Next$Entry:
0089 23              INX     H                       ;Update to next (or first) entry
             Check$Next$Entry1:
008A 23              INX     H
             Check$Next$Entry2:
008B 23              INX     H

008C 04              INR     B                       ;Update count

                     IF      Sector$Sparing
                                                     ;If sparing is used, the
                                                     ;  end of the table is indicated
                                                     ;  by an entry with the track number
                                                     ;  = to maximum track number
                     LXI     D,Maximum$Track         ;Get maximum track number
                     CALL    CMPM                    ;Compare DE to (HL), (HL+1)
                     JZ      Not$Bad$Sector          ;End of table reached
                     ENDIF

                                                     ;Note: For sector skipping
                                                     ;  the following search loop will
                                                     ;  terminate when the requested track
                                                     ;  is less than that in the table.
                                                     ;This will always happen when the
                                                     ;  maximum track number is encountered
                                                     ;  at the end of the table.
008D EB              XCHG                            ;DE -> table entry
008E 2A0C00          LHLD    Disk$Track              ;Get requested track
0091 EB              XCHG                            ;DE = req. track, HL -> table entry
0092 CDCD00          CALL    CMPM                    ;Compare req. track to table entry


                     IF      Sector$Sparing          ;Use the following code for
                                                     ;  sector sparing
                     JNZ     Check$Next$Entry        ;Track does not match
                     INX     H                       ;HL -> MS byte of track
                     INX     H                       ;HL -> sector
                     MOV     A,C                     ;Get requested sector
                     CMP     M                       ;Compare to table entry
                     JNZ     Check$Next$Entry2       ;Sector does not match

                                                     ;Track and sector match, so
                                                     ;  substitute spare track and
                                                     ;  appropriate sector
```

Figure 9-4. (Continued)

```
                        LXI     H,Spare$Track           ;Get track number used for spare
                                                        ;  sectors
                        SHLD    Disk$Track              ;Substitute track

                        MVI     A,First$Spare$Sector    ;Get first sector number
                        ADD     B                       ;Add on matched directory
                                                        ;  entry number
                        STA     Disk$Sector             ;Substitute sector
                        ENDIF

                        IF      Sector$Skipping         ;Use the following code for
                                                        ;  sector skipping
                                                        ;The object is to find the
                                                        ;  entry in the table which
                                                        ;  is greater or equal to the
                                                        ;  requested sector/track

0095 CA9E00             JZ      Tracks$Match            ;Possible match of track and sector
0098 D2AC00             JNC     Compute$Increment       ;Requested track < table entry
009B C38900             JMP     Check$Next$Entry        ;Requested track > table entry

                Tracks$Match:
009E 23                 INX     H                       ;HL -> MS byte of track
009F 23                 INX     H                       ;HL -> sector
00A0 77                 MOV     M,A                     ;Get sector from table

00A1 B9                 CMP     C                       ;Compare with requested sector
00A2 CAAB00             JZ      Sectors$Match           ;Track/sector matches
00A5 D2AC00             JNC     Compute$Increment       ;Req. trk/sec < spare trk/sec
00A8 C38B00             JMP     Check$Next$Entry2       ;Move to next table entry

                Sectors$Match:
00AB 04                 INR     B                       ;If track and sectors match with
                                                        ;  a table entry, then an additional
                                                        ;  sector must be skipped
                Compute$Increment:
                                                        ;B contains number of cumulative
                                                        ;  number of sectors to skip
00AC 79                 MOV     A,C                     ;Get requested sector
00AD 80                 ADD     B                       ;Skip required number
00AE 0612               MVI     B,Sectors$Per$Track     ;Determine final sector number
                                                        ;  and track increment
00B0 CDC300             CALL    DIV$A$BY$B              ;Returns C = quotient, A = remainder
00B3 320E00             STA     Disk$Sector             ;A = new sector number

00B6 59                 MOV     E,C                     ;Make track increment a word
00B7 1600               MVI     D,0
00B9 2A0C00             LHLD    Disk$Track              ;Get requested track
00BC 19                 DAD     D                       ;Add on increment
00BD 220C00             SHLD    Disk$Track              ;Save updated track
                        ENDIF

                Not$Bad$Sector:
                                                        ;Either track/sector were not bad,
                                                        ;  or requested track and sector have
                                                        ;  been updated.
00C0 C3D500             JMP     Read$Write$Disk         ;Go to physical disk read/write
                ;
                        IF      Sector$Skipping
                                                        ;Subroutine required for skipping
                                                        ;  routine
                ;
                ;
                ;       DIV$A$BY$B
                ;       Divide A by B
                ;
                ;       This routine divides A by B, returning the quotient in C
                ;       and the remainder in A.
                ;
                ;       Entry parameters
                ;
                ;               A = dividend
                ;               B = divisor
                ;
                ;       Exit parameters
```

Figure 9-4. (Continued)

```
                        ;
                        ;               A = remainder
                        ;               C = quotient
                        ;
                        DIV$A$BY$B:
00C3 0E00                       MVI     C,0             ;Initialize quotient
                        DIV$A$BY$B$Loop:
00C5 0C                         INR     C               ;Increment quotient
00C6 90                         SUB     B               ;Subtract divisor
00C7 F2C500                     JP      DIV$A$BY$B$Loop ;Repeat if result still +ve
00CA 0D                         DCR     C               ;Correct quotient
00CB 80                         ADD     B               ;Correct remainder
00CC C9                         RET
                                ENDIF
                        ;
                        ;       CMPM
                        ;       Compare memory
                        ;
                        ;       This subroutine compares the contents of DE to (HL) and (HL+1)
                        ;       returning with the flags as though the subtraction (HL) - DE
                        ;       were performed.
                        ;
                        ;       Entry parameters
                        ;
                        ;               HL -> word in memory
                        ;               DE = value to be compared
                        ;
                        ;       Exit parameters
                        ;
                        ;               Flags set for (HL) - DE
                        ;
                        CMPM:
00CD 7E                         MOV     A,M                     ;Get MS byte
00CE BA                         CMP     D
00CF C0                         RNZ                             ;Return now if MS bytes unequal
00D0 23                         INX     H                       ;HL -> LS byte
00D1 7E                         MOV     A,M                     ;Get LS byte
00D2 BB                         CMP     E
00D3 2B                         DCX     H                       ;Return with HL unchanged
00D4 C9                         RET

                        ;
                        ;
                        Absolute$Read:
                                                ;The absolute read (and write) routines
                                                ;  access the specified sector and track
                                                ;  without using bad sector management.
                        ;       Entry parameters
                        ;
                        ;               HL -> Buffer
                        ;               DE = Track
                        ;               A = Sector
                        ;               B = Physical disk drive number
                        ;               C = Number of bytes to read / 8
                        ;
                        ;       Set up disk controller command block with parameters in
                        ;       registers, then initiate read operation by falling through
                        ;       into Read$Write$Disk code below.
                        ;
                        ;
                        Read$Write$Disk:

                                ;-------------------------------------------------
                                ;The remainder of the low level disk drivers follow,
                                ;  reading the required sector and track.
                                ;-------------------------------------------------

                        ;
                        ;       Spare directory declarations
                        ;
                        ;       Note: The disk format utility creates an initial spare
                        ;       directory with track/sector entries for those track/sectors
                        ;       that it finds are bad. It fills the remainder of the
                        ;       directory with OFFH's (these serve to terminate the
                        ;       searching of the directory).
```

Figure 9-4. (Continued)

```
                    ;
                    ;
00C0 =              Spare$Length    EQU     64 * 3          ;64 Entries, 3 bytes each
                                                            ;  Byte 0,1 = track
                                                            ;  Byte 2 = sector

                    Spare$Directory$0:
00D5                        DS      Spare$Length        ;Spare directory itself
0195                        DS      2                   ;Set to maximum track number by SELDSK as
                                                        ;  a safety precaution. The FORMAT utility
                                                        ;  puts the maximum track number into all
                                                        ;  unused entries in the spare directory.

                    Spare$Directory$1:
0197                        DS      Spare$Length        ;Spare directory itself
0257                        DS      2                   ;End marker
```

Figure 9-4. Bad sector management (continued)

Improving Error Messages

The final extension to BIOS error handling discussed here is in disk-driver error-message handling. The subroutine shown in the example BIOS in Figure 8-10, although a significant improvement on the messages normally output by the BDOS, did not advise the user of the most suitable course of action for each error. Figure 9-5 shows an improved version of the error message processor.

```
                    ;       This shows slightly more user-friendly error processor
                    ;       for disk errors than that shown in the enhanced BIOS
                    ;       in Figure 8-10.
                    ;       This version outputs a recommended course of action
                    ;       depending on the nature of the error detected.
                    ;       Code that remains unchanged from Figure 8-10 has been
                    ;       abbreviated.
                    ;
                    ;       Dummy equates and data declarations needed to get
                    ;       an error free assembly of this example.
                    ;
0001 =              Floppy$Read$Code        EQU     01H     ;Read command for controller
0002 =              Floppy$Write$Code       EQU     02H     ;Write command for controller
                    ;
0000 00             Disk$Hung$Flag:         DB      0       ;Set NZ when watchdog timer times
                                                            ;  out
0258 =              Disk$Timer              EQU     600     ;10-second delay (16.66ms tick)
                    ;
0043 =              Disk$Status$Block       EQU     43H     ;Address in memory where controller
                                                            ;  returns status
                                                            ;Values from controller command table
0001 00             Floppy$Command:         DB      0
0002 00             Floppy$Head:            DB      0
0003 00             Floppy$Track:           DB      0
0004 00             Floppy$Sector:          DB      0
```

Figure 9-5. User-friendly disk-error processor

```
0005 00         Deblocking$Required:    DB      0       ;Flag set by SELDSK according
                                                        ;  to selected disk type

0006 00         Disk$Error$Flag:        DB      0       ;Error flag returned to BDOS
                ;
0007 00         In$Buffer$Disk:         DB      0       ;Logical disk Id. relating to current
                                                        ;  disk sector in deblocking buffer
                ;
                ;       Equates for Messages
                ;
0007 =          BELL    EQU     07H     ;Sound terminal bell
000D =          CR      EQU     0DH     ;Carriage return
000A =          LF      EQU     0AH     ;Line feed
                ;
0005 =          BDOS    EQU     5       ;BDOS entry point (for system reset)
                ;
                ;
                ;
                ;
                No$Deblock$Retry:
                        ;---------------------------------------------------
                        ; Omitted code to set up disk controller command table
                        ; and initiate the disk operation
                        ;---------------------------------------------------
0008 C31500             JMP     Wait$For$Disk$Complete
                ;
                ;
                Write$Physical:                 ;Write contents of disk buffer to
                                                ;  correct sector
000B 3E02               MVI     A,Floppy$Write$Code     ;Get write function code

000D C31200             JMP     Common$Physical ;Go to common code
                Read$Physical:                  ;Read previously selected sector
                                                ;  into disk buffer
0010 3E01               MVI     A,Floppy$Read$Code      ;Get read function code
                Common$Physical:
0012 320100             STA     Floppy$Command  ;Set command table

                ;
                Deblock$Retry:                  ;Re-entry point to retry after error
                        ;--------------------------------------------------
                        ; Omitted code sets up disk controller command block
                        ; and initiates the disk operation
                        ;--------------------------------------------------
                ;
                Wait$For$Disk$Complete:         ;Wait until disk status block indicates
                                                ;  operation has completed, then check
                                                ;  if any errors occurred
                                                ;On entry HL -> disk control byte
0015 AF                 XRA     A                       ;Ensure hung flag clear
0016 320000             STA     Disk$Hung$Flag

0019 213100             LXI     H,Disk$Timed$Out        ;Set up watchdog timer
001C 015802             LXI     B,Disk$Timer            ;Time delay
001F CD3B03             CALL    Set$Watchdog
                Disk$Wait$Loop:
0022 7E                 MOV     A,M                     ;Get control byte
0023 B7                 ORA     A
0024 CA3700             JZ      Disk$Complete           ;Operation done

0027 3A0000             LDA     Disk$Hung$Flag          ;Also check if timed out
002A B7                 ORA     A
002B C29F02             JNZ     Disk$Error              ;Will be set to 40H

002E C32200             JMP     Disk$Wait$Loop

                Disk$Timed$Out:                 ;Control arrives here from watchdog
                                                ;  routine itself -- so this is effectively
                                                ;  part of the interrupt service routine.
0031 3E40               MVI     A,40H                   ;Set disk hung error code
0033 320000             STA     Disk$Hung$Flag          ;  into error flag to pull
                                                        ;  control out of loop
0036 C9                 RET                             ;Return to watchdog routine
```

Figure 9-5. (Continued)

```
                        Disk$Complete:
0037 010000                     LXI     B,0                     ;Reset watchdog timer
                                                                ;HL is irrelevant here
003A CD3B03                     CALL    Set$Watchdog

003D 3A4300                     LDA     Disk$Status$Block       ;Complete -- now check status
0040 FE80                       CPI     80H                     ;Check if any errors occurred
0042 DA9F02                     JC      Disk$Error              ;Yes
                        ;
                        Disk$Error$Ignore:
0045 AF                         XRA     A                       ;No
0046 320600                     STA     Disk$Error$Flag         ;Clear error flag
0049 C9                         RET

                        ;
                        ;       Disk error message handling
                        ;
                        ;
                        Disk$Error$Messages:            ;This table is scanned, comparing the
                                                        ; disk error status with those in the
                                                        ; table. Given a match, or even when
                                                        ; the end of the table is reached, the
                                                        ; address following the status value
                                                        ; points to the correct advisory message text.
                                                        ; Following this is the address of an
                                                        ; error description message.
004A 40                         DB      40H
004B B0019500                   DW      Disk$Advice1,Disk$Msg$40
004F 41                         DB      41H
0050 C9019A00                   DW      Disk$Advice2,Disk$Msg$41
0054 42                         DB      42H
0055 E301A400                   DW      Disk$Advice3,Disk$Msg$42
0059 21                         DB      21H
005A 0702B400                   DW      Disk$Advice4,Disk$Msg$21
005E 22                         DB      22H
005F 1B02B900                   DW      Disk$Advice5,Disk$Msg$22
0063 23                         DB      23H
0064 1B02C000                   DW      Disk$Advice5,Disk$Msg$23
0068 24                         DB      24H
0069 3D02D200                   DW      Disk$Advice6,Disk$Msg$24
006D 25                         DB      25H
006E 3D02DE00                   DW      Disk$Advice6,Disk$Msg$25
0072 11                         DB      11H
0073 5302F100                   DW      Disk$Advice7,Disk$Msg$11
0077 12                         DB      12H
0078 5302FF00                   DW      Disk$Advice7,Disk$Msg$12
007C 13                         DB      13H
007D 53020C01                   DW      Disk$Advice7,Disk$Msg$13
0081 14                         DB      14H
0082 53021A01                   DW      Disk$Advice7,Disk$Msg$14
0086 15                         DB      15H
0087 53022901                   DW      Disk$Advice7,Disk$Msg$15
008B 16                         DB      16H
008C 53023501                   DW      Disk$Advice7,Disk$Msg$16
0090 00                         DB      0                       ;<== Terminator
0091 53024501                   DW      Disk$Advice7,Disk$Msg$Unknown   ;Unmatched code
                        ;
0005 =                  DEM$Entry$Size  EQU     5       ;Entry size in error message table
                        ;
                        ;
                        ;       Message texts
                        ;
0095 48756E6700Disk$Msg$40:     DB      'Hung',0                ;Timeout message
009A 4E6F742052Disk$Msg$41:     DB      'Not Ready',0
00A4 5772697465Disk$Msg$42:     DB      'Write Protected',0
00B4 4461746100Disk$Msg$21:     DB      'Data',0
00B9 466F726D61Disk$Msg$22:     DB      'Format',0
00C0 4D69737369Disk$Msg$23:     DB      'Missing Data Mark',0
00D2 4275732054Disk$Msg$24:     DB      'Bus Timeout',0
00DE 436F6E7472Disk$Msg$25:     DB      'Controller Timeout',0
00F1 4472697665Disk$Msg$11:     DB      'Drive Address',0
00FF 4865616420Disk$Msg$12:     DB      'Head Address',0
010C 547261636BDisk$Msg$13:     DB      'Track Address',0
```

Figure 9-5. (Continued)

```
011A 536563746FDisk$Msg$14:    DB      'Sector Address',0
0129 4275732041Disk$Msg$15:    DB      'Bus Address',0
0135 496C6C6567Disk$Msg$16:    DB      'Illegal Command',0
0145 556E6B6E6FDisk$Msg$Unknown:       DB      'Unknown',0
                ;
                Disk$EM$1:                      ;Main disk error message -- part 1
014D 070D0A                    DB      BELL,CR,LF
0150 4469736B20                DB      'Disk ',0
                ;
                                                ;Error text output next
                ;
                Disk$EM$2:                      ;Main disk error message -- part 2
0156 204572726F                DB      ' Error ('
015E 0000       Disk$EM$Status: DB      0,0     ;Status code in hex
0160 290D0A2020                DB      ')',CR,LF,'    Drive '
016E 00         Disk$EM$Drive:  DB      0       ;Disk drive code, A,B...
016F 2C20486561                DB      ', Head '
0176 00         Disk$EM$Head:   DB      0       ;Head number
0177 2C20547261                DB      ', Track '
017F 0000       Disk$EM$Track:  DB      0,0     ;Track number
0181 2C20536563                DB      ', Sector '
018A 0000       Disk$EM$Sector: DB      0,0     ;Sector number
018C 2C204F7065                DB      ', Operation - '
019A 00                        DB      0               ;Terminator
                ;
019B 526561642EDisk$EM$Read:   DB      'Read.',0       ;Operation names
01A1 5772697465Disk$EM$Write:  DB      'Write.',0
                ;
01A8 0D0A202020Disk$Advice0:   DB      CR,LF,'     ',0
01B0 436865636BDisk$Advice1:   DB      'Check disk loaded, Retry',0
01C9 506F737369Disk$Advice2:   DB      'Possible hardware problem',0
01E3 5772697465Disk$Advice3:   DB      'Write enable if correct disk, Retry',0
0207 5265747279Disk$Advice4:   DB      'Retry several times',0
021B 5265666F72Disk$Advice5:   DB      'Reformat disk or use another disk',0
023D 4861726477Disk$Advice6:   DB      'Hardware error, Retry',0
0253 4861726477Disk$Advice7:   DB      'Hardware or Software error, Retry',0
                ;
0275 2C206F7220Disk$Advice9:   DB      ', or call for help if error persists',CR,LF
                ;
                Disk$Action$Confirm:
029B 00                        DB      0       ;Set to character entered by user
029C 0D0A00                    DB      CR,LF,0
                ;
                ;       Disk error processor
                ;
                ;       This routine builds and outputs an error message.
                ;       The user is then given the opportunity to:
                ;
                ;               R -- retry the operation that caused the error
                ;               I -- ignore the error and attempt to continue
                ;               A -- abort the program and return to CP/M
                ;
                Disk$Error:
029F F5                 PUSH    PSW             ;Preserve error code from controller
02A0 215E01             LXI     H,Disk$EM$Status        ;Convert code for message
02A3 CD3B03             CALL    CAH                     ;Converts A to hex

02A6 3A0700             LDA     In$Buffer$Disk          ;Convert disk id. for message
02A9 C641               ADI     'A'                     ;Make into letter
02AB 326E01             STA     Disk$EM$Drive

02AE 3A0200             LDA     Floppy$Head             ;Convert head number
02B1 C630               ADI     '0'
02B3 327601             STA     Disk$EM$Head

02B6 3A0300             LDA     Floppy$Track            ;Convert track number
02B9 217F01             LXI     H,Disk$EM$Track
02BC CD3B03             CALL    CAH

02BF 3A0400             LDA     Floppy$Sector           ;Convert sector number
02C2 218A01             LXI     H,Disk$EM$Sector
02C5 CD3B03             CALL    CAH

02C8 214D01             LXI     H,Disk$EM$1             ;Output first part of message
02CB CD3B03             CALL    Output$Error$Message
```

Figure 9-5. (Continued)

```
02CE F1                 POP     PSW                     ;Recover error status code
02CF 47                 MOV     B,A                     ;For comparisons
02D0 214500             LXI     H,Disk$Error$Messages - DEM$Entry$Size
                                                        ;HL -> table -- one entry
02D3 110500             LXI     D,DEM$Entry$Size        ;For loop below
                Disk$Error$Next$Code:
02D6 19                 DAD     D                       ;Move to next (or first) entry

02D7 7E                 MOV     A,M                     ;Get code number from table
02D8 B7                 ORA     A                       ;Check if end of table
02D9 CAE302             JZ      Disk$Error$Matched      ;Yes, pretend a match occurred
02DC B8                 CMP     B                       ;Compare to actual code
02DD CAE302             JZ      Disk$Error$Matched      ;Yes, exit from loop
02E0 C3D602             JMP     Disk$Error$Next$Code    ;Check next code
                ;
                Disk$Error$Matched:
02E3 23                 INX     H                       ;HL -> advisory text address
02E4 5E                 MOV     E,M
02E5 23                 INX     H
02E6 56                 MOV     D,M                     ;DE -> advisory test
02E7 D5                 PUSH    D                       ;Save for later

02E8 23                 INX     H                       ;HL -> message text address
02E9 5E                 MOV     E,M                     ;Get address into DE
02EA 23                 INX     H
02EB 56                 MOV     D,M

02EC EB                 XCHG                            ;HL -> text
02ED CD3B03             CALL    Output$Error$Message    ;Display explanatory text

02F0 215601             LXI     H,Disk$EM$2             ;Display second part of message
02F3 CD3B03             CALL    Output$Error$Message

02F6 219B01             LXI     H,Disk$EM$Read          ;Choose operation text
                                                        ;  (assume a read)
02F9 3A0100             LDA     Floppy$Command          ;Get controller command
02FC FE01               CPI,    Floppy$Read$Code
02FE CA0403             JZ      Disk$Error$Read         ;Yes
0301 21A101             LXI     H,Disk$EM$Write         ;No, change address in HL
                Disk$Error$Read:
0304 CD3B03             CALL    Output$Error$Message    ;Display operation type

0307 21A801             LXI     H,Disk$Advice0          ;Display leading blanks
030A CD3B03             CALL    Output$Error$Message

030D E1                 POP     H                       ;Recover advisory text pointer
030E CD3B03             CALL    Output$Error$Message

0311 217502             LXI     H,Disk$Advice9          ;Display trailing component
0314 CD3B03             CALL    Output$Error$Message
                ;
                Disk$Error$Request$Action:              ;Ask the user what to do next
0317 CD3B03             CALL    Request$User$Choice     ;Display prompt and get single
                                                        ;  character response (folded to
                                                        ;  uppercase)
031A FE52               CPI     'R'                     ;Retry
031C CA2C03             JZ      Disk$Error$Retry
031F FE41               CPI     'A'                     ;Abort?
0321 CA3603             JZ      System$Reset
0324 FE49               CPI     'I'                     ;Ignore?
0326 CA4500             JZ      Disk$Error$Ignore
0329 C31703             JMP     Disk$Error$Request$Action
                ;
                Disk$Error$Retry:                       ;The decision on where to return to
                                                        ;  depends on whether the operation
                                                        ;  failed on a deblocked or
                                                        ;  nondeblocked drive
032C 3A0500             LDA     Deblocking$Required
032F B7                 ORA     A
0330 C21500             JNZ     Deblock$Retry
0333 C30800             JMP     No$Deblock$Retry
```

Figure 9-5. (Continued)

```
                ;
                System$Reset:                           ;This is a radical approach, but
                                                        ;  it does cause CP/M to restart
0336 0E00               MVI     C,0                     ;System reset
0338 CD0500             CALL    BDOS

                ;
                ;       Omitted subroutines (listed in full in Figure 8-10)
                ;
                Set$Watchdog:           ;Set watchdog timer (to number of "ticks" in BC, and
                                        ;  to transfer control to (HL) if timer hits zero).
                CAH:                    ;Convert A to two ASCII hex characters, storing
                                        ;  the output in (HL) and (HL+1)
                Output$Error$Message:   ;Display the 00-byte terminated error message
                                        ;  pointed to by HL. Output is directed only to
                                        ;  those console devices not being used for list
                                        ;  output as well.
                Request$User$Choice:    ;Display prompt "Enter R, A, I..." and return
                                        ;  single keyboard character (uppercase) in A
033B C9                 RET             ;Dummy
```

Figure 9-5. User-friendly disk-error processor (continued)

10

Debugging A New CP/M System

This chapter deals with some of the problems you will face bringing up CP/M on a computer system for the first time or enhancing it once it is up and running on your system.

In the first case, when CP/M does not yet run on your computer, you may be writing the complete BIOS yourself, although you can model what you do on the example BIOS provided on the CP/M release diskette and the example code from Chapter 6.

In the second case, you can extend the existing BIOS by adding code—from the examples in Chapters 8 and 9, code from computer magazines, or code you create yourself. To do this, you will need access to the BIOS source code—a problem if the manufacturer of your computer does not make it available. In general, however, the BIOS source code is included with the system or can be obtained at nominal or no cost. If you cannot obtain the source code, you can, of course, take the bull by the horns and reimplement CP/M on your system. This may require many hours of disassembling the current BIOS machine code to find out how to access all the various ports and how to control the devices to which they are connected.

Although the BIOS is the major component of a new CP/M implementation, remember that it is only the beginning—you can spend the same amount of time and effort getting the bootstrap loader and all the utilities to function.

Basic Debugging Techniques

Before getting involved in the details of how to debug a CP/M implementation, it is worth considering the nature of the task. Some quotations that are appropriate here:

> "Program testing can be used to show the presence of bugs, but never to show their absence." —Dijkstra

> "We call them bugs because to call them mistakes would be psychologically unacceptable." —Hopkins

> "Constants aren't, variables won't." —Osborne

Debugging is the name we give to the process of executing programs and ascertaining whether the programs are running correctly. "Correctly" means in accordance with the mental model we have built of how the program should behave, subject to the constraints imposed by the physical hardware. Therein lies the first of the problems; you and the hardware are the arbiters of correct performance. The hardware is usually unforgiving; if there is a flaw in the way you program it, it will either be dramatically "uncooperative" or not work at all. As for how you perceive the system, several fairly simple tests, along with attempts to use the system for useful work for a few days, will shake the system down fairly well. The most difficult problems will be with intermittent failures or logical contradictions.

Computers are deterministic. That is, if you start from a known state and perform a known series of operations, the computer will always yield the same results. To achieve a known state is not so difficult—resetting the system and clearing memory will do it. Performing a known series of operations just means running the program again, although if you are using interrupts, you cannot truthfully say that exactly the same operations are being performed, because the interrupts will not happen at *exactly* the same time as before.

The "Orville Wright" Approach

Your role in debugging a new CP/M system is comparable to the popular, though untrue, idea of the way the Wright brothers developed flying machines:

build a machine, take it to the top of a hill, throw it off, and, when it crashes, examine the debris to discover what went wrong.

Each time you do an assembly and test, you are building the aircraft and lobbing it off the edge of a cliff. Each time it crashes, you examine the wreckage and try to determine the possible cause.

This is a highly inferential process. With the wreckage as a starting point, you use inference and intuition to extrapolate the real problem and the correction for it.

Built-In Debug Code

The single most important concept that you will need in testing CP/M systems is the same as that used in the modern day "black box" flight recorder. This device is essentially a multi-channel tape recorder that records all of the relevant conditions of the aircraft, its height, altitude, throttle settings, flap settings, and even the voice communications among crew members. If the airplane crashes, investigators can replay the information and understand what happened during the flight.

Applying this concept to debugging CP/M means that you must build into your code some method for recording what it is doing, so that if the system crashes, you can see what it was doing. Make the code tell you what went wrong.

The debug code should be designed at the same time as the rest of the program. Plan the debugging code while the design is still on the drawing board. The source code for debugging should be a permanent part of the BIOS. Use conditional assembly to "IF" out most of the debug code from the final version, or make the code sensitive to a flag in the configuration block so that you can re-enable the debug code at a moment's notice if the system begins to behave strangely.

The more meaningful the debug output data, the less you will have to guess at what is wrong, and therefore the less painful and time-consuming the debugging process will be. Make the output intelligible to others who may use it or yourself several months hence. Data that tells you what is happening is more useful than internal hexadecimal values, particularly if someone else must interpret it or relay it to you over the telephone.

Debug Subroutines

Many programmers do their debugging on a casual "catch as catch can" basis because they are overwhelmed by the task of building the necessary tools. Others are too eager to start on a new program to take a few extra hours or days to build debug subroutines.

To help solve this problem, the following section provides some ready-made debugging tools that can be used "as is." Each of these routines has been thoroughly debugged (there's nothing worse than debug code with bugs in it!) and has been used in actual program testing.

Overall Design Philosophy

Some common methods run through the examples that follow. These include displaying meaningful "captions" (including the specific address that called the debug routine), grouping all debugging code together, preserving the contents of all registers, and setting up the stack area in a standard way.

Debug Code Captions When the contents of registers or memory are output as part of a debugging process, a caption of explanatory text describing the values should be displayed. For example, rather than displaying the contents of the A register like this,

```
A = 1F
```

you can use a meaningful caption such as:

```
Transaction Code A = 1F.
```

When you write additional debugging code, especially if you need to add it to an existing routine, it is cumbersome to have to write the call to the debug routine and then search through the source code to find a convenient place to put an ASCII caption string. A caption string several pages removed from the point where it is referenced makes for problems when you want to relate the debug output on the screen or listing to the source code itself. Therefore, all of the routines that follow allow you to declare the caption strings "in-line" like this:

```
IF          DEBUG
CALL        Debug$Routine
DB          'Caption string here',CR,LF,0
ENDIF

MVI         .....             ;Next instruction
```

All of the following routines that output a caption recognize one specific 8-bit value in the caption string. If they encounter a value of 0ADH (mnemonic for ADdress), they will output the address of the byte following the call to the debug routine. For example,

```
0210        CALL        Debug$Routine
0213        DB          0ADH,'Caption string',0
```

will cause the routine to display the following:

```
0213 Caption string
```

This identifies the point in your program from which the debug routine was called, and thus avoids any possible ambiguity between different calls to the same debug routine with similar captions.

Grouping Debug Code Grouping all the debug code together lends itself to using conditional assembly with IF/ENDIF statements.

Setting Up the Stack Area All of the following routines preserve the CPU registers so that there are no side effects from using them. All of them assume that they can use the stack pointer and that there is sufficient room in the stack area. Hence you will need to declare adequate stack space for your main code and for the debug routines. Fill the stack area with a known pattern like this:

```
        DW    9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
        DW    9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
        DW    9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
Stack$Area:                    ;Label the upper end of the area
```

Then, during debugging, you can examine the stack area and determine how much of it is unused. For example, if you looked at the stack area you might see something like this:

```
                                  "Low-water mark"
                                         V
99 99 99 99 99 99 99 99 99 99 99 99 09 15 43 42
01 29 00 00 1A 2B 10 FF FF 39 02 ED 11 01 37 44
DD 00 00 11 1A 23 31 00 41 AE FE 00 01 10 70 C9
```

Stack area overflow can give arcane bugs; the program seems to leap off into space in a nondeterministic way. By setting up the stack area in this way, you can recognize an overflow condition easily.

Debug Initialization Before you can execute any of the debug subroutines in this chapter, you must make a call to the initialization subroutine, DB$Init. The DB$Init routine sets up some of the internal variables needed by the debug package. You may need to add some of your own initialization code here.

Console Output

Normally, you can use the CONOUT functions either via the BDOS (Function 2), or via the BIOS by calling the jump vector directly. You cannot do this when you need to debug console routines themselves, nor when you need to debug interrupt service routines. In the latter case, if an interrupt pulled control out of the CONOUT routine in the BIOS, you would get unwanted re-entrancy if the debug code again entered the CONOUT driver to display a caption. Therefore, the debug routines have been written to call their own local CONOUT routine, which is called DB$CONOUT. DB$CONOUT can be changed to call the BDOS, the BIOS, or a "private" polled output routine.

A counterpart DB$CONIN routine for console input is provided for essentially the same reasons.

Controlling Debug Output

All output of debug routines in this chapter is controlled by a single master flag, DB$Flag. If this flag is nonzero, debug output will occur; if zero, all output is suppressed.

This flag can be set and cleared from any part of the program you are testing. It is especially useful when you need to debug a subroutine that is called many times from many different places. You can write additional code to enable debug output when certain conditions prevail; for example, when a particular track or sector is about to be written or when a character input buffer is almost full.

Two subroutines, DB$On and DB$Off, are shown that access the debug control flag. These, as their names suggest, turn debug output on and off.

Turning the debug output on and off from within the program can create a confusing display of debug output, lacking any apparent continuity. DB$Off gives you the option of outputting a character string indicating that debug output has been turned off.

Pass Counters

Another method of controlling debug output is to use a *pass counter,* enabling debug output only after control has passed through a particular point in the code a specific number of times.

Two subroutines are provided for this purpose. DBSetPass sets the pass counter to a specific value. DB$Pass decrements this pass count each time control is transferred to it. When the pass count hits zero, the debug control flag DB$Flag is nonzero and debug output begins.

Using pass counter techniques can save you time and effort in tracking down a problem that occurs only after the code has been running for several minutes.

Displaying Contents of Registers and Memory

Figure 10-2 shows a series of display subroutines, the primary one of which is DB$Display. It takes several parameters, depending on the information you want displayed. The generic call to DB$Display is as follows:

```
CALL    DB$Display
DB      Code      <- Indicates the data to be
                     displayed
{DW     Optional additional parameters}
 DB     'Caption string',0
```

The codes that can be used in this call are shown in Table 10-1.

The only function that uses additional parameters is DB$Memory. This displays bytes from memory in hexadecimal and ASCII, using the start and finish addresses following the call. Here is an example:

```
CALL        DB$Display
DB          DB$Memory
DW          Start$Address,End$Address
DB          'Caption string',0
```

Table 10-1. Codes for DB$Display

Code	Value displayed
8-bit registers	
DB$F	Condition Flags
DB$A	Register A
DB$B	Register B
DB$C	Register C
DB$D	Register D
DB$E	Register E
DB$H	Register H
DB$L	Register L
Memory	
DB$Memory	Bytes starting and ending at the addresses specified by the two word values following the code value.
16-bit registers	
DB$BC	Register pair BC
DB$DE	Register pair DE
DB$HL	Register pair HL
DB$SP	Stack Pointer
Byte values	
DBBBC	Byte addressed by BC
DBBDE	Byte addressed by DE
DBBHL	Byte addressed by HL
Word values	
DBWBC	Word addressed by BC
DBWDE	Word addressed by DE
DBWHL	Word addressed by HL

Debugging Program Logic

In addition to displaying the contents of registers and memory, you need to display the program's execution path, not in terms of addresses, but in terms of the *problem.* You can do this by displaying debug messages that indicate what decisions have been made by the program as it executes. For example, if your BIOS checks a particular value to see whether the system should read or write on a particular device, the debug routine should display a message like this:

```
Entering Disk Read Routine
```

This is more meaningful than just displaying the function code for the drivers — although you may want to display this as well, in case it has been set to some strange value.

Two subroutines are provided to display debug messages. They are DB$MSG and DB$MSGI. Both of these display text strings are terminated with a byte of 00H. You can see the difference between the two subroutines if you examine the way they are called.

DB$MSG is called like this:

```
LXI  H,Message$Text        ;HL -> text string
CALL DB$MSG
```

DB$MSGI is called like this:

```
CALL DB$MSG
DB   0DH,0AH,'Message Text',0 ;In-line
```

DB$MSGI is more convenient to use. If you decide that you need to add a message, you can declare the message immediately following the call. This also helps when you look at the listing, since you can see the complete text at a glance.

Use DB$MSG when the text of the message needs to be selected from a table. Get the address of the text into HL and then call DB$MSG to display it.

Creating Your Own Debug Displays

If you need to build your own special debug display routines, you may find it helpful to incorporate some of the small subroutines in the debug package. The following are the subroutines you may want to use:

DB$CONOUT
Displays the character in the C register.

DB$CONIN
Returns the next keyboard character in A.

DB$CONINU
Returns the next keyboard character in A, converting lowercase letters to uppercase.

DB$DHLH
Displays contents of HL in hexadecimal.

DB$DAH
Displays contents of A in hexadecimal.

DB$CAH
Converts contents of A to hexadecimal and stores in memory pointed at by HL.

DB$Nibble$To$Hex
Converts the least significant four bits of A into an ASCII hexadecimal character in A.

DB$CRLF
Displays a CARRIAGE RETURN/LINE FEED.

DB$Colon
Displays the string " : ".

DB$Blank
Displays a single space character.

DB$Flag$Save$On
Saves the current state of the debug output control flag and then sets the flag "on" to enable debug output.

DB$Flag$Restore
Restores the debug output control flag to the state it was in when the DB$Flag$Save$On routine was last called.

DB$GHV
Gets a hexadecimal value from the keyboard, displaying a prompt message first. From one to four characters can be specified as the maximum number of characters to be input.

DBATo$Upper
If the A register contains a lowercase letter, this converts it to an uppercase letter.

Debugging I/O Drivers

Debugging low-level device drivers creates special problems. The major one is that you do not normally want to read and write via actual hardware ports while you are debugging the code—either because doing so would cause strange things to happen to the hardware during the debugging, or because you are developing and debugging the drivers on a system different from the target hardware on which the drivers are to execute.

Before considering the solution, remember that the input and output instructions (IN and OUT) are each two bytes long. The first byte is the operation code

(0DBH for input, 0D3H for output), and the second byte is the port number to "input from" or "output to."

Debug subroutines are provided here to intercept all IN and OUT instructions, displaying the port number and either accepting a hexadecimal value from the console and putting it into the A register (in the case of IN), or displaying the contents of the A register (for the OUT instruction).

IN and OUT instructions can be "trapped" by changing the operation code to one of two RST (restart) instructions. An RST is effectively a single-byte CALL instruction, calling down to a predetermined address in low memory. The debug routines arrange for JMP instructions in low memory to receive control when the correct RST is executed. The code that receives control can pick up the port number, display it, and then accept a hex value for the A register (for IN) or display the current contents of the A register (for OUT). The example subroutines shown later in this chapter use RST 4 in place of IN instructions, RST 5 for OUT.

Wherever you plan to use IN, use the following code:

```
IF          Debug
RST         4
ENDIF
IF          NOT Debug
DB          IN
ENDIF
DB          Port$Number
```

Note that you can use the IN operation code as the operand of a DB statement. The assembler substitutes the correct operation code.

Use the following code wherever you need to use an OUT instruction:

```
IF          Debug
RST         5
ENDIF
IF          NOT Debug
DB          OUT
ENDIF
DB          Port$Number
```

When the RST 4 (IN) instruction is executed, the debug subroutine displays

```
1AB3 : Input from Port 01 : _
```

The "1AB3" is the address in memory of the byte containing the port number. It serves to pinpoint the IN instruction in memory. You can then enter one or two hexadecimal digits. These will be converted and put into the A register before control returns to the main program at the instruction following the byte containing the port number.

When the RST 5 (OUT) instruction is encountered, the debug subroutine displays

```
1AB5 :  Output to Port 01 : FF
```

This identifies where the OUT instruction would normally be as well as the port number and the contents of the A register when the RST 5 (OUT) is executed.

Debugging Interrupt Service Routines

You can use a technique similar to that of the RST instruction just described to "fake" an interrupt. You preset the low-memory address for the RST instruction you have chosen for the jump into the interrupt service routine under test.

When the RST instruction is executed, control will be transferred into the interrupt service routine just as though an interrupt had occurred. You will need to intercept any IN or OUT instructions as described above—otherwise the code probably will go into an endless loop.

Before executing the RST instruction to fake the interrupt, load all the registers with known values. For example:

```
MVI     A,0AAH
LXI     B,0BBCCH
LXI     D,0DDEEH
LXI     H,01122H
RST     6          ;Fake interrupt
NOP
```

When control returns from the service routine, you can check to see that it restored all of the registers to their correct values. An interrupt service routine that does not restore all the registers can produce bugs that are very hard to find.

Check, too, that the stack pointer register has been restored and that the service routine did not require too many bytes on the stack.

You also can use the CALL instruction to transfer control to the interrupt service routine in order to fake an interrupt. RST and CALL achieve the same effect, but RST is closer to what happens when a real interrupt occurs. As it is a single-byte instruction, it also is easier to patch in.

Subroutine Listings

Figure 10-1 is a functional index to the source code listing for the debug subroutines shown in Figure 10-2. The listing's commentary defines precisely how each debug subroutine is called.

Figure 10-3 shows the output from the debug testbed.

Software Tools for Debugging

In addition to building in debugging subroutines, you will need one of the following proprietary debug programs:

DDT (Dynamic Debugging Tool)

This program, included with the standard CP/M release, allows you to load programs, set and display memory and registers, trace through your program instruction by instruction, or execute it at full speed, but stopping

Start Line	Functional Component or Routines
00001	Debug subroutine's Testbed
00100	Test register display
00200	Test memory dump display
00300	Test register pair display
00400	Test byte indirect display
00500	Test DB$On/Off
00600	Test DBSetPass and DB$Pass
00700	Test debug input/output
00800	Debug subroutines themselves
01100	DB$Init - initialization
01200	DB$CONINU - get uppercase keyboard character
01300	DB$CONIN - get keyboard character
01400	DB$CONOUT - display character in C
01500	DB$On - enable debug output
01600	DB$Off - disable debug output
01700	DBSetPass - set pass counter
01800	DB$Pass - execute pass point
01900	DB$Display - main debug display routine
02200	Main display processing subroutines
02500	DB$Display$CALLA - display CALL's address
02600	DB$DHLH - display HL in hexadecimal
02700	DB$DAH - display A in hexadecimal
02800	DB$CAH - convert A to hexadecimal in memory
02900	DB$Nibble$To$Hex - convert LS 4 bits of A to hex.
02930	DB$CRLF - display Carriage Return, Line Feed
02938	DB$Colon - display " : "
02946	DB$Blank - display " "
03100	DB$MSGI - display in-line message
03147	DB$MSG - display message addressed by HL
03300	DB$Input - debug INput routine
03500	DB$Output - debug OUTput routine
03700	DB$Flag$Save$On - save debug flag and enable
03800	DB$Flag$Restore - restore debug control flag
03900	DB$GHV - get hexadecimal value from keyboard
04100	DBATo$Upper - convert A to upper case

Figure 10-1. Functional index for Figure 10-2

at certain addresses (called breakpoints). It also has a built-in mini-assembler and disassembler so you do not have to hand assemble any temporary code "patches" you add.

SID (Symbolic Interactive Debug)

Similar to DDT in many ways, SID has enhancements that are helpful if you use Digital Research's MAC (Macro Assembler) or RMAC (Relocating Macro Assembler). Both of these assemblers can be told to output a file

```
                        ;
                        ;       Debug Subroutines
                        ;
                        ;<----  NOTE:
                        ;       The line numbers at the extreme left are included purely
                        ;       to reference the code from the text.
                        ;       There are deliberately induced discontinuities
                        ;       in the numbers in order to allow space for expansion.
                        ;
                        ;       Because of the need to test these routines thoroughly,
                        ;       and in case you wish to make any changes, the testbed
                        ;       routine for the debug package itself has been left in
                        ;       in this figure.
                        ;
                        ;       Debug testbed
                        ;
0100                            ORG     100H
                        START:
0100 316B03                     LXI     SP,Test$Stack           ;Set up local stack
0103 CDEA04                     CALL    DB$Init                 ;Initialize the debug package
0106 CD1505                     CALL    DB$On                   ;Enable debug output
                                                                ;Simple test of A register display
0109 3EAA                       MVI     A,0AAH                  ;Preset a value in the A register
010B 01CCBB                     LXI     B,0BBCCH                ;Prefill all other registers, partly
010E 11EEDD                     LXI     D,0DDEEH                ;  to check the debug display, but
0111 2111FF                     LXI     H,0FF11H                ;  also to check register save/restore
                        ;#
                        ;       Test register display
                        ;
0114 B7                         ORA     A                       ;Set M-flag, clear Z-flag, set E-flag
0115 37                         STC                             ;Set carry
0116 CD5205                     CALL    DB$Display              ;Call the debug routine
0119 00                         DB      DB$F
011A 466C616773                 DB      'Flags',0
                        ;
0120 CD5205                     CALL    DB$Display              ;Call the debug routine
0123 02                         DB      DB$A
0124 4120526567                 DB      'A Register',0
                        ;
012F CD5205                     CALL    DB$Display              ;Call the debug routine
0132 04                         DB      DB$B
0133 4220526567                 DB      'B Register',0
                        ;
013E CD5205                     CALL    DB$Display              ;Call the debug routine
0141 06                         DB      DB$C
0142 4320526567                 DB      'C Register',0
                        ;
014D CD5205                     CALL    DB$Display              ;Call the debug routine
0150 08                         DB      DB$D
0151 4420526567                 DB      'D Register',0
                        ;
015C CD5205                     CALL    DB$Display              ;Call the debug routine
015F 0A                         DB      DB$E
0160 4520526567                 DB      'E Register',0
                        ;
016B CD5205                     CALL    DB$Display              ;Call the debug routine
016E 0C                         DB      DB$H
016F 4820526567                 DB      'H Register',0
                        ;
017A CD5205                     CALL    DB$Display              ;Call the debug routine
017D 0E                         DB      DB$L
017E 4C20526567                 DB      'L Register',0
                        ;#
                        ;       Test Memory Dump Display
                        ;
0189 CD5205                     CALL    DB$Display
018C 18                         DB      DB$M                    ;Dump memory
018D 08012801                   DW      108H,128H               ;Check start/end at nonmultiples
0191 4D656D6F72                 DB      'Memory Dump #1',0      ;  of 10H
                        ;
01A0 CD5205                     CALL    DB$Display
01A3 18                         DB      DB$M                    ;Dump memory
01A4 00011F01                   DW      100H,11FH               ;Check start and end on displayed
01A8 4D656D6F72                 DB      'Memory Dump #2',0      ;  line boundaries
                        ;
```

Figure 10-2. Debug subroutines

```
 01B7 CD5205          CALL   DB$Display
 01BA 18              DB     DB$M                  ;Dump memory
 01BB 01010001        DW     101H,100H             ;Check error handling where
 01BF 4D656D6F72      DB     'Memory Dump #3',0    ;  start > end address
                  ;
 01CE CD5205          CALL   DB$Display
 01D1 18              DB     DB$M                  ;Dump memory
 01D2 00010001        DW     100H,100H             ;Check end-case of single byte
 01D6 4D656D6F72      DB     'Memory Dump #4',0    ;  output
                  ;#
                  ;       Test register pair display
                  ;
 01E5 CD5205          CALL   DB$Display            ;Call the debug routine
 01E8 10              DB     DB$BC
 01E9 4243205265      DB     'BC Register',0
                  ;
 01F5 CD5205          CALL   DB$Display            ;Call the debug routine
 01F8 12              DB     DB$DE
 01F9 4445205265      DB     'DE Register',0
                  ;
 0205 CD5205          CALL   DB$Display            ;Call the debug routine
 0208 14              DB     DB$HL
 0209 484C205265      DB     'HL Register',0
                  ;
 0215 CD5205          CALL   DB$Display            ;Call the debug routine
 0218 16              DB     DB$SP
 0219 5350205265      DB     'SP Register',0
                  ;
 0225 013203          LXI    B,Byte$BC             ;Set up registers for byte tests
 0228 113303          LXI    D,Byte$DE
 022B 213403          LXI    H,Byte$HL
                  ;#
                  ;       Test byte indirect display
                  ;
 022E CD5205          CALL   DB$Display            ;Call the debug routine
 0231 1A              DB     DB$B$BC
 0232 4279746520      DB     'Byte at (BC)',0
                  ;
 023F CD5205          CALL   DB$Display            ;Call the debug routine
 0242 1C              DB     DB$B$DE
 0243 4279746520      DB     'Byte at (DE)',0


                  ;
 0250 CD5205          CALL   DB$Display            ;Call the debug routine
 0253 1E              DB     DB$B$HL
 0254 4279746520      DB     'Byte at (HL)',0
                  ;
 0261 013503          LXI    B,Word$BC             ;Set up the registers for word tests
 0264 113703          LXI    D,Word$DE
 0267 213903          LXI    H,Word$HL

 026A CD5205          CALL   DB$Display            ;Call the debug routine
 026D 20              DB     DB$W$BC
 026E 576F726420      DB     'Word at (BC)',0
                  ;
 027B CD5205          CALL   DB$Display            ;Call the debug routine
 027E 22              DB     DB$W$DE
 027F 576F726420      DB     'Word at (DE)',0
                  ;
 028C CD5205          CALL   DB$Display            ;Call the debug routine
 028F 24              DB     DB$W$HL
 0290 576F726420      DB     'Word at (HL)',0
                  ;#
                  ;       Test DB$On/Off
                  ;
 029D CD1D05          CALL   DB$Off                ;Disable debug output
 02A0 CDD607          CALL   DB$MSGI               ;Display in-line message
 02A3 0D0A546869      DB     0DH,0AH,'This message should NOT appear',0

 02C4 CD1505          CALL   DB$On
 02C7 CDD607          CALL   DB$MSGI
 02CA 0D0A446562      DB     0DH,0AH,'Debug output has been re-enabled.',0
                  ;#
                  ;       Test pass count logic
                  ;
```

Figure 10-2. (Continued)

```
02EE CD1D05              CALL    DB$Off                  ;Disable debug output
02F1 CD2405              CALL    DB$Set$Pass             ;Set pass count
02F4 1E00                DW      30
                ;
02F6 3E22                MVI     A,34                    ;Set loop counter greater than pass
                                                         ;  counter
                Test$Pass$Loop:
02F8 CD3505              CALL    DB$Pass                 ;Decrement pass count
02FB CDD607              CALL    DB$MSGI                 ;Display in-line message
02FE 0D0A546869          DB      0DH,0AH,'This message should display 5 times',0
0324 3D                  DCR     A
0325 C2F802              JNZ     Test$Pass$Loop
                ;#
                ;        Test debug input/output
                ;
0328 CD1D05              CALL    DB$Off                  ;Check that debug IN/OUT
                                                         ;  must still occur when debug
                                                         ;  output is disabled.
032B E7                  RST     4                       ;Debug input
032C 11                  DB      11H                     ;Port number
032D EF                  RST     5                       ;Debug output (value return from input)
032E 22                  DB      22H                     ;Port number

032F C30000              JMP     0                       ;Warm boot at end of testbed
                ;
                ;
                ;        Dummy values for byte and word displays
0332 BC         Byte$BC:         DB      0BCH
0333 DE         Byte$DE:         DB      0DEH
0334 F1         Byte$HL:         DB      0F1H
                ;
0335 0C0B       Word$BC:         DW      0B0CH
0337 0E0D       Word$DE:         DW      0D0EH
0339 010F       Word$HL:         DW      0F01H
                ;
033B 9999999999                  DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
034B 9999999999                  DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
035B 9999999999                  DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
                Test$Stack:
                ;
                ;
                ;
0400                     ORG     400H                    ;To avoid unnecessary listings
                                                         ;  when only the testbed changes
                ;-----------------------------------------------------------------
                ;#
                ;
                ;        Debug subroutines
                ;
                ;
                ;        Equates for DB$Display codes
                ;        These equates are the offsets down the table of addresses
                ;        for various subroutines to be used.
                ;
0000 =          DB$F     EQU     00      ;Flags
0002 =          DB$A     EQU     02      ;A register
0004 =          DB$B     EQU     04      ;B
0006 =          DB$C     EQU     06      ;C
0008 =          DB$D     EQU     08      ;D
000A =          DB$E     EQU     10      ;E
000C =          DB$H     EQU     12      ;H
000E =          DB$L     EQU     14      ;L
0010 =          DB$BC    EQU     16      ;BC
0012 =          DB$DE    EQU     18      ;DE
0014 =          DB$HL    EQU     20      ;HL
0016 =          DB$SP    EQU     22      ;Stack pointer
0018 =          DB$M     EQU     24      ;Memory
001A =          DB$B$BC  EQU     26      ;(BC)
001C =          DB$B$DE  EQU     28      ;(DE)
001E =          DB$B$HL  EQU     30      ;(HL)
0020 =          DB$W$BC  EQU     32      ;(BC+1),(BC)
0022 =          DB$W$DE  EQU     34      ;(DE+1),(DE)
0024 =          DB$W$HL  EQU     36      ;(HL+1),(HL)
                ;
                ;
                ;        Equates
0020 =          RST4             EQU     20H     ;Address for RST 4 - IN instruction
```

Figure 10-2. (Continued)

```
0028 =          RST5            EQU     28H     ;Address for RST 5 - OUT instruction
                ;
0001 =          B$CONIN         EQU     1       ;BDOS CONIN function code
0002 =          B$CONOUT        EQU     2       ;BDOS CONOUT function code
000A =          B$READCONS      EQU     10      ;BDOS read console function code
0005 =          BDOS            EQU     5       ;BDOS entry point
                ;
0000 =          False           EQU     0
FFFF =          True            EQU     NOT False
                ;
                                                        ;Equates to specify how DB$CONOUT
                                                        ;  and DB$CONIN should perform
                                                        ;  their input/output
0000 =          DB$Polled$IO    EQU     False           ;)
0000 =          DB$BIOS$IO      EQU     False           ;) Only one must be true
FFFF =          DB$BDOS$IO      EQU     True            ;)
                ;
                                                        ;Equates for polled I/O
0001 =          DB$Status$Port  EQU     01H             ;Console status port
0002 =          DB$Data$Port    EQU     02H             ;Console data port
                ;
0002 =          DB$Input$Ready  EQU     0000$0010B      ;Incoming data ready
0001 =          DB$Output$Ready EQU     0000$0001B      ;Ready for output
                ;
                                                        ;Data for BIOS I/O
0400 C3         BIOS$CONIN:     DB      JMP             ;The initialization routine sets these
0401 0000                       DW      0               ;  two JMP addresses into the BIOS
0403 C3         BIOS$CONOUT:    DB      JMP
0404 0000                       DW      0
                ;
                ;       Main debug variables and constants
                ;
0406 00         DB$Flag:        DB      0       ;Main debug control flag
                                                ;  When this flag is nonzero, all debug
                                                ;  output will be made. When zero, all
                                                ;  debug output will be suppressed.
                                                ;  It is altered either directly by the user
                                                ;  or using the routines DB$On, DB$Off and
                                                ;  DB$Pass.
                ;
0407 0000       DB$Pass$Count:  DW      0       ;Pass counter
                                                ;  When this is nonzero, calls to DB$Pass
                                                ;  decrement it by one. When it reaches
                                                ;  zero, the debug control flag, DB$Flag,
                                                ;  is set nonzero, thereby enabling
                                                ;  debug output.
                ;
                DB$Save$HL:                     ;Save area for HL
0409 00         DB$Save$L:      DB      0
040A 00         DB$Save$H:      DB      0

040B 0000       DB$Save$SP:     DW      0       ;Save area for stack pointer
040D 0000       DB$Save$RA:     DW      0       ;Save area for return address
040F 0000       DB$Call$Address:        DW      0       ;Starts out the same as DB$Save$RA
                                                ;  but DB$Save$RA gets updated during
                                                ;  debug processing. This value is
                                                ;  output ahead of the caption
                DB$Start$Address:               ;Start address for memory display
0411 0000                       DW      0
                DB$End$Address:                 ;End address for memory display
0413 0000                       DW      0
                DB$Display$Code:                ;Display code requested
0415 00                         DB      0
                ;
                ;
                                                ;Stack area
0416 9999999999                 DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
0426 9999999999                 DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
0436 9999999999                 DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
0446 00         DB$Save$E:      DB      0       ;E register
0447 00         DB$Save$D:      DB      0       ;D register
0448 00         DB$Save$C:      DB      0       ;C register
0449 00         DB$Save$B:      DB      0       ;B register
044A 00         DB$Save$F:      DB      0       ;Flags
044B 00         DB$Save$A:      DB      0       ;A register
                DB$Stack:                       ;Debug stack area
                                                ;  The registers in the stack area are PUSHed
                                                ;  onto the stack and accessed directly.
```

Figure 10-2. (Continued)

```
                        ;
                        ;       Register caption messages
                        ;
                        ;       The table below, indexed by the Display$Code is used to access
                        ;       the register caption string.
                        ;
                        DB$Register$Captions:
044C 7204                       DW      DB$F$RC         ;Flags
044E 7804                       DW      DB$A$RC         ;A register
0450 7A04                       DW      DB$B$RC         ;B
0452 7C04                       DW      DB$C$RC         ;C
0454 7E04                       DW      DB$D$RC         ;D
0456 8004                       DW      DB$E$RC         ;E
0458 8204                       DW      DB$H$RC         ;H
045A 8404                       DW      DB$L$RC         ;L
045C 8604                       DW      DB$BC$RC        ;BC
045E 8904                       DW      DB$DE$RC        ;DE
0460 8C04                       DW      DB$HL$RC        ;HL
0462 8F04                       DW      DB$SP$RC        ;Stack pointer
0464 9204                       DW      DB$M$RC         ;Memory
0466 A604                       DW      DB$B$BC$RC      ;(BC)
0468 AB04                       DW      DB$B$DE$RC      ;(DE)
046A B004                       DW      DB$B$HL$RC      ;(HL)
046C B504                       DW      DB$W$BC$RC      ;(BC+1),(BC)
046E C104                       DW      DB$W$DE$RC      ;(DE+1),(DE)
0470 CD04                       DW      DB$W$HL$RC      ;(HL+1),(HL)
                        ;
0472 466C616773DB$F$RC:         DB      'Flags',0       ;Flags
0478 4100       DB$A$RC:        DB      'A',0           ;A register
047A 4200       DB$B$RC:        DB      'B',0           ;B
047C 4300       DB$C$RC:        DB      'C',0           ;C
047E 4400       DB$D$RC:        DB      'D',0           ;D
0480 4500       DB$E$RC:        DB      'E',0           ;E
0482 4800       DB$H$RC:        DB      'H',0           ;H
0484 4C00       DB$L$RC:        DB      'L',0           ;L
0486 424300     DB$BC$RC:       DB      'BC',0          ;BC
0489 444500     DB$DE$RC:       DB      'DE',0          ;DE
048C 484C00     DB$HL$RC:       DB      'HL',0          ;HL
048F 535000     DB$SP$RC:       DB      'SP',0          ;Stack pointer
0492 5374617274DB$M$RC:         DB      'Start, End Address ',0 ;Memory
04A6 2842432900DB$B$BC$RC:      DB      '(BC)',0        ;(BC)
04AB 2844452900DB$B$DE$RC:      DB      '(DE)',0        ;(DE)
04B0 28484C2900DB$B$HL$RC:      DB      '(HL)',0        ;(HL)
04B5 2842432B31DB$W$BC$RC:      DB      '(BC+1),(BC)',0 ;(BC+1),(BC)
04C1 2844452B31DB$W$DE$RC:      DB      '(DE+1),(DE)',0 ;(DE+1),(DE)
04CD 28484C2B31DB$W$HL$RC:      DB      '(HL+1),(HL)',0 ;(HL+1),(HL)
                        ;
                        ;       Flags message
                        ;
04D9 43785A784DDB$Flags$Msg:    DB      'CxZxMxExIx',0  ;Compatible with DDT's display
                        ;
                        ;       Flags masks used to test user's flag byte
                        ;
                        DB$Flag$Masks:
04E4 01                         DB      0000$0001B      ;Carry
04E5 40                         DB      0100$0000B      ;Zero
04E6 80                         DB      1000$0000B      ;Minus
04E7 04                         DB      0000$0100B      ;Even parity
04E8 10                         DB      0001$0000B      ;Interdigit carry (aux carry)
04E9 00                         DB      0               ;Terminator
                        ;#
                        ;       DB$Init
                        ;       This routine initializes the debug package.
                        ;
                        DB$Init:
                                IF      DB$BIOS$IO              ;Use BIOS for CONIN/CONOUT
                                LHLD    1                       ;Get warm boot address from base
                                                                ;  page. H = BIOS jump vector page
                                MVI     L,09H                   ;Get CONIN offset in jump vector
                                SHLD    BIOS$CONIN + 1          ;Set up address
                                MVI     L,0CH                   ;Get CONOUT offset in jump vector
                                SHLD    BIOS$CONOUT + 1
                                ENDIF

                                        ;Set up JMP instructions to receive control
                                        ;  when an RST instruction is executed
04EA 3EC3                       MVI     A,JMP           ;Set JMP instructions at RST points
```

Figure 10-2. (Continued)

```
04EC 322000             STA     RST4
04EF 322800             STA     RST5
04F2 211A08             LXI     H,DB$Input      ;Address of fake input routine
04F5 222100             SHLD    RST4 + 1
04F8 216C08             LXI     H,DB$Output     ;Address of fake output routine
04FB 222900             SHLD    RST5 + 1

04FE C9                 RET
                ;#
                ;       DB$CONINU
                ;       This routine returns the next character from the console,
                ;       but converting "a" to "z" to uppercase letters.
                ;
                DB$CONINU:
04FF CD0505             CALL    DB$CONIN                ;Get character from keyboard
0502 C31B09             JMP     DB$A$To$Upper           ;Fold to upper and return
                ;#
                ;       DB$CONIN
                ;       This routine returns the next character from the console.
                ;       According to the setting of equates, it uses simple
                ;       polled I/O, the BDOS (function 2) or the BIOS.
                ;
                ;       Exit parameters
                ;
                ;               A = character from console
                ;
                DB$CONIN:
                        IF      DB$Polled$IO            ;Simple polled input
                        IN      DB$Status$Port          ;Check if incoming data
                        ANI     DB$Input$Ready
                        JZ      DB$CONIN                ;No
                        IN      DB$Data$Port            ;Input data character
                        PUSH    PSW                     ;Save data character
                        MOV     C,A                     ;Ready for output
                        CALL    DB$CONOUT               ;Echo it back
                        POP     PSW                     ;Recover data character
                        RET
                        ENDIF

                        IF      DB$BDOS$IO              ;Use BDOS for input
0505 0E01               MVI     C,B$CONIN               ;Read console
0507 C30500             JMP     BDOS                    ;BDOS returns to our caller
                        ENDIF

                        IF      DB$BIOS$IO              ;Use BIOS for input
                        JMP     BIOS$CONIN              ;This was set up during BIOS
                                                        ;  initialization
                        ENDIF

                ;#
                ;       DB$CONOUT
                ;       This routine outputs the character in the C register to the
                ;       console, using simple polled I/O, the BDOS or the BIOS.
                ;
                ;       Entry parameters
                ;               A = byte to be output
                ;
                DB$CONOUT:
050A 3A0604             LDA     DB$Flag                 ;Check if debug output enabled
050D B7                 ORA     A
050E C8                 RZ                              ;Ignore output if disabled

                        IF      DB$Polled$IO            ;Use simple polled output
                        IN      DB$Status$Port          ;Check if ready for output
                        ANI     DB$Output$Ready
                        JZ      DB$CONOUT               ;No
                        MOV     A,C                     ;Get data byte
                        OUT     DB$Data$Port
                        RET
                        ENDIF

                        IF      DB$BDOS$IO              ;Use BDOS for output
050F 59                 MOV     E,C                     ;Move into correct register
0510 0E02               MVI     C,B$CONOUT
0512 C30500             JMP     BDOS                    ;BDOS returns to our caller
                        ENDIF

                        IF      DB$BIOS$IO              ;Use BIOS for output
```

Figure 10-2. (Continued)

```
                                MOV     A,C                     ;Move into correct register
                                JMP     BIOS$CONOUT             ;Set up during debug initialization
                                ENDIF
                        ;#
                        ;
                        ;       DB$On
                        ;       This routine enables all debug output by setting the
                        ;       DB$Flag nonzero.
                        ;
                        DB$On:
        0515 F5                 PUSH    PSW                     ;Preserve registers
        0516 3EFF               MVI     A,0FFH
        0518 320604             STA     DB$Flag                 ;Set control flag on
        051B F1                 POP     PSW
        051C C9                 RET
                        ;#
                        ;
                        ;       DB$Off
                        ;       This routine disables all debug output by setting the
                        ;       DB$Flag to zero.
                        ;
                        DB$Off:
        051D F5                 PUSH    PSW                     ;Preserve registers
        051E AF                 XRA     A
        051F 320604             STA     DB$Flag                 ;Clear control flag
        0522 F1                 POP     PSW
        0523 C9                 RET
                        ;#
                        ;
                        ;       DB$Set$Pass
                        ;       This routine sets the pass counter. Subsequent calls to DB$Pass
                        ;       decrement the count, and when it reaches 0, debug output
                        ;       is enabled.
                        ;
                        ;       Calling sequence
                        ;
                        ;               CALL    DB$Set$Pass
                        ;               DW      Pass$Count$Value
                        ;
                        DB$Set$Pass:
        0524 220904             SHLD    DB$Save$HL              ;Preserve user's HL
        0527 E1                 POP     H                       ;Recover return address
        0528 D5                 PUSH    D                       ;Preserve user's DE
        0529 5E                 MOV     E,M                     ;Get LS byte of count
        052A 23                 INX     H                       ;Update pointer
        052B 56                 MOV     D,M                     ;Get MS byte
        052C 23                 INX     H                       ;HL points to return address
        052D EB                 XCHG                            ;HL = pass counter
        052E 220704             SHLD    DB$Pass$Count           ;Set debug pass counter
        0531 EB                 XCHG                            ;HL points to return address
        0532 D1                 POP     D                       ;Recover user's DE
        0533 E3                 XTHL                            ;Recover user's HL and set
                                                                ;  return address on top of stack
        0534 C9                 RET
                        ;#
                        ;
                        ;       DB$Pass
                        ;       This routine decrements the debug pass counter -
                        ;       if the result is negative, it takes no further action.
                        ;       If the result is zero, it sets the debug control flag nonzero
                        ;       to enable debug output.
                        ;
                        DB$Pass:
        0535 F5                 PUSH    PSW                     ;Save user's registers
        0536 E5                 PUSH    H
        0537 2A0704             LHLD    DB$Pass$Count           ;Get pass count
        053A 2B                 DCX     H
        053B 7C                 MOV     A,H                     ;Check if count now negative
        053C B7                 ORA     A
        053D FA4705             JM      DB$Pass$x               ;Yes, take no further action
        0540 220704             SHLD    DB$Pass$Count           ;Save downdated count
        0543 B5                 ORA     L                       ;Check if count now zero
        0544 CA4A05             JZ      DB$Pass$ED              ;Yes, enable debug
                        DB$Pass$x:                              ;
        0547 E1                 POP     H                       ;Recover user's registers
        0548 F1                 POP     PSW
        0549 C9                 RET
```

Figure 10-2. (Continued)

```
                        ;
                        DB$Pass$Ed:                                     ;Enable debug
054A 3EFF                       MVI     A,0FFH
054C 320604                     STA     DB$Flag                         ;Set debug control flag
054F C34705                     JMP     DB$Pass$x
                        ;#
                        ;
                        ;       DB$Display
                        ;       This is the primary debug display routine.
                        ;
                        ;       Calling sequence
                        ;
                        ;               CALL    DB$Display
                        ;               DB      Display$Code
                        ;               DB      'Caption String',0
                        ;
                        ;               Display code identifies which register(s) are to be
                        ;               displayed.
                        ;
                        ;               When the display code specifies a block of memory
                        ;               the sequence is:
                        ;
                        ;               CALL    DB$Display
                        ;               DB      Display$Code
                        ;               DW      Start$Address,End$Address
                        ;               DB      'Caption String',0
                        ;
                        DB$Display:
                        ;
                        DB$Display$Enabled:
0552 220904                     SHLD    DB$Save$HL                      ;Save user's HL

0555 E3                         XTHL                                    ;Get return address from stack
0556 220D04                     SHLD    DB$Save$RA                      ;This gets updated by debug code
0559 E5                         PUSH    H                               ;Save return address temporarily
055A 2B                         DCX     H                               ;Subtract 3 to address call instruction
055B 2B                         DCX     H                               ;  itself
055C 2B                         DCX     H
055D 220F04                     SHLD    DB$Call$Address                 ;Save actual address of CALL
0560 E1                         POP     H                               ;Recover return address

0561 F5                         PUSH    PSW                             ;Temporarily save flags to avoid
                                                                        ;  them being changed by DAD SP
0562 210000                     LXI     H,0                             ;Preserve stack pointer
0565 39                         DAD     SP
0566 23                         INX     H                               ;Correct for extra PUSH PSW needed
0567 23                         INX     H                               ;  to save the flags
0568 220B04                     SHLD    DB$Save$SP
056B F1                         POP     PSW                             ;Recover flags

056C 314C04                     LXI     SP,DB$Stack                     ;Switch to local stack

056F F5                         PUSH    PSW                             ;Save other user's registers
0570 C5                         PUSH    B                               ;The stack area is specially laid
0571 D5                         PUSH    D                               ;  out to access these registers

0572 2A0D04                     LHLD    DB$Save$RA                      ;Get return address
0575 7E                         MOV     A,M                             ;Get display code
0576 321504                     STA     DB$Display$Code
0579 23                         INX     H                               ;Update return address

057A FE18                       CPI     DB$M                            ;Check if memory to be displayed
057C C29105                     JNZ     DB$Not$Memory
057F 5E                         MOV     E,M                             ;Get DE = start address
0580 23                         INX     H
0581 56                         MOV     D,M
0582 23                         INX     H
0583 EB                         XCHG                                    ;HL = start address
0584 221104                     SHLD    DB$Start$Address
0587 EB                         XCHG                                    ;HL -> end address
0588 5E                         MOV     E,M                             ;Get DE = end address
0589 23                         INX     H
058A 56                         MOV     D,M
058B 23                         INX     H
058C EB                         XCHG                                    ;HL = end address, DE -> caption
058D 221304                     SHLD    DB$End$Address
0590 EB                         XCHG                                    ;HL -> caption string
```

Figure 10-2. (Continued)

```
                   DB$Not$Memory:
                   ;
                   ;       Output preamble and caption string
                   ;       The format for everything except memory display is:
                   ;
                   ;       nnnn : Caption String : RC = vvvv
                   ;         ^                      ^    ^
                   ;       Call Address             :    Value
                   ;                      Register Caption (A, B, C...)
                   ;
                   ;       A carriage return, line feed is output at the start of the
                   ;       message - but NOT at the end.
                   ;
                   ;       Memory displays look like :
                   ;
                   ; nnnn : Caption String : Start, End  ssss, eeee
                   ; ssss : hh hh hh hh hh hh hh hh hh hh hh hh hh hh hh hh : cccc cccc cccc cccc
                   ;

0591 E5                    PUSH    H                       ;Save pointer to caption string
0592 CDC107                CALL    DB$CRLF                 ;Display carriage return, line feed
0595 CD7C07                CALL    DB$Display$CALLA        ;Display DB$Call$Address in hex.

0598 E1                    POP     H                       ;Recover pointer to caption string
                   DB$Display$Caption:                     ;HL -> caption string
0599 7E                    MOV     A,M                     ;Get character
059A 23                    INX     H
059B B7                    ORA     A                       ;Check if end of string
059C CAA805                JZ      DB$End$Caption          ;Yes

059F E5                    PUSH    H                       ;Save string pointer
05A0 4F                    MOV     C,A                     ;Ready for output
05A1 CD0A05                CALL    DB$CONOUT               ;Display character
05A4 E1                    POP     H                       ;Recover string pointer
05A5 C39905                JMP     DB$Display$Caption      ;Go back for next character
                   ;
                   DB$End$Caption:
05A8 220D04                SHLD    DB$Save$RA              ;Save updated return address

05AB CDC807                CALL    DB$Colon                ;Display ' : '

                                                           ;Display register caption
05AE 3A1504                LDA     DB$Display$Code         ;Get user's display code
05B1 5F                    MOV     E,A                     ;Make display code into word
05B2 1600                  MVI     D,0
05B4 D5                    PUSH    D                       ;Save word value for later

05B5 FE18                  CPI     DB$M                    ;Memory display is a special case
05B7 CACF05                JZ      DB$Display$Mem$Caption  ;Yes

05BA 214C04                LXI     H,DB$Register$Captions  ;Make pointer to address in table
05BD 19                    DAD     D                       ;HL -> word containing address of
                                                           ;  register caption
05BE 5E                    MOV     E,M                     ;Get LS byte of address
05BF 23                    INX     H
05C0 56                    MOV     D,M                     ;DE -> register caption string
05C1 EB                    XCHG                            ;HL -> register caption string.
05C2 CDEE07                CALL    DB$MSG                  ;Display message addressed by HL

05C5 CDD607                CALL    DB$MSGI                 ;Display in-line message
05C8 203D2000              DB      ' = ',0
05CC C3ED05                JMP     DB$Select$Routine       ;Go to correct processor
                   ;
                   DB$Display$Mem$Caption:                 ;The memory display requires a special
                                                           ;  caption with the start and end
                                                           ;  addresses
05CF 219204                LXI     H,DB$M$RC               ;Display specific caption
05D2 CDEE07                CALL    DB$MSG
05D5 CDC807                CALL    DB$Colon                ;Display ' : '

05D8 2A1104                LHLD    DB$Start$Address        ;Display start address
05DB CD8707                CALL    DB$DHLH                 ;Display HL in hex.

05DE CDD607                CALL    DB$MSGI                 ;Display in-line message
05E1 2C2000                DB      ', ',0

05E4 2A1304                LHLD    DB$End$Address          ;Get end address
```

Figure 10-2. (Continued)

```
05E7 CD8707          CALL    DB$DHLH                   ;Display HL in hex.
05EA CDC107          CALL    DB$CRLF                   ;Display carriage return, line feed
                                                       ;Drop into select routine
                DB$Select$Routine:
05ED D1              POP     D                         ;Recover word value Display$Code
05EE 210A06          LXI     H,DB$Display$Table
05F1 19              DAD     D                         ;HL -> address of code to process
                                                       ;  display requirements
05F2 5E              MOV     E,M                       ;Get LS byte of address
05F3 23              INX     H                         ;Update pointer
05F4 56              MOV     D,M                       ;Get MS byte of address
05F5 EB              XCHG                              ;HL -> code

05F6 11FB05          LXI     D,DB$Exit                 ;Fake link on stack
05F9 D5              PUSH    D
05FA E9              PCHL                              ;"CALL" display processor
                ;
                DB$Exit:                               ;Return to the user
05FB D1              POP     D                         ;Recover user's registers saved
05FC C1              POP     B                         ;  on local debug stack
05FD F1              POP     PSW
05FE 2A0B04          LHLD    DB$Save$SP                ;Revert to user's stack
0601 F9              SPHL
0602 2A0D04          LHLD    DB$Save$RA                ;Get updated return address (bypasses
                                                       ;  in-line parameters)
0605 E3              XTHL                              ;Replace on top of user's stack
0606 2A0904          LHLD    DB$Save$HL                ;Get user's HL
0609 C9              RET                               ;Transfer to correct return address



                DB$Display$Table:
060A 3006            DW      DP$F           ;Flags
060C 5406            DW      DP$A           ;A register
060E 5A06            DW      DP$B           ;B
0610 6006            DW      DP$C           ;C
0612 6606            DW      DP$D           ;D
0614 6C06            DW      DP$E           ;E
0616 7206            DW      DP$H           ;H
0618 7806            DW      DP$L           ;L
061A 7E06            DW      DP$BC          ;BC
061C 8406            DW      DP$DE          ;DE
061E 8A06            DW      DP$HL          ;HL
0620 9006            DW      DP$SP          ;Stack pointer
0622 9606            DW      DP$M           ;Memory
0624 4907            DW      DP$B$BC        ;(BC)
0626 5007            DW      DP$B$DE        ;(DE)
0628 5707            DW      DP$B$HL        ;(HL)
062A 5E07            DW      DP$W$BC        ;(BC+1),(BC)
062C 6807            DW      DP$W$DE        ;(DE+1),(DE)
062E 7207            DW      DP$W$HL        ;(HL+1),(HL)
                ;#
                ;    Debug display processing routines
                ;
                DP$F:                               ;Flags
                                                    ;The flags are displayed in the same way that
                                                    ;  DDT uses: C1Z0M0E0I0
0630 3A4A04          LDA     DB$Save$F      ;Get flags
0633 47              MOV     B,A            ;Preserve copy
0634 21DA04          LXI     H,DB$Flags$Msg + 1     ;HL -> first 0/1 in message
0637 11E404          LXI     D,DB$Flag$Masks ;DE -> table of flag mask values
                DB$F$Next:
063A 1A              LDAX    D              ;Get next flag mask
063B B7              ORA     A              ;Check if end of table
063C CA4E06          JZ      DB$F$Display   ;Yes, display the results

063F A0              ANA     B              ;Check if this flag is set
0640 3E31            MVI     A,'1'          ;Assume yes
0642 C24706          JNZ     DB$F$NZ        ;Yes,it is set
0645 3E30            MVI     A,'0'          ;No,it is clear
                DB$F$NZ:
0647 77              MOV     M,A            ;Store '0' or '1' in message text
0648 23              INX     H              ;Update pointer to next 0/1
0649 23              INX     H
064A 13              INX     D              ;Update flag mask pointer
064B C33A06          JMP     DB$F$Next
                DB$F$Display:                       ;Display results
064E 21D904          LXI     H,DB$Flags$Msg
```

Figure 10-2. (Continued)

```
0651 C3EE07            JMP     DB$MSG          ;Display message and return
                  ;
                  DP$A:                 ;A register
0654 3A4B04            LDA     DB$Save$A       ;Get saved value
0657 C39107            JMP     DB$DAH          ;Display it and return
                  ;
                  DP$B:                 ;B
065A 3A4904            LDA     DB$Save$B       ;Get saved value
065D C39107            JMP     DB$DAH          ;Display it and return
                  ;
                  DP$C:                 ;C
0660 3A4804            LDA     DB$Save$C       ;Get saved value
0663 C39107            JMP     DB$DAH          ;Display it and return
                  ;
                  DP$D:                 ;D
0666 3A4704            LDA     DB$Save$D       ;Get saved value
0669 C39107            JMP     DB$DAH          ;Display it and return
                  ;
                  DP$E:                 ;E
066C 3A4604            LDA     DB$Save$E       ;Get saved value
066F C39107            JMP     DB$DAH          ;Display it and return
                  ;
                  DP$H:                 ;H
0672 3A0A04            LDA     DB$Save$H       ;Get saved value
0675 C39107            JMP     DB$DAH          ;Display it and return
                  ;
                  DP$L:                 ;L
0678 3A0904            LDA     DB$Save$L       ;Get saved value
067B C39107            JMP     DB$DAH          ;Display it and return
                  ;
                  DP$BC:                ;BC
067E 2A4804            LHLD    DB$Save$C       ;Get saved word value
0681 C38707            JMP     DB$DHLH         ;Display it and return
                  ;
                  DP$DE:                ;DE
0684 2A4604            LHLD    DB$Save$E       ;Get saved word value
0687 C38707            JMP     DB$DHLH         ;Display it and return
                  ;
                  DP$HL:                ;HL
068A 2A0904            LHLD    DB$Save$HL      ;Get saved word value
068D C38707            JMP     DB$DHLH         ;Display it and return
                  ;
                  DP$SP:                ;Stack Pointer
0690 2A0B04            LHLD    DB$Save$SP      ;Get saved word value
0693 C38707            JMP     DB$DHLH         ;Display it and return
                  ;
                  DP$M:                 ;Memory
0696 2A1304            LHLD    DB$End$Address  ;Increment end address to make
0699 23                INX     H               ;  arithmetic easier
069A 221304            SHLD    DB$End$Address

069D 2A1104            LHLD    DB$Start$Address
06A0 CD3A07            CALL    DB$M$Check$End  ;Compare HL to End$Address
06A3 DAD106            JC      DB$M$Address$OK ;End > start
06A6 CDD607            CALL    DB$MSGI         ;Error start > end
06A9 0D0A2A2A20        DB      0DH,0AH,'** ERROR - Start Address > End **',0
06CD C9                RET
                  ;
                  DB$M$Next$Line:
06CE CDC107            CALL    DB$CRLF                 ;Output carriage return, line feed
                  DB$M$Address$OK:                     ;Bypass CR,LF for first line
06D1 CDD607            CALL    DB$MSGI                 ;Indent line
06D4 202000            DB      '  ',0
06D7 2A1104            LHLD    DB$Start$Address        ;Get start of line address
06DA CD8707            CALL    DB$DHLH                 ;Display in hex

06DD CDC807            CALL    DB$Colon                ;Display ' : '

06E0 2A1104            LHLD    DB$Start$Address
                  DB$M$Next$Hex$Byte:
06E3 E5                PUSH    H                       ;Save memory address
06E4 CDD007            CALL    DB$Blank                ;Output a blank
06E7 E1                POP     H                       ;Recover current byte address
06E8 7E                MOV     A,M                     ;Get byte from memory
06E9 23                INX     H                       ;Update memory pointer
06EA E5                PUSH    H                       ;Save for later
06EB CD9107            CALL    DB$DAH                  ;Display in hex.
06EE E1                POP     H                       ;Recover memory updated address
```

Figure 10-2. (Continued)

```
06EF CD3A07            CALL    DB$M$Check$End          ;Compare HL vs.end address
06F2 CAFE06            JZ      DB$M$Display$ASCII      ;Yes, end of area
06F5 7D                MOV     A,L                     ;Check if at start of new line,
06F6 E60F              ANI     0000$1111B              ;  (is address XXX0H?)
06F8 CAFE06            JZ      DB$M$Display$ASCII      ;Yes
06FB C3E306            JMP     DB$M$Next$Hex$Byte      ;No, loop back for another

                   ;
                   DB$M$Display$ASCII:                 ;Display bytes in ASCII
06FE CDC807            CALL    DB$Colon                ;Display ' : '
0701 2A1104            LHLD    DB$Start$Address        ;Start ASCII as beginning of line
                   DB$M$Next$ASCII$Byte:
0704 7E                MOV     A,M                     ;Get byte from memory
0705 E5                PUSH    H                       ;Save memory address
0706 E67F              ANI     0111$1111B              ;Remove parity
0708 4F                MOV     C,A                     ;Prepare for output
0709 FE20              CPI     ' '                     ;Check if non-graphic
070B D21007            JNC     DB$M$Display$Char       ;Char >= space
070E 0E2E              MVI     C,'.'                   ;Display non-graphic as '.'
                   DB$M$Display$Char:
0710 FE7F              CPI     7FH                     ;Check if DEL (may be non-graphic)
0712 C21707            JNZ     DB$M$Not$DEL            ;No, it is graphic
0715 0E2E              MVI     C,'.'                   ;Force to '.'
                   ;
                   DB$M$Not$DEL:
0717 CD0A05            CALL    DB$CONOUT               ;Display character
071A E1                POP     H                       ;Recover memory address
071B 23                INX     H                       ;Update memory pointer
071C 221104            SHLD    DB$Start$Address        ;Update memory copy
071F CD3A07            CALL    DB$M$Check$End          ;Check if end of memory dump
0722 CA3707            JZ      DB$M$Exit               ;Yes, done
0725 7D                MOV     A,L                     ;Check if end of line
0726 E60F              ANI     0000$1111B              ;  by checking address = XXX0H
0728 CACE06            JZ      DB$M$Next$Line          ;Yes, start next line
072B 7D                MOV     A,L                     ;Check if extra blank needed
072C E603              ANI     0000$0011B              ;  if address is multiple of 4
072E C20407            JNZ     DB$M$Next$ASCII$Byte    ;No -- go back for next character
0731 CDD007            CALL    DB$Blank                ;Yes, output blank
0734 C30407            JMP     DB$M$Next$ASCII$Byte    ;Go back for next character

                   ;
                   DB$M$Exit:
0737 C3C107            JMP     DB$CRLF                 ;Output carriage return, line feed
                                                       ;  and return
                   ;
                   DB$M$Check$End:                     ;Compares HL vs End$Address
073A D5                PUSH    D                       ;Save DE (defensive programming)
073B EB                XCHG                            ;DE = current address
073C 2A1304            LHLD    DB$End$Address          ;Get end address
073F 7A                MOV     A,D                     ;Compare MS bytes
0740 BC                CMP     H
0741 C24607            JNZ     DB$M$Check$End$X        ;Exit now as they are unequal
0744 7B                MOV     A,E                     ;Compare LS bytes
0745 BD                CMP     L
                   DB$M$Check$End$X:
0746 EB                XCHG                            ;HL = current address
0747 D1                POP     D                       ;Recover DE
0748 C9                RET                             ;Return with condition flags set
                   ;
                   DP$B$BC:                ;(BC)
0749 2A4804            LHLD    DB$Save$C       ;Get saved word value
074C 7E                MOV     A,M             ;Get byte addressed by it
074D C39107            JMP     DB$DAH          ;Display it and return
                   ;
                   DP$B$DE:                ;(DE)
0750 2A4604            LHLD    DB$Save$E       ;Get saved word value
0753 7E                MOV     A,M             ;Get byte addressed by it
0754 C39107            JMP     DB$DAH          ;Display it and return
                   ;
                   DP$B$HL:                ;(HL)
0757 2A0904            LHLD    DB$Save$HL      ;Get saved word value
075A 7E                MOV     A,M             ;Get byte addressed by it
075B C39107            JMP     DB$DAH          ;Display it and return
                   ;
                   DP$W$BC:                ;(BC+1),(BC)
075E 2A4804            LHLD    DB$Save$C       ;Get saved word value
0761 5E                MOV     E,M             ;Get word addressed by it
0762 23                INX     H
```

Figure 10-2. (Continued)

```
0763 56                 MOV     D,M
0764 EB                 XCHG                    ;HL = word to be displayed
0765 C38707             JMP     DB$DHLH         ;Display it and return
                ;
                DP$W$DE:                ;(DE+1),(DE)
0768 2A4604             LHLD    DB$Save$E       ;Get saved word value
076B 5E                 MOV     E,M             ;Get word addressed by it
076C 23                 INX     H
076D 56                 MOV     D,M
076E EB                 XCHG                    ;HL = word to be displayed
076F C38707             JMP     DB$DHLH         ;Display it and return
                ;
                DP$W$HL:                ;(HL+1),(HL)
0772 2A0904             LHLD    DB$Save$HL      ;Get saved word value
0775 5E                 MOV     E,M             ;Get word addressed by it
0776 23                 INX     H
0777 56                 MOV     D,M
0778 EB                 XCHG                    ;HL = word to be displayed
0779 C38707             JMP     DB$DHLH         ;Display it and return
                ;
                ;#
                ;       DB$Display$CALLA
                ;       This routine displays the DB$Call$Address in hexadecimal,
                ;       followed by " : ".
                ;
                DB$Display$CALLA:
077C E5                 PUSH    H               ;Save caller's HL
077D 2A0F04             LHLD    DB$Call$Address ;Get the call address
0780 CD8707             CALL    DB$DHLH         ;Display HL in hex.
0783 E1                 POP     H               ;Recover caller's HL
0784 C3C807             JMP     DB$Colon        ;Display " : " and return
                ;
                ;#
                ;
                ;       DB$DHLH
                ;       Display HL in hex.
                ;
                ;       Entry parameters
                ;
                ;               HL = value to be displayed
                ;
                DB$DHLH:
0787 E5                 PUSH    H               ;Save input value
0788 7C                 MOV     A,H             ;Get MS byte first
0789 CD9107             CALL    DB$DAH          ;Display A in hex.
078C E1                 POP     H               ;Recover input value
078D 7D                 MOV     A,L             ;Get LS byte
078E C39107             JMP     DB$DAH          ;Display it and return
                ;
                ;#
                ;
                ;       DB$DAH
                ;       Display A register in hexadecimal
                ;
                ;       Entry parameters
                ;
                ;               A = value to be converted and output
                ;
                DB$DAH:
0791 F5                 PUSH    PSW             ;Take a copy of the value to be converted
0792 0F                 RRC                     ;Shift A right four places
0793 0F                 RRC
0794 0F                 RRC
0795 0F                 RRC
0796 CDB407             CALL    DB$Nibble$To$Hex        ;Convert LS 4 bits to ASCII
0799 CD0A05             CALL    DB$CONOUT               ;Display the character
079C F1                 POP     PSW                     ;Get original value again
079D CDB407             CALL    DB$Nibble$To$Hex        ;Convert LS 4 bits to ASCII
07A0 C30A05             JMP     DB$CONOUT               ;Display and return to caller
                ;#
                ;
                ;       DB$CAH
                ;       Convert A register to hexadecimal ASCII and store in
                ;       specified address.
                ;
                ;       Entry parameters
                ;
```

Figure 10-2. (Continued)

```
                        ;               A = value to be converted and output
                        ;               HL -> buffer area to receive two characters of output
                        ;
                        ;       Exit parameters
                        ;
                        ;               HL -> byte following last hex.byte output
                        ;
                        DB$CAH:
07A3 F5                         PUSH    PSW             ;Take a copy of the value to be converted
07A4 0F                         RRC                     ;Shift A right four places
07A5 0F                         RRC
07A6 0F                         RRC
07A7 0F                         RRC
07A8 CDB407                     CALL    DB$Nibble$To$Hex        ;Convert to ASCII hex.
07AB 77                         MOV     M,A                     ;Save in memory
07AC 23                         INX     H                       ;Update pointer
07AD F1                         POP     PSW                     ;Get original value again
07AE CDB407                     CALL    DB$Nibble$To$Hex        ;Convert to ASCII hex.
07B1 77                         MOV     M,A                     ;Save in memory
07B2 23                         INX     H                       ;Update pointer
07B3 C9                         RET
                        ;#
                        ;
                        ;       Minor subroutines
                        ;
                        ;
                        ;       DB$Nibble$To$Hex
                        ;       This is a minor subroutine that converts the least
                        ;       significant four bits of the A register into an ASCII
                        ;       hex. character in A and C
                        ;
                        ;       Entry parameters
                        ;
                        ;               A = nibble to be converted in LS 4 bits
                        ;
                        ;       Exit parameters
                        ;
                        ;               A,C  = ASCII hex. character
                        ;
                        DB$Nibble$To$Hex:
07B4 E60F                       ANI     0000$1111B      ;Isolate LS four bits
07B6 C630                       ADI     '0'             ;Convert to ASCII
07B8 FE3A                       CPI     '9' + 1         ;Compare to maximum
07BA DABF07                     JC      DB$NTH$Numeric  ;No need to convert to A -> F
07BD C607                       ADI     7               ;Convert to a letter
                        DB$NTH$Numeric:
07BF 4F                         MOV     C,A             ;For convenience of other routines
07C0 C9                         RET


                        ;
                        ;       DB$CRLF
                        ;       Simple routine to display carriage return, line feed.
                        ;
                        DB$CRLF:
07C1 CDD607                     CALL    DB$MSGI         ;Display in-line message
07C4 0D0A00                     DB      0DH,0AH,0
07C7 C9                         RET
                        ;
                        ;       DB$Colon
                        ;       Simple routine to display ' : '.
                        ;
                        DB$Colon:
07C8 CDD607                     CALL    DB$MSGI         ;Display in-line message
07CB 203A2000                   DB      ' : ',0
07CF C9                         RET
                        ;
                        ;       DB$Blank
                        ;       Simple routine to display ' '.
                        ;
                        DB$Blank:
07D0 CDD607                     CALL    DB$MSGI         ;Display in-line message
07D3 2000                       DB      ' ',0
07D5 C9                         RET
                        ;#
                        ;
                        ;       Message processing subroutines
```

Figure 10-2. (Continued)

```
                        ;
                        ;       DB$MSGI (message in-line)
                        ;       Output null-byte terminated message that follows the
                        ;       CALL to MSGOUTI
                        ;
                        ;       Calling sequence
                        ;
                        ;               CALL    DB$MSGI
                        ;               DB      'Message',0
                        ;               ... next instruction
                        ;
                        ;       Exit parameters
                        ;               HL -> instruction following message
                        ;
                        ;
                        DB$MSGI:
                                                        ;Get return address of stack, save
                                                        ;  user's HL on top of stack
        07D6 E3                 XTHL                    ;HL -> message

        07D7 F5                 PUSH    PSW             ;Save all user's registers
        07D8 C5                 PUSH    B
        07D9 D5                 PUSH    D
                        DB$MSGI$Next:
        07DA 7E                 MOV     A,M             ;Get next data byte
        07DB 23                 INX     H               ;Update message pointer
        07DC B7                 ORA     A               ;Check if null byte
        07DD C2E507             JNZ     DB$MSGIC        ;No, continue

        07E0 D1                 POP     D               ;Recover user's registers
        07E1 C1                 POP     B
        07E2 F1                 POP     PSW
        07E3 E3                 XTHL                    ;Recover user's HL from stack, replacing
                                                        ;  it with updated return address
        07E4 C9                 RET                     ;Return to address after 00-byte
                                                        ;  after in-line message
                        DB$MSGIC:
        07E5 E5                 PUSH    H               ;Save message pointer
        07E6 4F                 MOV     C,A             ;Ready for output
        07E7 CD0A05             CALL    DB$CONOUT
        07EA E1                 POP     H               ;Recover message pointer
        07EB C3DA07             JMP     DB$MSGI$Next    ;Go back for next char.

                        ;
                        ;       DB$MSG
                        ;       Output null-byte terminated message
                        ;
                        ;       Calling sequence
                        ;
                        ;       MESSAGE:        DB      'Message',0
                        ;               :
                        ;               LXI     H,MESSAGE
                        ;               CALL    DB$MSG
                        ;
                        ;       Exit parameters
                        ;               HL -> null byte terminator
                        ;
                        ;
                        DB$MSG:
        07EE F5                 PUSH    PSW             ;Save user's registers
        07EF C5                 PUSH    B
        07F0 D5                 PUSH    D
                        DB$MSG$Next:
        07F1 7E                 MOV     A,M             ;Get next byte for output
        07F2 B7                 ORA     A               ;Check if 00-byte terminator
        07F3 CA0008             JZ      DB$MSG$X        ;Exit
        07F6 23                 INX     H               ;Update message pointer
        07F7 E5                 PUSH    H               ;Save updated pointer
        07F8 4F                 MOV     C,A             ;Ready for output
        07F9 CD0A05             CALL    DB$CONOUT
        07FC E1                 POP     H               ;Recover message pointer
        07FD C3F107             JMP     DB$MSG$Next     ;Go back for next character
                        ;
                        DB$MSG$X:
        0800 D1                 POP     D               ;Recover user's registers
        0801 C1                 POP     B
        0802 F1                 POP     PSW
```

Figure 10-2. (Continued)

```
0803 C9                        RET
                    ;#
                    ;
                    ;          Debug input routine
                    ;
                    ;          This routine helps debug code in which input instructions
                    ;          would normally occur. The opcode of the IN instruction
                    ;          must be replaced by a value of 0E7H (RST 4).
                    ;
                    ;          This routine picks up the port number contained in the byte
                    ;          following the RST 4, converts it to hexadecimal, and
                    ;          displays the message:
                    ;
                    ;                  Input from port XX :
                    ;
                    ;          It then accepts two characters (in hex.) from the keyboard,
                    ;          converts these to binary in A, and then returns control
                    ;          to the byte following the port number
                    ;
                    ;          *******
                    ;          WARNING - This routine uses both DB$CONOUT and BDOS calls
                    ;          *******
                    ;
0804 496E707574DBIN$Message:   DB      'Input from Port '
0814 5858203A20DBIN$Port:      DB      'XX : ',0
                    ;
                    ;
                    DB$Input:
081A 220904                    SHLD    DB$Save$HL      ;Save user's HL
081D E1                        POP     H               ;Recover address of port number
081E 2B                        DCX     H               ;Backup to point to RST
081F 220F04                    SHLD    DB$Call$Address ;Save for later display
0822 23                        INX     H               ;Restore to point to port number
                                                       ;Note: A need not be preserved
0823 7E                        MOV     A,M             ;Get port number
0824 23                        INX     H               ;Update return address to bypass port number
0825 220D04                    SHLD    DB$Save$RA      ;Save return address
0828 C5                        PUSH    B               ;Save remaining registers
0829 D5                        PUSH    D
082A F5                        PUSH    PSW             ;Save port number for later


082B CDB108                    CALL    DB$Flag$Save$On ;Save current state of debug flag
                                                       ;  and enable debug output

082E CDC107                    CALL    DB$CRLF         ;Display carriage return, line feed
0831 CD7C07                    CALL    DB$Display$CALLA;Display call address
0834 F1                        POP     PSW             ;Recover port number
0835 211408                    LXI     H,DBIN$Port
0838 CDA307                    CALL    DB$CAH          ;Convert to hex.and store in message
083B 210408                    LXI     H,DBIN$Message  ;Output prompting message
083E CDEE07                    CALL    DB$MSG
0841 0E02                      MVI     C,2             ;Get 2 digit hex. value
0843 CDCF08                    CALL    DB$GHV          ;Returns value in HL
0846 7D                        MOV     A,L             ;Get just single byte

0847 CDBF08                    CALL    DB$Flag$Restore ;Restore debug output to previous state

084A D1                        POP     D               ;Recover registers
084B C1                        POP     B
084C 2A0904                    LHLD    DB$Save$HL      ;Get previous HL
084F E5                        PUSH    H               ;Put on top of stack
0850 2A0D04                    LHLD    DB$Save$RA      ;Get return address
0853 E3                        XTHL                    ;TOS = return address, HL = previous value
0854 C9                        RET
                    ;#
                    ;
                    ;          Debug output routine
                    ;
                    ;          This routine helps debug code in which output instructions
                    ;          would normally occur. The opcode of the OUT instruction
                    ;          must be replaced by a value of 0EFH (RST 5).
                    ;
                    ;          This routine picks up the port number contained in the byte
                    ;          following the RST 5, converts it to hexadecimal, and
                    ;          displays the message:
                    ;
```

Figure 10-2. (Continued)

```
                        ;               Output to port XX : AA
                        ;
                        ;       where AA is the contents of the A register prior to the
                        ;       RST 5 being executed.
                        ;       Control is then returned to the byte following the port number.
                        ;
                        ;       *******
                        ;       WARNING - This routine uses both DB$CONOUT and BDOS calls
                        ;       *******
                        ;
                        ;
0855 4F75747075DBO$Message:     DB      'Output to Port '
0864 5858203A20DBO$Port:        DB      'XX : '
0869 414100     DBO$Value:      DB      'AA',0
                        ;
                        ;
                        DB$Output:
086C 220904             SHLD    DB$Save$HL      ;Save user's HL
086F E1                 POP     H               ;Recover address of port number
0870 2B                 DCX     H               ;Backup to point to RST
0871 220F04             SHLD    DB$Call$Address ;Save for later display
0874 23                 INX     H               ;Restore to point at port number
0875 324B04             STA     DB$Save$A       ;Preserve value to be output
0878 7E                 MOV     A,M             ;Get port number
0879 23                 INX     H               ;Update return address to bypass port number
087A 220D04             SHLD    DB$Save$RA      ;Save return address
087D C5                 PUSH    B               ;Save remaining registers
087E D5                 PUSH    D
087F F5                 PUSH    PSW             ;Save port number for later

0880 CDB108             CALL    DB$Flag$Save$On ;Save current state of debug flag
                                                ;  and enable debug output

0883 CDC107             CALL    DB$CRLF         ;Display carriage return, line feed
0886 CD7C07             CALL    DB$Display$CALLA;Display call address
0889 F1                 POP     PSW             ;Recover port number
088A 216408             LXI     H,DBO$Port
088D CDA307             CALL    DB$CAH          ;Convert to hex.and store in message

0890 3A4B04             LDA     DB$Save$A
0893 216908             LXI     H,DBO$Value     ;Convert value to be output
0896 CDA307             CALL    DB$CAH          ;Convert to hex.and store in message

0899 215508             LXI     H,DBO$Message   ;Output prompting message
089C CDEE07             CALL    DB$MSG

089F CDBF08             CALL    DB$Flag$Restore ;Restore debug flag to previous state

08A2 D1                 POP     D               ;Recover registers
08A3 C1                 POP     B
08A4 2A0904             LHLD    DB$Save$HL      ;Get previous HL
08A7 E5                 PUSH    H               ;Put on top of stack
08A8 2A0D04             LHLD    DB$Save$RA      ;Get return address
08AB E3                 XTHL                    ;TOS = return address, HL = previous value
08AC 3A4B04             LDA     DB$Save$A       ;Recover A (NOTE: FLAG NOT RESTORED)
08AF C9                 RET
                        ;#
                        ;
                        ;       DB$Flag$Save$On
                        ;       This routine is only used for DB$IN/OUT.
                        ;       It saves the current state of the debug control flag,
                        ;       D$Flag, and then enables it to make sure that
                        ;       DB$IN/OUT output always goes out.
                        ;
08B0 00                 DB$Flag$Previous:       DB      0       ;Previous flag value
                        ;
                        DB$Flag$Save$On:
08B1 F5                 PUSH    PSW                     ;Save caller's registers
08B2 3A0604             LDA     DB$Flag                 ;Get current value
08B5 32B008             STA     DB$Flag$Previous        ;Save it
08B8 3EFF               MVI     A,0FFH                  ;Set flag
08BA 320604             STA     DB$Flag
08BD F1                 POP     PSW
08BE C9                 RET
                        ;#
                        ;
```

Figure 10-2. (Continued)

```
                        ;       DB$Flag$Restore
                        ;       This routine is only used for DB$IN/OUT.
                        ;       It restores the debug control flag, DB$Flag, to
                        ;       its former state.
                        ;
                        DB$Flag$Restore:
08BF F5                         PUSH    PSW
08C0 3AB008                     LDA     DB$Flag$Previous        ;Get previous setting
08C3 320604                     STA     DB$Flag                 ;Set debug control flag
08C6 F1                         POP     PSW
08C7 C9                         RET

                        ;
                        ;#
                        ;
                        ;       Get hex. value
                        ;
                        ;       This subroutine outputs a prompting message, and then reads
                        ;       the keyboard in order to get a hexadecimal value.
                        ;       It is somewhat simplistic in that the first non-hex value
                        ;       terminates the input. The maximum number of digits to be
                        ;       converted is specified as an input parameter. If more than the
                        ;       maximum number is entered, only the last four are significant.
                        ;
                        ;*****************************************************************
                        ;                       W A R N I N G
                        ;       DB$GHV will always use the BDOS to perform a read console
                        ;       function (#10). Be careful if you use this routine from
                        ;       within an executing BIOS.
                        ;*****************************************************************
                        ;
                        ;       Entry parameters
                        ;
                        ;               HL -> 00-byte terminated message to be output
                        ;               C = number of hexadecimal digits to be input
                        ;
                        ;
                        DB$GHV$Buffer:                  ;Input buffer for console characters
                        DB$GHV$Max$Count:
08C8 00                         DB      0               ;Set to the maximum number of chars.
                                                        ;  to be input
                        DB$GHV$Input$Count:
08C9 00                         DB      0               ;Set by the BDOS to the actual number
                                                        ;  of chars. entered
                        DB$GHV$Data$Bytes
08CA                            DS      5               ;Buffer space for the characters
                        ;
                        ;
                        DB$GHV:
08CF 79                         MOV     A,C                     ;Get maximum characters to be input
08D0 FE05                       CPI     5                       ;Check against maximum count
08D2 DAD708                     JC      DB$GHV$Count$OK         ;Carry set if A < 5
08D5 3E04                       MVI     A,4                     ;Force to only four characters
                        DB$GHV$Count$OK:
08D7 32C808                     STA     DB$GHV$Max$Count        ;Set up maximum count in input buffer
08DA CDEE07                     CALL    DB$MSG                  ;Output prompting message
08DD 11C808                     LXI     D,DB$GHV$Buffer         ;Accept characters from console
08E0 0E0A                       MVI     C,B$READCONS            ;Function code
08E2 CD0500                     CALL    BDOS

08E5 0E02                       MVI     C,B$CONOUT              ;Output a line feed
08E7 1E0A                       MVI     E,0AH
08E9 CD0500                     CALL    BDOS

08EC 210000                     LXI     H,0                     ;Initial value
08EF 11CA08                     LXI     D,DB$GHV$Data$Bytes     ;DE -> data characters
08F2 3AC908                     LDA     DB$GHV$Input$Count      ;Get count of characters input
08F5 4F                         MOV     C,A                     ;Keep count in C
                        DB$GHV$Loop:
08F6 0D                         DCR     C                       ;Downdate count
08F7 F8                         RM                              ;Return when all done (HL has value)
08F8 1A                         LDAX    D                       ;Get next character from buffer
08F9 13                         INX     D                       ;Update buffer pointer
08FA CD1B09                     CALL    DB$A$To$Upper           ;Convert A to uppercase if need be
08FD FE30                       CPI     '0'                     ;Check if less than 0
08FF D8                         RC                              ;Yes, terminate
0900 FE3A                       CPI     '9' + 1                 ;Check if > 9
0902 DA1009                     JC      DB$GHV$Hex$Digit        ;No, it must be numeric
```

Figure 10-2. (Continued)

```
0905 FE41           CPI     'A'                         ;Check if < 'A'
0907 D8             RC                                  ;Yes, terminate
0908 FE47           CPI     'F' + 1                     ;Check if > 'F'
090A D0             RNC                                 ;Yes, terminate
090B D637           SUI     'A' - 10                    ;Convert A through F to numeric
090D C31209         JMP     DB$GHV$Shift$Left$4         ;Combine with current result
              ;
              DB$GHV$Hex$Digit:
0910 D630           SUI     '0'                         ;Convert to binary
              DB$GHV$Shift$Left$4:
0912 29             DAD     H                           ;Shift HL left four bits
0913 29             DAD     H
0914 29             DAD     H
0915 29             DAD     H
0916 85             ADD     L                           ;Add binary value in LS 4 bits of A
0917 6F             MOV     L,A                         ;Put back into HL total
0918 C3F608         JMP     DB$GHV$Loop                 ;Loop back for next character
              ;#
              ;
              ;     A to upper
              ;     Converts the contents of the A register to an uppercase
              ;     letter if it is currently a lowercase letter
              ;
              ;     Entry parameters
              ;
              ;             A = character to be converted
              ;
              ;     Exit parameters
              ;
              ;             A = converted character
              ;
              DB$A$To$Upper:
091B FE61           CPI     'a'             ;Compare to lower limit
091D D8             RC                      ;No need to convert
091E FE7B           CPI     'z' + 1         ;Compare to upper limit
0920 D0             RNC                     ;No need to convert
0921 E65F           ANI     5FH             ;Convert to uppercase
0923 C9             RET
```

Figure 10-2. Debug subroutines (continued)

```
B>ddt fig10-2.hex<cr>
DDT VERS 2.0
NEXT  PC
0924 0000
-g100<cr>

0116 : Flags : Flags = C1Z0M1E1I0
0120 : A Register : A = AA
012F : B Register : B = BB
013E : C Register : C = CC
014D : D Register : D = DD
015C : E Register : E = EE
016B : H Register : H = FF
017A : L Register : L = 11
0189 : Memory Dump #1 : Start, End Address  : 0108, 0128
  0108 :  05 3E AA 01 CC BB 11 EE : .>*. L;.n
  0110 :  DD 21 11 FF B7 37 CD 52 05 00 46 6C 61 67 73 00 : ]!.. 77MR ..Fl ags.
  0120 :  CD 52 05 02 41 20 52 65 67 : MR.. A Re g

01A0 : Memory Dump #2 : Start, End Address  : 0100, 011F
  0100 :  31 6B 03 CD EA 04 CD 15 05 3E AA 01 CC BB 11 EE : 1k.M j.M. .>*. L;.n
  0110 :  DD 21 11 FF B7 37 CD 52 05 00 46 6C 61 67 73 00 : ]!.. 77MR ..Fl ags.

01B7 : Memory Dump #3 : Start, End Address  : 0101, 0100

** ERROR - Start Address > End **
01CE : Memory Dump #4 : Start, End Address  : 0100, 0100
  0100 :  31 : 1
```

Figure 10-3. Console output from debug testbed run

```
01E5 : BC Register : BC = BBCC
01F5 : DE Register : DE = DDEE
0205 : HL Register : HL = FF11
0215 : SP Register : SP = 0369
022E : Byte at (BC) : (BC) = BC
023F : Byte at (DE) : (DE) = DE
0250 : Byte at (HL) : (HL) = F1
026A : Word at (BC) : (BC+1),(BC) = 0B0C
027B : Word at (DE) : (DE+1),(DE) = 0D0E
028C : Word at (HL) : (HL+1),(HL) = 0F01
Debug output has been re-enabled.
This message should display 5 times
This message should display 5 times
This message should display 5 times
This message should display 5 times
This message should display 5 times
032B : Input from Port 11 : aa

032D : Output to Port 22 : AA
```

Figure 10-3. Console output from debug tested run (continued)

containing all of the symbols in your program, along with their respective addresses. Once the program has been loaded by SID, you can refer to the memory image of your program not by address, but by the actual symbol name from your source code. SID also supports the "pass count" concept when using breakpoints.

ZSID (Z80 Symbolic Debug)

This is the Z80 CPU's version of SID. The mini-assembler/disassembler uses Zilog instruction mnemonics rather than those used by Intel.

Bringing Up CP/M for the First Time

It is much harder to bring up CP/M on a new computer system than to debug an enhanced version on a system already running CP/M. You will often find yourself staring at a programmatic "brick wall" with no adequate debugging tools to assist you.

For example, you install the CP/M system on a diskette (using another CP/M-based computer system), put the diskette into the new computer, and press the RESET button. The disk head loads on the disk, and then—nothing! You cannot use any programs such as DDT or SID because you do not yet have CP/M up and running on the new computer. Or can you?

The answer is, wherever possible, debug the code for the new machine on an existing CP/M system. You may have to "fake" some aspects of the new bootstrap or BIOS so that the act of testing it on the host machine does not interact with the CP/M already running on it.

This scheme permits you to be fairly sure of your program logic before loading the diskette into the new machine. It will help pin down problems caused by hardware problems on the new computer.

The hardest situation of all is if you have only the new computer and the release diskettes from Digital Research. Your only option is to find a way of reading the CP/M image on the release diskette into memory, hand patch in new console and disk drivers (not a trivial task), write the patched image back onto a diskette, and resort to Orville Wright testing.

If you value your time, it is always more cost-effective to use another system with CP/M already installed. This is true even if the two systems do not have the same diskette format. You can still do the bootstrap and build the CP/M image on the host machine. Then download the image directly into the memory of the new machine and write it out to a diskette.

This *downloading* process does require, however, that the new computer have a read-only memory (ROM) monitor program. Depending on the capability of this ROM monitor program, you may have to hand patch into the new machine's memory a primitive "download" program that reads 8-bit characters from a serial port, stacking them up in memory and returning control to the monitor program when you press a keyboard character on the new machine's console. In fact, some ROM monitor programs have a downloading program built in.

Debugging the CP/M Bootstrap Loader

The CP/M bootstrap loader, as you may recall, is written on one of the outermost tracks on a diskette or hard disk. On a standard 8-inch single-sided, single-density diskette, CP/M's bootstrap loader is stored on the first sector of the first track. The loader is brought into memory by firmware that gets control of the CPU when you turn your machine on or press the RESET button.

The bootstrap has to be compact, as the diskette space on which it is stored is limited: no more than 128 bytes for standard 8-inch diskettes. This tends to rule out the use of the debug subroutines already described, so you have to fall back to more primitive techniques.

Testing the Bootstrap Under CP/M

A bootstrap is best developed on a CP/M-based system. The task is easiest of all if you already have CP/M running on your new machine and are simply preparing an enhanced version of the bootstrap loader. In this case, you can test most of the code as though it were a user program running in the transient program area (TPA).

Most bootstraps get loaded into memory at location 0000H, so at the front of the code to be debugged you must put a temporary origin line that reads

```
ORG         100H
```

If you omit this and ask DDT to load the HEX file output by the assembler, it will load at the true origin, 0000H, and wipe out the contents of the base page for the version of CP/M that you are running. This will cause a system crash; you will have to press the RESET button and reload CP/M. When this happens, DDT does not tell you directly that anything is amiss; it just displays a "?" after your request to load the HEX file. You will discover that the system has "gone away" only when you try to do something else.

You also will need to adjust the addresses into which the bootstrap tries to load the CP/M image. If you do not, you will overwrite the version of CP/M presently running.

With these adjustments made, you can load the bootstrap under DDT and watch it execute, confirming that it does load the correct image into the correct addresses for debugging and transfer control to the BIOS jump vector. When everything appears to be functioning correctly, use the IF instruction to disable the debug code, reassemble the bootstrap, and write it onto a diskette. Then put the diskette into drive A and press RESET.

Was the Bootstrap Loaded?

At this point you must establish whether the bootstrap is being loaded into memory when the machine is turned on or RESET is pressed. The best way of doing this, and one that you can leave in place permanently, is to output a sign-on message as soon as the loader gets control. This requires hardware set up to prepare the USART (Universal Synchronous/Asynchronous Receive/Transmit) chip to output data, although some manufacturers write this initialization code into the firmware that loads the bootstrap. A suitable sign-on message would be the following:

```
CP/M Bootstrap Loader : Vn 1.0 11/18/82
```

If you do not see this message, assume that control is *not* being transferred to the bootstrap loader. This will be useful in the future if someone should call you with a complaint that CP/M cannot be loaded. If this message does not appear, they probably do not have CP/M on the disk.

Did the Bootstrap Load CP/M?

This is a harder question to answer than whether the bootstrap itself has been loaded, especially if the bootstrap loader sign-on is displayed and then the system crashes. A sign-on message early in the BIOS cold boot processing can confirm the correct transfer of control into the BIOS.

If the problems with the bootstrap program are severe, you may have to adapt the memory-dump debugging subroutine, dumping the contents of memory to the console in order to see what information the bootstrap loader is placing in memory. Display 100H bytes starting from the front of the BIOS jump vector. This table has an immediately recognizable pattern of 0C3H values every three bytes.

You should also check to see that the bootstrap is loading the correct number of sectors from the disk into memory. If it loads too few, CP/M may sign on only to crash a few moments later because it attempts either to execute code or access a constant at the end of the BIOS. If the bootstrap loads too many sectors from the disk, the excess may "wrap around" the top of memory and overwrite the bootstrap itself, down at location 0000H, before it has completed its task. In this case, you would see only the sign-on for the bootstrap, not for the BIOS.

Debugging the BIOS

Rather than try to debug the BIOS as a single piece of code, debug it as a series of separate functional modules.

Notwithstanding current "top-down" philosophies of dealing with overall structure first, it can be quicker to debug the low-level subroutines in a device driver first. This gives you a solid base on which to build.

The BIOS can be divided up into its constituent modules as follows:

- Character input
  - Interrupt service
  - Non-interrupt service
- Character output
- Interrupt routines
  - Real time clock
  - Watchdog timers
- Disk drivers
  - High-level (deblocking)
  - Low-level (physical I/O)

Plan to write a *testbed* program for each of these modules. This testbed code serves two purposes; first, it provides a means of transferring control into the module under test in a controlled way. Second, it includes the necessary modules or dummy modules to "fool" the module under test into responding as if it were running in a complete BIOS under CP/M.

Using the testbed, you can check every part of the module's logic except the part that may be time-critical. Problems caused by timing, such as interrupts disabled for too long or code that is too slow or too fast for a particular peripheral controller chip, tend to show up only when you are testing on the final hardware and when you are running your new BIOS under CP/M.

What You Should Test for in the BIOS

Describing fully how to debug each module in the BIOS ould fill several books. Remember that you are trying to establish the *absence* of errors using a technique that, by its very nature, tends to show only their *presence*.

There are two basic approaches to debugging. One is the plodding method, checking every aspect of the code to ensure that every feature really does work. The second is to try to do something useful with the code.

Plan to use both. Start with the plodding method, testing each feature under control of the testbed until you are sure that it is working *in vitro*. When all of the BIOS modules have been tested individually, build a CP/M system and try to do some useful work with it. Trying to use the system for actual work testing *in vitro* can be a good test.

Feature Checklist

Make a list of the specific features included in the various BIOS modules. Then devise specific test sequences that will show that each of the features is working correctly.

The same testbed code can often test all of the features of a driver module. If it cannot, create a new testbed for the more exotic features.

Keep the testbed routines. Experience shows that they are most often needed shortly after you have erased them. Even after you have tested the BIOS, the testbed routines will come in handy if you decide to enhance a particular driver later on. You can extract the driver code from the BIOS, glue it together with the testbed, and test the new feature code in isolation from the BIOS.

The following sections show example testbeds for the various drivers, along with example checklists. These checklists were used to test the example BIOS routines shown in earlier chapters.

Character Drivers

Figure 10-4 shows the code for an example testbed routine for character I/O drivers in the BIOS. This code would be followed by the actual character I/O drivers, exactly as they would appear in the BIOS except that all IN and OUT instructions would be replaced with RST 4's and 5's respectively (see Figure 10-2) so that you could enter input values and inspect output values on the console.

This example contains the initialization code for the debug package shown in Figure 10-2 and the code setting up an RST 6 used to "fake" incoming character interrupts.

The main testbed loop consists of a faked incoming character interrupt followed by optional calls to CONIN or CONOUT, the return of control to DDT, or a loop back to fake another character interrupt. You can only return control to DDT if you used DDT to load the testbed and driver programs in the first place.

```
                ;       Testbed for character I/O drivers in the BIOS
                ;
                ;       The complete source file consists of three components:
                ;
                ;               1. The testbed code shown here
                ;               2. The character I/O drivers destined for the BIOS
                ;               3. The debug package shown in Figure 10-2.
                ;
FFFF =          TRUE    EQU     0FFFFH
0000 =          FALSE   EQU     NOT TRUE

FFFF =          DEBUG   EQU     TRUE            ;For conditional assembly of RST
                                                ;  instructions in place of IN and
                                                ;  OUT instructions in the drivers
0030 =          RST6    EQU     30H             ;Use RST 6 for fake incoming character
                                                ;  interrupt
0100                    ORG     100H
                START:
0100 31D101             LXI     SP,Test$Stack   ;Use a local stack
0103 CDD101             CALL    DB$Init         ;Initialize the debug package
0106 3EC3               MVI     A,JMP           ;Set up RST 6 with JMP opcode
0108 323000             STA     RST6
010B 21D101             LXI     H,Character$Interrupt   ;Set up RST 6 JMP address
010E 223100             SHLD    RST6 + 1
                ;
                ;       Make repeated entry to character interrupt routine
                ;       to ensure that characters can be captured and stored in
                ;       an input buffer
                ;
                Testbed$Loop:
0111 3EAA               MVI     A,0AAH          ;Set registers to known pattern
0113 01CCBB             LXI     B,0BBCCH
0116 11EEDD             LXI     D,0DDEEH
0119 2111FF             LXI     H,0FF11H
011C F7                 RST     6               ;Fake interrupt for incoming character

011D CDD101             CALL    DB$MSGI         ;Display in-line message
0120 0D0A456E74         DB      0DH,0AH,'Enter I to Input Char., O to Output, D to enter
0152 444454203A         DB      'DDT : ',0

0159 CDD101             CALL    DB$CONINU       ;Get uppercase character
015C FE49               CPI     'I'             ;CONIN?
015E CA7201             JZ      Go$CONIN
0161 FE44               CPI     'D'             ;DDT?
0163 CA6E01             JZ      Go$DDT
0166 FE4F               CPI     'O'             ;CONOUT?
0168 CA9101             JZ      Go$CONOUT
016B C31101             JMP     Testbed$Loop    ;Loop back to interrupt again
                Go$DDT:
016E FF                 RST     7               ;Enter DDT (RST 7 set up by DDT)
016F C31101             JMP     Testbed$Loop
                Go$CONIN:
0172 CDD101             CALL    CONST           ;Get console status
0175 CA1101             JZ      Testbed$Loop    ;No data waiting
0178 CDD101             CALL    CONIN           ;Get data from buffer

017B CDD101             CALL    DB$Display      ;Display character returned
017E 02                 DB      DB$A            ;  in A register
017F 434F4E494E         DB      'CONIN returned',0

018E C37201             JMP     Go$CONIN        ;Repeat CONIN loop until no chars.
                                                ;  waiting
                ;
                Go$CONOUT:
0191 CDD101             CALL    CONST           ;Get console status
0194 CA1101             JZ      Testbed$Loop    ;No data waiting
0197 CDD101             CALL    CONIN
019A 4F                 MOV     C,A             ;Ready for output
019B CDD101             CALL    CONOUT          ;Output to console
019E C39101             JMP     Go$CONOUT       ;Repeat while there is still data
                ;
01A1 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
01B1 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
01C1 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
```

Figure 10-4. Testbed for character I/O drivers in the BIOS

```
                Test$Stack:
                ;
                ;       Dummy routines for those shown in other figures
                ;
                ;       BIOS routines (Figure 8-10)
                ;
                CONST:                  ;BIOS console status
                CONIN:                  ;BIOS console input
                CONOUT:                 ;BIOS console output;
                Character$Interrupt:    ;Interrupt service routine for incoming chars.
                ;
                ;       Debug routines (Figure 10-2)
                ;
                DB$Init:                ;Debug initialization
                DB$MSGI:                ;Display message in-line
                DB$CONINU:              ;Get uppercase character from keyboard
                DB$Display:             ;Main debug display routine
0002 =          DB$A    EQU     02      ;Display code for DB$Display
```

Figure 10-4. Testbed for character I/O drivers in the BIOS (continued)

Executing an RST 7 without using DDT will cause a system crash, as DDT sets up the necessary JMP instruction at location 0038H in the base page.

The faked incoming character interrupt transfers control directly to the interrupt service routine in the BIOS (see the example in Figure 8-10, line 04902, label Character$Interrupt). This reads the status ports of each of the character devices; you can enter the specific status byte values that you want. If you enter a value that indicates that a data character is "incoming," you will be prompted for the actual 8-bit data value to be "input." You can make the interrupt service routine appear to be inputting characters and stacking characters up in the input buffer. For debugging purposes, reduce the size of the input buffer to eight bytes. Making it larger means you will have to input more characters to test the buffer threshold logic. To check the interrupt service routine, you will pass through the main testbed loop doing nothing but faking incoming character interrupts and entering status and data values. The data characters will then be stacked up in the input buffer.

To check the correct functioning of the interrupt service routines, you can stay in control with DDT from the outset. Alternatively, you can just use DDT to load the testbed/driver HEX file, loop around inputting several characters, and then request that the testbed return control to DDT. Then you can use DDT to inspect the contents of the device table(s) and input buffers.

Another possibility is to create debugging routines that display the contents of the device table in a meaningful way, with each field captioned like this:

```
DEVICE TABLE 0
     Status Port     81    Data Port       80
     Output Ready    01    Input Ready     02
     DTR high        40
     Reset Int. Prt D8    Reset Int. Val. 20
     :
     :
     Status Byte 1
          Output Suspended
          Output Xon Enabled
     :
```

```
:
Buffer Base OE8C
Put Offset      05      Get Offset       01
Char. Count     04      Control Count    00
Data Buffer
41 42 43 44 45 00 00 00
```

This display device table routine will require a fair amount of effort to code and debug—but it will pay dividends. You can obtain a complete "snapshot" of the device table without having to decode hexadecimal memory dumps and individual bits. Constant values in the device tables are also displayed, so that if a bug in your code corrupts the table, you will know about it immediately.

The next section shows examples of the specific tests you need to make, along with a description of the strategy you can use.

Interrupt Service Routine Checklist In a functioning BIOS, control is transferred to the interrupt service module whenever an incoming character causes an interrupt. In the example BIOS in Figure 8-10 (line 4900), the code scans each character device in turn to determine which one is causing the interrupt.

When you are debugging the interrupt service routines using the "fake" input/output instructions, you will have to enter specific status byte values. Refer to the device table declarations in Figure 8-10, line 1500, to determine what values you must enter to make the service routine think that an incoming character is arriving or that data terminal ready (DTR) is high or low.

Start the debugging process using the first device table. Then repeat the tests on the other device tables.

The following is a checklist of features that should be checked in debugging the interrupt service routine:

Are all registers restored correctly on exit from the interrupt servicing?

Using DDT, start execution from the beginning of the testbed. Set a breakpoint (with the G100,nnnn command) to get control back immediately before the CALL Character$Interrupt. Use the X command to display all of the registers, and then, by using the G,nnnn command, you set a breakpoint at the instruction that immediately follows the CALL Character$Interrupt. The character drivers will prompt you for the status values. Enter 00 (which indicates that no character is incoming). Display the registers again—their values should be the same. Remember to check the value of the stack pointer and the amount of the stack area that has been used.

NOTE: Do not be too surprised if you lose control of the machine when you first try this test. You may have some fundamental logic errors initially. If the system crashes, reset it, reload CP/M, and then start the test again. This time, rather than setting the second breakpoint at the instruction following the CALL Character$Interrupt, venture down into the Character$Interrupt code and go through the code a few instructions at a time, setting breakpoints before any instructions that could cause a transfer of control. Find out how far you are getting into the driver before it either jumps off into space or settles into a loop.

Does the service routine push a significant number of bytes onto the stack after an interrupt has occurred?

When you get control back after the CALL Character$Interrupt, use the D (dump) command to dump the stack area's memory on the console. Check how far down the stack came by looking for the point where the constants that used to fill the stack area are overwritten by other data.

The example BIOS in Figure 8-10 saves only the contents of the HL register pair on the pre-interrupt stack. It then switches over to a private BIOS stack to save the contents of the rest of the registers and service the interrupt.

Are data characters added to the input buffer correctly?

"Input" a noncontrol character via the Character$Interrupt routine. Then check the contents of the appropriate device table. The character count and the put offset should both be set to one. Then check the contents of the input buffer itself; does it contain the character that you "input?"

Are control characters added to the input buffer correctly?

"Input" a control character such as 01H. Do not use ETX, ACK, XON, or XOFF (03H, 06H, 11H, and 13H, respectively); these may cause side effects if you have errors in the protocol handling logic. Check that the character is stored in the next byte of the input buffer and that the character and control counts are set to two and one, respectively. The put offset should also be set to two.

When the input buffer full threshold is reached, does the driver output the correct protocol character?

Set the first status byte in the first device table to enable input XON or RTS protocol, or both. Then go round the main testbed loop putting characters into the input buffer. Check the console display to see if the drivers output the correct values when the buffer is almost full (the default threshold is when five bytes remain). The driver should then drop the RTS line or output an XOFF character or both, according to the input protocol that you enabled.

When the input buffer is completely full, does the driver respond correctly?

This is an extension of the test above. Input one more character than can fit into the buffer. Check to see that the drivers do not stack the character into the input buffer and that a BELL character (07H) is output to the data port.

Are protocol characters XON/XOFF recognized and the necessary control flags set or reset?

Reload the testbed and drivers. Set the status byte to enable the output XON/XOFF protocol. Then use the Character$Interrupt routine to input an XOFF character (13H). Check to see that the XOFF character has not been put into the input buffer. Instead, the status byte should be set to indicate that output has indeed been suspended.

Input an XON and check to see that the output suspended flag has been reset.

Does the driver detect and reset hardware errors correctly?

Proceed as though you were going to input a character into the input buffer, but instead enter a status byte value that indicates that a hardware error has occurred (enter the value given in the device table for DT$Detect$Error$Value).

Check that the driver detects the error status and outputs the correct error-reset value to the appropriate control port.

Non-interrupt Service Routine Checklist In a "live" BIOS, non-interrupt service routines are accessed via the CONIN and CONST entry points in the BIOS jump vector. During debugging, the testbed can call the CONIN and CONST code directly.

Is input redirection functioning? Does control arrive in the driver with the correct device table selected?

This is best tested directly with DDT. Use the Gnnnn,bbbb command to transfer control into the CONIN code with a breakpoint at the RET instruction at the end of the Select$Device$Table routine (see Figure 8-10, line 04400). Check that the DE register pair is pointing at device table 0. If it is not, you will have to restart the test. Use the Tn command to make DDT trace through the Select$Device$Table subroutine to find the bug.

Are characters returned correctly from the buffer?

Use the testbed to "input" a character or two. Then use the testbed to make several entries into CONIN. Check the characters returned from the buffer.

Are the data character and control character counts correctly decremented?

After each character has been removed from the buffer by CONIN, use DDT to examine the device table and check that the data character and control character counts have been decremented correctly. Also check that the get pointer has moved up the input buffer.

When the buffer "almost empty" threshold is reached, does the driver emit the correct protocol character or manipulate the request to send (RTS) line correctly?

Use DDT to enable the input RTS or XON protocol or both. Then input characters into the input buffer until it reaches the buffer full threshold (the default is when only five spare bytes remain in the buffer). Confirm that "buffer almost full" processing occurs. Then make repetitive calls to CONIN to flush data out of the buffer. Check that the "buffer emptying" processing occurs when the correct threshold is reached. For RTS protocol, the driver should output a raise RTS value to the specified RTS control port. For XON, the driver should output an XON character to the data port (after first having read the status port to ensure that the hardware can output the character).

Does the driver handle buffer "wraparound" correctly?

Input characters to the input buffer until it becomes completely full. Then make a single CONIN call to remove the first character from the buffer. Follow this by inputting one more character to the buffer. Check that the get pointer is set to one and the put pointer set to zero.

Next, make successive CONIN calls to empty the buffer. Then input one more character to the buffer. Check that this last character is put into the first byte of the input buffer.

Can the driver handle "forced input" correctly?

Using DDT, set the forced input pointer to point to a 00-byte-terminated string; for example, use one of the function key decode default strings. (In Figure 8-10, the forced input pointer is initialized to point to a "startup string"—this is declared at the beginning of the configuration block at line 00400.)

Using DDT, call the CONST routine and check that it returns with A= 0FFH (indicating that there appears to be input data waiting).

Make successive calls to CONIN and confirm that the data bytes in the forced input string are returned. Check that the forcing of input ends when the 00H-byte is detected.

Does the console status routine operate correctly when it checks for data characters in the buffer, control characters in the buffer, and forced input?

Input a single noncontrol character, such as 41H, into the input buffer. Using DDT, check that the second status byte in the device table has the fake type-ahead flag set to zero. Call the CONST routine—it should return with A= 0FFH (meaning that there is data in the buffer). Then set the fake type-ahead bit in the second status byte and call CONST again. It should return with A = 00H (meaning that there is now "no data" in the buffer). Input a single control character into the buffer. Now CONST should return with A = 0FFH because there is a control character in the buffer.

Does the driver recognize escape sequences incoming from keyboard function keys?

This is a difficult feature to test when the real time clock routine is not running. The driver uses the watchdog timer to wait until all characters in the escape sequence have arrived. You will therefore have to modify the code in CONIN so that the watchdog timer appears to time out immediately, rather than waiting for the real time clock to tick. To make this change, refer to Figure 8-10, line 2200; this is the start of the CONIN routine. Look for the label CONIN$Wait$For$Delay. A few instructions later there is a JNZ CONIN$Wait$For$Delay. Using DDT, set all three bytes of this JNZ to 00H.

Then, using the testbed, input the complete escape sequence into the input buffer. For example, input hexadecimal values 1B, 4F, 51 (ESCAPE, O, P), which correspond to the characters emitted on a VT-100 terminal when FUNCTION KEY 1 (PF1) is pressed.

Next, use the testbed to make successive calls to CONIN. You should see the text associated with the function key (FUNCTION KEY 1, LINE FEED) being returned by CONIN.

Repeat this test using different function key sequences, including a sequence that does not correspond to any of the preset function keys. Check that the escape sequence itself is returned by CONIN without being changed into another string.

Can the driver differentiate between a function key and the same escape sequence generated by discrete key strokes?

This is almost the same test as above. Make the same patch to the CONIN code, only this time do not enter the complete escape sequence into the buffer. Enter only the hex characters 1B and 4F. Make sure that the CONIN routine does not substitute another string in place of this quasi-escape sequence.

This test only mimics the results of manually entering an escape sequence. You could not press the keys on a terminal fast enough to get all three characters into the input buffer within the time allowed by the watchdog timer.

Character Output Checklist

Can the driver output a character?

The CONOUT option in the testbed calls CONIN first to get a character. To start with, you may want to use DDT to set the C register to some graphic ASCII character such as 41H (A), and transfer control into CONOUT directly. Check that CONOUT reads the USART's status, waits for the output ready value, and then outputs the data to the data port. Note that the testbed will output all characters waiting in the input buffer (or forced input) when you select its CONOUT option. This is a convenience for advanced testing of the drivers—for initial testing you may want to modify the testbed to make only one call to CONIN and CONOUT and then return to the top of the testbed loop.

Does the driver suspend output when a protocol control flag indicates that output is to be suspended?

Using DDT, set the status byte in the device table to enable output XON/XOFF protocol. Then input an XOFF character and confirm that the output suspended bit in the status byte is set. Output a single character, and using DDT, confirm that the driver will remain in a status loop waiting for the output suspended bit to be cleared. Clear the bit using DDT and check that the character is output correctly.

When using ETX/ACK protocol, does the driver output an ETX after the specified number of characters have been output, then indicate that output is suspended?

For debugging purposes, alter the ETX message count value in the device table to three bytes. Then output three bytes of data via CONOUT. Check that the driver sends an ETX character (03H) after the three bytes have been output and that the output suspended flag in the status byte has been set.

Then input an ACK character (06H). Check that this character is not stored in the input buffer and that the output suspended flag is cleared.

Does the driver recognize and output escape sequences?

Input an ESCAPE, "t" (1BH, 74H) into the input buffer. Then output them via CONOUT. Using DDT, check that the CONOUT routine recognizes that an escape sequence is being output and selects the correct processing routine. In this case, the forced input pointer should be set to point at the ASCII time of day in the configuration block.

Does each of the escape sequence processors function correctly? Can the time and date be set to specified values using escape sequences?

Repeat the test above using all of the other escape sequences to make sure that they can be recognized and that they function correctly.

Real Time Clock Routines

A separate testbed program, shown in Figure 10-5, is used to check these routines. It calls the interrupt service routine directly to simulate a real time clock "tick," and then displays the time of day in ASCII on the console.

As you can see, the testbed makes a call into the debug package's initialization routine, DB$Init, and then uses an RST 6 to generate fake clock "ticks."

There is a JMP instruction in the testbed that bypasses a call to Set$Watchdog. Remove this JMP, either by editing it out or by using DDT to change it to NO OPERATIONs (NOP, 00H) when you are ready to test the watchdog routines.

Real Time Clock Test Checklist

Is the clock running at all?

Using DDT, trace through the interrupt service routine logic. Check that the seconds are being updated.

```
                ;       Testbed for real time clock driver in the BIOS.
                ;
                ;       The complete source file consists of three components:
                ;
                ;               1. The testbed code shown here
                ;               2. The real time clock driver destined for the BIOS.
                ;               3. The debug package shown in Figure 10-2.
                ;
FFFF =          TRUE    EQU     0FFFFH
0000 =          FALSE   EQU     NOT TRUE

FFFF =          DEBUG   EQU     TRUE            ;For conditional assembly of RST
                                                ;  instructions in place of IN and
                                                ;  OUT instructions in the drivers.
0030 =          RST6    EQU     30H             ;Use RST 6 for fake clock tick.

0100                    ORG     100H
                START:
0100 318B01             LXI     SP,Test$Stack   ;Use local stack
0103 CD8B01             CALL    DB$Init         ;Initialize the debug package
0106 3EC3               MVI     A,JMP           ;Set up RST 6 with JMP opcode
0108 323000             STA     RST6
010B 218B01             LXI     H,RTC$Interrupt ;Set up RST 6 JMP address
010E 223100             SHLD    RST6 + 1

0111 C31D01             JMP     Testbed$Loop    ;<=== REMOVE THIS JMP WHEN READY TO
                                                ;     TEST WATCHDOG ROUTINES

0114 013200             LXI     B,50            ;50 ticks before timeout
0117 214201             LXI     H,WD$Timeout    ;Address to transfer to
011A CD8B01             CALL    Set$Watchdog    ;Set the watchdog timer
                ;
                ;
                ;       Make repeated entry to RTC interrupt routine
                ;       to ensure that clock is correctly updated
                ;
                Testbed$Loop:
011D 3EAA               MVI     A,0AAH          ;Set registers to known pattern
011F 01CCBB             LXI     B,0BBCCH
0122 11EEDD             LXI     D,0DDEEH
0125 2111FF             LXI     H,0FF11H
0128 F7                 RST     6               ;Fake interrupt clock

0129 CD8B01             CALL    DB$MSGI         ;Display in-line message
012C 436C6F636B         DB      'Clock =',0

0134 218B01             LXI     H,Time$In$ASCII ;Get address of clock in driver
0137 CD8B01             CALL    DB$MSG          ;Display current clock value
                                                ; (Note: Time$In$ASCII already has
                                                ;  a line feed character in it)
013A CD8B01             CALL    DB$MSGI         ;Display in-line message
013D 0D00               DB      0DH,0           ;Carriage return

013F C31D01             JMP     Testbed$Loop
                ;
                ;       Control arrives here when the watchdog timer times
                ;       out
                WD$Timeout:
0142 CD8B01             CALL    DB$MSGI
0145 0D0A576174         DB      0DH,0AH,'Watchdog timed out',0
015A C9                 RET                     ;Return to watchdog routine
                ;
015B 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
016B 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
017B 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
                Test$Stack:
                ;
                ;       Dummy routines for those shown in other figures
                ;
                ;       BIOS routines (Figure 8-10)
                ;
                RTC$Interrupt:          ;Interrupt service routine for clock tick
                Set$Watchdog:           ;Set watchdog timer
                Time$In$ASCII:          ;ASCII string of HH:MM:SS, LF, 0
                ;
                ;       Debug routines (Figure 10-2)
                ;
                DB$Init:                ;Debug initialization
                DB$MSGI:                ;Display message in-line
                DB$MSG:                 ;Display message
```

Figure 10-5. Testbed for real-time-clock driver in the BIOS

Are the hours, minutes, and seconds carrying over correctly?

Let the testbed code run at full speed. You should see the time being updated on the console display—although it will be updated much more rapidly than real time.

Use DDT to set the minutes to 58 and then let the clock run again. Does it correctly show the hour and reset the minutes to 00? Then set the hours to 11 and the minutes to 58 and let the clock run. Do minutes carry over into hours and are hours reset to 0?

Repeat these tests with the clock update constants set for 24-hour format.

Is the clock interrupt service routine restoring the registers correctly?

Using DDT, check that the registers are still set correctly on return from the clock interrupt service routine.

How much of a load on the pre-interrupt stack is the service routine imposing?

Check the "low water mark" of the preset values remaining in the testbed stack area to see how much of a load the interrupt service routine is imposing on the stack.

Can the watchdog timer be set to a nonzero value? Can it be set back to zero?

Using the second part of the testbed, call the Set$Watchdog routine, and then monitor the testbed's execution as the watchdog timer times out. Check that the registers and stack pointer are set correctly when control is transferred to the timeout routine. Also check that control is returned properly from this routine, and thence from the interrupt service routine.

Disk Drivers

It is only feasible to check the low-level disk drivers in isolation from a real BIOS, as the BDOS interface to the deblocking code is very difficult to simulate. The testbed shown in Figure 10-6 serves only as a time-saver. It does not test the interface to the subroutines. Use DDT to set up the disk, track, and sector numbers, and then monitor the calls into SELDSK, SETTRK, SETSEC, SETDMA, and the read/write routines.

Unless you have the same disk controller on the host system as you do on the target machine, you will have to use the fake input/output system described earlier in this chapter, rather than attempt to read and write on real disks.

You can see that the testbed, after initializing the debugging package, makes calls to SELDSK, SETTRK, SETSEC, and SETDMA. It then calls a low-level read or write routine. The low-level routine called depends on which driver you wish to debug. For the standard floppy diskette driver shown in Figure 8-10, use ReadNoDeblock and WriteNoDeblock. For the 5 1/4-inch diskettes, use Read$Physical and Write$Physical. You will have to use DDT to set up some of the variables required by the low-level drivers that would normally be set up by the deblocking code.

```
                ;       Testbed for disk I/O drivers in the BIOS
                ;
                ;       The complete source file consists of three components:
                ;
                ;               1. The testbed code shown here
                ;               2. The Disk I/O drivers destined for the BIOS
                ;               3. The debug package shown in Figure 10-2.
                ;
FFFF =          TRUE    EQU     0FFFFH
0000 =          FALSE   EQU     NOT TRUE

FFFF =          DEBUG   EQU     TRUE            ;For conditional assembly of RST
                                                ;  instructions in place of IN and
                                                ;  OUT instructions in the drivers.
0100                    ORG     100H
                START:
0100 314704             LXI     SP,Test$Stack   ;Use a local stack
0103 CD4704             CALL    DB$Init         ;Initialize the debug package
                ;
                ;       Make calls to SELDSK, SETTRK, SETSEC and SETDMA,
                ;       then either a read or write routine.
                ;
                Testbed$Loop:
0106 314704             LXI     SP,Test$Stack   ;Use local stack

0109 3A1202             LDA     Logical$Disk    ;Set up for SELDSK call
010C 4F                 MOV     C,A
010D CD4704             CALL    SELDSK

0110 CD4704             CALL    DB$Display      ;Display return value in HL
0113 14                 DB      DB$HL
0114 53454C4453         DB      'SELDSK returned',0

0124 223201             SHLD    DPH$Start       ;Set up to display disk parameter header
0127 111000             LXI     D,16            ;Compute end address
012A 19                 DAD     D
012B 223401             SHLD    DPH$End         ;Store into debug call

012E CD4704             CALL    DB$Display      ;Display DPH
0131 18                 DB      DB$M            ;Memory
                DPH$Start:
0132 0000               DW      0
                DPH$End:
0134 0000               DW      0
0136 53656C6563         DB      'Selected DPH',0

0143 2A1302             LHLD    Track           ;Call SETTRK
0146 E5                 PUSH    H
0147 C1                 POP     B               ;SETTRK needs track in BC
0148 CD4704             CALL    SETTRK

014B 3A1502             LDA     Sector          ;Call SETSEC
014E 4F                 MOV     C,A             ;SETSEC need sector in C
014F CD4704             CALL    SETSEC

0152 011702             LXI     B,Test$Buffer   ;Set DMA address
0155 CD4704             CALL    SETDMA
0158 3A1602             LDA     Write$Disk      ;Check if reading or writing
015B B7                 ORA     A
015C C2D101             JNZ     Test$Write

015F CD4704             CALL    Read$No$Deblock ;*** or Read$Physical depending on which
                                                ;*** drivers you are testing
0162 CD4704             CALL    DB$Display      ;Display return code
0165 02                 DB      DB$A
0166 5465737420         DB      'Test Read returned',0

0179 CD0102             CALL    Check$Ripple    ;Check if ripple pattern in buffer
017C CA0601             JZ      Testbed$Loop    ;Yes, it is correct

017F CD4704             CALL    DB$MSGI         ;Indicate problem
0182 14                 DB      DB$HL           ;Display HL (points to offending byte)
0183 526970706C         DB      'Ripple pattern incorrect. HL -> failure.',0

01AC CD4704             CALL    DB$Display      ;Display test buffer
01AF CD1800             CALL    DB$M            ;Memory
01B2 1702               DW      Test$Buffer
```

Figure 10-6. Testbed for disk I/O drivers in the BIOS

```
01B4 0002               DW      Test$Buffer$Size
01B6 436F6E7465         DB      'Contents of Test$Buffer',0

01CE C30601             JMP     Testbed$Loop

               Test$Write:
01D1 CDF201             CALL    Fill$Ripple      ;Fill the test buffer with ripple pattern
01D4 CD4704             CALL    Write$No$Deblock;*** or Write$Physical depending on which
                                                 ;*** drivers you are testing

01D7 CD4704             CALL    DB$Display       ;Display return code
01DA 02                 DB      DB$A
01DB 5465737420         DB      'Test Write returned',0

01EF C30601             JMP     Testbed$Loop

               Fill$Ripple:                      ;Fills the Test$Buffer with a pattern
                                                 ;  formed by putting into each byte, the
                                                 ;  least significant 8-bits of the byte's
                                                 ;  address.
01F2 010002             LXI     B,Test$Buffer$Size
01F5 211702             LXI     H,Test$Buffer
               FR$Loop:
01F8 75                 MOV     M,L              ;Set pattern value into buffer
01F9 23                 INX     H                ;Update buffer pointer
01FA 0B                 DCX     B                ;Down date count
01FB 79                 MOV     A,C              ;Check if count zero
01FC B0                 ORA     B
01FD C2F801             JNZ     FR$Loop          ;Repeat until zero
0200 C9                 RET
               ;
               Check$Ripple:                     ;Check that the buffer is filled with the
                                                 ;  correct ripple pattern
                                                 ;  Returns with zero status if this is true,
                                                 ;  nonzero status if the ripple is not
                                                 ;  correct. HL point to the offending byte
                                                 ;  (which should  = L)
0201 010002             LXI     B,Test$Buffer$Size
0204 211702             LXI     H,Test$Buffer
               CR$Loop:
0207 7D                 MOV     A,L              ;Get correct value
0208 BE                 CMP     M                ;Compare to that in the buffer
0209 C0                 RNZ                      ;Mismatch, nonzero already indicated
020A 23                 INX     H                ;Update buffer pointer
020B 0B                 DCX     B                ;Downdate count
020C 79                 MOV     A,C              ;Check count zero
020D B0                 ORA     B
020E C20702             JNZ     CR$Loop          ;Repeat until zero
0211 C9                 RET                      ;Zero flag will already be set
               ;
               ;        Testbed variables
               ;
0212 00        Logical$Disk:    DB      0        ;A = 0, B = 1,...
0213 0000      Track:           DW      0        ;Disk track number
0215 00        Sector:          DB      0        ;Disk sector number
0216 00        Write$Disk:      DB      0        ;NZ to write to disk
               ;
0200 =         Test$Buffer$Size         EQU     512     ;<=== Alter as required
0217           Test$Buffer:     DS      Test$Buffer$Size
               ;
0417 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
0427 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
0437 9999999999         DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
               Test$Stack:
               ;
               ;        Dummy routines for those shown in other figures
               ;
               ;        BIOS routines (Figure 8-10)
               ;
               SELDSK:                  ;Select logical disk
               SETTRK:                  ;Set track number
               SETSEC:                  ;Set sector number
               SETDMA:                  ;Set DMA address
               Read$No$Deblock:         ;Driver read routines
               Read$Physical:
               Write$No$Deblock:        ;Driver write routines
               Write$Physical:
```

Figure 10-6. (Continued)

```
                ;
                ;       Debug routines (Figure 10-2)
                ;
                DB$Init:                        ;Debug initialization
                DB$MSGI:                        ;Display message in-line
                DB$Display:                     ;Main debug display routine
0002 =          DB$A    EQU     02              ;Display codes for DB$Display
0014 =          DB$HL   EQU     20
0018 =          DB$M    EQU     24
```

Figure 10-6. Testbed for disk I/O drivers in the BIOS (continued)

Before issuing the write call, the testbed fills the disk buffer with a known pattern. This pattern is checked on return from a read operation.

For both reading and writing, the testbed shows the contents of the A register. If you have added the enhanced disk error handling described in the previous chapter, the return value in A must *always* be zero.

Disk Driver Checklist *Does SELDSK return the correct address and set up the required system variables?*

Check that the correct disk parameter header address is returned for legitimate logical disks. Check, too, that it returns an address of 0000H for illegal disks.

Check that any custom processing, such as setting the disk type and deblocking requirements from extra bytes on the disk parameter blocks, is performed correctly.

Does the SETTRK and SETSEC processing function correctly?

Using DDT, check that the correct variables are set to the specified values.

Does the driver read in the spare-sector directory correctly?

Set up to execute a physical read and, using DDT, trace the logic of the READ entry point. Check that the spare-sector directory would be loaded into the correct buffer. If you are using fake input/output, use DDT to patch in a typical spare-sector directory with two or three "spared-out" sectors.

Does the driver produce the correct spare sector in place of a bad one?

Continuing with the physical read operation, check that, for "good" track/sectors, the sector-sparing logic returns the original track and sector number, and for "bad" track/sectors, it substitutes the correct spare track and sector. If you are using sector skipping, check that the correct number of sectors is skipped.

Can a sector be read in from the disk?

Continuing further with the physical read, check that the correct sector is read from the specified disk and track. If you are using real I/O (as opposed to faking it), the "ripple pattern" set by the testbed can be used, or you can fill the disk buffer area with some known pattern (using DDT's F command) so you can tell if any data gets read in.

Make sure you do not have any disks or diskettes in the computer system that are not write-protected — you may inadvertently write on a disk rather than read it during the early stages of testing.

Can a sector be written to the disk?

Using DDT, set up to write to a particular disk, track, and sector. Remove any write protection that you put on the target disk during earlier testing. You can either use the testbed's ripple pattern or fill the disk buffer area with a distinctive pattern. Write this data onto the disk, fill the buffer area with a *different* pattern, and read in the sector that you wrote. Check that the disk buffer gets changed back to the pattern written to the disk.

Does the driver display error messages correctly?

Rather than deliberately damaging a diskette to create errors, use DDT to temporarily sabotage the disk driver's logic. Make it return each of the possible error codes in turn, checking each time that the correct error message is displayed.

For each error condition in turn, check that the disk driver performs the correct recovery action, including interacting with the user and offering the choice of retrying, ignoring the error, or aborting the program.

Live Testing a New BIOS

Given that the drivers have passed all of the testing outlined above, you are ready to pull all of the BIOS pieces together and build a CP/M image.

For your initial testing, disable the real time clock, and use simple, polled I/O for the console driver if you can. It is important to get *something* up and running as soon as possible, and it is easier to do this without possible side effects from interrupts.

Prepare a complete listing of the BIOS and plan to spend at least an hour checking through it. Take a dry run through the console and disk driver — if there are any serious bugs left in these two drivers, CP/M may not start up. Remember that once the BIOS cold boot code has been executed and control is handed over to the CCP, the BDOS will be requested to log in the system disk, and this involves reading in the disk's directory.

Pay special attention to checking some of the major data structures. Make certain that everything is at a reasonable place in memory; for example, if the last address used by the BIOS is greater than 0FFFFH, you will need to move the entire CP/M image down in memory.

Then build a system disk, load it into the machine, and press the RESET button. You should see the bootstrap sign on, then the BIOS, and after a pause of about one second, the A> prompt (or 0A> if you have included the special feature that patches the CCP).

If you see both sign-on messages but do not get an A> prompt, a likely cause of the problem is in the disk drivers. Alternatively, the directory area on the disk may be full of random data rather than 0E5H's.

If you cannot see what is wrong with the system, you might try faking the disk drivers to return a 128-byte block of 0E5H's for each read operation. The CCP should then sign on.

Once you do have the A> prompt, you can proceed with the system checkout. Start by checking that the warm boot logic works. Type a CONTROL-C. There should be a slight pause, and the A> prompt should be output again.

Next, check that you can read the disk directory by using the DIR command. If you have an empty directory, you should get a NO FILE response. If you get strange characters instead, you either forgot to initialize the directory area or the disk parameter block is directing CP/M to the wrong part of the disk for the file directory. If the system crashes, there is a problem with the disk driver.

Check that you can write on the disk by entering the command SAVE 1 TEST. Then use the DIR command to confirm that file TEST shows up in the file directory. If it does, use the ERA command ERA TEST and do another DIR command to confirm that TEST has indeed been erased.

If TEST either does not show up on the disk or cannot be erased, then you have a problem with the disk driver WRITE routine.

Put a standard CP/M release diskette into drive B and use the DIR command to check that you can access the drive and display a disk directory. If you do, then load the DDT utility and exit from it by using a G0 (G, zero) command. This further tests if the disk drivers are functioning correctly.

To test the deblocking logic (if you are using disks that require deblocking), use the command:

```
PIP A:=B:*.*[V]
```

This copies all files from drive B to drive A using the verify option. It is a particularly good test of the system, and if you have any problems with the high-level disk drivers and deblocking code, you will get a Verify Error message from PIP. You can also get this message if you have hardware problems with the computer's memory, so run a memory test if you cannot find anything obviously wrong with the deblocking algorithm.

To completely test the deblocking code, you need to use PIP to copy a file of text larger than the amount of memory available. Thus, you may have to create a large text file using a text editor just to provide PIP with test data.

With the disk driver functioning correctly, rebuild the system with the real time clock enabled. Bring up the new system and check that the ASCII time of day is being updated in the configuration block; use DDT to inspect this in memory. Set the clock to the current time, let it run for five minutes, and see if it is still accurate. You may have to adjust one of the initialization time constants for the device that is providing the periodic interrupts for the clock.

Rebuild the system yet again, this time with the real interrupt-driven console input and the real console output routines. Check that the system comes up properly and that the initial forced-input startup string appears on the console.

Check that when you type characters on the keyboard they are displayed as you type them. If not, there could be a problem with either the CONIN or CONOUT routines. Experimentally type in enough characters to fill the input buffer. If the terminal's bell starts to sound, the interrupt service routine is probably not the culprit. Check the CONOUT routine again.

Check that the function key decode logic is working correctly. With the A> prompt displayed, press a function key. The CONIN driver should inject the correct function key string and it should appear on the terminal. For example, with the BIOS in Figure 8-10, pressing PF1 on the VT-100 terminal should produce this on the display:

```
A>Function Key1
Function?
A>
```

The CCP does not recognize "Function" as a legitimate command name, nor is there such a COM file — hence the question mark.

Using DDT, write a small program that outputs ESCAPE, "t" to the console, and check that the ASCII time of day string appears on the console. This checks that the escape sequence has been recognized.

11

Additional Utility Programs

This chapter contains the narrated source code for several useful utility programs. Two groups of such programs are included—those that supplement Digital Research's standard utility programs, and those that work in conjunction with features shown in the enhanced BIOS (Figure 8-10).

To avoid unnecessary detail, the programs shown in this chapter are all written in the C language. C is a good language to use for such purposes since it can show the overall logic of a program without the clutter of details common in assembly language.

In order to reuse as much source code as possible, this chapter includes a "library" of all the general-purpose C functions that can be called from within any of the utility programs. This file, called "LIBRARY.C", is shown in Figure 11-1. Once a utility program has been compiled, the necessary functions from the library can be linked with the utility's binary output to form the ".COM" file.

```
/* Library of commonly-used functions */

#include <LIBRARY.H>     /*  Standard defines and structures */

/*       Configuration block access       */

/*==============================================================*/
char
*get_cba(code)          /* Get configuration block address */
/*==============================================================*/
/* This function makes a call to a "private" entry in the BIOS
   jump vector to return the address of a specific data object in
   the BIOS. The code indicates which object is required.
   Each program using this function could make a direct call to
   the BIOS using the biosh() function provided by BDS C. This
   function provides a common point to which debugging code can
   be added to display the addresses returned. */

/* Entry parameters */                                                  a
int code;       /* Code that specifies the object
                   whose address is required */
/* Exit Parameters
   Address returned by the BIOS routine */

{
char *retval;           /* Value returned by the BIOS */

        retval = biosh(CBGADDR,code);
 /* printf("\nget_cba : code %d address %4x",code,retval); */
        return retval;
} /* End of get_cba(code) */



/*       Character manipulation functions        */


/*==============================================================*/
strscn(string,key)              /* String scan */
/*==============================================================*/
/* This function scans a 00-terminated character string looking
   for a key string in it. If the key string is found within the
   string, the function returns a pointer to it. Otherwise it
   returns a value of zero. */
/* Entry parameters */
char *string;           /* String to be searched */
char *key;              /* Key string to be searched for */

/* Exit parameters
   Pointer to key string within searched string, or
   zero if key not found
*/                                                                      b

{
while (*string)         /* For all non-null chars. in string */
        {
        if ((*string == *key) &&        /* First char. matches */
            (sstrcmp(string,key) == 0)  /* Perform substring
                                           compare on rest */
           )
                return string;          /* Substring matches,
                                           return pointer */
        string++;                       /* Move to next char. in string */
        }
return 0;                               /* Indicate no match found */
} /* End of strscn */


/*==============================================================*/
ustrcmp(string1,string2)        /* Uppercase string compare */
/*==============================================================*/
/* This function is similar to the normal strcmp function;              c
   it differs only in that the characters are compared as if they
   were all uppercase characters -- the strings are left
   unaltered. */
```

Figure 11-1. LIBRARY.C, commonly used functions, in C language

```
/* Entry Parameters */
char *string1;          /* Pointer to first string */
char *string2;          /* Pointer to second string */

/* Exit parameters
   0 - if string 1 = string 2
   -ve integer if string 1 > string 2
   +ve integer if string 1 < string 2
*/

{
int count;              /* Used to access chars. in both strings */

count = 0;              /* Start with the first character of both */

        /* While string 1 characters are non-null, and
           match their counterparts in string 2. */
while (string1[count] == string2[count])
        {
        if (string1[++count] == '\0')   /* Last char. in string 1 */
                return 0;               /* Indicate equality */
        }
return string2[count] - string1[count]; /* "Compare" chars. */

} /* End of sstrcmp */                                              c


/*============================================================*/
sstrcmp(string,substring)               /* Substring compare */
/*============================================================*/
/* This function compares two strings. The first, string, need not
   be 00-terminated. The second, substring, must be 00-terminated.
   It is similar to the standard function strcmp, except that the
   length of the substring controls how many characters are compared. */

/* Entry parameters */
char *string;           /* Pointer to main string */
char *substring;        /* Pointer to substring */

/* Exit parameters
   0 - substring matches corresponding characters in string
   -ve integer if char. in string is > char. in substring
   +ve integer if char. in string is < char. in substring
*/

{
int count;      /* Used to access chars. in string and substring */

count = 0;      /* Start with the first character of each */

        /* While substring characters are non-null, and
           match their counterparts in string. */
while (string[count] == substring[count])
        {
        if (substring[++count] == '\0') /* Last char in substring */
                return 0;               /* Indicate equality */
        }
return substring[count] - string[count];        /* "Compare" chars. */

} /* End of sstrcmp */                                              d


/*============================================================*/
usstrcmp(string,substring)      /* Uppercase substring compare */
/*============================================================*/
/* This function compares two strings. The first, string, need not
   be 00-terminated. The second, substring, must be 00-terminated.
   It is similar to the substring compare above except all
   characters are made uppercase. */

/* Entry parameters */
char *string;           /* Pointer to main string */
char *substring;        /* Pointer to substring */

/* Exit parameters
   0 -- substring matches corresponding characters in string      e
```

Figure 11-1. (Continued)

```
   -ve integer if char. in string is > char. in substring
   +ve integer if char. in string is < char. in substring
*/

{
int count;      /* Used to access chars in string and substring */

count = 0;      /* Start with the first character of each */

        /* While substring characters are non-null, and
           match their counterparts in string. */
while (toupper(string[count]) == toupper(substring[count]))
        {
        if (substring[++count] == '\0') /* Last char. in substring */
                return 0;               /* Indicate equality */
        }
return substring[count] - string[count];        /* "Compare" chars. */
} /* End of usstrcmp */                                              e


/*============================================================*/
comp_fname(scb,name)            /* Compare file names */
/*============================================================*/
/* This function compares a possibly ambiguous file name
   to the name in the specified character string. The number of
   bytes compared is determined by the number of characters in
   the mask.
   This function can be used to compare file names and types,
   or, by appending an extra byte to the mask, the file names,
   types, and extent numbers.
   For file directory entries, an extra byte can be prefixed to
   the mask and the function used to compare user number, file
   name, type, and extent.
   Note that a "?" in the first character of the mask will NOT
   match with a value of 0xE5 (this value is used to indicate
   an inactive directory entry). */

/* Entry parameters */
struct _scb *scb;       /* Pointer to search control block */
char *name;             /* Pointer to file name */

/* Exit parameter
   NAME_EQ if the names match the mask
   NAME_LT if the name is less than the mask
   NAME_GT if the name is greater than the mask
   NAME_NE if the name is not equal to the mask (but the outcome
        is ambiguous because of the wildcards in the mask)
*/                                                                   f

{
int count;              /* Count of the number of chars. processed */
short ambiguous;        /* NZ when the mask is ambiguous */
char *mask;             /* Pointer to bytes at front of SCB */

/* Set pointer to characters at beginning of search control block */
mask = scb;

        /* Ambiguous match on user number, matches
           only users 0 - 15, and not inactive entries */
if (mask[0] == '?')
        {
        if (name[0] == 0xE5)
                return NAME_NE; /* Indicate inequality */
        }
else    /* First char. of mask is not "?" */
        {
        if (mask[0] != name[0]) /* User numbers do not match */
                return NAME_NE; /* Indicate inequality */
        }

/* No, check the name (and, if the length is such, the extent) */
for (count = 1;                 /* Start with first name character */
     count <= scb -> scb_length; /* For all required characters */
     count++)                   /* Move to next character */
        {
        if (mask[count] == '?') /* Wildcard character in mask */
```

Figure 11-1. (Continued)

```
                {
                ambiguous = 1;  /* Indicate ambiguous name in mask */
                continue;       /* Do not make any comparisons */
                }
        if (mask[count] != (name[count] & 0x7F))
                {               /* Mask char. not equal to FCB char. */
                if (ambiguous)  /* If previous wildcard, indicate NE */
                        return NAME_NE;
                else
                        /* Compare chars. to determine relationship */
                        return (mask[count] > name[count] ?
                                NAME_LT : NAME_GT);
                }
        }
        /* If control reaches here, then all characters of the
        mask and name have been processed, and either there
        were wildcards in the mask, or they all matched. */
return NAME_EQ;         /* Indicate mask and name are "equal" */

} /* End of comp_fname */
```

f

```
/*==============================================================*/
conv_fname(fcb,fn)              /* Convert file name for output */
/*==============================================================*/
/* This function converts the contents of a file control
   block into a printable string "D:FILENAME.TYP." */

/* Entry parameters */
struct _fcb *fcb;               /* Pointer to file control block */
char *fn;                       /* Pointer to area to receive name */

{
        /* If the disk specification in the
           FCB is 0, use the current disk */
*fn++ = (fcb -> fcb_disk) ? (fcb -> fcb_disk + ('A'-1)) :
        (bdos(GETDISK) + 'A');

*fn++ = ':';                            /* Insert disk id. delimiter */

movmem(&fcb -> fcb_fname,fn,8);         /* Move file name  */
fn += 8;                                /* Update pointer */
*fn++ = '.';                            /* Insert file name/type delimiter */
movmem(&fcb -> fcb_fname+8,fn,3);       /* Move file type */
*fn++ &= 0x7F;                          /* Remove any attribute bits */
*fn++ &= 0x7F;                          /* Remove any attribute bits */
*fn++ &= 0x7F;                          /* Remove any attribute bits */
*fn = '\0';                             /* Terminator */

} /*  End of conv_fname  */
```

g

```
/*==============================================================*/
conv_dfname(disk,dir,fn)        /* Convert directory file name for output */
/*==============================================================*/
/* This function converts the contents of a file directory entry
   block into a printable string "D:FILENAME.TYP," */

/* Entry parameters */
short disk;                     /* Disk id. (A = 0, B = 1) */
struct _dir *dir;               /* Pointer to file control block */
char *fn;                       /* Pointer to area to receive name */

{
        /* Convert user number and disk id. */
sprintf(fn,"%2d/%c:",dir -> de_userno,disk + 'A');
fn += 5;                        /* Update pointer to file name */

movmem(&dir -> de_fname,fn,8);  /* Move file name  */
fn += 8;                        /* Update pointer */
*fn++ = '.';                    /* Insert file name/type delimiter */

movmem(&dir -> de_fname+8,fn,3); /* Move file type */
*fn++ &= 0x7F;                  /* Remove any attribute bits */
*fn++ &= 0x7F;                  /* Remove any attribute bits */
*fn++ &= 0x7F;                  /* Remove any attribute bits */
*fn = '\0';                     /* Terminator */
```

h

Figure 11-1. (Continued)

```
} /*  End of conv_dfname  */                                         h


/*============================================================*/
get_nfn(amb_fname,next_fname)   /* Get next file name */
/*============================================================*/
/* This function sets the FCB at "next_fname" to contain the
   directory entry found that matches the ambiguous file name
   in "amb_fname."
   On the first entry for a given file name, the most significant
   bit in the FCB's disk field must be set to one (this causes a
   search first BDOS call to be made). */

/* Entry parameters */
struct _fcb *amb_fname; /* Ambiguous file name */
struct _fcb *next_fname;/* First byte must have ms bit set for
                           first time entry)*/

/* Exit parameters
   0 = No further name found
   1 = Further name found (and set up in next_fname)
*/

{
char bdos_func;          /* Set to either search first or next */
char *pfname;            /* Pointer to file name in directory entry */

        /* Initialize tail-end of next file FCB to zero */
setmem(&next_fname -> fcb_extent,FCBSIZE-12,0);

bdos_func = SEARCHF;     /* Assume a search first must be given */    i

if (!(next_fname -> fcb_disk & 0x80))   /* If not first time */
        {
                /* search first on previous name */
        srch_file(next_fname,SEARCHF);
        bdos_func = SEARCHN;            /* Then do a search next */
        }
else    /* First time */
        next_fname -> fcb_disk &= 0x7F; /* Reset first-time flag */

        /* Refresh next_fname from ambiguous file name
           (move disk, name, type) */
movmem(amb_fname,next_fname,12);

        /* If first time, issue search first, otherwise
           issue a search next call. "srch_file" returns
           a pointer to the directory entry that matches
           the ambiguous file name, or 0 if no match */
if (!(pfname = srch_file(next_fname,bdos_func)) )
        {
        return 0;        /* Indicate no match */
        }
        /* Move file name and type */
movmem(pfname,&next_fname -> fcb_fname,11);
return 1;                /* Indicate match found */

} /*  End of get_nfn  */


/*============================================================*/
char *srch_file(fcb,bdos_code)  /* Search for file */
/*============================================================*/
/* This function issues either a search first or search next
   BDOS call. */

/* Entry Parameters */
struct _fcb *fcb;        /* pointer to file control block */          j
short bdos_code;         /* either SEARCHF or SEARCHN */

/* Exit parameters
   0 = no match found
   NZ = pointer to entry matched (currently in buffer)
*/
```

Figure 11-1. (Continued)

```
{
unsigned r_code;         /* Return code from search function
                            This is either 255 for no match, or 0, 1, 2, or 3
                            being the ordinal of the 32-byte entry in the
                            buffer that matched the name  */
char *dir_entry;         /* Pointer to directory entry */

        /* The BDOS C compiler always sets the BDOS DMA
           to location 0x80 */

r_code = bdos(bdos_code,fcb);   /* Issue the BDOS call */
if (r_code == 255)              /* No match found */                          j
        return 0;

        /* Set a pointer to the matching
           entry by multiplying return code by 128
           and adding onto the buffer address (0x80),
           also add 1 to point to first character of name */

return (r_code << 5) + 0x81;

}/* End of srch_file */


/*==============================================================*/
rd_disk(drb)             /* Read disk (via BIOS) */
/*==============================================================*/
/* This function uses the parameters previously set up in the
   incoming request block, and, using the BIOS directly,
   executes the disk read. */

/* Entry parameters */
struct _drb *drb;        /* Disk request block (disk, track, sector, buffer) */

/* Exit parameters
   0 = No data available
   1 = Data available                                                         k
*/

{
if (!set_disk(drb))      /* Call SELDSK, SETTRK, SETSEC */
        return 0;        /* If SELDSK fails, indicate
                            no data available */
if (bios(DREAD))         /* Execute BIOS read */
        return 0;        /* Indicate no data available if error returned */

return 1;                /* Indicate data available */

} /*  End of rd_disk  */


/*==============================================================*/
wrt_disk(drb)            /* Write disk (via BIOS) */
/*==============================================================*/
/* This function uses the parameters previously set up in the
   incoming request block, and, using the BIOS directly,
   executes the disk write. */

/* Entry parameters */
struct _drb *drb;        /* Disk request block (disk, track, sector, buffer) */

/* Exit parameters
   0 = Error during write
   1 = Data written OK                                                        l
*/

{
if (!set_disk(drb))      /* Call SELDSK, SETTRK, SETSEC, SETDMA */
        return 0;        /* If SELDSK fails, indicate no data written */
if (bios(DWRITE))        /* Execute BIOS write */
        return 0;        /* Indicate error returned */

return 1;                /* Indicate data written */

} /*  End of wrt_disk  */
```

Figure 11-1. (Continued)

```
/*==============================================================*/
short set_disk(drb)     /* Set disk parameters */
/*==============================================================*/
/* This function sets up the BIOS variables in anticipation of
   a subsequent disk read or write. */

/* Entry parameters */
struct _drb *drb;       /* Disk request block (disk, track, sector, buffer) */

/* Exit parameters
   0 = Invalid disk (do not perform read/write)
   1 = BIOS now set up for read/write
*/

{
        /* The sector in the disk request block contains a
           LOGICAL sector. If necessary (as determined by the
           value in the disk parameter header), this must be
           converted into the PHYSICAL sector.
           NOTE: skewtab is declared as a pointer to a pointer to
           a short integer (single byte). */
short **skewtab;        /* Skewtab -> disk parameter header -> skew table */
short phy_sec;          /* Physical sector */
                                                                          m
        /* Call the SELDSK BIOS entry point. If this returns
           a 0, then the disk is invalid. Otherwise, it returns
           a pointer to the pointer to the skew table */
if ( !(skewtab = biosh(SELDSK,drb -> dr_disk)) )
        return 0;               /*  Invalid disk  */

bios(SETTRK,drb -> dr_track);   /* Set track */

        /* Note that the biosh function puts the sector into
           registers BC, and a pointer to the skew table in
           registers HL. It returns the value in HL on exit
           from the BIOS */
phy_sec = biosh(SECTRN,drb -> dr_sector,*skewtab); /* Get physical sector */
bios(SETSEC,phy_sec);           /* Set sector */
bios(SETDMA,drb -> dr_buffer);  /* Set buffer address */

return 1;               /* Indicate no problems */

} /* End of setp_disk */



/*      Directory Management Functions          */


/*==============================================================*/
get_nde(dir_pb)         /* Get next directory entry */
/*==============================================================*/
/* This function returns a pointer to the next directory entry.
   If the directory has not been opened, it opens it.
   When necessary, the next directory sector is read in.
   If the current sector has been modified and needs to be written back
   onto the disk, this will be done before reading in the next sector. */

/* Entry parameters */
struct _dirpb *dir_pb;          /* Pointer to the disk parameter block */

/* Exit Parameters
   Returns a pointer to the next directory entry in the buffer.
   The directory open and write sector flags in the parameter             n
   block are reset as necessary.
*/

{
if(!dir_pb -> dp_open)          /* Directory not yet opened */
        {
        if (!open_dir(dir_pb))  /* Initialize and open directory */
                {
                err_dir(O_DIR,dir_pb);          /* Report error on open */
                exit();
                }
                /* Deliberately set the directory entry pointer to the end
                   of the buffer to force a read of a directory sector */
```

Figure 11-1. (Continued)

```
        dir_pb -> dp_entry = dir_pb -> dp_buffer + DIR_BSZ;
        dir_pb -> dp_write = 0;         /* Reset write-sector flag */
        }

        /* Update the directory entry pointer to the next entry in
           the buffer. Check if the pointer is now "off the end"
           of the buffer and another sector needs to be read. */
if (++dir_pb -> dp_entry < dir_pb -> dp_buffer + DIR_BSZ)
        {
        return dir_pb -> dp_entry;      /* Return pointer to next entry */
        }

        /* Need to move to next sector and read it in */

        /* Do not check if at end of directory or move to
           the next sector if the directory has just been
           opened (but the opened flag has not yet been set) */
if (!dir_pb -> dp_open)
        dir_pb -> dp_open = 1;  /* Indicate that the directory is now open */
else
        {
        /* Check if the sector currently in the buffer needs to be
           written back out to the disk (having been changed) */
        if (dir_pb -> dp_write)
                {
                dir_pb -> dp_write = 0;         /* Reset the flag */
                if(!rw_dir(W_DIR,dir_pb))       /* Write the directory sector */
                        {
                        err_dir(W_DIR,dir_pb);  /* Report error on writing */
                        exit();
                        }
                }

                /* Count down on number of directory entries left to process,
                   always four 32-byte entries per 128-byte sector */
        dir_pb -> dp_entrem -= 4;

                /* Set directory-end flag true if number of entries now < 0 */
        if (dir_pb -> dp_entrem == 0)           /* now at end of directory */
                {
                dir_pb -> dp_end = 1;           /* Indicate end */
                dir_pb -> dp_open = 0;          /* Indicate directory now closed */
                return 0;                       /* Indicate no more entries */
                }

                /* Update sector (and if need be track and sector) */
        if (++dir_pb -> dp_sector == dir_pb -> dp_sptrk)
                {
                ++dir_pb -> dp_track;           /* Update track */
                dir_pb -> dp_sector = 0;        /* Reset sector */
                }
        }

if(!rw_dir(R_DIR,dir_pb))       /* Read next directory sector */
        {
        err_dir(R_DIR,dir_pb);  /* Report error on reading */
        exit();
        }

        /* Reset directory-entry pointer to first entry in buffer */
return dir_pb -> dp_entry = dir_pb -> dp_buffer;

} /* End of get_nde */
```
n

```
/*==============================================================*/
open_dir(dir_pb)        /* Open directory */
/*==============================================================*/
/* This function "opens" up the file directory
   on a specified disk for subsequent processing
   by rw_dir, next_dir functions. */

/* Entry parameters */
struct _dirpb *dir_pb;  /* Pointer to directory parameter block */
```
o

Figure 11-1. (Continued)

```
/* Exit parameters
   0 = Error, directory not opened
   1 = Directory open for processing
*/

{
struct _dpb *dpb;                /* CP/M disk parameter block */

        /* Get disk parameter block address for the disk specified in
           the directory parameter block */
if ((dpb = get_dpb(dir_pb -> dp_disk)) == 0)
        return 0;       /* Return indicating no DPB for this disk */

        /* Set the remaining fields in the parameter block */
dir_pb -> dp_sptrk = dpb -> dpb_sptrk;  /* Sectors per track  */
dir_pb -> dp_track = dpb -> dpb_trkoff; /* Track offset of the directory  */
dir_pb -> dp_sector = 0;                /* Beginning of directory  */
dir_pb -> dp_nument = dpb -> dpb_maxden+1; /*  No. of directory entries  */    o
dir_pb -> dp_entrem = dir_pb -> dp_nument; /*  Entries remaining to process */
dir_pb -> dp_end = 0;                   /* Indicate not at end  */

        /* Set number of allocation blocks per directory entry to
           8 or 16 depending on the number of allocation blocks */
dir_pb -> dp_nabpde = (dpb -> dpb_maxabn > 255 ? 8 : 16);
        /* Set number of allocation blocks (one more than number of
           highest block) */
dir_pb -> dp_nab = dpb -> dpb_maxabn;

        /* Set the allocation block size based on the block shift.
           The possible values are: 3 = 1k, 4 = 2K, 5 = 4K, 6 = 8K, 7 = 16K.
           So a value of 16 is shifted right by (7 - bshift) bits. */
dir_pb -> dp_absize = 16 >> (7 - dpb -> dpb_bshift);

return 1;               /* Indicate that directory now opened */

} /*  End of open_dir  */


/*============================================================*/
rw_dir(read_op,dir_pb)  /* Read/write directory */
/*============================================================*/
/* This function reads/writes the next 128-byte
   sector from/to the currently open directory. */

/* Entry parameters */
short read_op;          /* True to read, false (0) to write */
struct _dirpb *dir_pb;  /* Directory parameter block */

/* Exit parameters
   0 = error -- operation not performed
   1 = operation completed
*/

{
struct _drb drb;                /* Disk request (for BIOS read/write) */
                                                                              p

drb.dr_disk = dir_pb -> dp_disk;        /* Set up disk request */
drb.dr_track = dir_pb -> dp_track;
drb.dr_sector = dir_pb -> dp_sector;
drb.dr_buffer = dir_pb -> dp_buffer;

if (read_op)
        {
        if (!rd_disk(&drb))     /* Issue read command */
                return 0;       /* Indicate error -- no data available */
        }
else
        {
        if (!wrt_disk(&drb))    /* Issue write command */
                return 0;       /* Indicate error -- no data written */
        }
return 1;                       /* Indicate operation complete */

} /* End of rd_dir */
```

Figure 11-1. (Continued)

```
/*==========================================================*/
err_dir(opcode,dir_pb)          /* Display directory error
/*==========================================================*/
/* This function displays an error message to report an error
   detected in the directory management functions open_dir and rw_dir. */
/* Entry parameters */
short opcode;                   /* Operation being attempted */
struct _dirpb *dir_pb;  /* Pointer to directory parameter block */

{
printf("\n\007Error during ");

switch(opcode)
        {
        case R_DIR:                                                         q
                printf("Reading");
                break;
        case W_DIR:
                printf("Writing");
                break;
        case O_DIR:
                printf("Opening");
                break;
        default:
                printf("Unknown Operation (%d) on",opcode);
        }

printf(" Directory on disk %c:. ",dir_pb -> dp_disk + 'A');

} /* End of err_dir */


/*==========================================================*/
setscb(scb,fname,user,extent,length)    /* Set search control block */
/*==========================================================*/
/* This function sets up a search control block according
   to the file name specified. The file name can take the
   following forms:

        filename
        filename.typ
        d:filename.typ
        *:filename.typ (meaning "all disks")
        ABCD...NOP:filename.typ (meaning "just the specified disks")

   The function sets the bit map according to which disks should be
   searched . For each selected disk, it checks to see if an error is
   generated when selecting the disk (i.e. if there are disk tables
   in the BIOS for the disk). */

/* Entry parameters */
struct _scb *scb;           /* Pointer to search control block */
char *fname;                /* Pointer to the file name */
short user;                 /* User number to search for */                 r
short extent;               /* Extent number to search for */
int length;                 /* Number of bytes to compare */

/* Exit parameters
   None.
*/

{
int disk;                   /* Disk number currently being checked */
unsigned adisks;            /* Bit map for active disks */

adisks = 0;                 /* Assume no disks to search */

if (strscn(fname,":"))          /* Check if ":" in file name */
        {
        if (*fname == '*')      /* Check if "all disks" */
                {
                adisks = 0xFFFF;        /* Set all bits */
                }
        else                    /* Set specific disks */
                {
                while(*fname != ':')    /* Until ":" reached */
```

Figure 11-1. (Continued)

```
                        {
                        /* Build the bit map by getting the next disk
                           id. (A - P), converting it to a number in
                           the range 0 - 15, shifting a 1-bit left
                           that many places, and OR-ing it into the
                           current active disks. */
                        adisks |= 1 << (toupper(*fname) - 'A');
                        ++fname;            /* Move to next character */
                        }
                ++fname;                    /* Bypass colon */
                }
        }
else    /* Use only current default disk */
        {
                /* Set just the bit corresponding to the current disk */
        adisks = 1 << bdos(GETDISK);
        }

setfcb(scb,fname);      /* Set search control block as though it
                           were a file control block. */

/* Make calls to the BIOS SELDSK routine to make sure that
   all of the active disk drives have disk tables for them
   in the BIOS. If they don't, turn off the corresponding                  r
   bits in the bit map. */

for (disk = 0;          /* Start with disk A: */
     disk < 16;         /* Until disk P: */
     disk++)            /* Use next disk */
        {
        if ( !((1 << disk) & adisks))
                continue;           /* Avoid selecting unspecified disks */
        if (biosh(SELDSK,disk) == 0)    /* Make BIOS SELDSK call */
                {                   /* Returns 0 if invalid disk */
                /* Turn OFF corresponding bit in mask
                   by AND-ing it with bit mask having
                   all the other bits set = 1  */
                adisks &= ((1 << disk) ^ 0xFFFF);
                }
        }

scb -> scb_adisks = adisks;     /* Set bit map in SCB */
scb -> scb_userno = user;       /* Set user number */
scb -> scb_extent = extent;     /* Set extent number */
scb -> scb_length = length;     /* Set number of bytes to compare */

} /* End setscb */


/*==============================================================*/
dm_clr(disk_map)                /* Disk map clear (to zeros) */
/*==============================================================*/
/* This function clears all elements of the disk map to zero. */

/* Entry Parameters */
unsigned disk_map[16][18];      /* Address of array of unsigned integers */

/* Exit parameters
   None.                                                                   s
*/

{
        /* WARNING -- The 576 in the setmem call below is based on
           the disk map array being [16][18] -- i.e. 288 unsigned
           integers, hence 576 bytes. */
setmem(disk_map,576,'\0');      /* Fill array with zeros */

} /* End of dm_clr */

/*==============================================================*/
dm_disp(disk_map,adisks)                /* Disk map display */
/*==============================================================*/
/* This function displays the elements of the disk map, showing            t
   the count in each element. A zero value-element is shown as
   blanks. For example:
```

Figure 11-1. (Continued)

```
      0   1   2   3   4   5   6   7   8   9  10  11  12  13  14  15  Used Free
A: 123     20  98         202     199 101 211                          954   70

   Lines will only be printed for active disks (as indicated by
   the bit map). */

/* Entry parameters */
unsigned disk_map[16][18];      /* Pointer to disk map array */
unsigned adisks;                /* Bit map of active disks */

{
#define USED_COUNT 16           /* "User" number for used entities */
#define FREE_COUNT 17           /* "User" number for free entities */

int disk;                       /* Current disk number */
int userno;                     /* Current user number */
unsigned dsum;                  /* Sum of entries for given disk */

printf("\n      0   1   2   3   4   5   6   7   8   9  10  11  12  13  14  15  Used Free");

for (disk = 0;                  /* Start with disk A: */
     disk < 16;                 /* Until disk P: */
     disk++)                    /* Next disk */
        {
        if (!(adisks & (1 << disk)))    /* Check if disk is active */
                continue;       /* No -- so bypass this one */

        printf("\n%c: ",disk + 'A');    /* Display disk number */

        dsum = 0;               /* Reset sum for this disk */
        for (userno = 0;        /* Start with user 0 */
             userno < 16;       /* Until user 15 */
             userno++)          /* Next user number */
                {
                dsum += disk_map[disk][userno]; /* Build sum */
                }

        if (dsum)       /* Check if any output for this disk,
                           and if not, display d: None */
                {
                /* Print either number or blanks */
                for (userno = 0;        /* Start with user 0 */
                     userno < 16;       /* Until user 15 */
                     userno++)          /* Next user number */
                        {
                        if (disk_map[disk][userno])
                                printf("%4d",disk_map[disk][userno]);
                        else
                                printf("    ");
                        }
                }
        else            /* No output for this disk */
                {
                printf(  " -- None --
                }
        printf("  %4d %4d",disk_map[disk][USED_COUNT],disk_map[disk][FREE_COUNT]);
        }

} /* End dm_disp */

/*============================================================*/
get_dpb(disk)           /* Get disk parameter block address */
/*============================================================*/
/* This function returns the address of the disk parameter
   block (located in the BIOS). */

/* Entry parameters */
char disk;              /* Logical disk for which DPB address is needed */

/* Exit parameters
        0 = Invalid logical disk
        NZ = Pointer to disk parameter block
*/

{
if (biosh(SELDSK,disk) == 0)            /* Make BIOS SELDSK call */
        return 0;                       /* Invalid disk */
```

Figure 11-1. (Continued)

```
    bdos(SETDISK,disk);                       /* Use BDOS SETDISK function */
    return bdos(GETDPARM);                    /* Get the disk parameter block */      u

    } /*  End of get_dpb  */



    /*       Code table functions     */

    /* Most programs that interact with a user must
       accept parameters from the user by name and translate
       the name into some internal code value.
       They also must be able to work in reverse, examining
       the setting of a variable, and determing what (ASCII
       name) it has been set to.

       An example is setting baud rates. The user may want to
       enter "19200," and have this translated into a number
       to be output to a chip. Alternatively, a previously
       set baud rate variable may have to be examined and the
       string "19200" generated to display its current
       setting to the user.

       A code table is used to make this task easier.
       Each element in the table logically consists of:
            A code value (unsigned integer)
            An ASCII character string (actually a pointer to it) */                   v


    /*============================================================*/
    ct_init(entry,code,string)     /* Initialize code table */
    /*============================================================*/
    /* This function initializes a specific entry in a code table
       with a code value and string pointer.

");    NOTE: By convention, the last entry in a given
       code table will have a code value of CT_SNF (string not found). */

    /* Entry parameters */
    struct _ct *entry;             /* Pointer to code table entry */
    int code;                      /* Code value to store in entry */
    char *string;                  /* Pointer to string for entry */

    /* Exit parameters
       None.
    */

    {
    entry -> _ct_code = code;                 /* Set _ct_code */
    entry -> _ct_sp = string;                 /* Set string pointer */
    } /*  end of ct_inti  */

    /*============================================================*/
    unsigned
    ct_parc(table,string)          /* Parameter - return code */
    /*============================================================*/
    /* This function searches the specified table for a
       matching string, and returns the code value that corresponds to it.
       If only one match is found in the table, then this function returns
       that code value. If no match or more than one match is found,
       it returns the error value, CT_SNF (string not found).
       This function is specifically designed for processing
       parameters on a command tail.
       Note that the comparison is done after conversion to uppercase
       (i.e. "STRING" matches "string"). A substring compare is used so               w
       that only the minimum number of characters for an unambiguous
       response need be entered. For example, if the table contained:

                    Code    Value
                    1       "APPLES"
                    2       "ORANGES"
                    3       "APRICOTS"

       A response of "O" would return code = 2, but "A" or "AP" would
       be ambiguous. "APR" or "APP" would be required. */

    struct _ct *table;             /* Pointer to table */
    char *string;                  /* Pointer to key string */
```

Figure 11-1. (Continued)

```
{
int mcode;                      /* Matched code to return */
int mcount;                     /* Count of number of matches found */

mcode = CT_SNF;                 /* Assume error */
mcount = 0;                     /* Reset match count */

while(table -> _ct_code != CT_SNF) /* Not at end of table */
        {
        /* Compare keyboard response to table entry using
           uppercase substring compare. */
        if (usstrcmp(table -> _ct_sp,string) == 0)
                {
                mcount++;       /* Update match count */
                mcode = table -> _ct_code;      /* Save code */
                }
        table++;                /* Move to next entry */
        }

if (mcount == 1)                /* Only one match found */
        return mcode;           /* Return matched code */
else                            /* Illegal or ambiguous */
        return CT_SNF;

} /* End ct_parc */


/*==============================================================*/
unsigned
ct_code(table,string)   /* Return code for string */
/*==============================================================*/
/* This function searches the specified table for the
   specified string. If a match occurs, it returns the
   corresponding code value. Otherwise it returns CT_SNF
   (string not found).
   Unlike ct_parc, this function compares every character in the
   key string, and will return the code on the first match found. */

/* Entry parameters */
struct _ct *table;      /* Pointer to table */
char *string;           /* Pointer to string */

/* Exit parameters
   Code value -- if string found
   CT_SNF -- if string not found
*/

{
while(table -> _ct_code != CT_SNF)      /* For all entries in table */
        {
        if (ustrcmp(table -> _ct_sp,string) == 0) /* Compare strings */
                return table -> _ct_code;       /* Return code */
        table++;                        /* Move to next entry */
        }
return CT_SNF;                          /* String not found */

} /* End ct_code */


/*==============================================================*/
ct_disps(table) /*   Displays all strings in specified table */
/*==============================================================*/
/* This function displays all of the strings in a given table.
   It is used to indicate valid responses for operator input. */

/* Entry parameters */
struct _ct *table;              /* Pointer to table */

/* Exit Parameters
        None.
*/

{
while(table -> _ct_code != CT_SNF)      /* Not end of table */
        {
        printf("\n\t\t%s",table -> _ct_sp);     /* Print string */
        table++;                        /* Move to next entry */
        }
```

Figure 11-1. (Continued)

```
putchar('\n');                         /* Add final return */          y

} /* End of ct_disps */


/*============================================================*/
ct_index(table,string) /* Returns index for a given string */
/*============================================================*/
/* This function searches the specified table, and returns
   the INDEX of the entry containing a matching string.
   All characters of the string are used for the comparison,
   after they have been made uppercase. */


/* Entry parameters */
struct _ct *table;               /* Pointer to table */
char *string;                    /* Pointer to string */

/* Exit parameters
   Index of entry matching string, or
   CT_SNF if string not found.
*/                                                                     z

{
int index;                       /* Current value of index */

index = 0;                       /* Initialize index */

while(table -> _ct_code != CT_SNF)       /* Not at end of table */

        {
        if (ustrcmp(table -> _ct_sp,string) == 0)
                return index;    /* Return index */
        table++;                 /* Move to next table entry */
        index++;                 /* Update index */
        }
return CT_SNF;          /* String not found */

}


/*============================================================*/
char *ct_stri(table,index)       /* Get string according to index */
/*============================================================*/
/* This function returns a pointer to the string in the
   table entry specified by the index. */

/* Entry parameters */
struct _ct *table;               /* Pointer to table */             aa
int index;                       /* Index into table */

{
struct _ct *entry;               /* Entry pointer */
        entry = table[index];    /* Point to entry */
        return entry -> _ct_sp;  /* Return pointer to string */

} /* End of ct_stri */


/*============================================================*/
char *ct_strc(table,code)  /* Get string according to code value */
/*============================================================*/
/* This function searches the specified table and returns a
   pointer to the character string in the entry with the
   matching code value or a pointer to a string of "unknown"
   if the code value is not found. */

/* Entry parameters */                                               bb
struct _ct *table;               /* Pointer to table */
unsigned  code;                  /* Code value */

{
while(table -> _ct_code != CT_SNF)       /* Until end of table */
        {
        if (table -> _ct_code == code)   /* Check code matches */
                return table -> _ct_sp;  /* Yes, return ptr. to str. */
        table++;                         /* No, move to next entry */
```

Figure 11-1. (Continued)

```
        }
return "Unknown";
}



/*      Bit vector functions    */                                          bb

/* These functions manipulate bit vectors. A bit vector is a group
   of adjacent bits, packed eight per byte. Each bit vector has the
   structure defined in the LIBRARY.H file.

   Bit vectors are used primarily to manipulate the operating
   system's allocation vectors and other values that can best
   be represented as a series of bits. */


/*===========================================================*/
bv_make(bv,bytes)       /* Make a bit vector and clear to zeros */
/*===========================================================*/
/* This function uses C's built-in memory allocation, alloc,
   to allocate the necessary amount of memory, and then
   sets the vector to zero-bits. */

/* Entry parameters */
struct _bv *bv;         /* Pointer to a bit vector */
unsigned bytes;         /* Number of bytes in bit vector */

/* Exit parameter
   NZ = vector created
   0 = insufficient memory to create vector                                 cc
*/

{
if(!(bv -> bv_bits = alloc(bytes)))     /* Request memory */
        return 0;                       /* Request failed */

bv -> bv_bytes = bytes;                 /* Set length */
bv -> bv_end = bv -> bv_bits + bytes;   /* Set pointer to end */

bv_fill(bv,0);                          /* Fill with 0's */
return 1;

} /* End bv_make */


/*===========================================================*/
bv_fill(bv,value)       /* Fill bit vector with value */
/*===========================================================*/
/* This function fills the specified bit vector with the
   specified value.
   This function exist only for consistency's sake and
   to isolate the main body of code from standard
   functions like setmem. */

/* Entry parameters */
struct _bv *bv;         /* Pointer to bit vector */                         dd
char value;             /* Value to fill vector with */

/* Exit parameters
   None.
*/


{
/*      address      length         value */
setmem(bv -> bv_bits,bv -> bv_bytes,value);
}

/*===========================================================*/
bv_set(bv,bitnum)               /* Set the specified bit number */
/*===========================================================*/
/* This function sets the specified bit number in the bit vector
   to one-bit. */                                                          ee

/* Entry parameters */
struct _bv *bv;                 /* Pointer to bit vector */
unsigned bitnum;                /* Bit number to be set */
```

Figure 11-1. (Continued)

```
/* Exit parameters
   None.
*/

{
unsigned byte_offset;          /* Byte offset into the bit vector */

if ((byte_offset = bitnum >> 3) > bv -> bv_bytes)
        return 0;       /* Bitnum is "off the end" of the vector */

/* Set the appropriate bit in the vector. The byte offset
   has already been calculated. The bit number in the byte
   is calculated by AND ing the bit number with 0x07.
   The specified bit is then OR ed into the vector */

bv -> bv_bits[byte_offset] |= (1 << (bitnum & 0x7));

return 1;               /* Indicate completion */

/* End of bv_set  */                                                    ee

/*============================================================*/
bv_test(bv,bitnum)              /* Test the specified bit number */
/*============================================================*/
/* This function returns a value that reflects the current
   setting of the specified bit. */

/* Entry parameters */
struct _bv *bv;                 /* Pointer to bit vector */
unsigned bitnum;                /* Bit number to be set */

/* Exit parameters
   None.
*/

{
unsigned byte_offset;           /* Byte offset into the bit vector */   ff

if ((byte_offset = bitnum >> 3) > bv -> bv_bytes)
        return 0;       /* Bitnum is "off the end" of the vector */

/* Set the appropriate bit in the vector. The byte offset
   has already been calculated. The bit number in the byte
   is calculated by AND ing the bit number with 0x07.
   The specified bit is then OR ed into the vector */

return bv -> bv_bits[byte_offset] & (1 << (bitnum & 0x7));

} /* End of bv_tests */


/*============================================================*/
bv_nz(bv)               /* Test bit vector nonzero */
/*============================================================*/
/* This function tests each byte in the specified vector,
   and returns indicating whether any bits are set in
   the vector. */

/* Entry parameters */
struct _bv *bv;         /* Pointer to bit vector */

/* Exit Parameters
   NZ = one or more bits are set in the vector                          gg
   0 = all bits are off
*/

{
char *bits;             /* Pointer to bits in bit vector */

bits = bv -> bv_bits;           /* Set working pointer */

while (bits != bv -> bv_end)    /* For entire bit vector */
        {
        if (*bits++)            /* If nonzero */
                return bits--;  /* Return pointer to NZ byte */
```

Figure 11-1. (Continued)

```
        }
return 0;                       /* Indicate vector is zero */

} /* End of by_nz */


/*==========================================================*/
bv_and(bv3,bv1,bv2)             /* bv3 = bv1 & bv2 */
/*==========================================================*/
/* This function performs a boolean AND between the bytes
   of bit vector 1 and 2, storing the result in bit vector 3. */

/* Entry parameters */
struct _bv *bv1;                /* Pointer to input bit vector */
struct _bv *bv2;                /* Pointer to input bit vector */

/* Exit parameters */
struct _bv *bv3;                /* Pointer to output bit vector */

{
char *bits1, *bits2, *bits3;    /* Working pointers to bit vectors */

bits1 = bv1 -> bv_bits;         /* Initialize working pointers */
bits2 = bv2 -> bv_bits;
bits3 = bv3 -> bv_bits;

        /* AND ing will proceed until the end of any one of the bit
           vectors is reached  */
while (bits1 != bv1 -> bv_end &&
       bits2 != bv2 -> bv_end &&
       bits3 != bv3 -> bv_end)
        {
                *bits3++ = *bits1++ & *bits2++; /* bv3 = bv1 & bv2 */
        }
} /* End of bv_and */


/*==========================================================*/
bv_or(bv3,bv1,bv2)              /* bv3 = bv1 or bv2 */
/*==========================================================*/
/* This function performs a boolean inclusive OR between the bytes
   of bit vectors 1 and 2, storing the result in bit vector 3. */

/* Entry parameters */
struct _bv *bv1;                /* Pointer to input bit vector */
struct _bv *bv2;                /* Pointer to input bit vector */

/* Exit parameters */
struct _bv *bv3;                /* Pointer to output bit vector */

{
char *bits1, *bits2, *bits3;    /* Working pointers to bit vectors */

bits1 = bv1 -> bv_bits;         /* Initialize working pointers */
bits2 = bv2 -> bv_bits;
bits3 = bv3 -> bv_bits;

        /* The OR ing will proceed until the end of any one of the bit
           vectors is reached. */
while (bits1 != bv1 -> bv_end &&
       bits2 != bv2 -> bv_end &&
       bits3 != bv3 -> bv_end)
        {
                *bits3++ = *bits1++ | *bits2++; /* bv3 = bv1 or bv2 */
        }
} /* End of bv_or */


/*==========================================================*/
bv_disp(title,bv)               /* Bit vector display */
/*==========================================================*/
/* This function displays the contents of the specified bit vector
   in hexadecimal. It is normally only used for debugging. */

/* Entry parameters */
char *title;                    /* Title for the display */
struct _bv *bv;                 /* Pointer to the bit vector */
```

gg

hh

ii

jj

Figure 11-1. (Continued)

```
/* Exit parameters
   None.
*/

{
char *bits;                        /* Working pointer */
unsigned byte_count;               /* Count used for formatting display */
unsigned bit_count;                /* Count for processing bits in a byte */
char byte_value;                          /* Value to be displayed */

printf("\nBit Vector : %s",title);        /* Display title */

bits = bv -> bv_bits;                     /* Set working pointer */
byte_count = 0;                           /* Initialize count */

while (bits != bv -> bv_end)              /* For the entire vector */
        {
        if (byte_count % 5 == 0)          /* Check if new line */
                                          /* Display bit number */          jj
                printf("\n%4d : ",byte_count << 3);

        byte_value = *bits++;    /* Get the next byte from the vector */

        for (bit_count = 0; bit_count < 8; bit_count++)
                {
                /* Display the leftmost bit, then shift the value
                   left one bit */
                if (bit_count == 4) putchar(' '); /* Separator */
                putchar((byte_value & 0x80) ? '1' : '0');
                byte_value <<= 1;       /* Shift value left */
                }
        printf("  ");                                   /* Separator */

        byte_count++;    /* Update byte count */
        }
} /* End of bv_disp */

/* End of LIBRARY.C */
```

Figure 11-1. (Continued)

Associated with the library of functions is another section of source code called "LIBRARY.H", shown in Figure 11-2. This "header" file must be included at the beginning of each program that calls any of the library functions.

For reasons of clarity, this chapter describes the simplest functions first, followed by the more complex, and finally by the utility programs that use the functions.

Several functions in the library and some definitions in the library header are not used by the utilities shown in this chapter. They have been included to illustrate techniques and because they might be useful in other utilities you could write.

```
#define LIBVN "1.0"     /* Library version number */

/* This file contains groups of useful definitions.
   It should be included at the beginning of any program
   that uses the functions in LIBRARY.C */


/* Definition to make minor language modification to C. */
#define short char              /* Short is not supported directly */     a
```

Figure 11-2. LIBRARY.H, code to be included at the beginning of any program that calls LIBRARY functions in Figure 11-1

```
/* One of the functions (bv_make) in the library uses the BDS C
   function, alloc, to allocate memory. The following definitions
   are provided for alloc. */

struct _header                  /* Header for block of memory allocated */
        {
        struct _header *_ptr;   /* Pointer to the next header in the chain */
        unsigned _size;         /* Number of bytes in the allocated block */
        };
struct _header _base;           /* Declare the first header of the chain */
struct _header *_allocp;        /* Used by alloc() and free() functions */
```
b

```
/* BDOS function call numbers  */

#define SETDISK 14      /* Set (select) disk */
#define SEARCHF 17      /* Search first */
#define SEARCHN 18      /* Search next */
#define DELETEF 19      /* Delete file */
#define GETDISK 25      /* Get default disk (currently logged in) */
#define SETDMA  26      /* Set DMA (Read/Write) Address */
#define GETDPARM 31     /* Get disk parameter block address */
#define GETUSER 32      /* Get current user number */
#define SETUSER 32      /* Set current user number */
```
c

```
/* Direct BIOS calls
   These definitions are for direct calls to the BIOS.
   WARNING:  Using these makes program less transportable.
   Each symbol is related to its corresponding jump in the
   BIOS jump vector.
   Only the more useful entries are defined. */

#define CONST   2       /* Console status */
#define CONIN   3       /* Console input */
#define CONOUT  4       /* Console output */
#define LIST    5       /* List output */
#define AUXOUT  6       /* Auxiliary output */
#define AUXIN   7       /* Auxiliary input */

#define HOME    8       /* Home disk */
#define SELDSK  9       /* Select logical disk */
#define SETTRK  10      /* Set track */
#define SETSEC  11      /* Set sector */
#define SETDMA  12      /* Set DMA address */
#define DREAD   13      /* Disk read */
#define DWRITE  14      /* Disk write */
#define LISTST  15      /* List status */
#define SECTRN  16      /* Sector translate */
#define AUXIST  17      /* Auxiliary input status */
#define AUXOST  18      /* Auxiliary output status */

                        /* "Private" entries in jump vector */
#define CIOINIT 19      /* Specific character I/O initialization */
#define SETDOG  20      /* Set watchdog timer */
#define CBGADDR 21      /* Configuration block, get address */
```
d

```
/* Definitions for accessing the configuration block */

#define CB_GET 21               /* BIOS jump number to access routine */
#define DEV_INIT 19             /* BIOS jump to initialize device */

#define CB_DATE 0               /* Date in ASCII */
#define CB_TIMEA 1              /* Time in ASCII */
#define CB_DTFLAGS 2            /* Date, time flags */
#define TIME_SET 0x01           /* This bit NZ means date has been set */
#define DATE_SET 0x02           /* This bit NZ means time has been set */

#define CB_FIP 3                /* Forced input pointer */
#define CB_SUM 4                /* System start-up message */

#define CB_CI 5                 /* Console input */
#define CB_CO 6                 /* Console output */
#define CB_AI 7                 /* Auxiliary input */
#define CB_AO 8                 /* Auxiliary output */
```
e

Figure 11-2. (Continued)

```
#define CB_LI 9                 /* List input */
#define CB_LO 10                /* List output */

#define CB_DTA 11               /* Device table addresses */
#define CB_C1224 12             /* Clock 12/24 format flag */
#define CB_RTCTR 13             /* Real time clock tick rate (per second) */

#define CB_WDC 14               /* Watchdog count */
#define CB_WDA 15               /* Watchdog address */

#define CB_FKT 16               /* Function key table */
#define CB_COET 17              /* Console output escape table */

#define CB_D0_IS 18             /* Device 0 initialization stream */
#define CB_D0_BRC 19            /* Device 0 baud rate constant */          e

#define CB_D1_IS 20             /* Device 1 initialization stream */
#define CB_D1_BRC 21            /* Device 1 baud rate constant */

#define CB_D2_IS 22             /* Device 2 initialization stream */
#define CB_D2_BRC 23            /* Device 2 baud rate constant */

#define CB_IV 24                /* Interrupt vector */
#define CB_LTCBO 25             /* Long term config. block offset */
#define CB_LTCBL 26             /* Long term config. block length */

#define CB_PUBF 27              /* Public files flag */
#define CB_MCBUF 28             /* Multi-command buffer */
#define CB_POLLC 29             /* Polled console flag */


        /* Device numbers and names for physical devices */
        /* NOTE: Change these definitions for your computer system */

#define T_DEVN 0                /* Terminal */                             f
#define M_DEVN 1                /* Modem */
#define P_DEVN 2                /* Printer */

#define MAXPDEV 2               /* Maximum physical device number */

        /* Names for the physical devices */

#define PN_T "TERMINAL"                                                    g
#define PN_M "MODEM"
#define PN_P "PRINTER"

        /* Structure and definitions for function keys */

#define FK_ILENGTH 2            /* No. of chars. input when func. key pressed
                                   NOTE: This does NOT include the ESCAPE. */
#define FK_LENGTH 16            /* Length of string  (not including fk_term) */
#define FK_ENTRIES 18           /* Number of function key entries in table */   h

struct _fkt                     /* Function key table */
        {
        char fk_input[FK_ILENGTH];      /* Lead-in character is not in table */
        char fk_output[FK_LENGTH];      /* Output character string */
        char fk_term;                   /* Safety terminating character */
        };


/* Definitions and structure for device tables */

        /* Protocol bits */
        /* Note: if the most significant bit is
           set = 1, then the set_proto function
           will logically OR in the value. This
           permits Input DTR to co-exist with
           XON or ETX protocol. */                                         i

#define DT_ODTR 0x8004          /* Output DTR high to send (OR ed in) */
#define DT_OXON 0x0008          /* Output XON */
#define DT_OETX 0x0010          /* Output ETX/ACK */

#define DT_IRTS 0x8040          /* Input RTS (OR-ed in) */
#define DT_IXON 0x0080          /* Input XON */
```

Figure 11-2. (Continued)

```
#define ALLPROTO 0xDC          /* All protocols combined */

struct _dt                     /* Device table */
        {
        char dt_f1[14];        /* Filler */
        char dt_st1;           /* Status byte 1 -- has protocol flags */          i
        char dt_st2;           /* Status byte 2 */
        unsigned dt_f2;        /* Filler */
        unsigned dt_etxml;     /* ETX/ACK message length */
        char dt_f3[12];        /* Filler */
        } ;


/* Values returned by the comp_fname (compare file name) */

#define NAME_EQ 0       /* Names equal */
#define NAME_LT 1       /* Name less than mask */                                 j
#define NAME_GT 2       /* Name greater than mask */
#define NAME_NE 3       /* Name not equal (and comparison ambiguous) */


/* Structure for standard CP/M file control block */

#define FCBSIZE 36             /* Define the overall length of an FCB */

struct _fcb
        {
        short fcb_disk;        /* Logical disk (0 = default) */
        char fcb_fname[11];    /* File name, type (with attributes) */
        short fcb_extent;      /* Current extent */
        unsigned fcb_s12;      /* Reserved for CP/M */                            k
        short fcb_reccnt;      /* Record count used in current extent */
        union                  /* Allocation blocks can be either */
                {              /* Single or double bytes */
                short fcbab_short[16];
                unsigned fcbab_long[8];
                } _fcbab;
        short fcb_currec;      /* Current record within extent */
        char fcb_ranrec[3];    /* Record for random read/write */
        };

/*  Parameter block used for calls to the directory management routines  */

#define DIR_BSZ 128            /* Directory buffer size */

struct _dirpb
        {
        short dp_open;         /* 0 to request directory to be opened */
        short dp_end;          /* NZ when at end of directory */
        short dp_write;        /* NZ to write current sector to disk */
        struct _dir *dp_entry; /* Pointer to directory entry in buffer */         l
        char dp_buffer [DIR_BSZ];      /* Directory sector buffer */
        char dp_disk;          /* Current logical disk */
        int dp_track;          /* Start track */
        int dp_sector;         /* Start sector */
        int dp_nument;         /* Number of directory entries */
        int dp_entrem;         /* Entries remaining to process */
        int dp_sptrk;          /* Number of sectors per track */
        int dp_nabpde;         /* Number of allocation blocks per dir. entry */
        unsigned dp_nab;       /* Number of allocation blocks */
        int dp_absize;         /* Allocation block size (in Kbytes) */
        };


/* The err_dir function is used to report errors found by the
   directory management routines, open_dir and rw_dir.
   Err_dir needs a parameter to define the operation being
   performed when the error occurred. The following definitions
   represent the operations possible. */                                         m

#define W_DIR   0       /* Writing directory */
#define R_DIR   1       /* Reading directory */
#define O_DIR   2       /* Opening directory */
```

Figure 11-2. (Continued)

```
/*  Disk parameter block maintained by CPM  */                              n

struct _dpb
        {
        unsigned dpb_sptrk;      /* Sectors per track */
        short dpb_bshift;        /* Block shift */
        short dpb_bmask;         /* Block mask */
        short dpb_emask;         /* Extent mask */
        unsigned dpb_maxabn;     /* Maximum allocation block number */
        unsigned dpb_maxden;     /* Maximum directory entry number */
        short dpb_rab0;          /* Allocation blocks reserved for */
        short dpb_rab1;          /*   directory blocks */
        unsigned dpb_diskca;     /* Disk changed workarea */
        unsigned dpb_trkoff;     /* Track offset */
        };


/*  Disk directory entry format  */                                         o

struct _dir {
        char de_userno;          /* User number or 0xE5 if free entry */
        char de_fname[11];       /* File name [8] and type [3] */
        int de_extent;           /* Extent number of this entry */
        int de_reccnt;           /* Number of 128-byte records used in last
                                      allocation block */
        union                    /* Allocation blocks can be either */
                {                /*   single or double bytes */
                short de_short[16];
                unsigned de_long[8];
                } _dirab;
        };


/*  Disk request parameters for BIOS-level read/writes */                   p

struct _drb
        {
        short dr_disk;           /* Logical disk A = 0, B = 1... */
        unsigned dr_track;       /* Track (for SETTRK) */
        unsigned dr_sector;      /* Sector (for SETSEC) */
        char *dr_buffer;         /* Buffer address (for SETDMA) */
        } ;


/* Search control block used by directory scanning functions */             q

struct _scb
        {
        short scb_userno;        /* User number(s) to match */
        char scb_fname[11];      /* File name and type */
        short scb_extent;        /* Extent number */
        char unused[19];         /* Dummy bytes to make this look like
                                      a file control block */
        short scb_length;        /* Number of bytes to compare */
        short scb_disk;          /* Current disk to be searched */
        unsigned scb_adisks;     /* Bit map of disks to be searched.
                                      the rightmost bit is for disk A:. */
        } ;


/* Code table related definitions */                                        r

#define CT_SNF 0xFFFF   /* String not found */

struct _ct              /* Define structure of code table */
        {
        unsigned _ct_code;       /* Code value */
        char *_ct_sp;            /* String pointer */
        };
```

Figure 11-2. (Continued)

```
/* Structure for bitvectors */

struct _bv
        {
        unsigned bv_bytes;          /* Number of bytes in the vector */
        char *bv_bits;              /* Pointer to the first byte in the vector */        s
        char *bv_end;               /* Pointer to byte following bit vector */
        } ;


/* End of LIBRARY.H */
```

Figure 11-2. (Continued)

Library Functions

This section describes the library functions and the sections from the header file that must be included at the beginning of each utility program.

A Minor Change to C Language

One minor problem with the BDS C Compiler is that it does not support "short" integers, or integers that are only a single byte long. It is convenient to declare certain values as short to serve as a reminder of the standard type definition. Therefore, the BDS C compiler must be "fooled" by declaring these values to be single characters. To do this, the library header file contains the declaration

```
#define short char.
```

shown in Figure 11-2, section a.

The "#define" tells the first part of the C compiler, the preprocessor, to substitute the string "char" (which declares a character variable) whenever it encounters the string "short" (which would ordinarily declare a short integer in standard C).

Note that character strings enclosed in "/*" and "*/" are regarded as comments and are ignored by the compiler.

BDOS Calls

The standard library of functions that comes with the BDS C compiler includes a function to make BDOS calls, called "bdos." It takes two parameters, and a typical call is of the following form:

```
bdos(c,de);
```

The "c" parameter represents the value that will be placed into the C register. This is the BDOS function code number. The "de" is the value that will be placed in the DE register pair.

The library header contains definitions (#define declarations) for BDOS functions 14 through 32, making these functions easier to use (Figure 11-2, c). Function 32 (Get/Set Current User Number) has two definitions; the "de" parameter is used to differentiate whether a get or a set function is to be performed.

BIOS Calls

The BDS C standard library also contains two functions that make direct BIOS calls. These are "bios" and "biosh." They differ only in that the bios function returns the value in the A register on return from the BIOS routine, whereas biosh, as its name implies, returns the value in the HL register pair. Examples of their use are

```
bios(jump_number,bc);
```

and

```
biosh(jump_number,bc,de);
```

Both functions take as their first parameter the number of the jump instruction in the BIOS jump vector to which control is to be transferred. For example, the console-status entry point is the third JMP in the vector. Numbering from 0, this would be jump number 2.

The library header file contains #defines for BIOS jumps 2 through 21 (Figure 11-2, d). The last group of these #defines (19 through 21) is for the "private" additions to the standard BIOS jump vectors described in Chapter 8.

Remember, though, that using direct BIOS calls makes programs more difficult to move from one system to another.

BIOS Configuration Block Access

As you may recall, the configuration block is a collection of data structures in the BIOS. These structures are used either to store the current settings of certain user-selectable options, or to point to other important data structures in the BIOS.

One of the "private" jumps appended to the standard BIOS jump vector transfers control to a routine that returns the address in memory of a specified data structure. For example, if a utility program needs to locate the word in the BIOS that determines from which physical device the console input is to read, it can transfer control to jump 21 in the BIOS jump vector (actually the 22nd jump) with a code value of 5 in the C register. This jump transfers control to the CBGet-Address code, which on its return will set HL to the address of the console input redirection vector. The utility program can then read from or write into this variable. The library header file contains #define declarations relating the code values to mnemonic names (Figure 11-2, e).

You will need to refer to the source code in Figure 8-10 to determine whether the address returned by the BIOS function is the address of the data element or the address of a higher-level table that in turn points to the data element.

In order to access the current system date, for example, you would include the following code:

```
char *ptr_to_date;              /* declare date pointer*/
ptr_to_date = biosh(CB_DATE);   /* get address */
```

The ptr_to_date can then be used to access the date directly.

During initial debugging of a utility, it is useful to be able to intercept all such accesses to the configuration block, partly to reassure yourself that the utility program is working as it should, and partly to ensure that the BIOS routine is returning the correct addresses to the data structures. Therefore, the utility library contains a function, "get_cba," that gets a configuration block address (Figure 11-1, a).

At first, it appears that get_cba is declared as a function that returns a pointer to characters. This is not strictly true. Sometimes the address it returns will point to characters, sometimes to integers, and sometimes to structures (such as the function key table).

The "printf" instruction has been left in the function in anticipation of debugging a utility. If you need to see some debug output whenever the get_cba function is used, delete the "/*" and "*/" surrounding the "printf" and recompile the library.

BIOS Function Key Table Access

The BIOS shown in Figure 8-10 contains code to recognize when an incoming escape sequence indicates that one of the terminal's function keys has been pressed. Instead of returning just the escape sequence, the console driver injects a previously programmed string of characters into the console input stream. For example, on a DEC VT-100 terminal, when the PF1 function key is pressed, the terminal emits the following character sequence: ESCAPE, "O", "P". The function key table contains the "OP" and a 00H-byte-terminated string of characters to be injected into the console input stream. In Figure 8-10, the example string is "FUNCTION KEY 1", LINE FEED. The library header file contains a declaration for the structure of the function key table (Figure 11-2, h).

Note the use of "#define" to declare the length of the incoming characters emitted by the terminal as well as the length of the output string.

In order to access a function key table entry, you must declare a pointer to a "_fkt" structure like this:

```
struct _fkt *ptr_to_fkt;      /* Declare pointer */
ptr_to_fkt = get_cba(CB_FKT); /* Set pointer */
printf("Display the first string : %s",
     ptr_to_fkt -> fk_output);
++ptr_to_fkt;                 /* Move to next entry */
```

The get_cba function is used to return the address of the first entry in the function key table and set a pointer to it. Then the printf function (part of the standard BDS C library) is used to print out the first string, which gets substituted for the "%s" in the quoted string. Note that the statement

```
++ptr_to_fkt
```

does not just add one to the pointer to the function key table—it adds whatever it takes to move the pointer to the next *entry* in the table.

BIOS Device Table Access

The device tables are important structures for the serial devices served by the console, auxiliary, and list device drivers in the BIOS. They are declared at line 1500 in Figure 8-10.

The get_cba function does not return a pointer to a specific device table, but a pointer to a table of device table addresses. Each entry in the address table corresponds to a specific device number. If there is no device table for a specific device number, then the corresponding entry in the table will be set to zero. the library header file contains definitions for the device table (Figure 11-2, i).

The device tables contain, among other things, the current serial line protocols used to synchronize the transmission and reception of data by the device drivers and the physical devices. An example utility, PROTOCOL, is shown later in the chapter. The example #define declarations and structure definition shown here are modeled on the requirements of this utility. The only relevant bytes are the two status bytes dt_st1 and dt_st2 and the message length used with the ETX/ACK protocol, dt_etxml. The #defines shown are for the specific bits in the device table's status bytes. The PROTOCOL utility uses the most significant bit to indicate whether a given protocol setting can coexist with others.

To access these fields, use the following code:

```
struct _ppdt
     {
     char *pdt[16];         /* Array of 16 pointers to device tables */
     } *ppdt;               /* Pointer to array of 16 pointers */
struct _dt *dt;             /* Pointer to device table */

ppdt = get_cba(CB_DTA);        /* Set pointer to array of pointers */
dt = ppdt -> pdt[device_no];   /* Set pointer to specified device
                                  table */

if (!dt)
     printf("\nError - no device table for this device.");

dt -> dt_etxml = 0;            /* Clear ETX message length */
```

BIOS Disk Parameter Block Access

Several of the utility programs shown in this chapter must access the file directory on a given logical disk. The disk parameter block (DPB) indicates the size and location of the file directory. The library header contains a structure definition that describes the DPB (Figure 11-2, n).

To locate the DPB, you can make a direct BIOS call to the SELDSK routine, which returns the address of the disk parameter header (DPH). You then can access the DPB pointer in the DPH. Alternatively, using the BDOS, you can make the required disk the default disk and then request the address of its DPB. The code for the latter method is shown in the get_dpb function included in the utility library (Figure 11-1, u).

The get_dpb function uses a BIOS SELDSK function first to see if the specified disk is legitimate. Only then does it use the BDOS.

Reading or Writing a Disk Using the BIOS

When you write a program that uses direct BIOS calls, you increase the possibility of problems in moving the program from one system to another. However, in certain circumstances it is necessary to use the BIOS. Reading and writing the file directory is one of these; the BDOS cannot be used to access the directory directly. The library header contains a structure declaration for a parameter block that contains the details of an "absolute" disk read or write (Figure 11-2, p).

Note the pointer to the 128-byte data buffer used to hold one of CP/M's "records."

The disk read and write functions are rd_disk (Figure 11-1, k) and wrt_disk (Figure 11-1, l). Both of them take a _drb as an input parameter, and both call the set_disk function to make the individual BIOS calls to SELDSK, SETTRK, and SETSEC.

Of special note is the code in set_disk (Figure 11-1, m) that converts a logical sector into a physical sector using the sector translation table and the SECTRAN entry point in the BIOS.

File Directory Entry Access

All of the utility programs that access a disk directory share the same basic logic regardless of their specific task. This logic can be described best in pseudo-code:

```
while (not at the end of the directory)
    {
    access the next directory entry
    if (this entry matches the current search criteria)
        {
        process the entry
        }
    }
```

There are two ways of implementing this logic. The first uses the BIOS to read the directory. Entries are presented to the utility exactly as they occur in the file directory. The second uses the BDOS functions Search First and Search Next and accesses the directory file-by-file rather than by entry. This latter method is more suited to utilities that process files rather than entries. The ERASE utility, described later in this chapter, illustrates this second method.

Three groups of functions are provided in the library: to access the next entry in the directory, to match the name in the current entry against a search key, and to assist with processing the directory.

Directory Accessing Functions

A number of functions involve access to the file directory. The first group of such functions performs the following:

get_nde (get next directory entry; Figure 11-1, n)
This function returns a pointer to the next directory entry, or returns zero if the end of the directory has been reached.

open_dir (open directory; Figure 11-1, o)
This function is called by get_nde to open up a directory for processing.

rw_dir (read/write directory; Figure 11-1, p)
This function reads or writes the current directory sector.

err_dir (error on directory; Figure 11-1, q)
This general-purpose routine displays an error message if the BIOS indicates that it had problems either reading or writing the directory.

All of these functions use a directory parameter block to coordinate their activity. The library header contains the definitions for this structure (Figure 11-2, l), as well as #define declarations for operation codes used by the directory-accessing functions (Figure 11-2, m).

Before calling get_nde, the calling program needs to set dp_open to zero (forcing a call to open_dir) and the dp_disk field to the correct logical disk. The open_dir function sets up all of the remaining fields, using get_dpb to access the disk parameter block for the disk specified in dp_disk.

Of the remaining flags, dp_end will be set to true, when the end of the directory is reached, and dp_write must be nonzero for rw_dir to write the current sector back onto the disk.

The get_nde function includes all of the necessary logic to move from one directory entry to the next, reading in the next sector when necessary, and writing out the previous sector if the dp_write flag has been set to a nonzero value by the calling program. It also counts down on the number of directory entries processed, detecting and indicating the end of the directory.

The code at the beginning of the function calls open_dir if the dp_open flag is false. Note the code at the end of open_dir that sets the number of allocation blocks per directory entry (dp_nabpde). This number is computed from the maximum allocation block number in the disk parameter block. If it is larger than 255, each allocation block must occupy a word, and there will be eight blocks per directory entry. If there are 255 or fewer allocation blocks, each will be one byte long and there will be 16 per entry. The allocation block size, in Kbytes, is computed from a simple formula.

In the early stages of debugging utilities, comment out the line that makes the call to wrt_disk. This will prevent the directory from being overwritten. You then can test even those utilities that attempt to erase entries from the directory without any risk of damaging any data on the disk.

The last function in this group, err_dir, is a common error handling function for taking care of errors while reading or writing the directory.

Directory Matching Functions

The second group of functions that access the file directory matches each directory entry against specific search criteria. These include the following functions:

setscb (set search control block; Figure 11-1, r)
A search control block (SCB) is a structure that defines the entries in the directory that are to be selected for processing.

comp_fname (compare file name; Figure 11-1, f)
This function compares the file name in the current directory entry with the one specified in the search control block.

The library header contains the structure definition for the search control block (Figure 11-2, q). This SCB is a hybrid structure. The first part of it is a cross between a file control block (FCB) and a directory entry. The last three fields, scb_length, scb_disk, and scb_adisks, are peculiar to the search control block. Note that its overall length is the same as an FCB's so that the standard BDS C function set_fcb can be used. This function sets the file name and type into an FCB, replacing "*" with as many "?" characters as are required, and clears all unused bytes to zero.

The scb_length field indicates to the comp_fname (compare file name) function how many bytes of the structure are to be compared. This field will be set to 12 to compare the user number, file name, and type, or to 13 to include the extent number.

Note that scb_disk is the *current* disk to be searched, whereas scb_adisks is a bit map with a 1 bit corresponding to each of the 16 possible logical disks that must be searched.

The search control block is initialized by the setscb function.

Note the form of the file name that setscb expects to receive. This is described in the comments at the beginning of the function.

Several of the utility programs use their own special versions of setscb, renaming it ssetscb (special setscb) to avoid the library version being linked into the programs.

The complementary function comp_fname is used to compare the first few bytes of the current directory entry to the corresponding bytes of the SCB.

The comp_fname function performs a specialized string match of the user number, the file name, the file type, and, optionally, the extent number. A "?" character in the search control block file name, type, and extent will match with any character in the file directory entry. However, in the SCB user number, a "?" will only match a number in the range 0 to 15; it will not match a directory entry that has the user number byte set to E5H (or 0xE5, as hexadecimal notation in C).

This function also returns one of several values to indicate the result of the comparison. These values are defined in the library header file (Figure 11-2, j).

Directory Processing Functions

The final group of functions that access the directory are those that help process the directory entries themselves. These functions use a structure definition to access each directory entry (Figure 11-2, o).

A union statement is used for the allocation block numbers. These can be single- or two-byte entries, depending on the maximum number of allocation blocks that must be represented. The union statement tells the BDS C compiler whether there will be a 16-byte array of short integers (characters) or an array of eight unsigned two-byte integers.

The functions contained in this group can be divided into three subgroups:

- Those that deal with converting directory entries for display on the console.
- Those that deal with a "disk map"—a convenient array for representing logical disks and the user numbers they contain.
- Those that deal with "bit vectors"—a convenient representation of which allocation blocks on a logical disk are in use or available.

The library contains only one function to convert a directory-entry file name into a suitable form for display on the console. This is the conv_dfname function (Figure 11-1, h). It takes the information from the specified directory entry (or, as a convenience, a search control block) and formats it into a string of the form

```
uu/d:filename.typ
```

The "uu" specifies the user number and the "d" specifies the disk identification.

The repetitive code at the end of the function is necessary to make sure that the characters in the file type do not have their high-order bits set. These bits are the file attributes. If they are set, they can render the characters nondisplayable on some terminals.

The second subgroup of functions, those that manipulate a "disk map," produce an array that looks like this:

```
Disks
  :
  v  User Numbers -->                                          -Totals-
  A  0  1  2  3  4  5  6  7  8  9 10 11 12 13 14 15 Used Free
  B
  :
  :
  P
```

This disk map is used by several utility programs. For example, the SPACE utility displays a disk map that shows, for each logical disk in the system, and for each user on each logical disk, how many Kbytes of disk space are in use. The totals at the right show the total of used and free space. In another example, the FIND utility shows how many files on each disk and in each user number match the search name.

Each utility program that uses a disk map is coded:

```
unsigned disk_map[16][18];
```

Two functions are provided in the library to deal with the disk map:

dm_clr (disk map clear; Figure 11-1, s)
: This function fills the entire disk map with zeros.

dm_disp (disk map display; Figure 11-1, t)
: This function displays the horizontal and vertical caption lines for the disk map and then converts each element of the disk map to a decimal number.

The first function, dm_clr, uses one of the standard BDS C functions to set a block of memory to a specific value. It presumes that the disk map is 16 × 18 elements, each two bytes long.

The second function, dm_disp, prints horizontal lines only for those disks specified in the bit map parameter. Here is an example of its output:

```
     0   1   2   3   4   ...   10  11  12  13  14  15  Used Free
A:   1   1                                               15  241
B:  66  20  74  50   3                                  245  779
C:  -- None --                                            0 1024
(NOTE: All user groups would be shown on the terminal.)
```

The final subgroup deals with processing "bit vectors." A bit vector is a string of bits packed eight bits per byte. Each bit is addressed by its relative number along the vector; the first bit is number 0.

An example of why bit vectors are used is a utility program that needs to scan the directory of a disk and build a structure showing which allocation blocks are in use. It can do this by accessing each active directory element and, for each nonzero allocation block number, setting the corresponding bit number in a bit vector.

The library header has a structure definition for a bit vector (Figure 11-2, s).

This vector contains the overall length of the bit vector in bytes, and two pointers. The first points to the start of the vector, the second to the end. The bytes that contain the vector bits themselves are allocated by the alloc function — one of the standard BDS C functions.

The following bit vector functions are provided in the library:

bv_make (bit vector make; Figure 11-1, cc)
: This function allocates memory for the bit vector (using the standard mechanism provided by BDS C) and sets all of the bits to zero.

bv_fill (bit vector fill; Figure 11-1, dd)
: This fills a specified vector, setting each byte to a specified value.

bv_set (bit vector set; Figure 11-1, ee)
: This sets the specified bit of a vector to one.

bv_test (bit vector test; Figure 11-1, ff)
: This function returns a value of zero or one, reflecting the setting of the specified bit in a bit vector.

bv_nz (bit vector nonzero; Figure 11-1, gg)
: This returns zero or a nonzero value to reflect whether *any* bits are set in the specified bit vector.

bv_and (bit vector AND; Figure 11-1, hh)
: This function performs a Boolean AND between two bit vectors and places the result into a third vector.

bv_or (bit vector OR; Figure 11-1, ii)
: This is similar to bv_and, except that it performs an inclusive OR on the two input vectors.

bv_disp (bit vector display; Figure 11-1, jj)
: This function displays a caption line and then prints out the contents of the specified bit vector as a series of zeros and ones. Each byte is formatted to make the output easier to read.

The bv_make function uses the alloc function to allocate a block from the unused part of memory between the end of a program and the base of the BDOS. It requires that two data structures be declared at the beginning of the program. These structures are declared in the library header file (Figure 11-2, b).

The bv_fill function uses the standard BDS C setmem function.

The bv_set function converts the bit number into a byte offset by shifting the bit number right three places. The least significant three bits of the original bit number specify which bit in the appropriate byte needs to be ORed in.

The bv_test function is effectively the reverse of bv_set. It accesses the specified bit and returns its value to the calling program.

The bv_nz function scans the entire bit vector looking for the first nonzero byte. If the entire vector is zero, it returns a value of zero. Otherwise, it returns a pointer to the first nonzero byte.

Both bv_and and bv_or functions take three bit vectors as parameters. The first vector is used to hold the result of either ANDing or ORing the second and third vectors together. Both of these functions assume that the output vector has already been created using bv_make. The shortest of the three vectors will terminate the bv_and or bv_or function; that is, these functions will terminate when they reach the end of the first (shortest) vector.

The final function, bv_disp, displays the title line specified by the calling program, and then displays all of the bits in the vector, with the bit number of the first bit on each line shown on the left.

None of the utility programs uses bv_disp—it has been left in the library purely as an aid to debugging.

Here is an example of bv_disp's output:

```
Bit Vector : Allocation Blocks in Use
   0 : 0000 0000  0001 1000  1000 0001  1111 1111  1111 1111
  40 : 1111 1111  1111 1111  1111 1111  1110 1011  0000 0000
  80 : 1100 0000  1111 1100  1111 1001  1100 0000  1001 1111
 120 : 1110 1100  0001 1111  0000 0000  1101 1000  0001 1110
 160 : 1111 1111  1110 1111  1110 1111  0000 0111  0000 0111
 200 : 1111 0010
```

Checking User-Specified Parameters

The C language provides a mechanism for accessing the parameters specified in the "command tail." It provides a count of the number of parameters entered, "argc" (argument count), and an array of pointers to each of the character strings, "argv" (argument vector). At the beginning of the main function of each program you must define these two variables like this:

```
main(argc,argv)
{
int argc;       /* Argument count */
char *argv[];   /* Array of pointers to char. strings */
:
: /* Remainder of main function */
:
}
```

Consider the minimum case—a command line with just the program name on it:

```
A>command
```

The convention is that the first argument on the line is the name of the program itself. Hence argc would be set to one, and argv[0] would be a pointer to the program name, "command."

Next consider a more complex case — a command line with parameters like the following:

```
A>command param1 123
```

In this case, argc will be three; argv[1] will be a pointer to param1; and argv[1][0] will access the 0 (the first) character of argv[1]—in this case the character "p."

To detect whether the second parameter is present and numeric, the code will be

```
if (isdigit(argv[1][0]))
     {
          /* Process digit */
     }
else
     {
          /* Parameter either not present or has
             alpha character at the front */
     }
```

In most of the utilities, you will get a much "friendlier" program if the user need only specify enough characters of a parameter to distinguish the value entered from the other possible values. For example, consider a program that can have as a parameter one of the following values: 300, 600, 1200, 2400, 4800, 9600, or 19200. It would be convenient if the user needed to type only the first digit, rather than having to enter redundant keystrokes. However, the values 1200 and 19200 would then be ambiguous. The user would have to enter 12 or 19. Novice users often prefer to specify the entire parameter for clarity and security.

The standard C library provides a character string comparison function, strcmp. Unfortunately, this function does not provide for the partial matching just described. Therefore, the library includes two special functions that do make this possible: sstrcmp (substring compare, Figure 11-1, d) and usstrcmp (uppercase substring compare, Figure 11-1, e). The latter function is necessary when you need to compare a substring that could contain lowercase characters; it converts characters to uppercase before the comparison.

To assist with character string manipulation, two additional functions have been included in the library. These are strscn (string scan, Figure 11-1, b) and ustrcmp (uppercase string compare, Figure 11-1, c).

Using Code Tables

A code table is a simple structure used by all of the utility programs that accept parameters that can have any of several values. The library header contains a structure definition for a code table (Figure 11-2, r).

A code table entry contains an unsigned code value and a pointer to a character string. It is used in the utility programs wherever there is a need to relate some arbitrary code number or bit pattern to an ASCII character string. For example, to program a serial port baud-rate-generator chip to various baud rates requires different time constants for each rate. Users do not need to know what these numbers are; they only need to be able to specify the baud rate as an ASCII string.

Thus, a code table is set up as follows:

Baud Rate Constant	User's Name
0x35	"300"
0x36	"600"
0x37	"1200"
0x3A	"2400"
0x3C	"4800"
0x3E	"9600"
0x3F	"19200"

A utility program now needs to be able to perform various operations using the code table:

- Given the input parameter on the command tail, the utility must check whether the ASCII string is in the code table, display all of the legal options on the console if it is not, and return the code value for subsequent processing if it is.
- Given the current baud rate constant (held in the BIOS), the utility must scan the code table and display the corresponding ASCII string to tell the user the current baud rate setting.

The library includes specialized functions to do this, plus some additional functions to make code tables more generally usable. These functions are

ct_init (code table initialize; Figure 11-1, v)
: This function initializes a specific entry in a code table, setting the code value and the pointer to the character string.

ct_parc (code table parameter return code; Figure 11-1, w)
: This performs an uppercase substring match on the specified key string, returning either an error (the value CT_SNF—string not found) or a code value.

ct_code (code table return code; Figure 11-1, x)
: This function is similar to ct_parc in that it scans a code table and returns the corresponding code. It differs in the way that the comparison is done. The entire search string is compared with the string in the code table entry. A match only occurs when all characters are the same.

ct_disps (code table display strings; Figure 11-1, y)
: This function displays all strings in a given code table. It is used either when the user has entered an invalid string, or when the utility program is requested to show what options are available for a parameter.

ct_index (code table return index; Figure 11-1, z)
: This function, given a string, searches the code table and returns the *index*

of the entry that has a string matching the search string. The index is not the code value; it is the number of the entry in the table.

ct_stri (code table string index; Figure 11-1, aa)
This function, given an entry index number, returns a pointer to the string in that entry.

ct_strc (code table string code; Figure 11-1, bb)
This function, given a code number, returns a pointer to the string in the entry that has a matching code number.

Accessing a Directory via the BDOS

One problem associated with accessing the file directory directly, as illustrated by earlier functions, is that the program is presented with directory entries in exactly the order that they occur in the directory. For some programs, such as those that process groups of files, it is better to use the BDOS Search First and Search Next functions to access the directory.

Using the BDOS, the program can process the first file name to match an ambiguous search key, then go back to the BDOS to get the name of the next file, and so on. The library header contains a structure definition for a standard CP/M file control block (Figure 11-2, k).

Notice that the first byte of the FCB is a disk number rather than the user number of the directory entry. Note also the use of a union statement to describe the allocation block numbers.

The standard BDS C library contains a function, setfcb, that is given the address of an FCB and a pointer to a string containing a file name. It converts any "*" in the name to the appropriate number of "?", and fills the remainder of the FCB with zeros.

The example library contains the following functions designed for BDOS file directory access:

get_nfn (get next file name; Figure 11-1, i)
This function is given a pointer to an ambiguous file name and a pointer to an FCB. It returns with the FCB set up to access the next file that matches the ambiguous file name.

srch_file (search for file; Figure 11-1, j)
This function, used by get_nfn, issues either a Search First or a Search Next BDOS call.

conv_fname (convert file name; Figure 11-1, g)
This function converts a file name from an FCB into a form suitable for display on the console. It is similar to the conv_dfname function described earlier except that it outputs only the disk, file name, and type (not the user number) in the form

```
d:filename.typ
```

To signal the get_nfn function that you want the first file name, you must set the most significant bit of the first byte, the disk number.

Here is an example showing how to use the get_nfn function:

```
struct _fcb fcb;            /* Declare a file control block */

setmem(fcb,FCB_SIZE,0);     /* Clear FCB to zeros */
fcb.fcb_disk = 0x80;        /* Mark FCB for "first time" */

while (get_nfn(fcb,"B:XYZ*.*"))
                            /* Until get_nfn returns a zero */
     {
                            /* Open the file using FCB */
     while                 (/* Not at end of file */)
          {
                            /* Process next record or
                               Character in file*/
          }
                            /* Close the file */
     }
```

The quoted string "B:XYZ*.*" could also be just a pointer to a string, or a parameter on the command line, argv[n].

The last function for BDOS processing of the file directory, conv_fname, is used to convert a file name for output to a terminal. Again, the repetitive code at the end clears the file attribute bits to avoid any side effects from the terminal.

Utility Programs Enhancing Standard CP/M

This group of utilities is designed to enhance those supplied by Digital Research. They do not take advantage of any special features of the enhanced BIOS in Figure 8-10 and can be used on *any* CP/M Version 2.2 installation.

With the exception of the ERASE utility, all of the utilities scan down the file directory using BIOS calls, as described earlier in this chapter.

ERASE — A Safer Way to Erase Files

There are two disadvantages to the Console Command Processor's built-in ERA command. First, it will unquestioningly erase groups of files. Second, if you have a file name with nongraphic or lowercase characters, you cannot use the ERA command, as the CCP converts the command tail characters to uppercase and terminates a file name on encountering any strange character in the string.

The ERASE utility shown in Figure 11-3 erases groups of files, but it asks the user for confirmation before it erases each file.

Rather than use the BIOS to access each directory entry, it uses the get_nfn function, which then calls the BDOS. Thus ERASE functions equally well for files that have multiple entries in the directory. It can use the BDOS Delete File function to erase all extents of a given file.

Here is an example console dialog showing ERASE in operation:

```
P3A>erase<CR>
ERASE Version 1.0 02/23/83 (Library 1.0)
Usage :
        ERASE {d:}file_name.typ


P3A>erase *.com<CR>
ERASE Version 1.0 02/23/83 (Library 1.0)

Searching for file(s) matching A:????????.COM.
        Erase A:UNERASE .COM y/n? n
        Erase A:TEMP1   .COM y/n? y <== Will be Erased!
        Erase A:TEMP2   .COM y/n? n
        Erase A:TEMP3   .COM y/n? n
        Erase A:TEMP4   .COM y/n? y <== Will be Erased!
        Erase A:ERASE   .COM y/n? n

Erasing files now...
        File A:TEMP1   .COM erased.
        File A:TEMP4   .COM erased.
```

```
#define VN "1.0 02/24/83"

/* ERASE
   This utility erases the specified file(s) logically
   by using a BDOS delete function. */

#include <LIBRARY.H>

struct _fcb amb_fcb;            /* Ambiguous name file control block */
struct _fcb fcb;                /* Used for BDOS search functions */

char file_name[20];             /* Formatted for display: d:FILENAME.TYP */
short cur_disk;                 /* Current logical disk at start of program */
                                /* ERASE saves the FCB's of the all the
                                   files that need to be erased in the
                                   following array  */
#define MAXERA 1024
struct _fcb era_fcb[MAXERA];
int ecount;                     /* Count of number of files to be erased */
int count;                      /* Used to access era_fcb during erasing */


main(argc,argv)
short argc;             /* Argument count */
char *argv[];           /* Argument vector (pointer to an array of char.   */
{

printf("\nERASE Version %s (Library %s)",VN,LIBVN);
chk_use(argc);                  /* Check usage */
cur_disk = bdos(GETDISK);       /* Get current default disk */

ecount = 0;                     /* Initialize count of files to erase */

setfcb(amb_fcb,argv[1]);        /* Set ambiguous file name */
if (amb_fcb.fcb_disk)           /* Check if default disk to be used */
        {
        bdos(SETDISK,amb_fcb.fcb_disk + 1);     /* Set to specified disk */
        }
```

Figure 11-3. ERASE.C, a utility that requests confirmation before erasing

```
        /* Convert ambiguous file name for output */
conv_fname(amb_fcb,file_name);
printf("\n\nSearching for file(s) matching %s.",file_name);

        /* Set the file control block to indicate a "first" search */
fcb.fcb_disk |= 0x80;   /* OR in the ms bit */

        /* While not at the end of the directory, set the FCB
           to the next name that matches */
while(get_nfn(amb_fcb,fcb))
        {
        conv_fname(fcb,file_name);
                /* Ask whether to erase file or not */
        printf("\n\tErase %s y/n? ",file_name);
        if (toupper(getchar()) == 'Y')
                {
                printf(" <== Will be erased!");
                        /* add current fcb to array of FCB's */
                movmem(fcb,&era_fcb[ecount++],FCBSIZE);
                        /* Check that the table is not full */
                if (ecount == MAXERA)
                        {
                        printf("\nWarning : Internal table now full. No more files can be erased");
                        printf("\n   until those already specified have been erased.");
                        break;  /* Break out of while loop */
                        }
                }
        }       /* All directory entries processed */

if (ecount)
        printf("\n\nErasing files now...");

        /* now process each FCB in the array, erasing the files */
for (count = 0;         /* Starting with the first file in the array */
     count < ecount;    /* Until all active entries processed */
     count++)           /* Move to next FCB */
        {
        conv_fname(&era_fcb[count],file_name);
        if (bdos(DELETEF,&era_fcb[count]) == -1)        /* error? */
                printf("\n\007Error trying to erase %s",file_name);
        else            /* File erased */
                printf("\n\tFile %s erased.",file_name);
        }
bdos(SETDISK,cur_disk); /* reset to current disk */
}


chk_use(argc)           /* Check usage */
/* This function checks that the correct number of
   parameters has been specified, outputting instructions if not. */


/* Entry parameter */
int argc;       /* Count of the number of arguments on the command line */

{
        /* The minimum value of argc is 1 (for the program name itself),
           so argc is always one greater than the number of parameters
           on the command line */

if (argc != 2)
        {
        printf("\nUsage :");
        printf("\n\tERASE {d:}file_name.typ");
        exit();
        }

}
```

Figure 11-3. (Continued)

UNERASE — Restore Erased Files

UNERASE, as its name implies, can be used to "revive" an accidentally erased file. Only files whose allocation blocks have not been reallocated to other files can be revived. The UNERASE utility shown in Figure 11-4 builds a bit vector of all the allocation blocks used by active directory entries. Then it builds a bit vector for all the allocation blocks required by the file to be UNERASEd. If a Boolean AND between the two vectors yields a nonzero vector, then one or more blocks that originally belonged to the erased file are now allocated to other files on the disk.

```
#define VN "1.0 02/12/83"

/* UNERASE --
   This utility does the inverse of ERASE: it restores
   specified files to the directory by changing the first byte of
   their directory entries from 0xE5 back to the specified user
   number. */

#include <LIBRARY.H>

struct _dirpb dir_pb;          /* Directory management parameter block */
struct _dir *dir_entry;        /* Pointer to directory entry */
struct _scb scb;               /* Search control block */
struct _scb scba;              /* SCB set up to match all files */
struct _dpb dpb;               /* CP/M's disk parameter block */
struct _bv inuse_bv;           /* Bit vector for blocks in use */
struct _bv file_bv;            /* Bit vector for file to be unerased */
struct _bv extents;            /* Bit vector for those extents unerased */

char file_name[20];            /* Formatted for display : un/d:FILENAME.TYP */

short cur_disk;                /* Current logical disk at start of program
                                  NZ = show map of number of files */
int count;                     /* Used to access the allocation block numbers
                                  in each directory entry */
int user;                      /* User in which the file is to be revived */


main(argc,argv)
short argc;               /* Argument count */
char *argv[];             /* Argument vector (pointer to an array of chars.) */

{
printf("\nUNERASE Version %s (Library %s)",VN,LIBVN);
chk_use(argc);                 /* Check usage */
cur_disk = bdos(GETDISK);      /* Get current default disk */

        /* Using a special version of the set search-control-block utility,
           set the disk, name, type (no ambiguous names), the user number
           to match only erased entries, and the length to compare
           the user, name, and type.
           This special version also returns the disk_id taken from
           the file name on the command line.  */
if ((dir_pb.dp_disk = ssetscb(scb,argv[1],0xE5,12)) == 0)
        {       /* Use default disk */
        dir_pb.dp_disk = cur_disk;
        }
else
        {       /* make disk A = 0, B = 1 (for SELDSK) */
        dir_pb.dp_disk--;
        }
printf("\nSearching disk %d.",dir_pb.dp_disk);

if(strscn(scb,"?"))       /* Check if ambiguous name */
        {
        printf("\nError -- UNERASE can only revive a single file at a time.");
        exit();
```

Figure 11-4. UNERASE.C, a utility program that "revives" erased files

```
        }

        /* Set up a special search control block that will match with
           all existing files. */

ssetscb(scba,"*.*",'?',12);     /* Set file name and initialize SCB */

if (argc == 2)                  /* No user number specified */
        user = bdos(GETUSER,0xFF);      /* Get current user number */
else
{
        user = atoi(argv[2]);           /* Get specified number */
        if (user > 15)
                {
                printf("\nUser number can only be 0 - 15.");
                exit();
                }
}

/* Build a bit vector that shows the allocation blocks
   currently in use. SCBA has been set up to match all
   active directory entries on the disk. */
build_bv(inuse_bv,scba);

/* Build a bit vector for the file to be restored showing
   which allocation blocks will be needed for the file. */
if (!build_bv(file_bv,scb))
        {
        printf("\nNo directory entries found for file %s.",
                argv[1]);
        exit();
        }
/* Perform a boolean AND of the two bit vectors. */
bv_and(file_bv,inuse_bv,file_bv);

/* Check if the result is nonzero -- if so, then one or more
   of the allocation blocks required by the erased file is
   already in use for an existing file and the file cannot
   be restored. */
if (bv_nz(file_bv))
{
        printf("\n--- This file cannot be restored as some parts of it");
        printf("\n    have been re-used for other files! ---");
        exit();
}


/* Continue on to restore the file by changing all the entries
   in the directory to have the specified user number.
   Note: There may be several entries in the directory for
   the same file name and type, and even with the same extent
   number. For this reason, a bit map is kept of the extent
   numbers unerased -- duplicate extent numbers will not be
   unerased. */

/* Set up the bit vector for up to 127 unerased extents */
bv_make(extents,16);            /* 16 * 8 bits */


/* Set the directory to "closed", and force the get_nde
   function to open it. */
dir_pb.dp_open = 0;

/* While not at the end of the directory, return a pointer to
   the next entry in the directory. */
while(dir_entry = get_nde(dir_pb))
{

        /* Check if user = 0xE5 and name, type match */
if (comp_fname(scb,dir_entry) == NAME_EQ)
        {
                /* Test if this extent has already been
                   unerased */
        if (bv_test(extents,dir_entry -> de_extent))
                {               /* Yes it has */
                printf("\n\t\tExtent #%d of %s ignored.",
                        dir_entry -> de_extent,argv[1]);
                continue;       /* Do not unerase this one */
                }
```

Figure 11-4. (Continued)

```
        else                    /* Indicate this extent unerased */
                {
                bv_set(extents,dir_entry -> de_extent);
                dir_entry -> de_userno = user; /* Unerase entry */
                dir_pb.dp_write = 1;    /* Need to write sector back */
                printf("\n\tExtent #%d of %s unerased.",
                        dir_entry -> de_extent,argv[1]);
                }
        }
}

printf("\n\nFile %s unerased in User Number %d.",
argv[1],user);

bdos(SETDISK,cur_disk); /* Reset to current disk */
}


build_bv(bv,scb)        /* Build bit vector (from directory) */
/* This function scans the directory of the disk specified in
   the directory parameter block (declared as a global variable),
   and builds the specified bit vector, showing all the allocation
   blocks used by files matching the name in the search control
   block. */

/* Entry parameters */
struct _bv *bv;         /* Pointer to the bit vector */
struct _scb *scb;       /* Pointer to search control block */
/* Also uses : directory parameter block (dir_pb) */

/* Exit parameters
   The specified bit vector will be created, and will have 1-bits
   set wherever an allocation block is found in a directory
   entry that matches the search control block.
   It also returns the number of directory entries matched. */
{
unsigned abno;          /* Allocation block number */
struct _dpb *dpb;       /* Pointer to the disk parameter block in the BIOS */
int mcount;             /* Match count of dir. entries matched */

mcount = 0;             /* Initialize match count */
dpb = get_dpb(dir_pb.dp_disk);  /* Get disk parameter block address */

/* make the bit vector with one byte for each eight allocation
   blocks + 1 */
if (!(bv_make(bv,(dpb -> dpb_maxabn >>3)+1)))
        {
        printf("\nError -- Insufficient memory to make a bit vector.");
        exit();
        }


/* Set directory to "closed" to force the get_nde
   function to open it. */
dir_pb.dp_open = 0;

/* Now scan the directory building the bit vector */
while(dir_entry = get_nde(dir_pb))
        {
                /* Compare user number (which can legitimately be
                   0xE5), the file name and the type). */
        if (comp_fname(scb,dir_entry) == NAME_EQ)
                {
                ++mcount;               /* Update match count */
                for (count = 0;         /* Start with the first alloc. block */
                    count < dir_pb.dp_nabpde;   /* For number of alloc. blks. per dir. entry */
                    count++)
                        {
                                /* Set the appropriate bit number for
                                   each nonzero allocation block number */
                        if (dir_pb.dp_nabpde == 8)      /* assume 8 2-byte numbers */
                                {
                                abno = dir_entry -> _dirab.de_long[count];
                                }
                        else    /* Assume 16 1-byte numbers */
                                {
```

Figure 11-4. (Continued)

```
                                        abno = dir_entry -> _dirab.de_short[count];
                                        }
                                if (abno) bv_set(bv,abno); /* Set the bit */
                                }
                        }
                }
        return mcount;          /* Return number of dir. entries matched */
}


chk_use(argc)           /* Check usage */
/* This function checks that the correct number of
   parameters has been specified, outputting instructions
   if not. */

/* Entry parameter */
int argc;       /* Count of the number of arguments on the command line */
{

/* The minimum value of argc is 1 (for the program name itself),
   so argc is always one greater than the number of parameters
   on the command line */

if (argc == 1 || argc > 3)
        {
        printf("\nUsage :");
        printf("\n\tUNERASE {d:}filename.typ {user}");
                printf("\n\tOnly a single unambiguous file name can be used.)");
        exit();
        }
} /* end chk_use */


ssetscb(scb,fname,user,length)  /* Special version of set search control block */
/* This function sets up a search control block according
   to the file name, type, user number, and number of bytes
   to compare.
   The file name can take the following forms :
        filename
        filename.typ
        d:filename.typ

   It sets the bit map according to which disks should be searched.
   For each selected disk, it checks to see if an error is generated
   when selecting the disk (i.e. if there are disk tables in the BIOS
   for the disk). */

/* Entry parameters */
struct _scb *scb;       /* Pointer to search control block */
char *fname;            /* Pointer to the file name */
short user;             /* User number to be matched */
int length;             /* Number of bytes to compare */

/* Exit parameters
   Disk number to be searched. (A = 1, B = 2...)
*/
{
short disk_id;          /* Disk number to search */

setfcb(scb,fname);      /* Set search control block as though it
                                were a file control block. */
disk_id = scb -> scb_userno;    /* Set disk_id before it gets overwritten
                                        by the user number */
scb -> scb_userno = user;       /* Set user number */
scb -> scb_length = length;     /* Set number of bytes to compare */
return disk_id;
} /* end setscb */
```

Figure 11-4. (Continued)

A further complication occurs if two or more directory entries of the erased file have the same extent number. This can happen if the file has been created and erased several times. Under these circumstances, UNERASE revives the first entry with a given extent number that it encounters, and displays a message on the console both when an extent is revived and when one is ignored.

Because of the complicated nature of the UNERASE process, the utility can process only a single, unambiguous file name.

The following console dialog shows UNERASE in operation:

```
P3A>dir *.com<CR>
A: UNERASE  COM : TEMP2    COM : TEMP3    COM : ERASE    COM

P3A>unerase<CR>
UNERASE Version 1.0 02/12/83 (Library 1.0)
Usage :
        UNERASE {d:}filename.typ {user}
        Only a single unambiguous file name can be used.

P3A>unerase temp1.com<CR>
UNERASE Version 1.0 02/12/83 (Library 1.0)
Searching disk A.
        Extent #0 of TEMP1.COM unerased.
                Extent #0 of TEMP1.COM ignored.

File TEMP1.COM unerased in User Number 3.

P3A>dir *.com<CR>
A: UNERASE  COM : TEMP1    COM : TEMP2    COM : TEMP3    COM
A: ERASE    COM

P3A>unerase temp5.com<CR>
UNERASE Version 1.0 02/12/83 (Library 1.0)
Searching disk A.
No directory entries found for file TEMP5.COM.
```

FIND — Find "Lost" Files

The FIND utility shown in Figure 11-5 searches all user numbers on specified logical disks, matching each entry against an ambiguous file name. It can then display either a disk map showing how many matching files were found in each user number for each disk, or the user number, file name, and type for each matched directory entry.

You can use FIND to locate a specific file or group of files, as shown in the following console dialog:

```
P3B>find<CR>
FIND Version 1.0 02/11/83 (Library 1.0)
Usage :
        FIND d:filename.typ {NAMES}
             *:filename.typ (All disks)
             ABCD..OP:filename.typ (Selected Disks)
        NAMES option shows actual names rather than map.

P3B>find ab:*.*<CR>
FIND Version 1.0 02/11/83 (Library 1.0)
```

```
Searching disk : A
Searching disk : B
                Numbers show files in each User Number.
                          --- User Numbers ---          Dir. Entries
      0    1    2    3    4    5   ...   11   12   13   14   15  Used Free
A:    1    1         8                                            23  233
B:   66   20   74   55    3                                      252  772

P3B>find *:*.com<CR>
FIND Version 1.0 02/11/83 (Library 1.0)
Searching disk : A
Searching disk : B
Searching disk : C
                          --- User Numbers ---          Dir. Entries
      0    1    2    3    4    5   ...   11   12   13   14   15  Used Free
A:                   5                                            23  233
B:   61    5    4   13                                           252  772
C:   -- None --                                                   16  112


P3B>find *.com names<CR>
FIND Version 1.0 02/11/83 (Library 1.0)
Searching disk : B
 0/B:CC      .COM 0/B:CC2     .COM 0/B:CLINK   .COM 2/B:CLIB    .COM
 1/B:CPM61   .COM 1/B:MOVCPM  .COM 1/B:PSWX    .COM 0/B:SUBMIT  .COM
 2/B:CDB     .COM 1/B:CPM60   .COM 0/B:DDT     .COM 0/B:EREMOTE .COM
 0/B:SPEEDSP .COM 0/B:PIP     .COM 0/B:PROTOSP .COM 0/B:RX      .COM
 0/B:TXA     .COM 0/B:EPUB    .COM 0/B:EPRIV   .COM 0/B:WSC     .COM
 0/B:X       .COM 0/B:CRCK    .COM 0/B:XSUB    .COM 0/B:DU      .COM
 0/B:QERA    .COM 0/B:FINDALL .COM 0/B:MOVEF   .COM 0/B:REMOTE  .COM
 0/B:LOCAL   .COM 0/B:DUMP    .COM 0/B:MRESET  .COM 0/B:ELOCAL  .COM
 0/B:PUTCPMF5.COM 0/B:TEST    .COM 0/B:FDUMP   .COM 0/B:INVIS   .COM
 0/B:L80     .COM 0/B:LIST    .COM 0/B:PUB     .COM 0/B:LOAD    .COM
 0/B:MAC     .COM 0/B:SCRUB   .COM 0/B:RXA     .COM 0/B:STAT    .COM
 0/B:TX      .COM 0/B:ERASEALL.COM 0/B:WM      .COM 0/B:MSFORMAT.COM
 0/B:STATUS  .COM 0/B:UNERA   .COM 0/B:MSINIT  .COM 0/B:VIS     .COM
 0/B:WSVTIP  .COM 0/B:XD      .COM 0/B:NEWVE   .COM 0/B:DDUMP   .COM
 0/B:FORMATMA.COM 0/B:PRIV    .COM 0/B:FCOMP   .COM 0/B:DDUMPA  .COM
 0/B:PUTSYS1C.COM 0/B:DDUMPNI .COM 0/B:DSTAT   .COM 0/B:ASM     .COM
 2/B:CDBTEST .COM 0/B:OLDSYS  .COM 0/B:E       .COM 2/B:F/C     .COM
 3/B:ERASE   .COM 3/B:FUNKEY  .COM 3/B:DATE    .COM 3/B:FIND    .COM

Press Space Bar to continue....
 3/B:SPACE   .COM 3/B:UNERASE .COM 3/B:MAKE    .COM 3/B:MOVE    .COM
 1/B:PUTSYSWX.COM 3/B:TIME    .COM 3/B:ASSIGN  .COM 3/B:SPEED   .COM
 3/B:PROTOCOL.COM 0/B:PRINTC  .COM 3/B:T       .COM
```

```
#define VN "1.0 02/11/83"

/* FIND - This utility can display either a map showing on which disks
   and in which user numbers files matching the specified ambiguous
   file name are found, or the actual names matched. */

#include <LIBRARY.H>

struct _dirpb dir_pb;          /* Directory management parameter block */
struct _dir *dir_entry;        /* Pointer to directory entry (somewhere in
                                     dir_pb) */
struct _scb scb;               /* Search control block */

char file_name[20];            /* Formatted for display : un/d:FILENAME.TYP */
```

Figure 11-5. FIND.C, a utility program that locates specific files or groups of files

```
short cur_disk;                 /* Current logical disk at start of program */
int mcount;                     /* Match count (no. of file names matched) */
int dmcount;                    /* Per disk match count */
int lcount;                     /* Line count (for lines displayed) */

int map_flag;                   /* 0 = show file names of matched files,
                                   NZ = show map of number of files */

        /* The array below is used to tabulate the results for each
           disk drive, and for each user number on the drive.
           In addition, two extra "users" have been added for "free"
           and "used" values. */

unsigned disk_map[16][18];      /* Disk A -> P, users 0 -> 15, free, used */
#define USED_COUNT 16           /* "User" number for used entities */
#define FREE_COUNT 17           /* "User" number for free entities */


main(argc,argv)
short argc;             /* Argument count */
char *argv[];           /* Argument vector (pointer to an array of chars.) */
{

printf("\nFIND Version %s (Library %s)",VN,LIBVN);
chk_use(argc);                  /* Check usage */
cur_disk = bdos(GETDISK);       /* Get current default disk */

dm_clr(disk_map);               /* Reset disk map */

        /* Set search control block
           disks, name, type, user number, extent number,
           and number of bytes to compare -- in this case, match all users,
           but only extent 0 */
setscb(scb,argv[1],'?',0,13);   /* Set disks, name, type */

map_flag = usstrcmp("NAMES",argv[2]);   /* Set flag for map option */

lcount = dmcount = mcount = 0;          /* Initialize counts */

for (scb.scb_disk = 0;          /* Starting with logical disk A: */
     scb.scb_disk < 16;         /* Until logical disk P: */
     scb.scb_disk++)            /* Move to next logical disk */
{

        /* Check if current disk has been selected for search */
if (!(scb.scb_adisks & (1 << scb.scb_disk)))
        continue;       /* No,so bypass this disk */

printf("\nSearching disk : %c",(scb.scb_disk + 'A'));
lcount++;               /* Update line count */

dir_pb.dp_disk = scb.scb_disk;  /* Set to disk to be searched*/
dmcount = 0;                    /* Reset disk matched count */

if (!map_flag)          /* If file names are to be displayed */
        putchar('\n');  /* Move to column 1 */

/* Set the directory to "closed", and force the get_nde
   function to open it  */
dir_pb.dp_open = 0;

        /* While not at the end of the directory, set a pointer to the
           next directory entry  */
while(dir_entry = get_nde(dir_pb))
        {
        /* Check if entry in use, to update
           the free/used counts */

        if (dir_entry -> de_userno == 0xE5)     /* Unused */
                disk_map[scb.scb_disk][FREE_COUNT]++;
        else    /* In use */
                disk_map[scb.scb_disk][USED_COUNT]++;

        /* Select only those active entries that are the
           first extent (numbered 0) of a file that matches
           the name supplied by the user  */
```

Figure 11-5. (Continued)

```
        if (
            (dir_entry -> de_userno != 0xE5) &&
            (dir_entry -> de_extent == 0) &&
            (comp_fname(scb,dir_entry) == NAME_EQ)
           )
                {

                mcount++;       /* Update matched counts */
                dmcount++;      /* Per disk count */

                if (map_flag)   /* Check map option */
                        {
                                /* Update disk map */
                        disk_map[scb.scb_disk][dir_entry -> de_userno]++;
                        }
                else            /* Display names */
                        {
                        conv_dfname(scb.scb_disk,dir_entry,file_name);
                        printf("%s  ",file_name);

                                /* Check if need to start new line */
                        if (!(dmcount % 4))
                                {
                                putchar('\n');
                                        if (++lcount > 18)
                                                {
                                                lcount = 0;
                                                printf("\nPress Space Bar to continue....");
                                                getchar();
                                                putchar('\n');
                                                }
                                        }
                                }
                        }
                } /* End of directory */
        } /* All disks searched */

if (map_flag)
{
printf("\n                Numbers show files in each user number.");
printf("\n                          --- User Numbers ---                    Dir. Entries");

dm_disp(disk_map,scb.scb_adisks);       /* Display disk map */
}

if (mcount == 0)
printf("\n --- File Not Found --- ");

bdos(SETDISK,cur_disk); /* Reset to current disk */
}


chk_use(argc)           /* check usage */
/* This function checks that the correct number of
   parameters has been specified, outputting instructions
   if not.
*/

/* Entry parameter */
int argc;       /* Count of the number of arguments on the command line */
{

/* The minimum value of argc is 1 (for the program name itself),
   so argc is always one greater than the number of parameters
   on the command line */

if (argc == 1 || argc > 3)
{
printf("\nUsage :");
printf("\n\tFIND d:filename.typ {NAMES}");
printf("\n\t     *:filename.typ (All disks)");
printf("\n\t     ABCD..OP:filename.typ (Selected Disks)");
printf("\n\tNAMES option shows actual names rather than map.");
exit();
}

}
```

Figure 11-5. (Continued)

SPACE — Show Used Disk Space

The SPACE utility shown in Figure 11-6 scans the specified logical disks and displays a disk map that shows, for each user number on each logical disk, how many Kbytes of storage have been used. It also displays the total number of Kbytes used and free on each logical disk.

Here is an example console dialog showing SPACE in operation:

```
P3B>space<CR>
SPACE Version 1.0 02/11/83 (Library 1.0)
Usage :
        SPACE *        (All disks)
        SPACE ABCD..OP (Selected Disks)

P3B>space *<CR>
SPACE Version 1.0 02/11/83 (Library 1.0)
Searching disk : A
Searching disk : B
Searching disk : C
                  Numbers show space used in kilobytes.
                          --- User Numbers ---          Space (Kb)
   0   1   2   3   4   5   ...  10  11  12  13  14  15 Used Free
A:  18 202     38                                       258 1196
B: 692 432 656 548  36                                 2364  996
C: 140                                                  140  204
```

```
#define VN "1.0 02/11/83"

/* SPACE -- This utility displays a map showing on the amount of space
   (expressed as relative percentages) occupied in each user number
   for each logical disk.It also shows the relative amount of space
   free. */

#include <LIBRARY.H>

struct _dirpb dir_pb;          /* Directory management parameter block */
struct _dir *dir_entry;        /* Pointer to directory entry */
struct _scb scb;               /* Search control block */
struct _dpb dpb;               /* CP/M's disk parameter block */

char file_name[20];            /* Formatted for display : un/d:FILENAME.TYP */

short cur_disk;          /* Current logical disk at start of program
                            NZ = show map of number of files */
int count;               /* Used to access the allocation block numbers
                            in each directory entry */
int user;                /* Used to access the disk map when calculating */

/* The array below is used to tabulate the results for each
   disk drive, and for each user number on the drive.
   In addition, two extra "users" have been added for "free"
   and "used" values.
*/
unsigned disk_map[16][18];     /* Disk A -> P, users 0 -> 15, free, used */
#define USED_COUNT 16          /* "User" number for used entities */
#define FREE_COUNT 17          /* "User" number for free entities */


main(argc,argv)
short argc;              /* Argument count */
char *argv[];            /* Argument vector (pointer to an array of chars.) */
{
```

Figure 11-6. SPACE.C, a utility that displays how much disk storage is used or available

```
printf("\nSPACE Version %s (Library %s)",VN,LIBVN);
chk_use(argc);                  /* Check usage */
cur_disk = bdos(GETDISK);       /* Get current default disk */

dm_clr(disk_map);               /* Reset disk map */

ssetscb(scb,argv[1]);           /* Special version : set disks,
                                   name, type */

for (scb.scb_disk = 0;          /* Starting with logical disk A: */
     scb.scb_disk < 16;         /* Until logical disk P: */
     scb.scb_disk++)            /* Move to next logical disk */
        {

        /* Check if current disk has been selected for search */
        if (!(scb.scb_adisks & (1 << scb.scb_disk)))
                continue;       /* No, so bypass this disk */

        printf("\nSearching disk : %c",(scb.scb_disk + 'A'));
        dir_pb.dp_disk = scb.scb_disk;  /* Set to disk to be searched */

        /* Set the directory to "closed", and force the get_nde
           function to open it  */
        dir_pb.dp_open = 0;

        /* While not at the end of the directory, set a pointer
           to the next entry in the directory */
        while (dir_entry = get_nde(dir_pb))
                {
                if (dir_entry -> de_userno == 0xE5)
                        continue;       /* Bypass inactive entries */

                for (count = 0;         /* Start with the first alloc. block */
                     count < dir_pb.dp_nabpde;  /* For number of alloc. blks. per dir. entry */
                     count++)
                        {
                        if (dir_pb.dp_nabpde == 8)      /* Assume 8 2-byte numbers */
                                {
                                disk_map[scb.scb_disk][dir_entry -> de_userno]
                                        += (dir_entry -> _dirab.de_long[count] > 0 ? 1 : 0);
                                }
                        else    /* Assume 16 1-byte numbers */
                                {
                                disk_map[scb.scb_disk][dir_entry -> de_userno]
                                        += (dir_entry -> _dirab.de_short[count] > 0 ? 1 : 0);
                                }
                        }       /* All allocation blocks processed */
                }       /* End of directory for this disk */


        /* Compute the storage used by multiplying the number of
           allocation blocks counted by the number of Kbytes in
           each allocation block. */

        for (user = 0;  /* Start with user 0 */
             user < 16; /* End with user 15 */
             user ++)   /* Move to next user number */
                {
                        /* Compute size occupied in Kbytes */
                disk_map[scb.scb_disk][user] *= dir_pb.dp_absize;
                        /* Build up sum for this disk */
                disk_map[scb.scb_disk][USED_COUNT] += disk_map[scb.scb_disk][user];
                }

        /* Free space = (# of alloc. blks * # of kbyte per blk)
                - used Kbytes
                - (directory entries * 32) / 1024 ... or divide by 32 */
        disk_map[scb.scb_disk][FREE_COUNT] = (dir_pb.dp_nab * dir_pb.dp_absize)
                - disk_map[scb.scb_disk][USED_COUNT]
                - (dir_pb.dp_nument >> 5);      /* Same as / 32 */
        }       /* All disks processed */


printf("\n                      Numbers show space used in kilobytes.");
printf("\n                              --- User Numbers ---                    Space (Kb)");

dm_disp(disk_map,scb.scb_adisks);       /* Display disk map */
```

Figure 11-6. (Continued)

```
bdos(SETDISK,cur_disk); /* Reset to current disk */
}

ssetscb(scb,ldisks)     /* Special version of set search control block */

/* This function sets up a search control block according
   to just the logical disks specified. The disk are specified as
   a single string of characters without any separators. An
   asterisk means "all disks." For example --

        ABGH    (disks A:, B:, G: and H: )
        *       (all disks for which SELDSK has tables)

   It sets the bit map according to which disks should be searched.
   For each selected disk, it checks to see if an error is generated
   when selecting the disk (i.e. if there are disk tables in the BIOS
   for the disk).
   The file name, type, and extent number are all set to "?" to match
   all possible entries in the directory. */

/* Entry parameters */
struct _scb *scb;       /* Pointer to search control block */
char *ldisks;           /* Pointer to the logical disks */

/* Exit parameters
   None.
*/
{
int disk;               /* Disk number currently being checked */
unsigned adisks;        /* Bit map for active disks */

adisks = 0;             /* Assume no disks to search */

if (*ldisks)            /* Some values specified */
        {
        if (*ldisks == '*')     /* Check if "all disks" */
                {
                adisks = 0xFFFF;        /* Set all bits */
                }
        else                    /* Set specific disks */
                {
                while(*ldisks)  /* Until end of disks reached */
                        {
                        /* Build the bit map by getting the next disk
                           id. (A - P), converting it to a number
                           in the range 0 - 15, and shifting a 1-bit
                           left that many places and OR ing it into
                           the current active disks.
                        */
                        adisks |= 1 << (toupper(*ldisks) - 'A');
                        ++ldisks;       /* Move to next character */
                        }
                }
        }
else    /* Use only current default disk */
        {
        /* Set just the bit corresponding to the current disk */
        adisks = 1 << bdos(GETDISK);
        }

        /* Set the user number, file name, type, and extent to "?"
           so that all active directory entries will match */
                /*      0123456789012       */
strcpy(&scb -> scb_userno,"?????????????");

        /* Make calls to the BIOS SELDSK routine to make sure that
           all of the active disk drives have disk tables for them
           in the BIOS. If they don't, turn off the corresponding
           bits in the bit map. */


for (disk = 0;          /* Start with disk A: */
     disk < 16;         /* Until disk P: */
     disk++)            /* Use next disk */
        {
        if ( !((1 << disk) & adisks))
                continue;               /* Avoid selecting unspecified disks */
```

Figure 11-6. (Continued)

```
        if (biosh(SELDSK,disk) == 0)    /* Make BIOS SELDSK call */
                {                       /* Returns 0 if invalid disk */
                /* Turn OFF corresponding bit in mask
                   by AND-ing it with bit mask having
                   all the other bits set = 1. */
                adisks &= ((1 << disk) ^ 0xFFFF);
                }
        }

scb -> scb_adisks = adisks;     /* Set bit map in scb */

} /* End ssetscb */



chk_use(argc)           /* Check usage */
/* This function checks that the correct number of
   parameters has been specified, outputting instructions
   if not. */

/* Entry parameter */
int argc;       /* Count of the number of arguments on the command line */
{

        /* The minimum value of argc is 1 (for the program name itself),
           so argc is always one greater than the number of parameters
           on the command line */

if (argc != 2)
        {
        printf("\nUsage :");
        printf("\n\tSPACE *         (All disks)");
        printf("\n\tSPACE ABCD..OP (Selected Disks)");
        exit();
        }
} /* End chk_use */
```

Figure 11-6. (Continued)

MOVE — Move Files Between User Numbers

The MOVE utility shown in Figure 11-7 moves files from one user number to another on the same logical disk. The movement is achieved by changing the user number in all the relevant directory entries. This is much faster than copying the files. It also avoids having multiple copies of the same file on the disk.

Here is a console dialog showing MOVE in operation:

```
P3B>move<CR>
MOVE Version 1.0 02/10/83 (Library 1.0)
Usage :
        MOVE d:filename.typ to_user {from_user} {NAMES}
             *:filename.typ (All disks)
             ABCD..OP:filename.typ (Selected Disks)
        NAMES option shows names of files moved.

P3B>dir *.com<CR>
B: ERASE    COM : FUNKEY  COM : DATE     COM : FIND     COM
B: SPACE    COM : UNERASE COM : MAKE     COM : MOVE     COM
B: TIME     COM : ASSIGN  COM : SPEED    COM : PROTOCOL COM

P3B>move *.com 0 names<CR>
MOVE Version 1.0 02/10/83 (Library 1.0)

Moving file(s)  3/B:????????.COM -> User 0.
```

```
O/B:ERASE    .COM  O/B:FUNKEY  .COM  O/B:DATE    .COM  O/B:FIND    .COM
O/B:SPACE    .COM  O/B:UNERASE .COM  O/B:MAKE    .COM  O/B:MOVE    .COM
O/B:TIME     .COM  O/B:ASSIGN  .COM  O/B:SPEED   .COM  O/B:PROTOCOL.COM

P3B>user 0<CR>
P0B>dir
B: ERASE    COM : FUNKEY  COM : DATE    COM : FIND     COM
B: SPACE    COM : UNERASE COM : MAKE    COM : MOVE     COM
B: TIME     COM : ASSIGN  COM : SPEED   COM : PROTOCOL COM
```

```
#define VN "1.0 02/10/83"

/* MOVE -- This utility transfers file(s) from one user number to
   another, but on the SAME logical disk. Files are not actually
   copied -- rather, their directory entries are changed. */

#include <LIBRARY.H>                    /

struct _dirpb dir_pb;           /* Directory management parameter block */
struct _dir *dir_entry;         /* Pointer to directory entry */
struct _scb scb;                /* Search control block */


#define DIR_BSZ 128             /* Directory buffer size */
char dir_buffer[DIR_BSZ];       /* Directory buffer */

char file_name[20];             /* Formatted for display : un/d:FILENAME.TYP */
short name_flag;                /* NZ to display names of files moved */

short cur_disk;                 /* Current logical disk at start of program */
int from_user;                  /* User number from which to move files */
int to_user;                    /* User number to which files will be moved */

int mcount;                     /* Match count (no. of file names matched) */
int dmcount;                    /* Per-disk match count */
int lcount;                     /* Line count (for lines displayed) */


main(argc,argv)
short argc;             /* Argument count */
char *argv[];           /* Argument vector (pointer to an array of chars.) */
{

printf("\nMOVE Version %s (Library %s)",VN,LIBVN);

chk_use(argc);                  /* Check usage */

to_user = atoi(argv[2]);        /* Convert user no. to integer */
        /* Set and check destination user number */
if(to_user > 15)
        {
        printf("\nError -- the destination user number cannot be greater than 15.");
        }

        /* Set the current user number */
from_user = bdos(GETUSER,0xFF);

        /* Check if source user number specified */
if (isdigit(argv[3][0]))
        {
                /* Set and check source user number */
        if((from_user = atoi(argv[3])) > 15)
                {
                printf("\nError -- the source user number cannot be greater than 15.");
                exit();
                }
                /* Set name suppress flag from parameter #4 */
        name_flag = usstrcmp("NAMES",argv[4]);
        }
else            /* No source user specified */
        {
```

Figure 11-7. MOVE.C, a utility program that changes files' user numbers

```
                /* Set name suppress flag from parameter #3 */
        name_flag = usstrcmp("NAMES",argv[3]);
        }

        /* To simplify the logic below, name_flag must be made
           NZ if it is equal to NAME_EQ, 0 if it is any other value */
name_flag = (name_flag == NAME_EQ ? 1 : 0);

if (to_user == from_user)       /* To = from */
        {
        printf("\nError - 'to' user number is the same as 'from' user number.");
        exit();
        }

        /* Set the search control block file name, type, user number,
           extent number, and length -- length matches user number, file
           name, and type. As the extent number does not enter into the
           comparison, all extents of a given file will be found. */
setscb(scb,argv[1],from_user,'?',13);

cur_disk = bdos(GETDISK);       /* Get current default disk */
lcount = dmcount = mcount = 0;  /* Initialize counts */

for (scb.scb_disk = 0;          /* Starting with logical disk A: */
     scb.scb_disk < 16;         /* Until logical disk P: */
     scb.scb_disk++)            /* Move to next logical disk */
        {
                /* Check if current disk has been selected for search */
        if (!(scb.scb_adisks & (1 << scb.scb_disk)))
                continue;       /* No, so bypass this disk */
                /* convert search user number and name for output */
        conv_dfname(scb.scb_disk,scb,file_name);
        printf("\n\nMoving file(s) %s -> User %d.",file_name,to_user);

        lcount++;               /* Update line count */

        dir_pb.dp_disk = scb.scb_disk;  /* Set to disk to be searched*/
        dmcount = 0;                    /* Reset disk matched count */

        if (name_flag)          /* If file names are to be displayed */
                putchar('\n');  /* Move to column 1 */

                /* Set the directory to "closed" to force the get_nde
                   function to open it. */
        dir_pb.dp_open = 0;

                /* While not at the end of the directory, set a pointer
                   to the next directory entry */
        while(dir_entry = get_nde(dir_pb))
                {
                        /* Match those entries that have the correct
                           user number, file name, type, and any
                           extent number. */
                if (
                    (dir_entry -> de_userno != 0xE5) &&
                    (comp_fname(scb,dir_entry) == NAME_EQ)
                   )
                        {

                        dir_entry -> de_userno = to_user;       /* Move to new user */
                                /* Request sector to be written back */
                        dir_pb.dp_write = 1;

                        mcount++;       /* Update matched counts */
                        dmcount++;      /* Per-disk count */

                        if (name_flag)  /* Check map option */
                                {
                                conv_dfname(scb.scb_disk,dir_entry,file_name);
                                        printf("%s  ",file_name);

                                /* Check if need to start new line */
                                if (!(dmcount % 4))
                                        {
                                        putchar('\n');
                                        if (++lcount > 18)
```

Figure 11-7. (Continued)

```
                                                        {
                                                        lcount = 0;
                                                        printf("\nPress Space Bar to continue....");
                                                        getchar();
                                                        putchar('\n');
                                                        }
                                                }
                                        }
                                }
                        }
                }

        if (mcount == 0)
                printf("\n --- No Files Moved --- ");

        bdos(SETDISK,cur_disk); /* Reset to current disk */
        }


        chk_use(argc)           /* Check usage */
        /* This function checks that the correct number of
           parameters has been specified, outputting instructions
           if not  */
        /* Entry parameter */
        int argc;       /* Count of the number of arguments on the command line */
        {

        /* The minimum value of argc is 1 (for the program name itself),
           so argc is always one greater than the number of parameters
           on the command line */

        if (argc == 1 || argc > 5)
                {
                printf("\nUsage :");
                printf("\n\tMOVE d:filename.typ to_user {from_user} {NAMES}");
                printf("\n\t     *:filename.typ (All disks)");
                printf("\n\t     ABCD..OP:filename.typ (Selected Disks)");
                printf("\n\tNAMES option shows names of files moved.");
                exit();
                }
        }
```

Figure 11-7. (Continued)

Other Utilities

The utility programs described in this section are by no means a complete set. You may want to develop many other specialized utility programs. Some possibilities are:

FILECOPY
: A more specialized version of PIP could copy ambiguously specified groups of files. Of special importance would be the ability to read a file containing the names of the files to be copied. A useful option would be the ability to detect the setting of the unused file attribute bit and copy only files that have been changed.

PROTECT/UNPROTECT
: This pair of utilities would allow you to "hide" files in user numbers greater than 15. Files so hidden could not be accessed other than by UNPROTECTing them, thereby moving them back into the normal user number range.

RECLAIM

This utility would read all sectors on a disk (using the BIOS). Any bad sectors encountered could then be logically removed by creating an entry in the file directory, with allocation block numbers that would effectively "reserve" the blocks containing the bad sectors.

OWNER

This utility, given a track or sector number, would access the directory and determine which file or files were using that part of the disk. This is useful if you have a bad sector or track on a disk. You then can determine which files have been damaged.

Utility Programs for the Enhanced BIOS

This section describes several utility programs that work with the enhanced BIOS shown in Figure 8-10. Several of these utilities work directly with the physical devices on the computer system, which can vary from computer to computer. The library header contains #define declarations for device numbers and names for physical devices (Figure 11-2, f and Figure 11-2, g).

These #define statements are used to build a physical-device code table. If you have more physical devices or want to change the names by which you refer to the devices, you will need to change these definitions.

All of these utilities share some common features in the way that they are invoked. If they are called without any parameters, they display instructions on the console regarding what parameters are available. If they are called with the word "SHOW" (or "S", "SH", and so forth) as a parameter, they display the current settings of whatever attribute the utility controls.

MAKE — Make Files "Invisible" or "Visible"

The MAKE utility shown in Figure 11-8 is designed to operate in conjunction with the public files option implemented in the enhanced BIOS of Figure 8-10. It has two modes of operation — making files "invisible" or "visible."

An invisible file is one in user 0 which has been set to Read-Only and System status. When the public files option is enabled, these files cannot be seen when you use the DIR command, nor can they be erased accidentally.

A visible file is one that has been set to Read/Write and Directory status.

When files are made invisible, they are transferred from the current user number to user 0. When files are made visible, they are transferred from user 0 to the current user number.

Here is an example console dialog showing MAKE in operation:

```
P3B>make<CR>
MAKE Version 1.0 02/12/83 (Library 1.0)
```

```
Usage :
        MAKE d:filename.typ INVISIBLE {NAMES}
                            VISIBLE
             *:filename.typ (All disks)
             ABCD..OP:filename.typ (Selected Disks)
        NAMES option shows names of files processed.

P3B>dir *.com<CR>
B: ERASE    COM : UNERASE  COM : ASSIGN   COM : PROTOCOL COM

P3B>make *.com invisible names<CR>
MAKE Version 1.0 02/12/83 (Library 1.0)

Moving files from User 3 to 0 and making them Invisible.
Searching disk : B

         0/B:ERASE   .COM made Invisible in User 0.
         0/B:UNERASE .COM made Invisible in User 0.
         0/B:ASSIGN  .COM made Invisible in User 0.
         0/B:PROTOCOL.COM made Invisible in User 0.

P3B>make erase.com visible names<CR>
MAKE Version 1.0 02/12/83 (Library 1.0)

Moving files from User 0 to 3 and making them Visible.
Searching disk : B

         3/B:ERASE   .COM made Visible in User 3.
```

```
#define VN "1.0 02/12/83"

/* MAKE - This utility is really two very similar programs;
   which one depends on the parameter specified on the command
   line.

   INVISIBLE finds all of the specified files, moves them
   to user number 0, and sets them to be System and Read Only
   status. These files can then be accessed from user numbers
   other than 0 when the public files feature is enabled in the
   BIOS.

   VISIBLE is the opposite in that the specified files are
   moved to the current user number and changed to Directory
   and Read/Write status. */

#include <LIBRARY.H>

struct _dirpb dir_pb;          /* Directory management parameter block */
struct _dir *dir_entry;        /* Pointer to directory entry */
struct _scb scb;               /* Search control block */
short to_user;                 /* User number to which files will be set */
short from_user;               /* User number from which files will be moved */

char file_name[20];            /* Formatted for display : un/d:FILENAME.TYP */
short name_flag;               /* NZ to display names of files moved */

short cur_disk;                /* Current logical disk at start of program */

int mcount;                    /* Match count (no. of file names matched) */

short invisible;               /* NZ when parameter specifies invisible */
char *operation;               /* Pointer to either "invisible" or "visible" */

main(argc,argv)
short argc;           /* Argument count */
char *argv[];         /* Argument vector (pointer to an array of chars.) */
```

Figure 11-8. MAKE.C, a utility that makes files "invisible" and protected or makes them "visible," accessible, and unprotected

```
{

printf("\nMAKE Version %s (Library %s)",VN,LIBVN);
chk_use(argc);                  /* Check usage */
cur_disk = bdos(GETDISK);       /* Get current default disk */
mcount = 0;                     /* Initialize count */


        /* Set the invisible flag according to the parameter */
invisible = usstrcmp("VISIBLE",argv[2]);

        /* Set the from_user and to_user numbers depending on which
           program is to be built, and the parameters specified. */

if (invisible)
        {
        from_user = bdos(GETUSER,0xFF); /* Get current user number */
        to_user = 0;    /* Always move files to user 0 */
        operation = "Invisible";        /* Set pointer to string */
        }
else    /* visible */
        {
        from_user = 0;                  /* Always move from user 0 */
        to_user = bdos(GETUSER,0xFF);   /* Get current user */
        operation = "Visible";          /* Set pointer to string */
        }

        /* Set search control block disks, name, type, user number,
           extent number, and number of bytes to compare -- in this
           case, match the "from" user, all extents. */
setscb(scb,argv[1],from_user,'?',13);   /* Set disks, name, type */

name_flag = usstrcmp("NAMES",argv[3]);  /* Set name-suppress flag from param. 3 */

        /* To simplify the logic below, name_flag must be made
           NZ if it is equal to NAME_EQ, 0 if it is any other value */
name_flag = (name_flag == NAME_EQ ? 1 : 0);


        /* Convert search user number and name for output */
conv_dfname(scb.scb_disk,scb,file_name);
printf("\n\nMoving files from User %d to %d and making them %s.",
        from_user,to_user,operation);

for (scb.scb_disk = 0;          /* Starting with logical disk A: */
     scb.scb_disk < 16;         /* Until logical disk P: */
     scb.scb_disk++)            /* Move to next logical disk */
        {

                /* Check if current disk has been selected for search */
        if (!(scb.scb_adisks & (1 << scb.scb_disk)))
                continue;       /* No -- so bypass this disk */

        printf("\nSearching disk : %c",(scb.scb_disk + 'A'));

        dir_pb.dp_disk = scb.scb_disk;  /* Set to disk to be searched*/

        if (name_flag)          /* If file names are to be displayed */
                putchar('\n');  /* Move to column 1 */


                /* Set the directory to "closed", and force the get_nde
                   function to open it. */
        dir_pb.dp_open = 0;

                /* While not at the end of the directory,
                   set a pointer to the next directory entry. */
        while(dir_entry = get_nde(dir_pb))
                {

                        /* Match those entries that have the correct
                           user number, file name, type, and any
                           extent number. */
                if (
                    (dir_entry -> de_userno != 0xE5) &&
                    (comp_fname(scb,dir_entry) == NAME_EQ)
                   )
                        {
```

Figure 11-8. (Continued)

```
                        mcount++;       /* Update matched counts */

                        if (invisible)
                                {       /* Set ms bits */
                                dir_entry -> de_fname[8] |= 0x80;
                                dir_entry -> de_fname[9] |= 0x80;
                                }
                        else    /* Visible */
                                {       /* Clear ms bits */
                                dir_entry -> de_fname[8] &= 0x7F;
                                dir_entry -> de_fname[9] &= 0x7F;
                                }

                                /* Move to correct user number */
                        dir_entry -> de_userno = to_user;

                                /* Indicate sector to be written back */
                        dir_pb.dp_write = 1;

                                /* Check if name to be displayed */
                        if (name_flag)
                                {
                                conv_dfname(scb.scb_disk,dir_entry,file_name);
                                printf("\n\t%s made %s in User %d.",
                                        file_name,operation,to_user);
                                }
                        }
                }       /* All directory entries processed */
        }       /* All disks processed */

if (mcount == 0)
        printf("\n --- No Files Processed --- ");

bdos(SETDISK,cur_disk); /* Reset to current disk */
}


chk_use(argc)           /* Check usage */
/* This function checks that the correct number of
   parameters has been specified, outputting instructions
   if not.
*/
/* Entry parameter */
int argc;       /* Count of the number of arguments on the command line */
{

        /* The minimum value of argc is 1 (for the program name itself),
           so argc is always one greater than the number of parameters
           on the command line */

if (argc == 3 || argc == 4)
        return;
else
        {
        printf("\nUsage :");
        printf("\n\tMAKE d:filename.typ INVISIBLE {NAMES}");
        printf("\n\t                    VISIBLE");
        printf("\n\t     *:filename.typ (All disks)");
        printf("\n\t     ABCD..OP:filename.typ (Selected Disks)");
        printf("\n\tNAMES option shows names of files processed.");
        exit();
        }

}
```

Figure 11-8. (Continued)

SPEED — Set Baud Rates

The SPEED utility shown in Figure 11-9 sets the baud rate for a specific serial device. Here is an example console dialog that shows several of the options:

```
P3B>speed<CR>
SPEED 1.0 02/17/83
The SPEED utility sets the baud rate speed for each physical device.
Usage is :  SPEED physical-device baud-rate, or
            SPEED SHOW     (to show current settings)

Valid physical devices are:
                TERMINAL
                PRINTER
                MODEM

Valid baud rates are:
                300
                600
                1200
                2400
                4800
                9600
                19200


P3B>speed show<CR>
SPEED 1.0 02/17/83
Current Baud Rate settings are :
        TERMINAL set to 9600 baud.
        PRINTER set to 9600 baud.
        MODEM set to 9600 baud.


P3B>speed m 19<CR>
SPEED 1.0 02/17/83
Current Baud Rate settings are :
        TERMINAL set to 9600 baud.
        PRINTER set to 9600 baud.
        MODEM set to 19200 baud.


P3B>speed xyz 12<CR>
SPEED 1.0 02/17/83
Physical Device 'XYZ' is invalid or ambiguous.
Legal Physical Devices are :
                TERMINAL
                PRINTER
                MODEM
```

```
#define VN "\nSPEED 1.0 02/17/83"

/* This utility sets the baud rate speed for each of the physical
   devices. */

#include <LIBRARY.H>

struct _ct ct_pdev[MAXPDEV + 2];        /* Physical device table */

        /* Hardware specific items */
```

Figure 11-9. SPEED.C, a utility that sets the baud rate for a specific device

```
                                        /* Baud rates for serial ports */
#define B300    0x35                    /* 300 baud */
#define B600    0x36                    /* 600 baud */
#define B1200   0x37                    /* 1200 baud */
#define B2400   0x3A                    /* 2400 baud */
#define B4800   0x3C                    /* 4800 baud */
#define B9600   0x3E                    /* 9600 baud */
#define B19200  0x3F                    /* 19200 baud */
struct _ct ct_br[10];   /* Code table for baud rates (+ spare entries) */


        /* Parameters on the command line */
#define PDEV argv[1]    /* Physical device */
#define BAUD argv[2]    /* Baud rate */


main(argc,argv)
int argc;
char *argv[];
{
printf(VN);     /* Display sign-on message */
setup();        /* Set up code tables */
chk_use(argc);  /* Check correct usage */

        /* Check if request to show current settings */
if (usstrcmp("SHOW",argv[1]))
        {               /* No -- assume setting is required */
        set_baud(get_pdev(PDEV),get_baud(BAUD)); /* Set baud rate */
        }

show_baud();            /* Display current settings */

} /* end of program */

setup()                 /* set up the code tables for this program */
{
        /* Initialize the physical device table */
ct_init(ct_pdev[0],T_DEVN,PN_T);        /* Terminal */
ct_init(ct_pdev[1],P_DEVN,PN_P);        /* Printer */
ct_init(ct_pdev[2],M_DEVN,PN_M);        /* Modem */
ct_init(ct_pdev[3],CT_SNF,"*"); /* Terminator */

        /* Initialize the baud rate table */
ct_init(ct_br[0],B300,"300");
ct_init(ct_br[1],B600,"600");
ct_init(ct_br[2],B1200,"1200");
ct_init(ct_br[3],B2400,"2400");
ct_init(ct_br[4],B4800,"4800");
ct_init(ct_br[5],B9600,"9600");
ct_init(ct_br[6],B19200,"19200");
ct_init(ct_br[7],CT_SNF,"*");   /* Terminator */
}

unsigned
get_pdev(ppdev)         /* Get physical device */
/* This function returns the physical device code
   specified by the user in the command line. */
char *ppdev;            /* Pointer to character string */
{
unsigned retval;                        /* Return value */

retval = ct_parc(ct_pdev,ppdev);        /* Get code for ASCII string */
if (retval == CT_SNF)           /* If string not found */
        {
        printf("\n\007Physical Device '%s' is invalid or ambiguous.",
                ppdev);
        printf("\nLegal Physical Devices are : ");
        ct_disps(ct_pdev);      /* Display all values */
        exit();
        }
return retval;                  /* Return code */
}

unsigned
get_baud(pbaud)
/* This function returns the baud rate time constant for
   the baud rate specified by the user in the command line   */
```

Figure 11-9. (Continued)

```
char *pbaud;            /* Pointer to character string */
{
unsigned retval;                        /* Return value */
retval = ct_parc(ct_br,pbaud);  /* Get code for ASCII string */
if (retval == CT_SNF)           /* If string not found */
        {
        printf("\n\007Baud Rate '%s' is invalid or ambiguous.",
                pbaud);
        printf("\nLegal Baud Rates are : ");
        ct_disps(ct_br);        /* Display all values */
        exit();
        }
return retval;          /* Return code */
}

set_baud(pdevc,baudc)   /* Set the baud rate of the specified device */
int pdevc;              /* Physical device code */
short baudc;            /* Baud rate code */
                        /* On some systems this may have to be a
                           two-byte (unsigned) value  */
{
short *baud_rc;         /* Pointer to the baud rate constant */
                        /* On some systems this may have to be a
                           two-byte (unsigned) value  */
/* Note: the respective codes for accessing the baud rate constants
   via the get_cba (get configuration block address) function are:
        Device #0 = 19,  #1 = 21, #2 = 23. This function uses this
   mathematical relationship  */

        /* Set up pointer to the baud rate constant */
baud_rc = get_cba(CB_DO_BRC + (pdevc << 1));

        /* Then set the baud rate constant */
*baud_rc = baudc;

        /* Then call the BIOS initialization routine */
bios(CIOINIT,pdevc);
}

show_baud()             /* Show current baud rate */
{

int pdevn;              /* Physical device number */
short baudc;            /* Baud rate code */
                        /* On some systems this may have to be a
                           two-byte (unsigned) value  */
short *baud_rc;         /* Pointer to the baud rate constant */
                        /* On some systems this may have to be a
                           two-byte (unsigned) value  */
/* Note: the respective codes for accessing the baud rate constants
   via the get_cba (get configuration block address) function are:
        Device #0 = 19,  #1 = 21, #2 = 23. This function uses this
   mathematical relationship  */

printf("\nCurrent baud rate settings are :");


for (pdevn = 0; pdevn <= MAXPDEV; pdevn ++)     /* All physical devices */
        {
                /* Set up pointer to the baud rate constant --
                   the code for the get_cba function is computed
                   by adding the physical device number *2 to
                   the Baud Rate code for device #0 */

        baud_rc = get_cba(CB_DO_BRC + (pdevn << 1));

                /* Then set the baud rate constant */
        baudc = *baud_rc;

        printf("\n\t%s set to %s baud.",
                ct_strc(ct_pdev,pdevn), /* Get ptr. to device name */
                ct_strc(ct_br,baudc) ); /* Get ptr. to baud rate */
        }
}

chk_use(argc)           /* Check correct usage */
int argc;               /* Argument count */
{
```

Figure 11-9. (Continued)

```
if (argc == 1)
        {
        printf("\nThe SPEED utility sets the baud rate speed for each physical device.");
        printf("\nUsage is :  SPEED physical-device baud rate, or");
        printf("\n             SPEED SHOW     (to show current settings)");
        printf("\n\nValid physical devices are: ");
        ct_disps(ct_pdev);
        printf("\nValid baud rates are: ");
        ct_disps(ct_br);
        exit();
        }
}
```

Figure 11-9. (Continued)

PROTOCOL — Set Serial Line Protocols

The PROTOCOL utility shown in Figure 11-10 is used to set the protocol for a specific serial device.

The drivers for each physical device can support several serial line protocols. The protocols are divided into two groups, depending on whether they apply to data output by or input to the computer.

Note that the output DTR and input RTS protocols can coexist with other protocols. The strategy is first to set the required character-based protocol and then to set the DTR/RTS protocol. There is an example of this in the following console dialog:

```
P3B>protocol<CR>
PROTOCOL Vn 1.0 02/17/83
PROTOCOL sets the physical device's serial protocols.
       PROTOCOL physical-device direction protocol {message-length}

Legal physical devices are :
              TERMINAL
              PRINTER
              MODEM

Legal direction/protocols are :
              Output DTR
              Output XON
              Output ETX
              Input RTS
              Input XON

       Message length can be specifed with Output ETX.


P3B>protocol show<CR>
PROTOCOL Vn 1.0 02/17/83
       Protocol for TERMINAL - None.
       Protocol for PRINTER - Output XON
       Protocol for MODEM - Input RTS


P3B>protocol m o e 128<CR>
PROTOCOL Vn 1.0 02/17/83
       Protocol for TERMINAL - None.
       Protocol for PRINTER - Output XON
```

```
                Protocol for MODEM - Output ETX  Message Length 128 bytes.


        P3B>protocol m o d<CR>
        PROTOCOL Vn 1.0 02/17/83
                Protocol for TERMINAL - None.
                Protocol for PRINTER - Output XON
                Protocol for MODEM - Output DTR Output ETX  Message Length
                  128 bytes.
```

```
#define VN "\nPROTOCOL Vn 1.0 02/17/83"
/* PROTOCOL -- This utility sets the serial port protocol for the
   specified physical device. Alternatively, it displays the
   current protocols for all of the serial devices. */

#include <LIBRARY.H>

        /* Code tables used to relate ASCII strings to code values */
struct _ct ct_iproto[3];        /* Code table for input protocols */
struct _ct ct_oproto[4];        /* Code table for output protocols */
struct _ct ct_dproto[7];        /* Code table for displaying protocols */
struct _ct ct_pdev[MAXPDEV + 2];/* Physical device table */
struct _ct ct_io[3];            /* Input, output */


        /* Parameters on the command line */
#define PDEV argv[1]    /* Physical device */
#define IO argv[2]      /* Input/output */
#define PROTO argv[3]   /* Protocol */
#define PROTOL argv[4]  /* Protocol message length */


main(argc,argv)
int argc;
char *argv[];
{
printf(VN);     /* Display sign-on message */
setup();        /* Set up code tables */
chk_use(argc);  /* Check correct usage */

        /* Check if request to show current settings */
if (usstrcmp("SHOW",argv[1]))
        {               /* No -- assume a set is required */
        set_proto(get_pdev(PDEV),       /* Physical device */
                        /* Input/output and protocol */
                get_proto(get_io(IO),PROTO),
                PROTOL);        /* Protocol message length */
        }
show_proto();

} /* end of program */

setup()                 /* Set up the code tables for this program */
{
        /* Initialize the physical device table */
ct_init(ct_pdev[0],0,PN_T);     /* Terminal */
ct_init(ct_pdev[1],1,PN_P);     /* Printer */
ct_init(ct_pdev[2],2,PN_M);     /* Modem */
ct_init(ct_pdev[3],CT_SNF,"*"); /* Terminator */

        /* Initialize the input/output table */
ct_init(ct_io[0],0,"INPUT");
ct_init(ct_io[1],1,"OUTPUT");
ct_init(ct_io[2],CT_SNF,"*");           /* Terminator */

        /* Initialize the output protocol table */
ct_init(ct_oproto[0],DT_ODTR,"DTR");
ct_init(ct_oproto[1],DT_OXON,"XON");
ct_init(ct_oproto[2],DT_OETX,"ETX");
```

Figure 11-10. PROTOCOL.C, a utility that sets the protocol governing input and output of a specified serial device

```
ct_init(ct_oproto[3],CT_SNF,"*");       /* Terminator */

        /* Initialize the input protocol table */
ct_init(ct_iproto[0],DT_IRTS,"RTS");
ct_init(ct_iproto[1],DT_IXON,"XON");
ct_init(ct_iproto[2],CT_SNF,"*");       /* Terminator */

        /* Initialize the display protocol */
ct_init(ct_dproto[0],DT_ODTR,"Output DTR");
ct_init(ct_dproto[1],DT_OXON,"Output XON");
ct_init(ct_dproto[2],DT_OETX,"Output ETX");
ct_init(ct_dproto[3],DT_IRTS,"Input RTS");
ct_init(ct_dproto[4],DT_IXON,"Input XON");
ct_init(ct_dproto[5],CT_SNF,"*");
}

unsigned
get_pdev(ppdev)         /* Get physical device */
/* This function returns the physical device code
   specified by the user in the command line. */
char *ppdev;            /* Pointer to character string */
{
unsigned retval;        /* Return value */

retval = ct_parc(ct_pdev,ppdev);/* Get code for ASCII string */
if (retval == CT_SNF)           /* If string not found */
        {
        printf("\n\007Physical Device '%s' is invalid or ambiguous.",
                ppdev);
        printf("\nLegal Physical Devices are : ");
        ct_disps(ct_pdev);      /* Display all values */
        exit();
        }
return retval;                  /* Return code */
}

unsigned
get_io(pio)             /* Get input/output parameter */
char *pio;              /* Pointer to character string */
{
unsigned retval;                /* Return value */

retval = ct_parc(ct_io,pio);    /* Get code for ASCII string */
if (retval == CT_SNF)           /* If string not found */
        {
        printf("\n\007Input/Output direction '%s' is invalid or ambiguous.",
                pio);
        printf("\nLegal values are : ");
        ct_disps(ct_io);        /* Display all values */
        exit();
        }
return retval;                  /* Return code */
}

unsigned

get_proto(output,pproto)
/* This function returns the protocol code for the
   protocol specified by the user in the command line. */
int output;             /* =1 for output, =0 for input */
char *pproto;           /* Pointer to character string */

{
unsigned retval;                /* Return value */

if (output)                     /* OUTPUT specified */
        {
                /* Get code for ASCII string */
        retval = ct_parc(ct_oproto,pproto);
        if (retval == CT_SNF)           /* If string not found */
                {
                printf("\n\007Output Protocol '%s' is invalid or ambiguous.",
        pproto);
                printf("\nLegal Output Protocols are : ");
                ct_disps(ct_oproto);    /* Display valid protocols */
                exit();
                }
```

Figure 11-10. (Continued)

```
        }
else                                /* INPUT specified */
        {
                /* Get code for ASCII string */
        retval = ct_parc(ct_iproto,pproto);
        if (retval == CT_SNF)           /* If string not found */
                {
                printf("\n\007Input Protocol '%s' is invalid or ambiguous.",
        pproto);
                printf("\nLegal Input Protocols are : ");
                ct_disps(ct_iproto);    /* Display valid protocols */
                exit();
                }
        }
return retval;                      /* Return code */
}


set_proto(pdevc,protoc,pplength)/* Set the protocol for physical device */
int pdevc;                          /* Physical device code */
unsigned protoc;                    /* Protocol byte */
char *pplength;                     /* Pointer to protocol length */
{
struct _ppdt
{
char *pdt[16];              /* Array of 16 pointers to the device tables */
} ;
struct _ppdt *ppdt;                 /* Pointer to the device table array */
struct _dt *dt;                     /* Pointer to a device table */

ppdt = get_cba(CB_DTA); /* Set pointer to array of pointers */
dt = ppdt -> pdt[pdevc];

if (!dt)                    /* Check if pointer in array is valid */
        {
        printf("\nError -- Array of Device Table Addresses is not set for device #%d.",
                pdevc);
        exit();
        }

if (protoc & 0x8000)        /* Check if protocol byte to be set
                               directly or to be OR ed in */
        {                   /* OR ed */
        dt -> dt_st1 |= (protoc & 0x7F);
        }
else
        {                   /* Set directly */
        dt -> dt_st1 = (protoc & 0x7F);
        }

if ((protoc & 0x7F) == DT_OETX) /* If ETX/ACK, check for message
                                   length */
        {
        if (isdigit(*pplength))         /* Check if length present */
                {
                        /* Convert length to binary and set device
                           table field. */
                dt -> dt_etxml = atoi(pplength);
                }
        }
}


show_proto()                /* Show the current protocol settings */
{
struct _ppdt
{
char *pdt[16];              /* Array of 16 pointers to the device tables */
} ;
struct _ppdt *ppdt;                 /* Pointer to the device table array */
struct _dt *dt;                     /* Pointer to a device table */
int pdevc;                          /* Physical device code */
struct _ct *dproto;                 /* Pointer to display protocols */

ppdt = get_cba(CB_DTA); /* Set pointer to array of pointers */

        /* For all physical devices */
```

Figure 11-10. (Continued)

```
for (pdevc = 0; pdevc <= MAXPDEV; pdevc++)
        {
                /* Set pointer to device table */
        dt = ppdt -> pdt[pdevc];

        if (dt) /* Check if pointer in array is valid */
                {
                printf("\n\tProtocol for %s - ",ct_strc(ct_pdev,pdevc));
                        /* Check if any protocols set */
                if (!(dt -> dt_st1 & ALLPROTO))
                        {
                        printf("None.");
                        continue;
                        }

                        /* Set pointer to display protocol table */
                dproto = ct_dproto;
                while (dproto -> _ct_code != CT_SNF)
                        {
                                /* Check if protocol bit set */
                        if (dproto -> _ct_code & dt -> dt_st1)
                                {       /* Display protocol */
                                printf("%s ",dproto -> _ct_sp);
                                }
                        ++dproto;       /* Move to next entry */
                        }
                        /* Check if ETX/ACK protocol and
                           message length to be displayed */
                if (dt -> dt_st1 & DT_OETX)
                        printf(" Message length %d bytes.",
                                dt -> dt_etxml);
                }
        }
}


chk_use(argc)           /* Check for correct usage */
int argc;               /* Argument count on commmand line */
{
if (argc == 1)
        {
        printf("\nPROTOCOL sets the physical device's serial protocols.");
        printf("\n\tPROTOCOL physical-device direction protocol {message-length}");
        printf("\n\nLegal physical devices are :");
        ct_disps(ct_pdev);
        printf("\nLegal direction/protocols are :");
        ct_disps(ct_dproto);
        printf("\n\tMessage length can be specifed with Output ETX.\n");
        exit();
        }
}
```

Figure 11-10. (Continued)

ASSIGN — Assign Physical to Logical Devices

The ASSIGN utility shown in Figure 11-11 sets the necessary bits in the physical input/output redirection bits in the BIOS. It assigns a logical device's input and output to physical devices. Input can only be derived from a single physical device, while output can be directed to multiple devices.

Here is an example console dialog showing ASSIGN in action:

```
P3B>assign<CR>
ASSIGN Vn 1.0 02/17/83
ASSIGN sets the Input/Output redirection.
        ASSIGN logical-device INPUT physical-device
        ASSIGN logical-device OUTPUT physical-dev1 {phy_dev2..}
        ASSIGN SHOW    (to show current assignments)
```

```
Legal logical devices are :
                CONSOLE
                AUXILIARY
                LIST

Legal physical devices are :
                TERMINAL
                PRINTER
                MODEM


P3B>assign show<CR>
ASSIGN Vn 1.0 02/17/83
Current Device Assignments are :
        CONSOLE INPUT is assigned to -  TERMINAL
        CONSOLE OUTPUT is assigned to -  TERMINAL
        AUXILIARY INPUT is assigned to -  MODEM
        AUXILIARY OUTPUT is assigned to -  MODEM
        LIST INPUT is assigned to -  PRINTER
        LIST OUTPUT is assigned to -  PRINTER

P3B>assign a o t m p<CR>
ASSIGN Vn 1.0 02/17/83
Current Device Assignments are :
        CONSOLE INPUT is assigned to -  TERMINAL
        CONSOLE OUTPUT is assigned to -  TERMINAL
        AUXILIARY INPUT is assigned to -  MODEM
        AUXILIARY OUTPUT is assigned to -  TERMINAL PRINTER MODEM
        LIST INPUT is assigned to -  PRINTER
        LIST OUTPUT is assigned to -  PRINTER
```

```
#define VN "\nASSIGN Vn 1.0 02/17/83"

#include <LIBRARY.H>

struct _ct ct_pdev[MAXPDEV + 2];        /* Physical device table */


        /* Names of logical devices */
#define LN_C    "CONSOLE"
#define LN_A    "AUXILIARY"
#define LN_L    "LIST"
struct _ct ct_ldev[4];          /* Logical device table */

struct _ct ct_io[3];            /* Input, output */

        /* Parameters on the command line */
#define LDEV argv[1]    /* Logical device */
#define IO argv[2]      /* Input/output */


main(argc,argv)
int argc;
char *argv[];
{

printf(VN);     /* Display sign-on message */
setup();        /* Set up code tables */
chk_use(argc);  /* Check correct usage */

        /* Check if request to show current settings */
if (usstrcmp("SHOW",argv[1]))
        {               /* No, assume a set is required */
```

Figure 11-11. ASSIGN.C, a utility that assigns a logical device's input and output to two physical devices

```
                /* NOTE : the number of physical devices to
                   process is given by argc - 3 */
        set_assign(get_ldev(LDEV),get_io(IO),argc - 3,argv);
        }
show_assign();

}

setup()                 /* Set up the code tables for this program */
{
        /* Initialize the physical device table */
ct_init(ct_pdev[0],0,PN_T);     /* Terminal */
ct_init(ct_pdev[1],1,PN_P);     /* Printer */
ct_init(ct_pdev[2],2,PN_M);     /* Modem */
ct_init(ct_pdev[3],CT_SNF,"*"); /* Terminator */

        /* Initialize the logical device table */
ct_init(ct_ldev[0],0,LN_C);     /* Terminal */
ct_init(ct_ldev[1],1,LN_A);     /* Auxiliary */
ct_init(ct_ldev[2],2,LN_L);     /* List */
ct_init(ct_ldev[3],CT_SNF,"*"); /* Terminator */

        /* Initialize the input/output table */
ct_init(ct_io[0],0,"INPUT");
ct_init(ct_io[1],1,"OUTPUT");
ct_init(ct_io[2],CT_SNF,"*");           /* Terminator */

}

unsigned
get_ldev(pldev)         /* Get logical device */
/* This function returns the logical device code
   specified by the user in the command line. */
char *pldev;            /* Pointer to character string */
{
unsigned retval;                        /* Return value */
retval = ct_parc(ct_ldev,pldev);        /* Get code for ASCII string */
if (retval == CT_SNF)           /* If string not found */
        {
        printf("\n\007Logical device '%s' is invalid or ambiguous.",
                pldev);
        printf("\nLegal logical devices are : ");
        ct_disps(ct_ldev);      /* Display all values */
        exit();
        }
return retval;                  /* Return code */
}

unsigned
get_io(pio)             /* Get input/output parameter */
char *pio;              /* Pointer to character string */
{
unsigned retval;                /* Return value */

retval = ct_parc(ct_io,pio);    /* Get code for ASCII string */
if (retval == CT_SNF)           /* If string not found */
        {
        printf("\n\007Input/output direction '%s' is invalid or ambiguous.",
                pio);
        printf("\nLegal values are : ");
        ct_disps(ct_io);        /* Display all values */
        exit();
        }
return retval;                  /* Return code */
}

set_assign(ldevc,output,argc,argv)      /* Set assignment (I/O redirection) */
int ldevc;                      /* Logical device code */
int output;                     /* I/O redirection code */
int argc;                       /* count of arguments to process */
char *argv[];                   /* Replica of parameter to main function */
{
unsigned *redir;                /* Pointer to redirection word */
int pdevc;                      /* Physical device code */
unsigned rd_val;                /* Redirection value */

        /* Get the address of the I/O redirection word.
```

Figure 11-11. (Continued)

```
            This code assumes that get_cba code values
            are ordered:
                 Device #0, input & output
                 Device #1, input & output
                 Device #2, input & putput

            The get_cba code is computed by multiplying the
            logical device code by 2 (that is, shift left 1)
            and added onto the code for Device #0, input
            Then the output variable (0 = input, 1 = output)
            is added on   */
redir = get_cba(CB_CI + (ldevc << 1) + output);

rd_val = 0;     /* Initialize redirection value */

        /* For output, assignment can be made to several physical
           devices, so this code may be executed several times  */
do
        {
                /* Get code for ASCII string */
                /* NOTE: the physical device parameters start
                   with parameter #3 (argv[3]). However argc
                   is a decreasing count of the number of physical
                   devices to be processed, Therefore, argc + 2
                   causes them to be processed in reverse order
                   (i.e. from right to left on the command line) */

        pdevc = ct_parc(ct_pdev,argv[argc + 2]);

        if (pdevc == CT_SNF)            /* If string not found */
                {
                printf("\n\007Physical device '%s' is invalid or ambiguous.",
                argv[argc + 2]);
                printf("\nLegal physical devices are : ");
                ct_disps(ct_pdev);      /* Display all values */
                exit();
                }
                /* Repeat this loop for as long as there are
                   more parameters (for output only) */
        else
                {
                /* Build new redirection value by OR ing in
                   a one-bit shifted left pdevc places. */
                rd_val |= (1 << pdevc);
                }
        } while (--argc && output);

*redir = rd_val;        /* Set the value into the config. block */
}

show_assign()                   /* Show current baud rate */
{
int rd_code;                    /* Redirection code for get_cba */
int ldevn;                      /* Logical device number */
int pdevn;                      /* Physical device number */
unsigned rd_val;                /* Redirection value */
unsigned *prd_val;              /* Pointer to the redirection value */

/* Note: the respective codes for accessing the redirection values
   via the get_cba (get configuration block address) function are:
        Device #0 console input -- 5
        Device #0 console putput -- 6
        Device #1 auxiliary input -- 7
        Device #1 auxiliary output -- 8
        Device #2 list input -- 9
        Device #2 list output -- 10

   This function uses this mathematical relationship  */

printf("\nCurrent device assignments are :");

        /* For all get_cba codes */
for (rd_code = CB_CI; rd_code <= CB_LO; rd_code++)
        {
                /* Set pointer to redirection value */
        prd_val = get_cba(rd_code);
                /* Get the input redirection value */
```

Figure 11-11. (Continued)

```
        rd_val = *prd_val;       /* This also performs byte reversal */

                /* Display device name. The rd_code is converted to a
                   device number by subtracting the first code number
                   from it and dividing by 2 (shift right one place).
                   The input/output direction is derived from the
                   least significant bit of the rd_code. */

        printf("\n\t%s %s is assigned to - ",
                ct_strc(ct_ldev,(rd_code - CB_CI) >> 1),
                ct_strc(ct_io,((rd_code & 0x01) ^ 1)));

                /* For all physical devices */
        for (pdevn = 0; pdevn < 16; pdevn++)
                {
                        /* Check if current physical device is assigned
                           by AND ing with a 1-bit shifted left pdevn times */
                if (rd_val & (1 << pdevn))      /* Is device active? */
                        {       /* Display physical device name */
                        printf(" %s",ct_strc(ct_pdev,pdevn) );
                        }
                }

        }
}


chk_use(argc)           /* Check for correct usage */
int argc;               /* Argument count on commmand line */
{
if (argc == 1)
        {
        printf("\nASSIGN sets the Input/Output redirection.");
        printf("\n\tASSIGN logical-device INPUT physical-device");
        printf("\n\tASSIGN logical-device OUTPUT physical-dev1 {phy_dev2..}");
        printf("\n\tASSIGN SHOW     (to show current assignments)");
        printf("\n\nLegal logical devices are :");
        ct_disps(ct_ldev);
        printf("\nLegal physical devices are :");
        ct_disps(ct_pdev);
        exit();
        }
}
```

Figure 11-11. (Continued)

DATE — Set the System Date

The DATE utility shown in Figure 11-12 sets the system date in the configuration block, along with a flag that indicates that the DATE utility has been used. Other utility programs can use this flag as a primitive test of whether the system date is current.

Here is an example console dialog:

```
P3B>date<CR>
DATE Vn 1.0 02/18/83
DATE sets the system date. Usage is :
       DATE mm/dd/yy
       DATE SHOW (to display current date)

P3B>date show<CR>
DATE Vn 1.0 02/18/83
       Current Date is 12/18/82

P3B>date 2/23/83<CR>
DATE Vn 1.0 02/18/83
       Current Date is 02/23/83
```

```
#define VN "\nDATE Vn 1.0 02/18/83"

/* This utility accepts the current date from the command tail,
   validates it, and set the internal system date in the BIOS.
   Alternatively, it can be requested just to display the current
   system date. */

#include <LIBRARY.H>

char *date;             /* Pointer to the date in the config. block */
char *date_flag;        /* Pointer to date-set flag */
int mm,dd,yy;           /* Variables to hold month, day, year */
int mcount;             /* Match count of numeric values entered */
int count;              /* Count used to add leading 0's to date */

main(argc,argv)
int argc;
char *argv[];
{
printf(VN);             /* Display sign-on message */
date = get_cba(CB_DATE);        /* Set pointer to date */
date_flag = get_cba(CB_DTFLAGS);/* Set pointer to date-set flag */

if (argc != 2)          /* Check if help requested (or needed) */
        show_use();     /* Display correct usage and exit */

if (usstrcmp("SHOW",argv[1]))   /* Check if not SHOW option */
        {
                /* Convert specified time into month, day, year */
        mcount = sscanf(argv[1],"%d/%d/%d",&mm,&dd,&yy);
        if (mcount != 3)                /* Input not numeric */
                show_use();     /* Display correct usage and exit */

                /* NOTE: The following validity checking is
                   simplistic, but could be expanded to accommodate
                   more context-sensitive checking: days in the month,
                   leap years, etc. */
        if (mm > 12 || mm < 1)  /* Check valid month, day, year */
                {
                printf("\nMonth = %d is illegal.",mm);
                show_use();     /* Display correct usage and exit */
                }
        if (dd > 31 || dd < 1)
                {
                printf("\nDay = %d is illegal.",dd);
                show_use();     /* Display correct usage and exit */
                }
        if (yy > 90 || yy < 83) /* <=== NOTE ! */
                {
                printf("\nYear = %d is illegal.",yy);
                show_use();     /* Display correct usage and exit */
                }

                /* Convert integers back into a formatted string */
        sprintf(date,"%2d/%2d/%2d",mm,dd,yy);
        date[8] = 0x0A;         /* Terminate with line feed */
        date[9] = '\0';         /* New string terminator */

                /* Change " 1/ 2/ 3" into "01/02/03" */
        for (count = 0; count < 7; count+=3)
                {
                if (date[count] == ' ')
                        date[count] = '0';
                }

                /* Turn flag on to indicate that user has set date */
        *date_flag |= DATE_SET;
        }
printf("\n\tCurrent Date is %s",date);
}

show_use()              /* Display correct usage and exit */
{
printf("\nDATE sets the system date. Usage is :");
printf("\n\tDATE mm/dd/yy");
printf("\n\tDATE SHOW (to display current date)\n");
exit();
}
```

Figure 11-12. DATE.C, a utility that makes the current date part of the system

TIME — Set the System Time

The TIME utility shown in Figure 11-13 sets the current system time. Like DATE, TIME sets a flag so that other utilities can test that the system time is likely to be current.

Here is an example console dialog:

```
P3B>time<CR>
TIME Vn 1.0 02/18/83
TIME sets the system time. Usage is :
        TIME hh{:mm{:ss}}
        TIME SHOW (to display current time)


P3B>time show<CR>
TIME Vn 1.0 02/18/83
        Current Time is 13:08:44


P3B>time 5:47<CR>
TIME Vn 1.0 02/18/83
        Current Time is 05:47:00
```

```
#define VN "\nTIME Vn 1.0 02/18/83"

/* This utility accepts the current time from the command tail,
   validates it, and sets the internal system time in the BIOS.
   Alternatively, it can just display the current system time. */

#include <LIBRARY.H>

char *time;             /* Pointer to the time in the config. block */
char *time_set;         /* Pointer to the time set flag */
int hh,mm,ss;           /* Variables to hold hours, minutes, seconds */
int mcount;             /* Match count of numeric values entered */
int count;              /* Count used to add leading zeros to time */

main(argc,argv)
int argc;
char *argv[];
{
printf(VN);             /* Display sign-on message */
time = get_cba(CB_TIMEA);       /* Set pointer to time */
time_flag = get_cba(CB_DTFLAGS);        /* Set pointer to the
                                           time-set flag */
hh = mm = ss = 0;       /* Initialize the time if seconds or
                           minutes are not specified */

if (argc != 2)          /* Check if help requested (or needed) */
        show_use();     /* Display correct usage and exit */

if (usstrcmp("SHOW",argv[1]))   /* Check if not SHOW option */
        {
                /* Convert time into hours, minutes, seconds */
        mcount = sscanf(argv[1],"%d:%d:%d",&hh,&mm,&ss);
        if (!mcount)            /* Input not numeric */
                show_use();     /* Display correct usage and exit */

        if (hh > 12)            /* Check valid hours, minutes, seconds */
                {
                printf("\n\007Hours = %d is illegal.",hh);
                show_use();     /* Display correct usage and exit */
                }
```

Figure 11-13. TIME.C, a utility that makes the current time part of the system

```
        if (mm > 59)
                {
                printf("\n\007Minutes = %d is illegal.",mm);
                show_use();     /* Display correct usage and exit */
                }
        if (ss > 59)
                {
                show_use();     /* Display correct usage and exit */
                printf("\n\007Seconds = %d is illegal.",ss);
                }

                /* Convert integers back into formatted string */
        sprintf(time,"%2d:%2d:%2d",hh,mm,ss);
        time[8] = 0x0A;         /* Terminate with line feed */
        time[9] = '\0';         /* New string terminator */

                /* Convert " 1: 2: 3" into "01:02:03" */
        for (count = 0; count < 7; count+=3)
                {
                if (time[count] == ' ')
                        time[count] = '0';
                }
                /* Turn bit on to indicate that the time has been set */
        *time_flag |= TIME_SET;
        }

printf("\n\tCurrent Time is %s",time);
}

show_use()              /* Display correct usage and exit */
{
printf("\nTIME sets the system time. Usage is :");
printf("\n\tTIME hh{:mm{:ss}}");
printf("\n\tTIME SHOW (to display current time)\n");
exit();
}
```

Figure 11-13. TIME.C, a utility that makes the current time part of the system (continued)

FUNKEY — Set the Function Keys

The FUNKEY utility shown in Figure 11-14 sets the character strings associated with specific function keys. In the specified character string, the character "<" is converted into a LINE FEED character. Here is an example console dialog:

```
P3B>funkey<CR>
FUNKEY sets a specific function key string.
        FUNKEY key-number "string to be programmed<"
                   (Note : '<' is changed to line feed.)
                   (       key-number is from 0 to 17.)
                   (       string can be up to 16 chars.)
        FUNKEY SHOW       (displays settings for all keys)


P3B>funkey show<CR>
FUNKEY Vn 1.0 02/18/83
        Key #0 = 'Function Key 1<'
        Key #1 = 'Function Key 2<'

P3B>funkey 0 "PIP B:=A:*.*[V]<"<CR>

P3B>funkey show<CR>
FUNKEY Vn 1.0 02/18/83
        Key #0 = 'PIP B:=A:*.*[V]<'
        Key #1 = 'Function Key 2<'
```

```
#define VN "\nFUNKEY Vn 1.0 02/18/83"

#include <LIBRARY.H>


int fnum;                       /* Function key number to be programmed */
char fstring[20];               /* String for function key */
struct _fkt *pfk;               /* Pointer to function key table */

main(argc,argv)
int argc;
char *argv[];
{

if (argc == 1 || argc > 3)
        show_use();

pfk = get_cba(CB_FKT);  /* Set pointer to function key table */

if (usstrcmp("SHOW",argv[1]))
        {
        if (!isdigit(argv[1][0]))
                {
                printf("\n\007'%s' is an illegal function key.",
                        argv[1]);
                show_use();
                }

        fnum = atoi(argv[1]);   /* Convert function key number */

        if (fnum > FK_ENTRIES)
                {
                printf("\n\007Function key number %d too large.",fnum);
                show_use();
                }

        if (get_fs(fstring) > FK_LENGTH)
                {
                printf("\n\007Function key string is too long.");
                show_use();
                }


        pfk += fnum;    /* Update pointer to string */
                /* Copy string into function key table */

                /* Check if function key input present */
        if (!(pfk -> fk_input[0]))
                {
                printf("\n\007Error : Function Key #%d is not set up to be programmed.",fnum);
                show_use();
                }
        strcpy(pfk -> fk_output,fstring);
        }
else            /* SHOW function specified */
        {
        printf(VN);             /* Display sign-on message */
        show_fun();
        }
}

get_fs(string)          /* Get function string from command tail */
char string[];          /* Pointer to character string */
{
char *tail;             /* Pointer to command tail */
short tcount;           /* Count of TOTAL characters in command tail */
int slen;               /* String length */

tail = 0x80;            /* Command line is in memory at 0080H */
tcount = *tail++;       /* Set TOTAL count of characters in command tail */
slen = 0;               /* Initialize string length */

while(tcount--)         /* For all characters in the command tail */
        {
        if (*tail++ == '"')     /* Scan for first quotes */
                break;
```

Figure 11-14. FUNKEY.C, a utility that sets the character strings associated with specific function keys

```
        }
if (!tcount)            /* No quotes found */
        {
        printf("\n\007No leading quotes found.");
        show_use();
        }

++tcount;               /* Adjust tail count */
while(tcount--)         /* For all remaining characters in tail */
        {
        if (*tail == '"')
                {
                string[slen] = '\0';    /* Add terminator */
                break;          /* Exit from loop */
                }
        string[slen] = *tail++; /* Move char. from tail into string */

        if (string[slen] == '<')
                string[slen] = 0x0A;
        ++slen;
        }
if (!tcount)            /* No terminating quotes found */
        {
        printf("\n\007No trailing quotes found.");
        show_use();
        }
return slen;            /* Return string length */
}


show_fun()              /* Display settings for all function keys */
{
struct _fkt *pfkt;      /* Local pointer to function keys */
int count;              /* Count to access function keys */
char *lf;               /* Pointer to "<" character (LINE FEED) */

pfkt = get_cba(CB_FKT); /* Set pointer to function key table */
for (count = 0; count <= FK_ENTRIES; count++)
        {
        if (pfkt -> fk_input[0])        /* Key is programmed */
                {
                        /* Check if at physical end of table */
                if (pfkt -> fk_input == 0xFF)
                        break;  /* Yes -- break out of for loop */
                strcpy(fstring,pfkt -> fk_output);
                        /* Convert all 0x0A chars to "<" */
                while (lf = strscn(fstring,"\012"))
                        {
                        *lf = '<';
                        }

                printf("\n\tKey #%d = '%s'",count,fstring);
                }
        ++pfkt;         /* Move to next entry */
        }
}


show_use()
{
printf("\nFUNKEY sets a specific function key string.");
printf("\n\tFUNKEY key-number \042string to be programmed<\042 ");
printf("\n\t            (Note : '<' is changed to line feed.)");
printf("\n\t            (        key-number is from 0 to %d.)",
        FK_ENTRIES-1);
printf("\n\t            (        string can be up to %d chars.)",
        FK_LENGTH);
printf("\n\tFUNKEY SHOW        (displays settings for all keys)");
exit();
}
```

Figure 11-14. (Continued)

Other Utilities

Because of space limitations, not all of the possible utility programs for the BIOS features can be shown in this chapter. Others that would need to be developed in order to have a complete set are

PUBLIC/PRIVATE
: This pair of utilities would turn the public files flag on or off, making the files in user 0 available from other user numbers or not, respectively.

SETTERM
: This program would program the CONOUT escape table, setting the various escape sequences as required. It could also program the characters in the function key table that match with those emitted by the terminal currently in use.

SAVESYS
: This utility would save the current settings in the long term configuration block.

LOADSYS
: This would load the long term configuration block from a previously saved image.

DO
: This utility would copy the command tail into the multi-command buffer, changing "\" into LINE FEED, and then set the forced input pointer to the multi-command buffer. As a result, characters from the multi-command buffer would be fed into the console input stream as though they had been typed one command at a time.

SPARE
: This utility would work in conjunction with the hard-disk bad-sector management in your disk drivers. It would spare out bad sectors or tracks on the hard disk. This done, all subsequent references to the sectors or tracks would be redirected to a different part of the disk.

Error Messages Displayed
Miscellaneous Errors

12

Error Messages

This chapter lists the error messages that emanate from standard CP/M and its utility programs. It does not include any error messages from the BIOS; these messages, if any, are the individualized product of the programmers who wrote the various versions of the BIOS.

The error messages are shown in alphabetical order, followed (in parentheses) by the name of the program or CP/M component outputting the message. Messages are shown in uppercase even if the actual message you will see contains lowercase letters. Additional characters that are displayed to "pretty up" the message have been omitted. For example, the message "** ABORTED **" will be listed as "ABORTED".

Following each message is an explanation and, where possible, some information to help you deal with the error.

The last section of the chapter deals with known errors or peculiarities in CP/M and its utilities. Read this section so that you will recognize these problems when they occur.

Error Messages Displayed

? (CCP)

The CCP displays a question mark if you enter a command name and there is no corresponding "command.COM" file on the disk.

It is also displayed if you omit the number of pages required as a parameter in the SAVE command.

? (DDT)

DDT outputs a question mark under several circumstances. You must use context (and some guesswork) to determine what has gone wrong. Here are some specific causes of problems:

- DDT cannot find the file that you have asked it to load into memory. Exit from DDT and investigate using DIR or STAT (the file may be set to System status and therefore invisible with DIR).
- There is a problem with the data in the HEX file that you have asked DDT to load. The problem could be a bad check-sum on a given line or an invalid field somewhere in the record. Try typing the HEX file out on a console, or use an editor to examine it. It is rare to have only one or two bad bits or bytes in a HEX file; large amounts of the file are more likely to have been corrupted. Therefore, you may be able to spot the trouble fairly readily. If you have the source code for the program, reassemble it to produce another copy of the HEX file. If you do not have the source code, there is no reliable way around this problem unless you are prepared to hand-create the HEX file—a difficult and tedious task.
- DDT does not recognize the instruction you have entered when using the "A" (assemble) command to convert a source code instruction into hexadecimal. Check the line that you entered. DDT does not like tabs in the line (although it appears to accept them) or hexadecimal numbers followed by "H". Check that the mnemonic and operands are valid, too.

??= (DDT)

This cryptic notation is used by DDT when you are using the "L" (list disassembled) command to display some part of memory in DDT's primitive assembly language form. DDT cannot translate all of the 256 possible values of a byte. Some of them are not used in the 8080 instruction set. When DDT encounters an untranslatable value, it displays this message as the instruction code, followed by the actual value of the byte in hexadecimal.

You will see this if you try to disassemble code written for the Z80 CPU, which uses unassigned 8080 instructions. You will also see it if you try to disassemble bytes that contain ASCII text strings rather than 8080 instructions.

ABORTED (STAT)

If you enter any keyboard character while STAT is working its way down the file directory setting files to $DIR (Directory), $SYS (System), $R/W (Read/Write), or $R/O (Read-Only) status, then it will display this message, stop what it is doing, and execute a warm boot.

By contrast, if you enter the command

```
A>stat *.*<cr>
```

to display all of the files on a disk, there is no way that the process can be aborted.

ABORTED (PIP)

This message is displayed if you press any keyboard character while PIP is copying a file to the list device.

BAD DELIMITER (STAT)

If your BIOS uses the normal IOBYTE method of assigning physical devices to logical devices, you use STAT to perform the assignment. The command has this format:

```
STAT RDR:=PTR:
```

STAT displays this message if it cannot find the "=" in the correct place.

BAD LOAD (CCP)

This is probably the most obscure error message that emanates from CP/M. You will get this message if you attempt to load a COM file that is larger than the transient program area. Your only recourse is to build a CP/M system that has a larger TPA.

BAD PARAMETER (PIP)

PIP accepts certain parameters in square brackets at the end of the command line. This message is displayed if you enter an invalid parameter or an illegal numeric value following a parameter letter.

BDOS ERROR ON d: BAD SECTOR (BDOS)

The BDOS displays this message if the READ and WRITE functions in your BIOS ever return indicating an error. The only safe response to this message is to type CONTROL-C. CP/M will then execute a warm boot. If you type CARRIAGE RETURN, the error will be ignored—with unpredictable results.

A well-implemented BIOS should include disk error recovery and control so that the error will never be communicated to the BDOS. If the BIOS gives you the option of ignoring an error, do so only when you are reasonably sure of the outcome or have adequate backup copies so that you can recreate your files.

BDOS ERROR ON d: FILE R/O (BDOS)

You will see this message if you attempt to erase (ERA) a file that has been set to Read-Only status. Typing any character on the keyboard causes the BDOS to perform a warm boot operation. Note that the BDOS does not tell you *which* file is creating the problem. This can be a problem when you use ambiguous file names in the ERA command. Use the STAT command to display all the files on the disk; it will tell you which files are Read-Only.

This message is also displayed if a program tries to delete a Read-Only file. Again, it can be difficult to determine which file is causing the problem. Your only recourse is to use STAT to try to infer which of the Read-Only files might be causing the problems.

BDOS ERROR ON d: R/O (BDOS)

This looks similar to the previous message, but it refers to an entire logical disk instead of a Read-Only file. However, it is rarely output because you have declared a disk to be Read-Only. Usually, it occurs because you changed diskettes without typing a CONTROL-C; CP/M will detect the new diskette and, without any external indication, will set the disk to Read-Only status.

If you or a program attempts to write any data to the disk, the attempt will be trapped by the BDOS and this message displayed. Typing any character on the keyboard causes a warm boot—then you can proceed.

BDOS ERROR ON d: SELECT (BDOS)

The BDOS displays this message if you or a program attempts to select a logical disk for which the BIOS lacks the necessary tables. The BDOS uses the value returned by SELDSK to determine whether a logical disk "exists" or not.

If you were trying to change the default disk to a nonexistent one, you will have to press the RESET button on your computer. There is no way out of this error.

However, if you were trying to execute a command that accessed the nonexistent disk, then you can type a CONTROL-C and CP/M will perform a warm boot.

BREAK x AT y (ED)

This is another cryptic message whose meaning you cannot guess. The list that follows explains the possible values of "x." The value "y" refers to the command ED was executing when the error occurred.

x	Meaning
#	Search failure. ED did not find the string you asked it to search for.
?	Unrecognized command.
0	File not found.
>	ED's internal buffer is full.
E	Command aborted.
F	Disk or directory full. You will have to determine which is causing the problem.

CANNOT CLOSE, READ/ONLY? (SUBMIT)

SUBMIT displays this message if the disk on which it is trying to write its output file, "$$$.SUB", is physically write protected. Do not confuse this with the disk being *logically* write protected.

The standard version of SUBMIT writes the output file onto the current default disk, so if your current default disk is other than drive A:, you may be able to avoid this problem if you switch the default to A: and then enter a command of the form

```
A>submit b:subfile<cr>
```

CANNOT CLOSE DESTINATION FILE (PIP)

PIP displays this message if the destination disk is physically write protected. Check the destination disk. If it is write protected, remove the protection and repeat the operation.

If the disk is not protected, you have a hardware problem. The directory data written to the disk is being written to the wrong place, even the wrong disk, or is not being recorded on the medium.

CANNOT CLOSE FILES (ASM)

ASM displays this message if it cannot close its output files because the disk is physically write protected, or if there is a hardware problem that prevents data being written to the disk. See the paragraph above.

CANNOT READ (PIP)

PIP displays this message if you attempt to read information from a logical device that can only output. For example:

```
A>pip diskfile=LST:<cr>
```

PIP also will display this message if you confuse it sufficiently, as with the following instruction:

```
A>pip file1=file2;file3<cr>
```

CANNOT WRITE (PIP)

PIP displays this message if you attempt to output (write) information to a logical device that can only be used for input, such as the RDR: (reader, the anachronistic name for the auxiliary input device).

CHECKSUM ERROR (LOAD)

LOAD displays this message if it encounters a line in the input HEX file that does not have the correct check sum for the data on the line.

LOAD also displays information helpful in pinpointing the problem:

```
CHECKSUM ERROR
LOAD ADDRESS  0110  <- First address on line in file
ERROR ADDRESS 0112  <- Address of next byte to be loaded
BYTES READ:
0110:
0110: 00 33 22 2B 02 21 27 02    <- Bytes preceding error
```

Note that LOAD does not display the check-sum value itself. Use TYPE or an editor to inspect the HEX file in order to see exactly what has gone wrong.

CHECKSUM ERROR (PIP)

If you ask PIP to copy a file of type HEX, it will check each line in the file, making sure that the line's check sum is valid. If it is not, PIP will display this message. Unfortunately, PIP does not tell you which line is in error—you must determine this by inspection or recreate the HEX file and try again.

COMMAND BUFFER OVERFLOW (SUBMIT)

SUBMIT displays this message if the SUB file you specified is too large to be processed. SUBMIT's internal buffer is only 2048 bytes. You must reduce the size of the SUB file; remove any comment lines, or split it into two files with the last line of the first file submitting the second to give a nested SUBMIT file.

COMMAND TOO LONG (SUBMIT)

The longest command line that SUBMIT can process is 125 characters. There is no way around this error other than reducing the length of the offending line. You will have to find this line by inspection—SUBMIT does not identify the line.

One way that you can remove a few characters from a command line is to rename the COM file you are invoking to a shorter name, or use abbreviated names for parameters if the program will accept these.

CORRECT ERROR, TYPE RETURN OR CTL-Z (PIP)

This message is a carryover from the days when PIP used to read hexadecimal data from a high-speed paper tape reader. If PIP detected the end of a physical roll of paper tape, it would display this message. The user could then check to see if the paper tape had torn or had really reached its end. If there was more tape to be read, the user could enter a CARRIAGE RETURN to resume reading tape or enter a CONTROL-Z to serve as the end-of-file character.

Needless to say, it is unlikely that you will see this message if you do not have a paper tape reader.

DESTINATION IS R/O, DELETE (Y/N)? (PIP)

PIP displays this message if you try to overwrite a disk file that has been set to Read-Only status. If you type "Y" or "y", PIP will overwrite the destination file. It leaves the destination file in Read/Write status with its Directory/System status unchanged. Typing any character other than "Y" or "y" makes PIP abandon the copy and display the message

```
** NOT DELETED**
```

You can avoid this message altogether if you specify the "w" option on PIP's command line. For example:

```
A>pip destfile=srcfile[w]<cr>
```

PIP will then overwrite Read-Only files without question.

DIRECTORY FULL (SUBMIT)

This message is displayed if the BDOS returns an error when SUBMIT tries to create its output file, "$$$.SUB". As a rough and ready approximation, use "STAT *.*" to see how many files and extents you have on the disk. Erase any unwanted ones. Then use "STAT DSK:" to find out the maximum number of directory entries possible for the disk.

You may also see this message if the file directory has become corrupted or if the disk formatting routine leaves the disk with the file directory full of some pattern other than E5H.

You can assess whether the directory has been corrupted by using "STAT USR:". STAT then displays which user numbers contain files. If the directory is corrupt, you will normally see user numbers greater than 15.

It is not easy to repair a corrupted directory. "ERA *.*" erases only the files for the current user number, so you will have to enter the command 16 times, once for each user number from 0 to 15. Alternatively, you can reformat the disk.

DISK OR DIRECTORY FULL (ED)

Self-explanatory.

DISK READ ERROR (PIP)
DISK WRITE ERROR (SUBMIT)
DISK WRITE ERROR (PIP)

These messages will normally be preceded by a BIOS error message. They will only be displayed if the BIOS returns indicating an error. As was described earlier, this is unlikely if the BIOS has any kind of error recovery logic.

END OF FILE, CTL-Z? (PIP)

PIP displays this message if, while copying a HEX file, it encounters a CONTROL-Z (end of file). Again, the underlying idea is based on the concept of physical paper tape. When you saw this message, you could look at the tape in the reader, and if it really was at the end of the roll, enter a CONTROL-Z on the keyboard to terminate the file. Given any other character, PIP would read the next piece of tape.

ERROR : CANNOT CLOSE FILES (LOAD)

LOAD displays this message if you have physically write protected the disk on which it is trying to write the output COM file.

ERROR : CANNOT OPEN SOURCE (LOAD)

LOAD displays this message if it cannot open the HEX file that you specified in the command tail.

ERROR : DISK READ (LOAD)
ERROR : DISK WRITE (LOAD)

These two messages would normally be preceded by a BIOS error message. If your BIOS includes disk error recovery, you would not normally see these messages; the error would have been handled by the BIOS.

ERROR : INVERTED LOAD ADDRESS (LOAD)

LOAD displays this message if it detects a load address less than 0100H in the input HEX file. It also displays the actual address input from the file, so you can examine the HEX file looking for this address to determine the likely cause of the problem.

Note that DDT, when asked to load the same HEX file, will do so without any error—and will probably damage the contents of the base page in so doing.

ERROR : NO MORE DIRECTORY SPACE (LOAD)

Self-explanatory.

ERROR ON LINE N (SUBMIT)

SUBMIT displays this message if it encounters a line in the SUB file that it does not know how to process. Most likely you have a file that has type .SUB but does not contain ASCII text.

The first line of the SUB file is number 001.

FILE EXISTS (CCP)

The CCP displays this message if you attempt to use the REN command to rename an existing file to a name already given to another file.

Use "STAT *.*" to display all of the files on the disk. DIR will show only those files that have Directory status, and you may not be able to see the file causing the problem.

FILE IS READ/ONLY (ED)

ED displays this message if you attempt to edit a file that has been set to Read-Only status.

FILE NOT FOUND (STAT)
FILENAME NOT FOUND (PIP)

STAT and PIP display their respective messages if you specify a nonexistent file. This applies to both specific and ambiguous file names.

INVALID ASSIGNMENT (STAT)

STAT can be used to assign physical devices to logical devices using the IOBYTE system described earlier. It will display this message if you enter an illogical assignment. Use the "STAT VAL:" command to display the valid assignments.

INVALID CONTROL CHARACTER (SUBMIT)

SUBMIT is supposed to be able to handle a control character in the SUB file—the notation being "^x", where "x" is the control letter. In fact, the standard release version of SUBMIT cannot handle this notation. A patch is available from Digital Research to correct this problem.

Given that this patch has been installed, SUBMIT will display this message if a character other than "A" to "Z" is specified after the circumflex character.

INVALID DIGIT (PIP)

PIP displays this message if it encounters non-numeric data where it expects a numeric value.

INVALID DISK ASSIGNMENT (STAT)

STAT displays this message if you try to set a logical disk to Read-Only status and you specify a parameter other than "R/O." Note that there is no leading "$" in this case (as there is when you want to set a file to Read-Only).

INVALID DRIVE NAME (USE A, B, C, OR D) (SYSGEN)

SYSGEN displays this message if you attempt to load the CP/M system from, or write the system to, a disk drive other than A, B, C, or D.

INVALID FILE INDICATOR (STAT)

STAT outputs this message if you specify an erroneous file attribute. File attributes can only be one of the following:

$DIR	Directory
$SYS	System
$R/O	Read-Only
$R/W	Read/Write

INVALID FORMAT (PIP)

PIP displays this message if you enter a badly formatted command; for example, a "+" character instead of an "=" (on some terminals these are on the same key).

INVALID HEX DIGIT (LOAD)

LOAD displays this message if it encounters a nonhexadecimal digit in the input HEX file, where only a hex digit can appear. LOAD then displays additional information to tell you where in the file the problem occurred:

```
INVALID HEX DIGIT
LOAD ADDRESS  0110  <- First address on line in file
ERROR ADDRESS 0112  <- Address of byte containing non-hex
BYTES READ:
0110:
0110: 00 33    <- Bytes preceding error
```

INVALID MEMORY SIZE (MOVCPM)

MOVCPM displays this message if you enter an invalid memory size for the CP/M system size you want to construct.

INVALID SEPARATOR (PIP)

PIP displays this message if you try to concatenate files using something other than a comma between file names.

INVALID USER NUMBER (PIP)

PIP displays this message if you enter a user number outside the range 0 to 15 with the "[gn]" option (where "n" is the user number).

NO 'SUB' FILE PRESENT (SUBMIT)

SUBMIT displays this message if it cannot find a file with the file name that you specified and with a type of .SUB.

NO DIRECTORY SPACE (ASM)
NO DIRECTORY SPACE (PIP)

Self-explanatory.

NO FILE (CCP)

The CCP displays this message if you use the REN (rename) command and it cannot find the file you wish to rename.

NO FILE (PIP)

PIP displays this message if it cannot find the file that you specified.

NO MEMORY (ED)

ED displays this message if it runs out of memory to use for storing the text that you are editing.

NO SOURCE FILE ON DISK (SYSGEN)

This error message is misleading. SYSGEN does not read source code files. The message should read "INPUT FILE NOT FOUND".

NO SOURCE FILE PRESENT (ASM)

In this case, ASM really does mean that the source code file cannot be found. Remember that ASM uses a strange form of specifying its parameters. ASM uses the file name that you enter and then searches for a file of that name, but with file type .ASM. The three characters of the file type that you specify are used to represent the logical disks on which the source, hex, and list files, respectively, are to be placed.

NO SPACE (CCP)

The CCP displays this message if you use the SAVE command and there is insufficient room on the disk to accommodate the file.

NOT A CHARACTER SOURCE (PIP)

PIP displays this message if you attempt to copy characters from a character output device, such as the auxiliary output device (known to PIP as PUN:).

OUTPUT FILE WRITE ERROR (ASM)

ASM will display this message if the BDOS returns an error from a disk write operation. If your BIOS has disk error recovery logic, you should never see this message.

PARAMETER ERROR (SUBMIT)

SUBMIT uses the "$" to mark points where parameter values are to be substituted. If you have a single "$" followed by an alphabetic character, SUBMIT will display this message. Use "$$" to represent a real "$".

PERMANENT ERROR, TYPE RETURN TO IGNORE (SYSGEN)

SYSGEN displays this message if the BIOS returns an error from a disk read or write operation. If your BIOS has disk error recovery logic, you should never see this message.

QUIT NOT FOUND (PIP)

PIP displays this message when it cannot find the string specified in the "[Qcharacter string^Z]" option, meaning "Quit copying when you encounter this string."

READ ERROR (CCP)

The CCP displays this message if the BIOS returns an error from a disk read or write operation. If your BIOS includes disk error recovery logic, you should not see this error message.

RECORD TOO LONG (PIP)

PIP displays this message if it encounters a line longer than 80 characters while copying a HEX file. Inspect the HEX file using the TYPE command or an editor.

REQUIRES CP/M 2.0 OR NEWER FOR OPERATION (PIP)
REQUIRES CP/M VERSION 2.0 OR LATER (XSUB)

Self-explanatory.

SOURCE FILE INCOMPLETE (SYSGEN)

SYSGEN displays this message if the file that you have asked it to read is too short. Use STAT to check the length of the file.

SOURCE FILE NAME ERROR (ASM)

ASM displays this message if you specify an ambiguous file name: that is, one that contains either "*" or "?".

SOURCE FILE READ ERROR (ASM)

ASM displays this message if it encounters problems reading the input source code file. Check the input file using the TYPE command or an editor.

START NOT FOUND (PIP)

PIP displays this message when it cannot find the string specified in the "[Scharacter string^Z]" option, meaning "Start copying when you encounter this string."

SYMBOL TABLE OVERFLOW (ASM)

ASM displays this message when you have too many symbols in the source code file. Your only recourse is to split the source file into several pieces and arrange for ORG (origin) statements to position the generated object code so that the pieces fit together.

SYNCRONIZATION ERROR (MOVCPM)

Apart from the spelling error, this message is designed to be cryptic. MOVCPM displays it when the Digital Research serial number embedded in MOVCPM does not match the serial number in the version of CP/M that you are currently running.

SYSTEM FILE NOT ACCESSIBLE (ED)

ED displays this message if you attempt to edit a file that has been set to System status. Use STAT to set the file to Directory status.

TOO MANY FILES (STAT)

STAT displays this message if there is insufficient memory available to sort and display all of the files on the specified disk. Try limiting the number of files it has to sort by judicious use of ambiguous file names.

UNRECOGNIZED DESTINATION (PIP)

PIP displays this message if you specify an "illegal" destination device.

VERIFY ERROR (PIP)

If you use the "[v]" (verify) option of PIP when copying to a disk file, PIP will write a sector to the disk, read it back, and compare the data. PIP displays this message if the data does not match.

If there is a problem with your disk system, you should have seen some form of disk error message preceding this one. If there is no preceding message, then you have a problem with the main memory on your system.

Wrong CP/M Version (Requires 2.0) (STAT)

Self-explanatory.

(XSUB ACTIVE) (XSUB)

This is not really an error message, but you may mistake it for one. XSUB is the eXtended SUBMIT program. Without it, SUBMIT can only feed command lines to the Console Command Processor. XSUB allows character-by-character input into any program that uses the BDOS to read console input.

XSUB is initiated by being the first command in a SUB file. Once initiated it stays in memory until the end of the SUB file has been reached. Until that happens, XSUB will output this message every time a warm boot occurs as a reminder that it is still in memory.

XSUB Already Present (XSUB)

XSUB will display this message if it is already active and you attempt to load it again.

Miscellaneous Errors

This section deals with errors that are not accompanied by any error message. It is included here to help you recognize a problem after it has already occurred.

The errors are shown grouped by product.

ASM: Fails to Detect Unterminated IF Clause

If you use the IF pseudo-operation, it must be followed by a matching ENDIF. ASM fails to detect the case that the end of the source file is encountered *before* the ENDIF.

If the condition specified on the IF line is false, you could have a situation in which ASM would ignore the majority of the source file without comment.

ASM: Creates HEX File That Cannot Be Loaded

If you omit the ORG statement at the front of a source file, ASM will assemble the code origined at location 0000H. This file will crash the system if you try to load it with DDT. The message "ERROR: INVERTED ADDRESS" will be shown from LOAD.

CP/M: Signs On and Then Dies Without A> Prompt

After the BIOS has signed on, it transfers control to the Console Command Processor. The CCP then attempts to log in the system disk, reading the file directory and building the allocation vector. If your file directory has been badly corrupted, it can cause the system to crash. Use another system disk and try to display the directory on the bad disk.

DDT: Loads HEX File and Then Crashes the System

DDT does not check the addresses specified in a HEX file. If you have forgotten to put an ORG statement at the front of the source file, or more subtly, if your source program has "wrapped around" by having addresses up at 0FFFFH and "above," the assembler will start assembling at 0000H again.

DIR: Shows Odd-Looking File Names

If you have odd-looking file names, or the vertical lines of ":" that DIR uses to separate the file names are misaligned, then the file directory has been corrupted. One strategy is to format a new disk, copy all of the valid files to it, and discard the corrupted disk.

DIR: Shows More than One Entry with the Same Name

This can happen if you use a program that creates a new file without asking the BDOS to delete any existing files of the same name. It can also happen if you use the custom MOVE utility carelessly.

To remedy the situation proceed as follows:

- Use PIP to copy the specific file to another disk. Do not use an ambiguous file name; specify the duplicated file name exactly. PIP will copy the first instance of the file it encounters in the directory.
- Use the ERA command to erase the duplicated file. *This will erase both copies of the file.*
- Use PIP to copy back the first instance of the file.

STAT: User Numbers > 15

If you use the "STAT USR:" command to display which user numbers contain active files, and user numbers greater than 15 are displayed, then the file directory on the disk has been corrupted.

Use PIP to copy the valid files from legitimate user numbers, and then discard the corrupted disk.

SUBMIT: Fails to Start Submit Procedure

There are several reasons why SUBMIT will not initiate a SUB file:

- You are using the standard release version of SUBMIT and your current default disk is other than drive A:. SUBMIT builds its "$$$.SUB" file on the default disk, but the CCP only looks on drive A: for "$$$.SUB". Use the following procedure to modify SUBMIT to build its "$$$.SUB" file on drive A:

```
A>DDT SUBMIT.COM<cr>
DDT VERS 2.2
NEXT  PC
0600 0100
-s5bb                <- Change 5bb
05BB 01 00<cr>       <- from 00 (default drive)
05BC 24 .<cr>           to 01 (drive A:)
-^C
A>SAVE 5 SUBMIT.COM<cr>
A>_
```

- If you forgot to terminate the last line of the SUB file with a CARRIAGE RETURN.
- If your SUB file contains a line with nothing but a CARRIAGE RETURN on it (that is, a blank line).

A

ASCII Character Set

The American Standard Code for Information Interchange (ASCII) consists of a set of 96 displayable characters and 32 nondisplayed characters. Most CP/M systems use at least a subset of the ASCII character set. When CP/M stores characters on a diskette as text, the ASCII definitions are used.

Several of the CP/M utility programs use the ASCII Character Code. Text created using ED is stored as ASCII characters on diskette. DDT, when displaying a "dump" of the contents of memory, displays both the hexadecimal and ASCII representations of memory's contents.

ASCII does not use an entire byte of information to represent a character. ASCII is a seven-bit code, and the eighth bit is often used for *parity.* Parity is an error-checking method which assures that the character received is the one transmitted. Many microcomputers and microcomputer devices ignore the *parity bit,* while others require one of the following two forms of parity:

Even Parity

The number of binary 1's in a byte is always an even number. If there is an odd number of 1's in the character, the parity bit will be a 1; if there is an even number of 1's in the character, the parity bit is made a 0.

Odd Parity

The number of binary 1's in a byte is always an odd number. If there is an even number of 1's in the character, the parity bit will be a 1; if there is an odd number of 1's in the character, the parity bit is made a 0.

Alternative ways of *coding* the information stored by the computer include the 8-bit EBCDIC (Extended Binary Coded Decimal Interchange Code), used by IBM, and a number of *packed binary* schemes, primarily used to represent numerical information.

Table A-1. ASCII Character Codes

				b7 →	0	0	0	0	1	1	1	1
				b6 →	**0**	**0**	**1**	**1**	**0**	**0**	**1**	**1**
				b5 →	**0**	**1**	**0**	**1**	**0**	**1**	**0**	**1**
b4	**b3**	**b2**	**b1**	**Row \ Col.**	**0**	**1**	**2**	**3**	**4**	**5**	**6**	**7**
0	0	0	0	0	NUL	DLE	SP	0	@	P	`	p
0	0	0	1	1	SOH	DC1	!	1	A	Q	a	q
0	0	1	0	2	STX	DC2	"	2	B	R	b	r
0	0	1	1	3	ETX	DC3	#	3	C	S	c	s
0	1	0	0	4	EOT	DC4	$	4	D	T	d	t
0	1	0	1	5	ENQ	NAK	%	5	E	U	e	u
0	1	1	0	6	ACK	SYN	&	6	F	V	f	v
0	1	1	1	7	BEL	ETB	'	7	G	W	g	w
1	0	0	0	8	BS	CAN	(	8	H	X	h	x
1	0	0	1	9	HT	EM	)	9	I	Y	i	y
1	0	1	0	10	LF	SUB	*	:	J	Z	j	z
1	0	1	1	11	VT	ESC	+	;	K	[	k	{
1	1	0	0	12	FF	FS	,	<	L	\	l	|
1	1	0	1	13	CR	GS	-	=	M	]	m	}
1	1	1	0	14	SO	RS	.	>	N	^	n	~
1	1	1	1	15	SI	US	/	?	O	_	o	DEL

NUL	Null	DC1	Device control 1
SOH	Start of heading	DC2	Device control 2
STX	Start of text	DC3	Device control 3
ETX	End of text	DC4	Device control 4
EOT	End of transmission	NAK	Negative acknowledge
ENQ	Enquiry	SYN	Synchronous idle
ACK	Acknowledge	ETB	End of transmission block
BEL	Bell or alarm	CAN	Cancel
BS	Backspace	EM	End of medium
HT	Horizontal tabulation	SUB	Substitute
LF	Line feed	ESC	Escape
VT	Vertical tabulation	FS	File separator
FF	Form feed	GS	Group separator
CR	Carriage return	RS	Record separator
SO	Shift out	US	Unit separator
SI	Shift in	SP	Space
DLE	Data link escape	DEL	Delete

Table A-2. ASCII Character Codes in Ascending Order

Hexadecimal	Binary	ASCII	Hexadecimal	Binary	ASCII
00	000 0000	NUL	30	011 0000	0
01	000 0001	SOH	31	011 0001	1
02	000 0010	STX	32	011 0010	2
03	000 0011	ETX	33	011 0011	3
04	000 0100	EOT	34	011 0100	4
05	000 0101	ENQ	35	011 0101	5
06	000 0110	ACK	36	011 0110	6
07	000 0111	BEL	37	011 0111	7
08	000 1000	BS	38	011 1000	8
09	000 1001	HT	39	011 1001	9
0A	000 1010	LF	3A	011 1010	:
0B	000 1011	VT	3B	011 1011	;
0C	000 1100	FF	3C	011 1100	<
0D	000 1101	CR	3D	011 1101	=
0E	000 1110	SO	3E	011 1110	>
0F	000 1111	SI	3F	011 1111	?
10	001 0000	DLE	40	100 0000	
11	001 0001	DC1	41	100 0001	A
12	001 0010	DC2	42	100 0010	B
13	001 0011	DC3	43	100 0011	C
14	001 0100	DC4	44	100 0100	D
15	001 0101	NAK	45	100 0101	E
16	001 0110	SYN	46	100 0110	F
17	001 0111	ETB	47	100 0111	G
18	001 1000	CAN	48	100 1000	H
19	001 1001	EM	49	100 1001	I
1A	001 1010	SUB	4A	100 1010	J
1B	001 1011	ESC	4B	100 1011	K
1C	001 1100	FS	4C	100 1100	L
1D	001 1101	GS	4D	100 1101	M
1E	001 1110	RS	4E	100 1110	N
1F	001 1111	US	4F	100 1111	O
20	010 0000	SP	50	101 0000	P
21	010 0001	!	51	101 0001	Q
22	010 0010	”	52	101 0010	R
23	010 0011	#	53	101 0011	S
24	010 0100	$	54	101 0100	T
25	010 0101	%	55	101 0101	U
26	010 0110	&	56	101 0110	V
27	010 0111	’	57	101 0111	W
28	010 1000	(	58	101 1000	X
29	010 1001	)	59	101 1001	Y
2A	010 1010	*	5A	101 1010	Z
2B	010 1011	+	5B	101 1011	[
2C	010 1100	,	5C	101 1100	\
2D	010 1101	-	5D	101 1101	]
2E	010 1110	.	5E	101 1110	^
2F	010 1111	/	5F	101 1111	_

Table A-2. ASCII Character Codes in Ascending Order (Continued)

Hexadecimal	Binary	ASCII	Hexadecimal	Binary	ASCII
60	110 0000		70	111 0000	p
61	110 0001	a	71	111 0001	q
62	110 0010	b	72	111 0010	r
63	110 0011	c	73	111 0011	s
64	110 0100	d	74	111 0100	t
65	110 0101	e	75	111 0101	u
66	110 0110	f	76	111 0110	v
67	110 0111	g	77	111 0111	w
68	110 1000	h	78	111 1000	x
69	110 1001	i	79	111 1001	y
6A	110 1010	j	7A	111 1010	z
6B	110 1011	k	7B	111 1011	{
6C	110 1100	l	7C	111 1100	|
6D	110 1101	m	7D	111 1101	}
6E	110 1110	n	7E	111 1110	~
6F	110 1111	o	7F	111 1111	DEL

B

CP/M Command Summary

This appendix summarizes the command line format and the function of each CP/M built-in and transient command. The commands are listed in alphabetical order.

ASM Command Lines

ASM filename<cr> Assembles the file filename.ASM; uses the currently logged disk for all files.

ASM filename.opt<cr> Assembles the file filename.ASM on drive o: (A:,B:,...,P:). Writes HEX file on drive p: (A:,B:,...,P:), or skips if p: is Z:.

Writes PRN file on drive t: (A:,B:,...,P:), sends to console if p: is X:, or skips if p: is Z:.

DDT Command Lines

DDT<cr> Loads DDT and waits for DDT commands.

DDT x:filename.typ<cr> Loads DDT into memory and also loads filename.typ from drive x: into memory for examination, modification, or execution.

DDT Command Summary

Assss Enters assembly language statements beginning at hexadecimal address ssss.

D Displays the contents of the next 192 bytes of memory.

Dssss,ffff Displays the contents of memory starting at hexadecimal address ssss and finishing at hexadecimal address ffff.

Fssss,ffff,cc Fills memory with the 8-bit hexadecimal constant cc starting at hexadecimal address ssss and finishing with hexadecimal address ffff.

G Begins execution at the address contained in the program counter.

G,bbbb Sets a breakpoint at hexadecimal address bbbb, then begins execution at the address contained in the program counter.

G,bbbb,cccc Sets breakpoints at hexadecimal addresses bbbb and cccc, then begins execution at the address contained in the program counter.

Gssss Begins execution at hexadecimal address ssss.

Gssss,bbbb Sets a breakpoint at hexadecimal address bbbb, then begins execution at hexadecimal address ssss.

Hx,y Hexadecimal sum and difference of x and y.

Ifilename.typ Sets up the default file control block using the name filename.typ.

L Lists the next eleven lines of assembly language program disassembled from memory.

Lssss Lists eleven lines of assembly language program disassembled from memory starting at hexadecimal address ssss.

Lssss,ffff Lists the assembly language program disassembled from memory starting at hexadecimal address ssss and finishing at hexadecimal address ffff.

Mssss,ffff,dddd Moves the contents of the memory block starting at hexadecimal address ssss and ending at hexadecimal address ffff to the block of memory starting at hexadecimal address dddd.

R Reads a file from disk into memory (use "I" command first).

Rnnnn Reads a file from disk into memory beginning at the hexadecimal address nnnn higher than normal (use "I" command first).

Sssss Displays the contents of memory at hexadecimal address ssss and optionally changes the contents.

Tnnnn Traces the execution of (hexadecimal) nnnn program instructions.

Unnnn Executes (hexadecimal) nnnn program instructions, then stops and displays the CPU register's contents.

X Displays the CPU register's contents.

Xr Displays the contents of CPU or Flag r and optionally changes them.

DIR Command Lines

DIR x:<cr> Displays directory of all files on drive x:. Drive x: is optional; if omitted, the currently logged drive is used.

DIR x:filename.typ<cr> Displays directory of all files on drive x: whose names match the ambiguous or unambiguous filename.typ. Drive x: is optional; if omitted, the currently logged drive is used.

DUMP Command Line

DUMP x:filename.typ <cr> Displays the hexadecimal representations of each byte stored in the file filename.typ on drive x:. If filename.typ is ambiguous, displays the first file which matches the ambiguous file name.

ED Command Line

ED x:filename.typ <cr> Invokes the editor, which then searches for filename.typ on drive x: and creates a temporary file x:filename.$$$ to store the edited text. The filename.typ is unambiguous. Drive x: is optional; if omitted, the currently logged drive is assumed.

ED Command Summary

NOTE: Non-alphabetic commands follow the "Z" command.

nA Append lines. Moves "n" lines from original file to edit buffer. 0A moves lines until edit buffer is at least half full.

+/−B Begin/Bottom. Moves CP.
+B moves CP to beginning of edit buffer
−B moves CP to end of edit buffer.

+/−nC Move by characters. Moves CP by "n" character positions.
+ moves forward
− moves backward.

+/−nD Delete characters. Deletes "n" characters before or after the CP in the edit buffer.
+ deletes before the CP
− deletes after the CP.

E End. Ends edit, closes files, and returns to CP/M; normal end.

nFstring^Z Find string. Finds the "n"th occurrence of string, beginning the search after the CP.

H Move to head of edited file. Ends edit, renames files, and then edits former temporary file.

I<cr> Enter insert mode. Text from keyboard goes into edit buffer after the CP; exit with CONTROL-Z.

Istring^Z Insert string. Inserts string in edit buffer after the CP.

Istring<cr> Insert line. Inserts string and CRLF in the edit buffer after the CP.

nJfindstring^Zinsertstring^Zendstring^Z Juxtaposition. Beginning after the CP, finds findstring, inserts insertstring after it, then deletes all following characters up to but not including endstring; repeats until performed "n" times.

+/−nK Kill lines. Deletes "n" lines.
+ deletes after the CP
− deletes before the CP.

+/−nL Move by lines. Moves the CP to the beginning of the line it is in, then moves the CP "n" lines forward or backward.
+ moves forward
− moves backward.

nMcommandstring^Z Macro command. Repeats execution of the ED commands in commandstring "n" times. "n" = 0, "n" = 1, or "n" absent repeats execution until error occurs.

nNstring^Z Find string with autoscan. Finds the "n"th occurrence of string, automatically appending from original file and writing to temporary file as necessary.

O Return to original file. Empties edit buffer, empties temporary file, returns to beginning of original file, ignores previous ED commands.

+/−nP Move CP and print pages. Moves the CP forward or backward one page, then displays the page following the CP. "nP" displays "n" pages, pausing after each.

Q Quit edit. Erases temporary file and block move file, if any, and returns to CP/M; original file is not changed.

R<cr> Read block move file. Copies the entire block move file X$$$$$$$.LIB from disk and inserts it in the edit buffer after the CP.

Rfilename<cr> Read library file. Copies the entire file filename with extension LIB from the disk and inserts it in the edit buffer after the CP.

nSfindstring^Zreplacestring^Z Substitute string. Starting at the CP, repeats "n" times: finds findstring and replaces it with replacestring.

+/−nT Type lines. Displays "n" lines.

+ displays the "n" lines after the CP

− displays the "n" lines before the CP.

If the CP is not at the beginning of a line

0T displays from the beginning of the line to the CP

T displays from the CP to the end of the line

0TT displays the entire line without moving the CP.

+/−U Uppercase translation. After +U command, alphabetic input to the edit buffer is translated from lowercase to uppercase; after −U, no translation occurs.

0V Edit buffer free space/size. Displays the decimal number of free (empty) bytes in the edit buffer and the total size of the edit buffer.

+/−V Verify line numbers. After +V, a line number is displayed with each line displayed; ED's prompt is then preceded by the number of the line containing the CP. After −V, line numbers are not displayed, and ED's prompt is "*".

nW Write lines. Writes first "n" lines from the edit buffer to the temporary file; deletes these lines from the edit buffer.

nX Block transfer (Xfer). Copies the "n" lines following the CP from the edit buffer to the temporary block move file X$$$$$$$.LIB; adds to previous contents of that file.

nZ Sleep. Delays execution of the command which follows it. Larger "n" gives longer delay, smaller "n" gives shorter delay.

n: Move CP to line number "n." Moves the CP to the beginning of the line number "n" (*see* "+/−V").

:m Continue through line number "m." A command prefix which gives the ending point for the command which follows it. The beginning point is the location of the CP (*see* "+/−V").

+/−n Move and display one line. Abbreviated form of +/−nLT.

ERA Command Lines

ERA x:filename.typ<cr> Erases the file filename.typ on the disk in drive x:. The filename and/or typ can be ambiguous. Drive x: is optional; if omitted, the currently logged drive is used.

ERA x:*.*<cr> Erases all files on the disk in drive x:. Drive x: is optional; if omitted, the currently logged drive is used.

Line Editing Commands

CONTROL-C Restarts CP/M if it is the first character in command line. Called *warm start.*

CONTROL-E Moves to the beginning of next line. Used for typing long commands.

CONTROL-H or BACKSPACE Deletes one character and erases it from the screen (CP/M version 2.0 and newer).

CONTROL-J or LINE FEED Same as CARRIAGE RETURN (CP/M version 2.0 and newer).

CONTROL-M Same as CARRIAGE RETURN (<cr>).

CONTROL-P Turns on the list device (usually your printer). Type it again to turn off the list device.

CONTROL-R Repeats current command line (useful with version 1.4); it verifies the line is corrected after you delete several characters (CP/M version 1.4 and newer).

CONTROL-S Temporarily stops display of data on the console. Press any key to continue.

CONTROL-U or CONTROL-X Cancels current command line (CP/M version 1.4 and newer).

RUBOUT (RUB) or DELETE (DEL) Deletes one character and echoes (repeats) it.

Load Command Line

LOAD x:filename<cr> Reads the file filename.HEX on drive x: and creates the executable program file filename.COM on drive x:.

MOVCPM Command Lines

MOVCPM<cr> Prepares a new copy of CP/M which uses all of memory; gives control to the new CP/M, but does not save it on disk.

MOVCPM nn<cr> Prepares a new copy of CP/M which uses "nn" K bytes of memory; gives control to the new CP/M, but does not save it on disk.

MOVCPM * * <cr> Prepares a new copy of CP/M that uses all of memory, to be saved with SYSGEN or SAVE.

MOVCPM nn * <cr> Prepares a new copy of CP/M that uses "nn" K bytes of memory, to be saved with SYSGEN or SAVE.

The "nn" is an integer decimal number. It can be 16 through 64 for CP/M 1.3 or 1.4. For CP/M 2.0 and newer "nn" can be 20 through 64.

PIP Command Lines

PIP<cr> Loads PIP into memory. PIP prompts for commands, executes them, then prompts again.

PIP pipcommandline<cr> Loads PIP into memory. PIP executes the command pipcommandline, then exits to CP/M.

PIP Command Summary

x:new.typ=y:old.typ[p]<cr> Copies the file old.typ on drive y: to the file new.typ on drive x:, using parameters p.

x:new.typ=y:old1.typ[p],z:old2.typ[q]<cr> Creates a file new.typ on drive x: that consists of the contents of file old1.typ on drive y: using parameters p followed by the contents of file old2.typ on drive z: using parameters q.

x:filename.typ=dev:[p]<cr> Copies data from device dev: to the file filename.typ on drive x:.

dev:=x:filename.typ[p]<cr> Copies data from filename.typ on drive x: to device dev:.

dst:=src:[p]<cr> Copies data to device dst: from device src:.

PIP Parameter Summary

B	Specifies block mode transfer.
Dn	Deletes all characters after the "n"th column.
E	Echoes the copying to the console as it is being performed.
F	Removes form feed characters during transfer.
Gn	Directs PIP to copy a file from user area "n."
H	Checks for proper Intel Hex File format.
I	Ignores any *:00* records in Intel Hex File transfers.
L	Translates uppercase letters to lowercase.
N	Adds a line number to each line transferred.
O	Object file transfer (ignores end-of-file markers).
Pn	Issues page feed after every "n"th line.
Qs^Z	Specifies quit of copying after the string "s" is encountered.
R	Directs PIP to copy from a system file.
Ss^Z	Specifies start of copying after the string "s" is encountered.
Tn	Sets tab stops to every "n"th column.
U	Translates lowercase letters to uppercase.
V	Verifies copy by comparison after copy finished.
W	Directs PIP to copy onto an R/O file.
Z	Zeroes the "parity" bit on ASCII characters.

PIP Destination Devices

CON:	PUN:	LST:	Logical devices
TTY:	PTP:	LPT:	
CRT:	UP1:	UL1:	
UC1:	UP2:		Physical devices
OUT:	PRN:		Special PIP devices

PIP Source Devices

CON:	RDR:		Logical devices
TTY:	PTR:		
CRT:	UR1:		
UC1:	UR2:		Physical devices
NUL:	EOF:	INP:	Special PIP devices

REN Command Line

REN newname.typ=oldname.typ<cr> Finds the file oldname.typ and renames it newname.typ.

SAVE Command Line

SAVE nnn x:filename.typ<cr> Saves a portion of the Transient Program Area of memory in the file filename.typ on drive x: where nnn is a decimal number representing the number of pages of memory. Drive x: is the option drive specifier.

STAT Command Lines

STAT<cr> Displays attributes and amount of free space for all diskette drives accessed since last warm or cold start.

STAT x:<cr> Displays amount of free space on the diskette in drive x:.

STAT x:filename.typ<cr>(CP/M 2.0 and newer) Displays size and attributes of file(s) filename.typ on drive x:. filename.typ may be ambiguous. x: is optional; if omitted, currently logged drive is assumed.

STAT x:filename.typ $atr<cr> Assigns the attribute atr to the file(s) filename.typ on drive x:. File filename.typ may be ambiguous. Drive x: is optional; if omitted, currently logged drive is assumed.

STAT DEV:<cr> Reports which physical devices are currently assigned to the four logical devices.

STAT VAL:<cr> Reports the possible device assignments and partial STAT command line summary.

STAT log:=phy:<cr> Assigns the physical device phy: to the logical device log: (may be more than one assignment on the line; each should be set off by a comma).

STAT USR:<cr> (CP/M 2.0 and newer) Reports the current user number as well as all user numbers for which there are files on currently logged disks.

STAT x:DSK<cr> (CP/M 1.4 and newer) Assigns a temporary write-protect status to drive x:.

SUBMIT Command Lines

SUBMIT filename<cr> Creates a file $$$.SUB which contains the commands listed in filename.SUB; CP/M then executes commands from this file rather than the keyboard.

SUBMIT filename parameters<cr> Creates a file $$$.SUB which contains commands from the file filename.SUB; certain parts of the command lines in filename. SUB are replaced by parameters during creation of $$$.SUB. CP/M then gets commands from this file rather than the keyboard.

SYSGEN Command Line

SYSGEN<cr> Loads the SYSGEN program to transfer CP/M from one diskette to another.

TYPE Command Line

TYPE x:filename.typ<cr> Displays the contents of file filename.typ from drive x: on the console.

USER Command Line

USER n<cr> Sets the User Number to "n," where "n" is an integer decimal number from 0 to 15, inclusive.

x: Command Line

x:<cr> Changes the currently logged disk drive to drive x:. Drive x: can be "A" through "P."

C

Summary of BDOS Calls

Table C-1. BDOS Function Definitions for CP/M-80 Version 2.2

Function No.	Function Name	Entry Parameter(s)	Exit Parameter(s)	Explanation
00	SYSTEM RESET	None	None	Restarts CP/M-80 by returning control to the the CCP after reinitializing the disk subsystem.
01	CONSOLE INPUT	None	A = ASCII character	Returns the next character typed to the character calling program. Any non-printable character is echoed to the screen (like BACKSPACE, TAB, or CARRIAGE RETURN). Execution does not return to the calling program until a character has been typed. Standard CCP control characters are recognized and their actions performed (CONTROL-P begins or ends printer echoing and so on).

Table C-1. (Continued)

Function		Entry Parameter(s)	Exit Parameter(s)	Explanation
No.	**Name**			
02	CONSOLE OUTPUT	E = ASCII character	None	Displays the character in the E register on the console device. Standard CCP control characters are recognized and their actions performed (CONTROL-P begins or ends printer echoing and so on.).
03	READER INPUT	None	A = ASCII character	Returns the next character received from the reader device to the calling program. Execution does not return to the calling program until a character is received.
04	PUNCH OUTPUT	E = ASCII character	None	Transmits the character in the E register to the punch device.
05	LIST OUTPUT	E = ASCII character	None	Transmits the character in the E register to the list device.
06	DIRECT CONSOLE IN DIRECT CONSOLE OUT	E = FF hex E = ASCII character	A = ASCII None	If register E contains an FF hex, the console device is interrogated to see if a character is ready. If no character is ready, a 00 is returned to the calling program in register A; otherwise the character detected is returned in register A. If register E contains any character other than an FF hex, that character is passed to the console display. All CCP control characters are ignored. The user must protect the program against nonsensical characters being sent from or received by the console device.
07	GET IOBYTE	None	A = IOBYTE	Places a copy of the byte stored at location 0003 hex in the A register before returning control to the calling program.
08	SET IOBYTE	E = IOBYTE	None	Places a copy of the value in register E into the memory location of 0003 hex before returning control to the calling program.
09	PRINT STRING	DE = String address	None	Sends the string of characters stored beginning at the address stored in the DE register pair to the console device. All characters in subsequent addresses are sent until BDOS encounters a memory location which contains a 24 hex (an ASCII "$"). The CCP control characters are checked for and performed if encountered.

NOTE: CP/M-80 always copies the contents of the H register in the A register if nothing is to be specifically returned in the A register. Some manufacturers, specifically Microsoft, make use of such information to reduce movement of information between the H and A registers.

Table C-1. (Continued)

Function		Entry Parameter(s)	Exit Parameter(s)	Explanation
No.	Name			
0A	READ CONSOLE BUFFER	DE = Buffer address	Data in buffer	This function performs essentially the same as the CCP would in that it takes the characters the user types and stores them into the buffer that begins at the address stored in the DE register pair. The first byte in the buffer pointed to by the DE pair must be the maximum length of the command; BDOS will place the number of characters encountered in the second byte, with the typed command beginning with the third byte pointed to by the DE pair. All standard CCP editing characters are recognized during the command entry.
0B	GET CONSOLE STATUS	None	A = Status	BDOS checks the status of the console device and returns a 00 hex if no character is ready, FF hex if a character has been typed.
0C	GET VERSION NUMBER	None	HL = Version	If the byte returned in the H register is 00 hex then CP/M is present, if 01, then MP/M is present. The byte returned in the L register is 00 if the version is previous to CP/M 2.0, 20 hex if the version is 2.0, 21 hex if 2.1 and so on.
0D	RESET DISK SYSTEM	None		Used to tell CP/M to reset the disk subsystem. Should be used any time diskettes are changed.
0E	SELECT DISK	E = Disk number	None	Selects the disk to be used for subsequent disk operations. A 00 hex in the E register indicates disk A, a 01 hex indicates disk B, etc.
0F	OPEN FILE	DE = FCB address	A = 'Found'/ not found code	Used to activate a file on the current disk drive and current user area. BDOS scans the first 14 bytes of the designated FCB block and attempts to find a match to the filename in the block. A 3F hex (ASCII "?") can be used in any of the filename positions to indicate a "don't care" character. If a match is found, the relevant information about that file is filled into the rest of the FCB by CP/M-80. A value of 00 hex to 03 in register A upon return indicates the open operation was successful, while an FF hex indicates that the file could not be found. If question marks are used to identify a file, the first matching entry is used.

NOTE: CP/M-80 always copies the contents of the H register in the A register if nothing is to be specifically returned in the A register. Some manufacturers, specifically Microsoft, make use of such information to reduce movement of information between the H and A registers.

Table C-1. (Continued)

Function		Entry Parameter(s)	Exit Parameter(s)	Explanation
No.	Name			
10	CLOSE FILE	DE = FCB address	A = 'Found'/ not found code	Performs the opposite of the open file function. A close file function must be performed upon completion of use of any file which has had information written into it.
11	SEARCH FOR FIRST	DE = FCB address	A= 'Found'/ not found code	Performs the same as the open file function with the difference being that the current disk buffer is filled with the 128-byte record which is the directory entry of the matched file.
12	SEARCH FOR NEXT	None	A= 'Found'/ not found code	Performs the same as search for first function except that the search continues on from the last matched entry.
13	DELETE FILE	DE = FCB address	A='Found'/ not found code	Changes a flag on the directory entry for the file pointed to by the FCB so that CP/M-80 no longer recognizes it as a valid file. No information is actually erased when this function is performed, although subsequent writes to diskette may use some of the area previously associated with the "deleted" file.
14	READ SEQUEN-TIAL	DE = FCB address	A = Error code	If a file has been activated for use by an open file or make file function, the read sequential function reads the next 128-byte block into memory at the current DMA address. The value of 00 hex is returned in the A register if the read was successful, while any nonzero value in the A register indicates failure.
15	WRITE SEQUEN-TIAL	DE = FCB address	A = Error code	If a file has been activated for use by an open file or make file function, the write sequential function writes the 128-byte block of memory at the current DMA address to the next 128-byte record of the named file.
16	MAKE FILE	DE = FCB address	A = DIR code	Creates a new file with the information (name) indicated by the FCB. CP/M-80 does not check to see if the file indicated already exists, so you must first check to see if the file exists (or delete it). A newly created file need not be opened, as the make file function also performs the necessary opening operations.
17	RENAME FILE	DE = FCB address	A = DIR code	Changes the name of the file referenced by the first 16 bytes of the FCB to the name in the second 16 bytes.

NOTE: CP/M-80 always copies the contents of the H register in the A register if nothing is to be specifically returned in the A register. Some manufacturers, specifically Microsoft, make use of such information to reduce movement of information between the H and A registers.

Table C-1. (Continued)

Function		Entry Parameter(s)	Exit Parameter(s)	Explanation
No.	Name			
18	RETURN LOGIN VECTOR	None	HL = Disk login	The bits in the HL register are used to specify which disk drives are active. The first bit in the L register refers to drive A, the last bit in the H register corresponds to drive P, the highest possible drive. A bit value of 1 indicates active status, a zero denotes an inactive drive.
19	RETURN CURRENT DISK	None	A = Current disk	The numbers 0 through 15 are used to represent the current default disk drive upon return from this function.
1A	SET DMA ADDRESS	DE = DMA	None	Used to select the 128-byte memory block to be used for buffering all disk transfers. Upon system or disk reset, cold or warm start, the buffer is reset to 0080 hex on a normal CP/M-80 system.
1B	GET ALLOC ADDRESS	None	HL= Allocation address	Returns the starting address of the allocation vector, a table which is maintained in memory for each on-line disk drive that indicates the portions of the diskette which are in use.
1C	WRITE PROTECT DISK	None	None	Provides temporary write protection for the diskette in the current default disk drive.
1D	GET R/O VECTOR	None	HL = Disk R/O	Returns a 16-bit value in the HL registers which indicate which drives on the system are write protected. The drives are assigned as in the LOGIN VECTOR, with a value 1 indicating write-protection.
1E	SET FILE ATTRIBUTES	DE = FCB address	A = DIR code	Sets the file attributes that indicate system/directory and R/O or R/W file status for the file pointed to by the FCB address.
1F	GET DISK PARMS	None	HL = DPB address	Retrieves the disk parameter block for the current active disk drive. These parameters can be used to determine space available on a diskette or to change the characteristics of the disk drive under user control.
20	GET USER CODE SET USER CODE	E = FF E = User code	A = Current User or None	If the E register contains an FF hex, the current user number is returned in the A register. To reset the user number, the appropriate user code is placed in the E register. While the USER command allows user numbers in the range 0-15, this BDOS function can set user numbers in the range of 0-31.

NOTE: CP/M-80 always copies the contents of the H register in the A register if nothing is to be specifically returned in the A register. Some manufacturers, specifically Microsoft, make use of such information to reduce movement of information between the H and A registers.

Table C-1. (Continued)

Function		Entry Parameter(s)	Exit Parameter(s)	Explanation
No.	Name			
21	READ RANDOM	DE = FCB address	A = Error code	Reads the random record number contained in the 33rd, 34th, and 35th byte (a 24-bit address) of the FCB pointed to.
22	WRITE RANDOM	DE = FCB address	A = Error code	Writes information from the current DMA address to the random record pointed to by the number contained in the 33rd, 34th, and 35th bytes of the indicated FCB.
23	COMPUTE FILE SIZE	DE = FCB address	RRF set	Returns the current size of the random record file in the three bytes that constitute the random record field of the FCB. If the third byte contains a 1, then the file contains the maximum record count of 65536, otherwise the value in the first two bytes is a 16-bit value that represents the file size.
24	SET RANDOM RECORD	DE = FCB address	RRF set	Returns the next random record (fills in the random record field of the FCB) after the last sequentially read record. Digital Research suggests that this function is most appropriate to file indexing.
25	RESET DRIVE	DE = Reset drive bits	A = Error code	Forces the specified drives to be reset to the drive bits initial non-logged status.
28	WRITE RANDOM (ZERO)	DE = FCB address	A = Error code	Writes a record of all zeros to diskette before a record is written; useful for identifying unused random records (an unused record would contain zeros instead of data).

NOTE: CP/M-80 always copies the contents of the H register in the A register if nothing is to be specifically returned in the A register. Some manufacturers, specifically Microsoft, make use of such information to reduce movement of information between the H and A registers.

D

Summary of BIOS Calls

Table D-1. CP/M-80 BIOS Routine Definitions

Label in Jump Table	Entry Parameter(s)	Exit Parameter(s)	Explanation
COLDSTART	None	C = 0	Your routine should perform all the necessary start-up operations, including initializing all the values in the base page. Before exiting, the C register must be set to zero.
WARMSTART	None	C = Drive	Your routine should perform all the necessary restart operations but does not need to reinitialize the base page. The C register, on exit, should contain the current drive number.
CONSOLE STATUS (CONST)	None	A = Status	
CONSOLE* INPUT	None	A = Character	

Table D-1. (Continued)

Label in Jump Table	Entry Parameter(s)	Exit Parameter(s)	Explanation
READER* INPUT	None	A = Character	Your routine should wait for a character to be entered at the appropriate device and then return the character in the A register.
CONSOLE* OUTPUT	C = Character	None	
LIST* OUTPUT	C = Character	None	
PUNCH* OUTPUT	C = Character	None	Your routine should take the character in the C register and display it on the appropriate device.
HOME DISK	None	None	The head of the disk drive should be returned to the home position (track 0, sector 0).
SELECT DISK	C = Drive	HL = DHA	Your routine should select the drive indicated by the number in the C register. The HL register on return should contain the address of the disk parameter header.
SET TRACK	C = Track	None	The track indicated by the C register value should be set as the next track to be accessed by the disk drive.
SET SECTOR	C = Sector	None	The sector indicated by the C register value should be set as the next track to be accessed by the disk drive.
SET DMA ADDRESS	BC = DMA address	None	The DMA address indicated by the BC register pair should be set as the address to use for all information transfers from memory to diskette and vice versa.
READ DISK	None	A = Status	Read the current track and sector and transfer the data to the DMA address already set. A 01 hex should be returned if there was an error during transfer.
WRITE DISK	None	A = Status	Write the current track and sector from the data at the DMA address.
SECTOR	BC = Logical sector	HL = Physical sector	
TRANSLATION	DE = Sector map address		A special routine used for systems which maintain data in other than 128-byte blocks. The logical sector on entry is changed to reflect the appropriate actual sector on the diskette.
LIST STATUS	None	A = Status	Your routine should interrogate the appropriate device to see if a character is ready and return a 00 hex in the A register if not ready, or a FF hex if ready.

*All console and device I/O should be done by first looking at the IOBYTE (0003 hex) to determine which device is selected.

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